

COMBINADICS ADJACENCY MATRIX

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COMBINADICS

The combinatorial number system, or combinadics, is a way to enumerate all m -combinations. Each combination of m natural numbers is given a ranking from 0 to $\binom{n}{m} - 1$. The ranking $N(J)$ of an m -combination $J = \{c_1, \dots, c_m\}$, arranged in lexicographical order, of natural numbers $\{0, \dots, n-1\}$ is equal to

$$N(J) = \binom{c_m}{m} + \dots + \binom{c_1}{1}.$$

ADJACENCY MATRIX

Definition 1 (Adjacency). *Two combinations J and K with the same cardinality are said to be adjacent if they differ by one element. That is,*

$$|K \cap J| = |K| - 1 = |J| - 1.$$

Let A_m^n be the adjacency matrix for m -combinations taken over n elements. The following lemma gives a recurrence relation to construct A_m^n .

Lemma 2. *For $n > m > 1$*

$$A_m^n = \begin{bmatrix} A_m^{n-1} & \tilde{A}_m^n \\ (\tilde{A}_m^n)^\top & A_{m-1}^{n-1} \end{bmatrix},$$
$$\tilde{A}_m^n = \begin{bmatrix} \tilde{A}_m^{n-1} & 0_{\binom{n-2}{m-2} \times \binom{n-2}{m-2}} \\ I_{\binom{n-2}{m-1}} & A_{m-1}^{n-1} \end{bmatrix},$$

with starting conditions

$$A_1^n = \begin{bmatrix} 0 & 1 & \dots & 1 \\ 1 & 0 & \ddots & \vdots \\ \vdots & \ddots & \ddots & 1 \\ 1 & \dots & 1 & 0 \end{bmatrix}, \quad \tilde{A}_2^4 = \begin{bmatrix} 1 & 1 & 0 \\ 1 & 0 & 1 \\ 0 & 1 & 1 \end{bmatrix}.$$

UNRANKING

Generating J from $N(J)$, called unranking, can be achieved through a greedy algorithm using the same coefficients. Let c be the largest number in $\{0, \dots, n-1\}$ such that $N(J) \geq \binom{c}{m}$. Then $c \in J$. Subtract $\binom{c}{m}$ from $N(J)$ and repeat, now considering the numbers $\{0, \dots, c-1\}$ and the $(m-1)$ -combination $J \setminus \{c\}$.

IMPORTANT LEMMA

Lemma 1. *For m -combination J*

$$N(J^c) + N(J) = \binom{n}{m} - 1$$

ANTI-TRANSPOSE

The anti-transpose of a matrix A , denoted A^τ , is its reflection over its northeast-to-southwest diagonal. That is, if $A \in \mathbb{R}^{m \times m}$ then $(A^\tau)_{i,j} = (A)_{m-j+1, m-i+1}$, where $i, j = \{1, \dots, m\}$.

PROOF

$$\begin{array}{c|cc} & J_i & J_j \\ \hline K_i & A_m^{n-1} & \tilde{A}_m^n \\ K_j & (\tilde{A}_m^n)^\tau & A_{m-1}^{n-1} \end{array}$$

The combinations J_i and K_i are m -combinations over $n-1$ elements, so the upper left block is A_m^{n-1} :

$$(A_m^{n-1})_{N(J_i)+1, N(K_i)+1} = (A_m^n)_{N(J_i)+1, N(K_i)+1}.$$

Both J_j and K_j have the element $n-1$, so

$$N(J_j) = \binom{n-1}{m} + N(J_j \setminus \{n-1\}),$$

$$N(K_j) = \binom{n-1}{m} + N(K_j \setminus \{n-1\}).$$

This makes the bottom right block A_{m-1}^{n-1} . Subdivide the combinations J_i and K_j into sets that do or do not contain $n-1$ or $n-2$.

