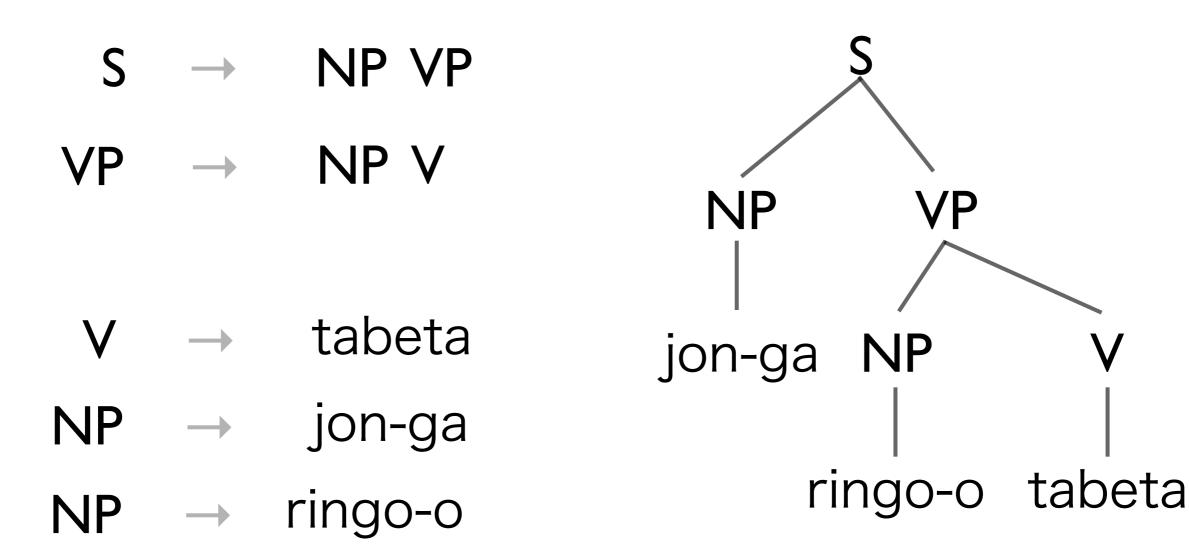
# Decoding and Inference with Syntactic Translation Models



April 8, 2014

### **CFGs**



Output: jon-ga ringo-o tabeta

# Synchronous CFGs

```
S \rightarrow NP VP
```

VP - NP V

V → tabeta

NP → jon-ga

NP → ringo-o

# Synchronous CFGs

```
V → tabeta : ate
```

NP → jon-ga : John

NP → ringo-o : an apple

# Synchronous CFGs

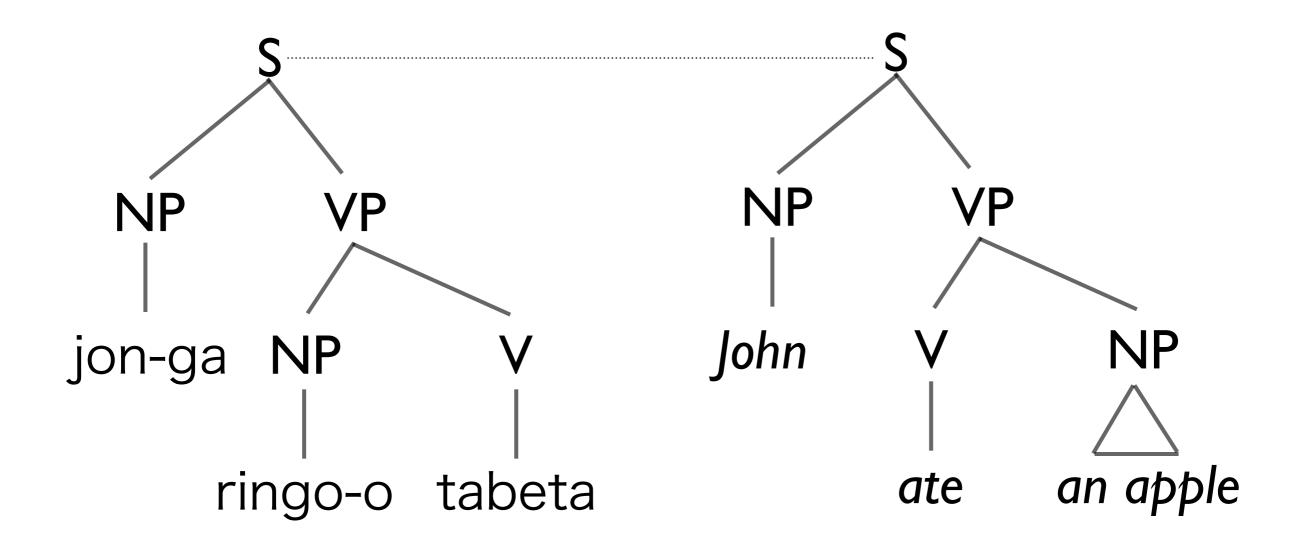
```
S \rightarrow NP VP : 1 2 (monotonic)
VP \rightarrow NP V : 2 1 (inverted)
```

```
V → tabeta : ate
```

