

# Assignment # 1

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Course = Discrete Structures (CSC - 102)

Q=1

## Prepositions

## Truth Values

|   |   |
|---|---|
| 1) $p \vee q$   | T |
| 2) $\sim p \vee \sim q$   | T |
| 3) $\sim p \vee q$  | T |
| 4) $\sim p \vee \sim (q \wedge r)$  | T |
| 5) $\sim (p \vee q) \wedge (\sim p \vee r)$                                   | F |
| 6) $(p \vee \sim r) \wedge \sim (l q \vee r) \leftrightarrow \sim (r \vee p)$ | F |
| 7) $\sim p \vee \sim (q \wedge r) \rightarrow (\sim p \vee q)$                | T |
| 8) $(q \vee \sim r) \oplus (p \wedge r)$                                      | T |



Q=2

- i) Exclusive or
- ii) Inclusive or
- iii) Inclusive or
- iv) Exclusive or
- v) Inclusive or



Q=3

1)  $\sim p \wedge (q \vee r)$

In words : You cannot use wireless network and either you pay the semester fees or you are

a subscriber to the service.

Truth Value : F

$$2) \sim(p \vee q) \wedge r$$

In Words : It is not the case that you can use the wireless network or you can pay the semester fee , and you are a subscriber to the service.

Truth Value : F

$$3) (p \wedge q) \wedge \sim(r \vee p)$$

In Words : You can use the wireless network and you can pay the semester fees , and it is not the case that you are a subscriber or you can use the wireless network.

Truth Value : F

$$4) (p \wedge q) \vee \sim r$$

In Words : You can use the wireless network and you pay the semester fee , or you are not a subscriber to the service.

Truth Value : T

$$5) \sim(p \vee q) \leftrightarrow r$$

In Words : It is not the case that you can use the wireless network or you pay the semester fee iff you are a subscriber to the service.

Truth Value : T

$$6) (p \wedge (q \vee r)) \rightarrow (r \vee p)$$

In words: If you can use the wireless network and either you pay the semester fee or you are a subscriber to the service, then you are a subscriber to the service or you can use the wireless network.

Truth Value : T



$\underline{Q=4}$

- |                       |                                 |
|-----------------------|---------------------------------|
| i) $q \rightarrow p$  | iii) $q \rightarrow p$          |
| ii) $q \wedge \sim p$ | iv) $\sim q \rightarrow \sim p$ |

Truth Table :-

| P | q | $\sim p$ | $\sim q$ | $q \rightarrow p$ | $q \wedge \sim p$ | $\sim q \rightarrow \sim p$ |
|---|---|----------|----------|-------------------|-------------------|-----------------------------|
| T | T | F        | F        | T                 | F                 | T                           |
| T | F | F        | T        | T                 | F                 | F                           |
| F | T | T        | F        | F                 | T                 | T                           |
| F | F | T        | T        | T                 | F                 | T                           |



$\underline{Q=5}$

$\Rightarrow$  Restatements

- 1) If Javed studies hard, then he will pass the Discrete Structure exam

- 2) If Ayesha passes 133 credit hours, then she may graduate.
- 3) If Farid buys a computer, then he has obtained \$2000.
- 4) If Kiran passes discrete structure, then can take algorithm course.
- 5) If Bruch will build them, then better cars are built.
- 6) If the chairperson gives the lecture, then the audience will go to sleep.
- 7) If the program is readable, then it is well structured.

⇒ Part B:-

- 1) Converse : If Javed passes the Discrete Structure exam, then he studied hard.  
Inverse : If Javed do not study hard, then he will not pass Discrete Structure exam.  
Contrapositive : If Javed didn't pass Discrete Structure exam, then he did not study hard.

2) Converse: If Ayesha may graduate, then she passes 133 credit hours.

Inverse: If Ayesha does not pass 133 credit hours, then she may not graduate.

Contrapositive: If Ayesha does not graduate, then she did not pass 133 credit hours.

3) Converse: If Farid has obtained \$ 2000, then he buys a computer.

Inverse: If Farid does not buy a computer, then he has not obtained \$2000.

Contrapositive: If Farid has not obtained \$2000, then he does not buy a computer.

4) Converse: If Kiran takes the algorithm course, then she passes discrete structure.

Inverse: If Kiran does not pass discrete structure, then she cannot take the algorithm course.

If Kiran does not take the algorithm course, then she did not pass discrete structure.

5) Converse: If better cars are built, then Buick will build them.

Inverse: If Buick will not build them, then better cars are not built.

Contrapositive: If better cars are not built, then Buick does not build them.

6) Converse: If audience goes to sleep, then the chairperson gives the lecture.

Inverse: If the chairperson does not give the lecture, then the audience will not go to sleep.

