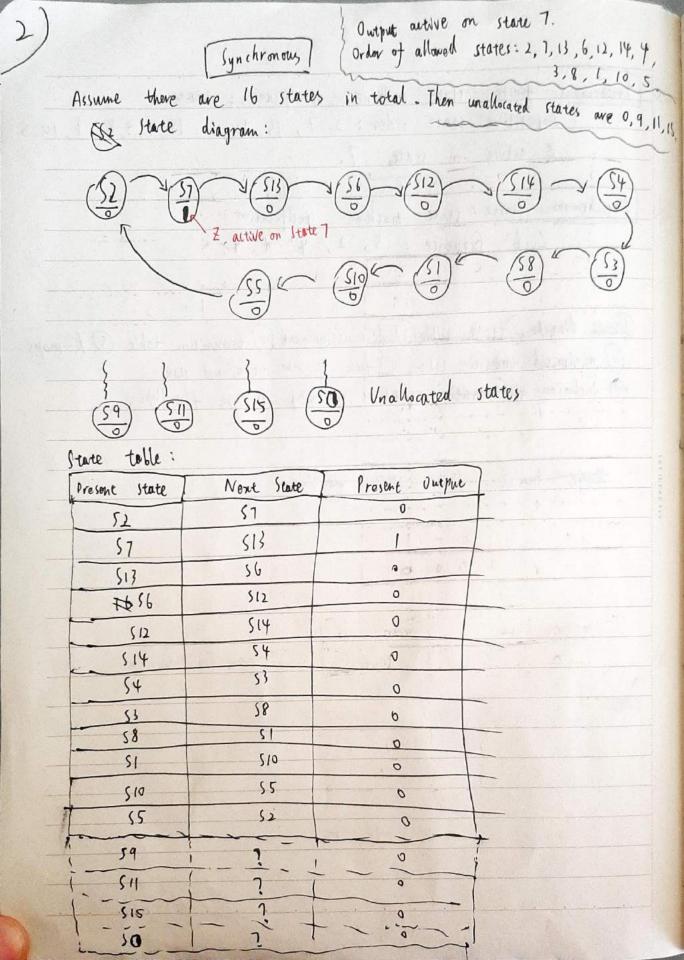
* Synchronous finite state machine specification: Repetitive state order: 2, 7, 13, 6, 12, 14, 4, 7,8, 1, 10, Z active on state: 7. X: [Asynchronous finite state machine specification: Lock sequence: 5, 2, 4, 1, 9. Ostate diagram Dstate table Dode assignment (4) transition table 5 K-max (6) minimised transition table of force it into not-allowed state 1 inclusion of transition equations assuming D-type + lip-flops Moore model limite state machine



1	Assign	inary Gray code to each of these states.
1	State	Assigned code
1		lada a a a a a a a a a a a a a a a a a a
1	So	0 0 0 0
	51	0001
1	Sz	OO 1 1 Assigned code Bits Output Bit
	53	We can reduce the code to
1	54	ollow bits.
1	22	0111
1	56	0101
1	57	0100
-	58	1100
	59	
	510	
1	511	11.10
	512	1010
	513	1011
	514	1001
	SIS	1000