NP → jon-ga : John

NP → ringo-o : an apple

# Synchronous generation

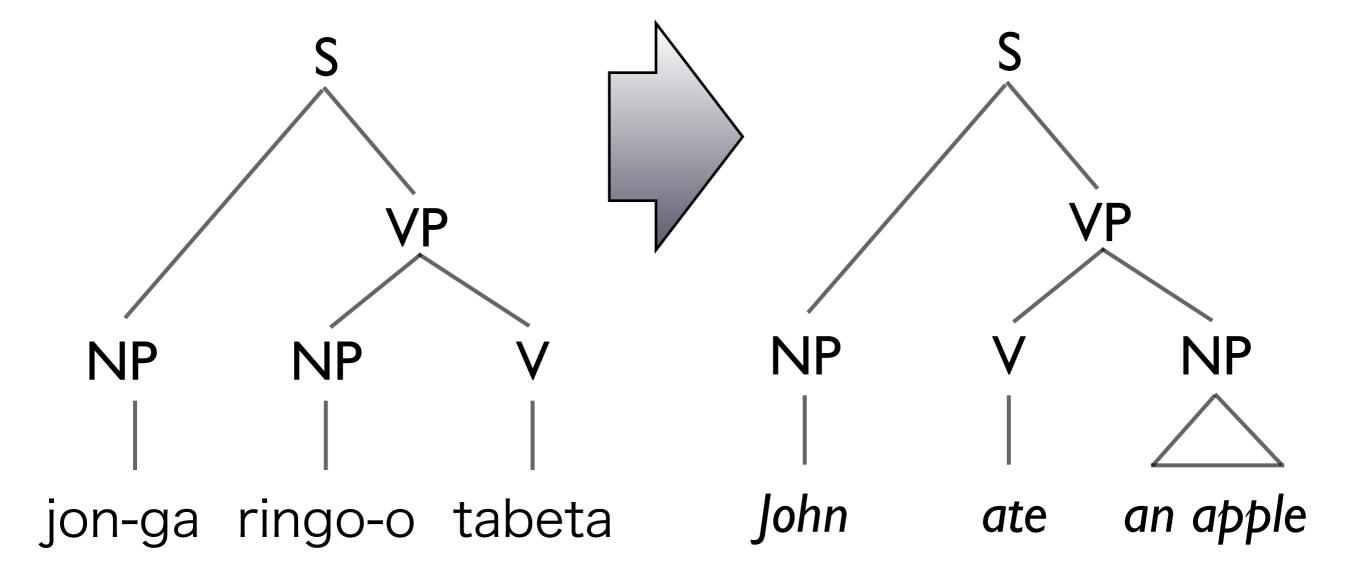


Output: (jon-ga ringo-o tabeta : John ate an apple)