Contrapositive: If the audience does not go to sleep, then the chairperson didn't give the lecture.

7) Converse: If the program is well structured, then it is readable.

Inverse: If the program is not readable, then it is not well structured.

Contrapositive: If the program is not well structured, then it is not readable.



Q = 6

i) The router can send packets to the edge system only if it supports the new address space  
=  $s \rightarrow a$

ii) For the router to support the new address space, it is necessary that the latest software release be installed. =  $a \rightarrow r$

iii) The router can send address packets to the edge system if the latest software release is installed. =  $r \rightarrow s$

iv) The router does not support the new address space =  $\sim a$

So, from iv statement, we came to know that  $a = F$ . By putting value of  $a$  in i & iii, we will get the truth values. Therefore,

$$i) s \rightarrow a$$

If  $a$  is  $F$ ,  $s$  is also  $F$ .

$$ii) r \rightarrow a$$

If  $a$  is  $F$ ,  $r$  is also  $F$ .

Truth Values :-

$$s = F$$

$$a = F$$

$$r = F$$

$$\text{Q} = 7 \quad \text{a) ii)}$$

Laws of logic

Start with L.H.S:-

$$\sim[r \vee (q \wedge (\sim r \rightarrow \sim p))]$$

Rewrite:-

$$\sim[r \vee (q \wedge (r \vee \sim p))]$$

Distribute:-

$$q \wedge (r \vee \sim p) \equiv (q \wedge r) \vee (q \wedge \sim p)$$

$$\Rightarrow \sim[r \vee ((q \wedge r) \vee (q \wedge \sim p))]$$

De-Morgan's law:-

$$\sim r \wedge \sim(q \wedge r) \wedge \sim(q \wedge \sim p)$$

Combiner

Since,  $\sim r$  &  $\sim q$  is common.

So,

$$[\sim r \wedge \sim(q \wedge \sim p)]$$

This shows that both expressions are equivalent.

De Morgan to Components:-

For  $\sim(q \wedge r)$ :

$$\sim q \vee \sim r$$

For  $\sim(q \wedge \sim p)$ :

$$\sim q \vee p$$

Substituting

$$\sim r \wedge (\sim q \vee \sim r) \wedge (\sim q \vee p)$$

$Q = \neg p$

$\neg Q$

$p \rightarrow r$

$\neg p \rightarrow r$

$p \vee r$

$\neg p \vee r$

$(p \vee r) \wedge (\neg p \vee r)$

$\neg (\neg p \wedge \neg r)$

| $\neg Q$ | $\neg p$ | $\neg r$ | $\neg (\neg p \wedge \neg r)$ |
|----------|----------|----------|-------------------------------|
| F        | F        | T        | T                             |
| F        | F        | T        | T                             |
| F        | T        | F        | T                             |
| T        | F        | F        | F                             |
| T        | T        | F        | F                             |
| T        | T        | T        | F                             |

| $\neg Q$ | $\neg p$ | $\neg r$ | $\neg (\neg p \wedge \neg r)$ |
|----------|----------|----------|-------------------------------|
| T        | T        | F        | T                             |
| T        | T        | T        | T                             |

| $\neg Q$ | $\neg p$ | $\neg r$ | $\neg (\neg p \wedge \neg r)$ |
|----------|----------|----------|-------------------------------|
| T        | T        | T        | T                             |

| $\neg Q$ | $\neg p$ | $\neg r$ | $\neg (\neg p \wedge \neg r)$ |
|----------|----------|----------|-------------------------------|
| T        | F        | T        | T                             |
| F        | F        | T        | T                             |
| F        | T        | F        | T                             |
| T        | F        | F        | F                             |
| T        | T        | F        | F                             |
| T        | T        | T        | F                             |

| $\neg Q$ | $\neg p$ | $\neg r$ | $\neg (\neg p \wedge \neg r)$ |
|----------|----------|----------|-------------------------------|
| T        | F        | F        | T                             |
| F        | F        | F        | T                             |
| F        | T        | F        | T                             |
| T        | F        | F        | F                             |
| T        | T        | F        | F                             |
| T        | T        | T        | F                             |

| $\neg Q$ | $\neg p$ | $\neg r$ | $\neg (\neg p \wedge \neg r)$ |
|----------|----------|----------|-------------------------------|
| T        | T        | F        | T                             |
| T        | T        | T        | T                             |

| $\neg Q$ | $\neg p$ | $\neg r$ | $\neg (\neg p \wedge \neg r)$ |
|----------|----------|----------|-------------------------------|
| T        | F        | T        | T                             |
| F        | F        | T        | T                             |
| F        | T        | F        | T                             |
| T        | F        | F        | F                             |
| T        | T        | F        | F                             |
| T        | T        | T        | F                             |

| $\neg Q$ | $\neg p$ | $\neg r$ | $\neg (\neg p \wedge \neg r)$ |
|----------|----------|----------|-------------------------------|
| T        | T        | T        | T                             |