Present T	ine	Future time
Dn	an	Onti
0	0	0
0	1	0
1	O	
	1	

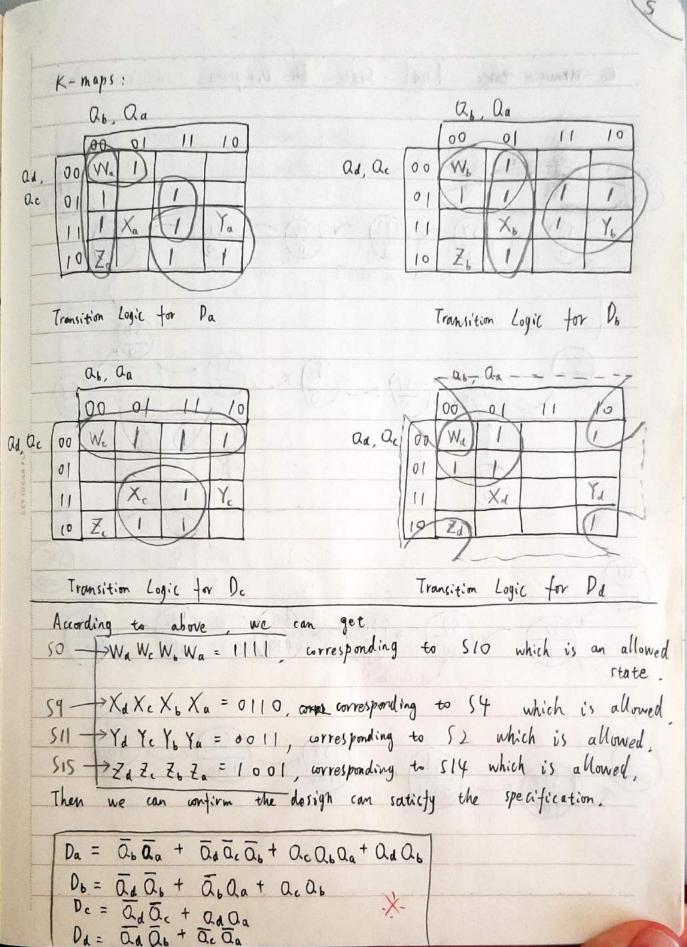
Full-Form Transition Table of D-type

D-type tlip-tlops will be used in this assignment.

The transition table can be derived: (ordered by specification)

Sto a O. O	T	0	Nex State 57 513 56	0	a. 1	ab o	Qa O	Nex O	1.	D _b	Da	
0 1 0	0	1	57 513	0	1			-	1000	0	0	
10	11/2	-	513	-		0	0		_	The second second		
1	11/2	-		1	0		1	1-1	0	1	1	1
1		1	56	-	0	1	1	0	11	0	1	
-	0		10	0	1	0	1		1	-		
-	-		\$12	1	0	1	0	1	0	1	0	
	11	0	514	1	0	0	1	1	0	0	1	
		1	54	0	1	1	0	0	1	1	0	
					0	1	0	0	0	1	0	
				-	-	-	-		1	0	0	
0	1	0	(diameter)	-			-		0	0	1	
11	0	0	51	0	0	0	-				1	
0	0	1	Slo	1	1	1	11					
1		1	55	0	1	1	1	0	1	1	1	
+ ;		1	52	0	0	1		O	0	1	11	
1		0		_		W	Wal	Wal	We	W	Wa	
10			T		-	-	-			and the same of th	-	
111	0	Lif		-	Appendix Committee	The second second		1		1		
111	110	0	3	Yd_	1c	16	la	- +	- 1	- (200	37
101	00	5	J 14	Zd	7	71	7	7 1	7	7	7	- 1
	0 1 0 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1	0 0 1 1 0 0 1 1 0 0 0 0 1 1 0 0 0 0 1 1 0 0 0 1 1 0 0 0 1 1 0 0 0 1 1 0 0 0 1 1 0 0 1 0 1 0 0 1	0 0 1	0 0 1 54 1 1 0 53 0 1 0 58 1 0 0 51 0 0 1 510 1 1 1 55 1 1 1 52 1 0 1 7	0 0 1 54 0 1 1 0 53 0 0 1 0 58 1 1 0 0 51 0 0 0 1 510 1 1 1 1 55 0 1 1 1 55 0 1 1 1 52 0 0 0 0 ? Wa 1 0 1 ? Xa	0 0 1 54 0 1 1 1 0 53 0 0 0 1 0 58 1 1 1 0 0 51 0 0 0 0 1 510 1 1 1 1 1 55 0 1 1 1 1 52 0 0 0 0 0 ? Wa Wa 1 0 1 ? Ya Ya 1 1 0 ? Ya Ya	0 0 1 54 0 1 1 1 1 0 53 0 0 1 0 1 0 58 1 1 0 0 0 51 0 0 0 1 510 1 1 1 1 1 55 0 1 1 1 1 52 0 0 1 1 1 1 52 0 0 1 1 0 0 ? Wd Wc Wb 1 0 1 ? Yd Yc Yb 1 1 0 ? Yd Yc Yb	0 0 1 54 0 1 1 0 1 1 0 53 0 0 1 0 0 1 0 58 1 1 0 0 1 0 0 51 0 0 0 1 0 0 1 510 1 1 1 1 1 1 1 55 0 1 1 1 1 1 1 52 0 0 1 1 0 0 0 ? Wd WL Wb Wa 1 0 1 ? Yd Yc Yb Ya 1 1 0 ? Yd Yc Yb Ya	0 0 1 54 0 1 1 0 0 1 1 0 53 0 0 1 0 0 0 1 0 58 1 1 0 0 1 1 0 0 51 0 0 0 1 0 0 0 1 510 1 1 1 1 1 1 1 1 55 0 1 1 1 0 1 1 1 52 0 0 1 1 0 1 0 0 ? Wd WL Wb Wa Wd 1 0 1 ? Yd Yc Yb Ya Yd	0 0 1 54 0 1 1 0 0 1 1 1 0 53 0 0 1 0 0 0 0 1 0 58 1 1 0 0 1 1 1 0 0 51 0 0 0 1 0 0 0 0 1 510 1 1 1 1 1 1 1 1 55 0 1 1 1 0 1 1 1 1 52 0 0 1 1 0 0 0 0 ? Wd Wc Wb Wa Wd Wc 1 0 1 7 Xd Xc Xb Xb Xd Xc 1 0 1 7 Yd Yc Yb Ya Yd Yc 1 1 0 ? Yd Yc Yb Ya Yd Yc 1 7 7 7 7 7 7 7	0 0 1 54 0 1 1 0 0 1 1 1 0 0 1 1 1 0 0 1 1 1 0 0 0 1 1 0 0 0 0 1 0 0 0 0 1 0	0 0 1 54 0 1 1 0 0 1 1 0 0 1 1 0 0 1 1 0 0 1 1 0 0 0 1 0 0 0 1 0 0 0 0 1 0 0 0 0 1 0 0 0 0 1 0 0 0 0 1 0 0 0 0 1 0 0 0 0 1 0 0 0 0 1 0 0 0 0 1 0 0 0 0 1 0 0 0 0 1 0 0 0 0 1 0 0 0 0 1 1 0 0 0 1