$$\begin{array}{c|cc} & n-1 \notin & n-1 \in \\ \hline n-2 \notin & \tilde{K}_i & \tilde{J}_i \\ n-2 \in & \tilde{K}_j & \tilde{J}_j \end{array}$$

\tilde{K}_i and \tilde{J}_j differ by at least two elements, so the relevant block is a zero matrix. \tilde{J}_i and \tilde{K}_j are adjacent if and only if $\tilde{J}_i \setminus \{n-1\} = \tilde{K}_j \setminus \{n-2\}$. This makes the relevant block an identity matrix. \tilde{J}_j and \tilde{K}_j are adjacent if and only if $\tilde{J}_j \setminus \{n-1\}$ and $\tilde{K}_j \setminus \{n-2\}$ are adjacent. The relevant block is then a copy of \tilde{A}_{m-1}^{n-1} .

For the final block, consider a new set $J_k = \tilde{J}_i \setminus \{n-1\} \cup \{n-2\}$. Then \tilde{J}_i is adjacent to \tilde{K}_i if and only if J_k is adjacent to \tilde{K}_i . Thus,

$$(A_m^n)_{N(\tilde{J}_i)+1, N(\tilde{K}_i)+1} = (A_m^{n-1})_{N(J_k)+1, N(\tilde{K}_i)+1}.$$

The block is then a copy of \tilde{A}_m^{n-1} .

Since adjacency is bidirectional, A_m^n is symmetric, and the bottom left block is the transpose of the upper right block.

The following corollary means only half of all A_m^n need to be found.

Corollary 1.

$$(A_m^n)^\tau = A_{n-m}^n.$$

UNION OF ADJACENT COMBINATIONS

Define a matrix B_m^n that stores the values of $N(J \cup K)$ for when J and K are adjacent:

$$(B_m^n)_{N(J)+1, N(K)+1} = \begin{cases} N(J \cup K) + 1 & J, K \text{ adjacent} \\ 0 & \text{otherwise} \end{cases}$$

Lemma 3.

$$B_1^n = \begin{bmatrix} 0 & 1 & 2 & \dots & \binom{n}{2} + 1 \\ & 0 & 3 & \dots & \binom{n}{2} + 2 \\ & & \ddots & \ddots & \vdots \\ & & & 0 & \binom{n}{2} + n \\ & & & & 0 \end{bmatrix}, \quad B_{n-1}^n = A_1^n.$$

Lemma 4.

$$B_m^n = \begin{bmatrix} B_m^{n-1} & \tilde{B}_m^n \\ (\tilde{B}_m^n)^\top & B_{m-1}^{n-1} \end{bmatrix} + \binom{n-1}{m+1} \begin{bmatrix} 0 & \tilde{A}_m^n \\ 0 & A_{m-1}^{n-1} \end{bmatrix},$$
$$\tilde{B}_m^n = \begin{bmatrix} \tilde{B}_m^{n-1} & 0 \\ \text{diag}(K) & \tilde{B}_{m-1}^{n-1} \end{bmatrix} - \binom{n-2}{m+1} \begin{bmatrix} \tilde{A}_m^{n-1} & 0 \\ 0 & 0 \end{bmatrix},$$

where $\text{diag}(K)$ is the diagonal matrix with entries equal to the row index.

INTERSECTIONS

Corollary 2. *Let J and K be two adjacent m -combinations, then*

$$(B_m^n)_{N(J)+1, N(K)+1}^\tau = \binom{n}{m} - N(J \cap K).$$

$$\begin{aligned} (B_m^n)_{N(J)+1, N(K)+1}^\tau &= (B_m^n)_{\binom{n}{m}-N(K), \binom{n}{m}-N(J)} \\ &= (B_m^n)_{N(K^c)+1, N(J^c)+1} \\ &= N(K^c \cup J^c) + 1 \\ &= N((J \cap K)^c) + 1 \\ &= \binom{n}{m} - N(J \cap K). \end{aligned}$$

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