Thus, both are  
equivalent

Q=7(B):-

a)  $(p \wedge q) \rightarrow p$

| i) | $p$ | $q$ | $p \wedge q$ | $(p \wedge q) \rightarrow p$ |
|----|-----|-----|--------------|------------------------------|
|    | T   | T   | T            | T                            |
|    | T   | F   | F            | T                            |
|    | F   | T   | F            | T                            |
|    | F   | F   | F            | T                            |

ii)  $(p \wedge q) \rightarrow p$

$\sim(p \wedge q) \vee p$

$\sim p \vee \sim q \vee p$

$p \vee \sim p \vee \sim q$  ( $p \vee \sim p$  is always true)

Therefore, " $(p \wedge q) \rightarrow p$  is a tautology"

b)  $p \rightarrow (p \vee q)$

| i) | $p$ | $q$ | $(p \vee q)$ | $p \rightarrow (p \vee q)$ |
|----|-----|-----|--------------|----------------------------|
|    | T   | T   | T            | T                          |
|    | T   | F   | T            | T                          |
|    | F   | T   | T            | T                          |
|    | F   | F   | F            | T                          |

ii)  $p \rightarrow (p \vee q)$

$\sim p \vee (p \vee q)$

$(\sim p \vee p) \vee q$  ( $\sim p \vee p$ ) is always true.

Therefore, " $p \rightarrow (p \vee q)$  is a tautology"

c)  $\sim p \rightarrow (p \rightarrow q)$

| i) | $p$ | $q$ | $\sim p$ | $p \rightarrow q$ | $\sim p \rightarrow (p \rightarrow q)$ |
|----|-----|-----|----------|-------------------|--|
|    | T   | T   | F        | T                 | T                                      |
|    | T   | F   | F        | F                 | T                                      |
|    | F   | T   | T        | T                 | T                                      |
|    | F   | F   | T        | T                 | T                                      |

ii)  $\sim p \rightarrow (p \rightarrow q)$

$\sim p \rightarrow (\sim p \vee q)$

$p \vee \sim p \vee q$  "( $p \vee \sim p$  is always true)".

Therefore, " $\sim p \rightarrow (p \rightarrow q)$  is a tautology".

d)  $(p \wedge q) \rightarrow (p \rightarrow q)$

| i) | $p$ | $q$ | $p \wedge q$ | $p \rightarrow q$ | $(p \wedge q) \rightarrow (p \rightarrow q)$ |
|----|-----|-----|--------------|-------------------|--|
|    | T   | T   | T            | T                 | T  |
|    | T   | F   | F            | F                 | T  |
|    | F   | T   | F            | T                 | T  |
|    | F   | F   | F            | T                 | T  |

ii)  $(p \wedge q) \rightarrow (p \rightarrow q)$

$\sim(p \wedge q) \vee (p \rightarrow q)$

$\sim p \vee \sim q \vee (\sim p \vee q)$

$\sim p \vee \sim q \vee q$  " $(\sim q \vee q)$  is always true".

Therefore, " $(p \wedge q) \rightarrow (p \rightarrow q)$  is a tautology".

$$e) \sim(p \rightarrow q) \rightarrow p$$

|    | $p$ | $q$ | $p \rightarrow q$ | $\sim(p \rightarrow q)$ | $\sim(\sim(p \rightarrow q)) \rightarrow p$ |
|----|-----|-----|-------------------|-------------------------|---|
| i) | T   | T   | T                 | F                       | T   |
|    | T   | F   | F                 | T                       | T   |
|    | F   | T   | T                 | F                       | T   |
|    | F   | F   | T                 | F                       | T   |

$$\text{ii)} \sim p \sim(p \rightarrow q) \rightarrow p$$

$$\sim(\sim p \vee q) \rightarrow p$$

$$(p \wedge \sim q) \rightarrow p$$

$$\sim(p \wedge \sim q) \vee p$$

$$\sim p \vee q \vee p$$

$$p \vee \sim p \vee q \quad "(\sim p \vee p \text{ is always true})"$$

Therefore,  $(\sim p \rightarrow q) \rightarrow p$  is a tautology

$$f) \sim(p \rightarrow q) \rightarrow \sim q$$

|    | $p$ | $q$ | $\sim q$ | $p \rightarrow q$ | $\sim(p \rightarrow q)$ | $\sim(\sim(p \rightarrow q)) \rightarrow \sim q$ |
|----|-----|-----|----------|-------------------|-------------------------|--|
| i) | T   | T   | F        | T                 | F                       | T  |
|    | T   | F   | T        | F                 | T                       | F  |
|    | F   | T   | F        | T                 | F                       | F  |
|    | F   | F   | T        | T                 | F                       | T  |

$$\text{ii)} \sim(p \rightarrow q) \rightarrow \sim q$$

$$\sim(\sim p \vee q) \rightarrow \sim q$$

$$(p \wedge \sim q) \rightarrow \sim q$$

$$\sim(p \wedge \sim q) \vee \sim q$$

$$\sim p \vee q \vee \sim q \quad "(\sim q \vee \sim q \text{ is always true})"$$