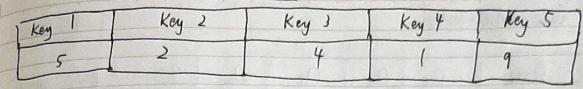
Z active on ST, Z= QaQcQbQa



Final transition Table:

1	Curren	t St	ate			Newe S	tate				Nea	t	D-Ir	purs	Z
13	tate	ad	ac	26	Qa	State	Qa	Qc	10,	Qa	Da	De	Do	Da	0
	82	0	0	1	1	57	0	1	0	0	0	1	0	0	0
	57	0	1	0	0	513	1	0	1	1	1	0	I	1	I
	513	1	0	1		56	0	1	0	1	0	1	0		0
	56	0	1	0	1	\$12	1	0	l	0		0	1	0	0
1	512		0	1	0	514	1	0	0	1	1	0	0		0
	514	1	0	0	1	54	0	1	1	0	D			0	9
1	54	0	1		0	53	0	0	1	0	0	0	1	0	0
1	53	0	0	1	0	58	1	1	0	0	1	1	0	0	0
1	58	1	11	0	0	SI	0	0	0		0	0	0	1	0
	SI	0	0	0	1	\$10	1	1	1	1	1	1	1	1	0
	510	1	1	1	1	55	0	1	1	1	0	1	ı	1	0
1	55	0	1	1	Is	52	0	0	1	1	0	0	1	1	0
1	So	0	0	0	0	510	1	1	1	1	1	1	1	1	0
1	59	1	1	0	1	154	0	1	1	0	0	1	1	0	0
t	511	1	1	1	0	52	0	0	1	1	0	0	1	1	0
	515	11	0	0	0	514	1	0	0	1	1	0	0	1	0

Asynchronous



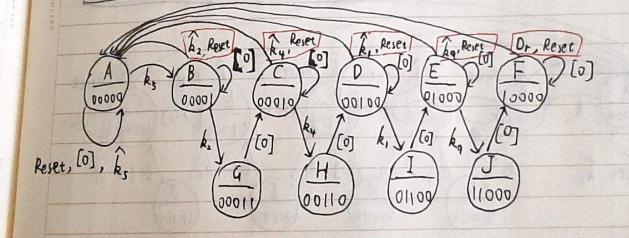
State Diagram:

Firse, we should define some pre-processed functions. As

the five-digit code sequence is 52419 in specification,

tex pre-prossed functions are as following:

	PX:	Tepresone	reg	MUMBO	N 3	pressed	~ (d na	1	ey-	Non-
4	kx:	represent	key	nun	nber X	. (0~	4)	pressed	and	N	other
	^ .	key. represent of key	anu	keu	other	than	ben	X,	or	a	combination
+	k.:	of key	s pr	essed.			~~				

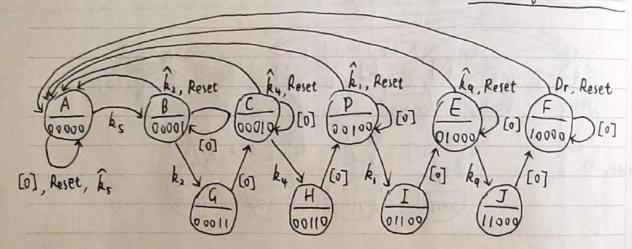


In order to add fraud-protection transitions and avoid race hazards we use one-hot take star and three-hot two-hot code strategy to assign states.

Code Assignment:

State	1 Qe	ad	Qc	Qb	aa	Z	
A	0	0	0	0	0	0	
β	0	0	0	0	1	0	
C	0	0	0	1	0	0	
P	0	0	1	0	0	0	
E	0	1	0	0	0	0	
F	150	0	0	0	0	1	-12
4	0	0	0	1	1	0	
Н	0	0	1	1	0	0	
1	0	1	1	0	0	0	
J	1	1	0	0	0	0	

Lock sequence: 52419



Flow Table :

Low				Next	[tate							
Current State	Reset	[0]	ks	\hat{k}_s	k,	lk.	ku	Â,	16.	k,	A.	la,	Dr
A	A	A	B	A			1			La.	22		-
B	A	В			G	A	-						
(A	C	1.00				H	A					
D	A	D							1	A			
Ε	A	E						-			J	A	
F	A	F		1				- 1					A
9		C											
Н		D										1.	
ī		E							* 1				
J		F	N. III					1				-m	

32 0 19 10 0 10

Julio la la la

19 12 1- 13

14)

Expand the pre-processed functions: (0) = ko k, ko ki b2 k3 k4 k5 k6 k7 k8 k9 ko ki R= 60 k1 k2 k3 k4 k5 k6 k7 k8 k9 ky = ko ki k2 k3 k4 k5 k6 k7 k8 k9 ko ki k2 k3 k4 k5 k6 k7 k8 k9 k, = 60 k1 b2 k3 ky k5 k6 k7 k8 k9 40 k1 k2 k3 k4 k5 k6 k7 k8 k9 ko kl k2 k3 k4 k5 kb k7 k8 ko ki [0] = k0 k1 k2 k3 k4 k5 kb k7 k8 k9 Note: Ten keys on key board

Extraction of implicants:

06 = \(\overline{a} \overline{Q}_{0} \o

= Qe Qa Qa Qa Qa ko ki k2 k3 k4 k5 kb k7 k8 k9 + Qe Qa Qa Qa Qa ko ki k2 k3 k4 k5 kb k7 k8 k9 + Qe Qa Qa Qa Qa ko ki k2 k3 k4 k5 kb k7 k8 k9 + Qe Qa Qa Qa Qa ko ki k2 k3 k4 k5 kb k7 k8 k9

 $Q_d = \overline{Q}_e \overline{Q}_a Q_c \overline{Q}_b \overline{Q}_o k_1 + \overline{Q}_e Q_d \overline{Q}_c \overline{Q}_b \overline{Q}_a [0] + \overline{Q}_e Q_d \overline{Q}_c \overline{Q}_b \overline{Q}_o k_q$ $+ \overline{Q}_e Q_d Q_c \overline{Q}_b \overline{Q}_a [0]$ $= \overline{Q}_e \overline{Q}_a Q_c \overline{Q}_b \overline{Q}_a \overline{k} 0 \quad k_1 \quad \overline{k}_2 \quad \overline{k}_3 \quad \overline{k}_4 \quad \overline{k}_5 \quad \overline{k}_6 \quad \overline{k}_7 \quad \overline{k}_8 \quad \overline{k}_9$