Therefore,  $\sim(p \rightarrow q) \rightarrow \sim q$  is a tautology

Q=8 :-

i)  $P(11)$

In Words : "11 divides 77"

T or F : True, because  $77 \div 11 = 7$

ii)  $P(1)$

In Words : "1 divides 77"

T or F : True, because any positive integer is divisible by 1.

iii)  $P(3)$

In Words : "3 divides 77"

T or F : False, because  $77 \div 3$  does not yield an integer.

iv)  $\forall n P(n)$

In Words : "For all positive integers  $n$ ,  $n$  divides 77."

T or F : False, since there are positive integers (like 3) that do not divide 77.

v)  $\exists n P(n)$

In Words : "There exists positive integers  $n$  such that  $n$  divides 77."

T or F : True, because integers like 1, 7, 11 divides 77.

vi)  $\forall n \sim P(n)$

In Words : "For all positive integers  $n$ ,  $n$  does not divide 77".

T or F : False, since some integers (like 1, 7, 11) do divide 77.

vii)  $\exists n \sim P(n)$

In Words : There exists a positive integer  $n$  such that  $n$  does not divide 77.

T or F : True, integers like 3 do not divide 77.

viii)  $\sim(\forall n P(n))$

In Words: "It is not true that all positive integers  $n$ ,  $n$  divides 77."

T or F: True.

ix)  $\sim(\exists n P(n))$

In Words: "It is not true that there exists a positive integer  $n$  such that  $n$  divides 77."

T or F: False.



$$\textcircled{Q} = 8 \quad \text{iii} :-$$

i)  $\forall x (P(x) \rightarrow Q(x))$

In Words: "For all people  $x$ , if  $x$  is a professional athlete, then  $x$  plays soccer."

T or F: False, not all athletes play soccer.

ii)  $\exists x (P(x) \rightarrow Q(x))$

In Words: "There exists people  $x$  such that if  $x$  is a professional athlete, then  $x$  plays soccer."

T or F: True.

iii)  $\forall x (Q(x) \rightarrow P(x))$

In Words: "For all people, if  $x$  plays soccer, then  $x$  is a professional athlete."

T or F: False. Many people play soccer recreationally.

$$\text{iv) } \exists n (Q(n) \rightarrow P(n))$$

In Words: "There exists a person  $x$  such that if  $x$  plays soccer, then  $x$  is a professional athlete."

T or F: True.

$$\text{v) } \forall n (P(n) \vee Q(n))$$

In Words: "For all people  $x$ ,  $x$  is a professional athlete or  $x$  plays soccer."

T or F: False. There are many who are neither professional athletes nor soccer players.

$$\text{vi) } \exists n (P(n) \vee Q(n))$$

In Words: "There exists a person  $n$  such that  $n$  is a professional athlete or  $n$  plays soccer."

T or F: True.

$$\text{vii) } \forall n (P(n) \wedge Q(n))$$

In Words: For all people  $x$ ,  $x$  is a professional athlete and  $x$  plays soccer.

T or F: False.

$$\text{viii) } \exists n (P(n) \wedge Q(n))$$

In Words: There exists a person  $n$  such that  $n$  is a professional athlete and  $n$  plays soccer.

T or F: True.



The End

Q-8 i :-

i)  $(2n+1)^2$  is an odd integer:

Propositional Function: Yes, it depends upon variable "n".

Domain : The set of integers ( $\mathbb{Z}$ ).

ii) Choose an integer b/w 1 & 10.

Propositional Function : Yes, it implies a variable (the chosen integer).

Domain : The integers (1-9)

iii) Let n be real number.

Propositional Function : Yes, it defines a variable "n".

Domain : The set of Real numbers ( $\mathbb{R}$ ).

iv) The movie won the Academy Award as the best picture of 1995.

Propositional Function : No, it is a statement about particular event, not dependent on a variable.

v)  $1+3=4$ .

P.F : No, it is a mathematical statement, that is always true and no mathematical statement is involved.

vi) There exists n such that  $n < y$  ( $x, y$  real numbers).

Propositional Function : Yes, it depends on y.

Domain : The set of Real numbers ( $\mathbb{R}$ ).