 $Q_e = \overline{Q}_e Q_d \overline{Q}_c \overline{Q}_b \overline{Q}_a k_q + Q_e \overline{Q}_d \overline{Q}_c \overline{Q}_b Q_a [0] + Q_e Q_d \overline{Q}_c \overline{Q}_b \overline{Q}_a [0]$ $= \overline{Q}_e Q_d \overline{Q}_c \overline{Q}_b \overline{Q}_a \overline{k}_0 \overline{k}_1 \overline{k}_2 \overline{k}_3 \overline{k}_4 \overline{k}_5 \overline{k}_5 \overline{k}_5 \overline{k}_5 \overline{k}_7 \overline{k}_8 \overline{k}_9$ $+ Q_e \overline{Q}_d \overline{Q}_c \overline{Q}_b \overline{Q}_a \overline{k}_0 \overline{k}_1 \overline{k}_2 \overline{k}_3 \overline{k}_4 \overline{k}_5 \overline{k}_5 \overline{k}_5 \overline{k}_7 \overline{k}_8 \overline{k}_9$ $+ Q_e Q_d \overline{Q}_c \overline{Q}_b \overline{Q}_a \overline{k}_0 \overline{k}_1 \overline{k}_2 \overline{k}_3 \overline{k}_4 \overline{k}_5 \overline{k}_5 \overline{k}_5 \overline{k}_7 \overline{k}_8 \overline{k}_9$

Extraction of Prime Implicants:

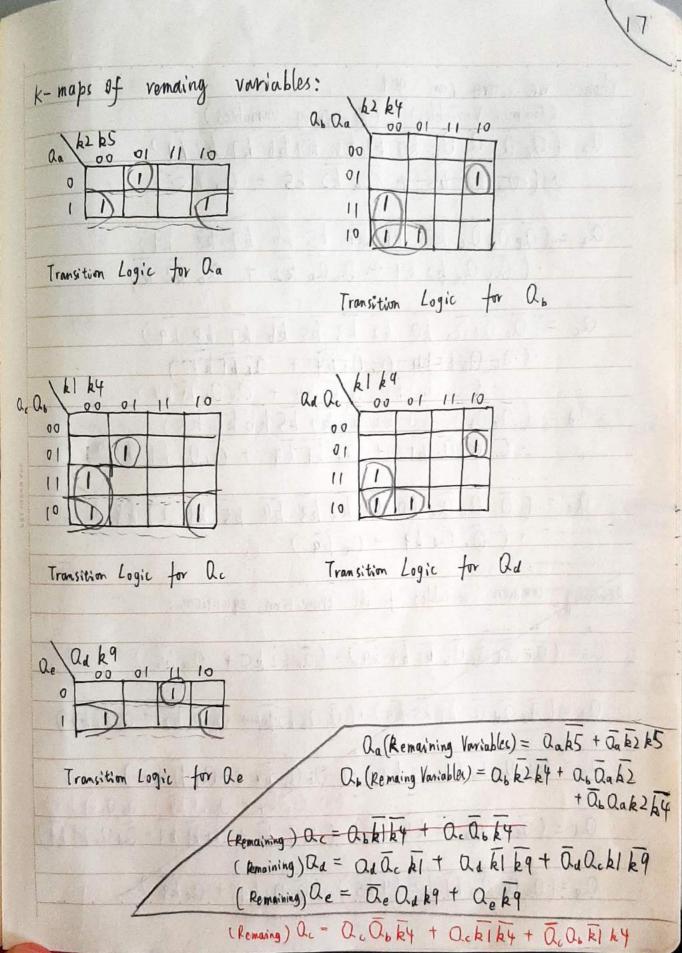
Oa = ((ommon Variables). (Remaining Variables)

Qa = Qe Qd Qc Qb ko kl k3 ky kb k7 k8 k9 (Qak5 + Qak5 · (Qa k2 k5 + Qa k2 k5 + Qa k2 k5)

a = (\overline \overline a \o

ac = (\bar{a} = \bar{a} = \bar{a} = \bar{k} =

Qd = (Oe Q, Qa ko k2 k3 k4 k5 k6 k7 k8) · (Qd Qc k1 k9 + Qd Qc k1 k9 + Qd Qc k1 k9 + Qd Qc k1 k9)



```
Then we can can get:
      (Common Variables) · (Remarking Variables)
 Qa = (Qe Qd Qc Qb RO EI R3 EY kb k7 k8 k9)
      · ( da k2 k5 + Qa k2 k5 + Qa k5)
 a. = (QeOd Qcko ki k3 k5 k6 k7 k8 k9)
       · ( Q6 Qa k2 k4 + Q6 Qa k2 + Q6 k2 k4)
  ac = ( ae ada ko k2 k3 k5 kb k7 k8 k9)
       · ( ac abby ac abby + ac abbiby)
· ( ac abby + ac ki ky + ac abbiby)
  an = ( ae a, a ko ki ki ki ky ks kb k7 k8)
        · ( Qa Qck1 + Qd kl k9 + Qd Qck1 k9)
   ae = ( ac ab aa ko kl k2 k3 k4 k5 kb k7 k8)
         · ( Qe Ouk9 + Qe kg)
Removing common variables of all transition equations:
 a= (ae aacab ki ky k9). (a. k2k5 + aak5)
  as = ( ae ad a. ki ks kg) · ( abaa bzky + abaakz + abkzky)
 Qu = ( Que Qua k2 k4 k5) - (QuQc ki + Onli k9 + Qua Qck k9)
```

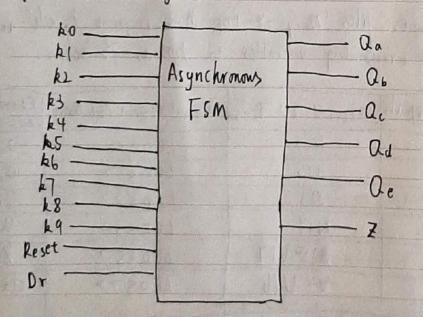
Implementation	: () ynchronou		et De (Clk)
JAM-A —	Synchronous	— Qa	2 Surport
JAM-B	Finite	— Q, kes	et De-bounced clock circuit
JAM_C -	- State -	— ac	De-Bounder Court Con Con
JAM-D -	- Machine -	- Qd	al a Mil
JAM-Grable-	- URBS	Z1/A	967-831
clk_Set	0408	- (lk_Q	clk = Reset + Set·clk
(Input)		(outpur)	

Editing the constraints files "Nexys-4-DDR_MSc_Fragment_1.xde" and map the name of corresponding variables to Jins as 76 below.

Artix-7 10 Pins	Original Digilent Net Name	Edited Digilant Net Name
Ht7	LED<0>	Qa
K15	LEDCIT	06
113	LEO < 27	ac
NIT	LED <3)	ad
R18 VI7	LEO<57	Z
₩ U16	LE0<7>	clk_a
֡֡֜֜֜֜֜֜֜֜֜֜֜֜֜֜֜֜֜֜֜֜֜֜֜֜֜֜֜֜֜֜֜֜֜֜֜	H17 K15 J13 N14 R18 V17	KIS LED < 1 >

(Laput) Syndronous FSM Signal	Artia I/O Pins	Original PhDigilant Nee Name	Edited Digilar Nee No.
JAM-A	JIS	SWCO>	JAM_A
JAM-B	L16	SWC 1>	JAM-B
JAM-C	WI3	SW(27	JAM_C
JAM-D	RIS	SW < 37	JAM-0
JAM_Enable	R17	SW (4>	JAM_Enabl
clk-see	W(8	BTNU	clk-ser
Clk-Reset	P18	BTND	Clk_Rese

Implementation: (Asynchronous)



Editing the constraints tiles as below

(netput) Aggrehmonous FSM signal	Artix 1/0 Pins	Original Digilene Net Name	Edited Digilent Net Name
Qd	HIT	LED <0>	- Qa
Q _b	K15	LED <1>	a
Oc	JI3	LEO <2>	Qe
ad	NI4	LEO (3)	0-4
he	R18	LED (4>	O.e
Z	U16	LE0 <7>	7

(Input) Asynchronous FSM signal	Artic 1/0 Pins	Original Pigillent Net Name	Edited Digileno Net Name
k0	JIS	SW<0>	ko
k	L16	SM < 1>	kl
k2	M13	5 W (2>	k2
k}	R15	(W<3>	k3
ky	RIT	SW <4>	kY
k5	T18	SM < 2>	ks
26	U18	SW (6>	kb
_ k7	RIS	SW <17	k7
k8	T8	SM < 8>	k8
19	V8	SW <9>	k9
Dr	713	SW < 11>	Dr
Reset	U12	2M<13>	Reset

Summary	
1) map numes of variables and pins in the constraints	files
D Ignore the combinational loop alert by adding see property Allow_CombINATIONAL_Loops TRUE [ge	et_nets] '
3) Wire — assign reg — always	
9 initial { sequential execution } —	Test
alway { parallel precution } —	Design
@ Edie tor source tiles properties, clear 'implementation' and 'synthesis' for	Test_Bencl
(b) Synchronous: clock	

Asychronous: race hazard