# Translation as parsing

Parse source

Project to target



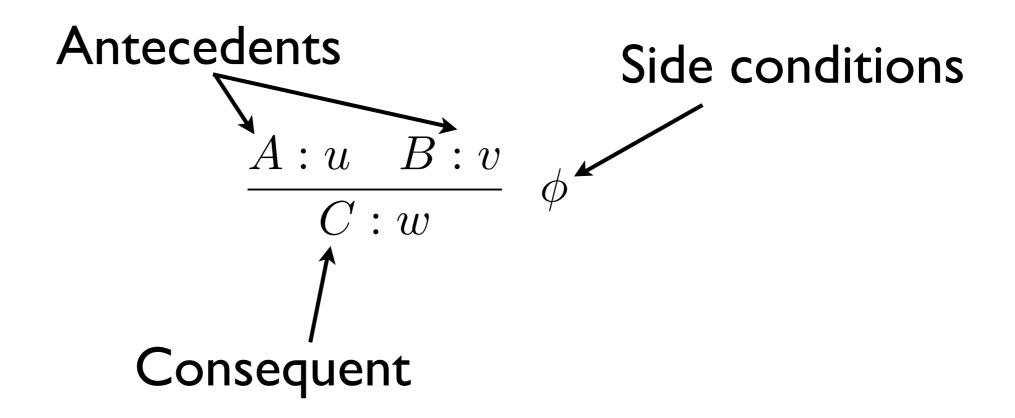
# A closer look at parsing

- Parsing is usually done with dynamic programming
  - Share common computations and structure
  - Represent exponential number of alternatives in polynomial space
  - With SCFGs there are two kinds of ambiguity
    - source parse ambiguity
    - translation ambiguity
    - parse forests can represent both

# A closer look at parsing

- Any monolingual parser can be used (most often: CKY / "dotted" CKY variants)
- Parsing complexity is  $O(|n^3|)$ 
  - cubic in the length of the sentence (n<sup>3</sup>)
  - cubic in the number of non-terminals ( $|G|^3$ )
    - adding nonterminal types increases parsing complexity substantially!
    - With few NTs, exhaustive parsing is tractable

# Parsing as deduction



"If A and B are true with weights u and v, and phi is also true, then C is true with weight w."

# Example: CKY

### Inputs:

$$\mathbf{f} = \langle f_1, f_2, \dots, f_\ell \rangle$$

Context-free grammar in Chomsky normal form.

#### Item form:

[X,i,j] A subtree rooted with NT type X spanning i to j has been recognized.

# Example: CKY

Goal:

$$[S,0,\ell]$$

Axioms:

$$\overline{[X, i-1, i] : w} \quad (X \xrightarrow{w} f_i) \in G$$

Inference rules:

$$\frac{[X,i,k]:u\quad [Y,k,j]:v}{[Z,i,j]:u\times v\times w} \quad (Z\xrightarrow{w} XY)\in G$$

S → PRP VP

 $VP \rightarrow V NP$ 

 $VP \rightarrow V SBAR$ 

SBAR → PRP V

 $NP \rightarrow PRP NN$ 

V → saw

NN → duck

V → duck

 $PRP \rightarrow I$ 

PRP → her





l saw her duck

 $VP \rightarrow V NP$ 

 $VP \rightarrow V SBAR$ 

SBAR → PRP V

 $NP \rightarrow PRP NN$ 

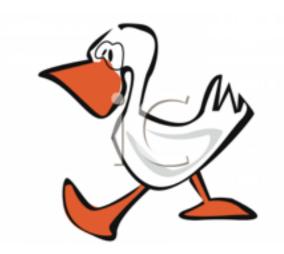
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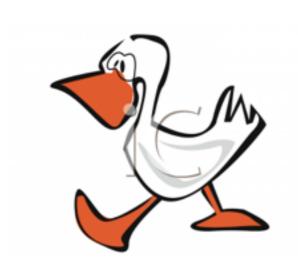


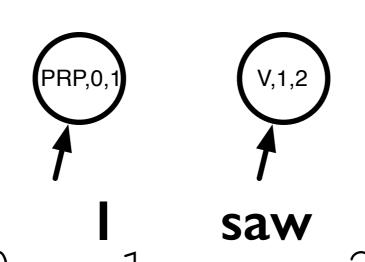
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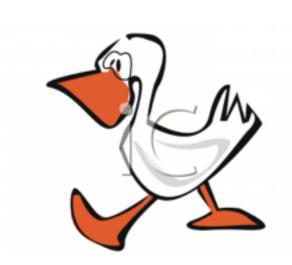
V → saw

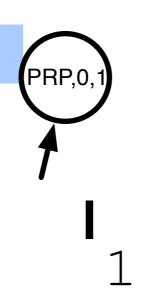
NN → duck

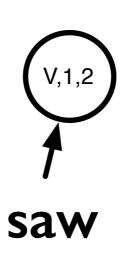
V → duck

 $PRP \rightarrow I$ 

PRP → her









her duck



 $VP \rightarrow V NP$ 

 $VP \rightarrow V SBAR$ 

SBAR → PRP V

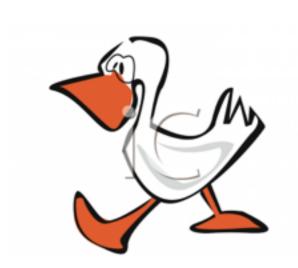
 $NP \rightarrow PRP NN$ 

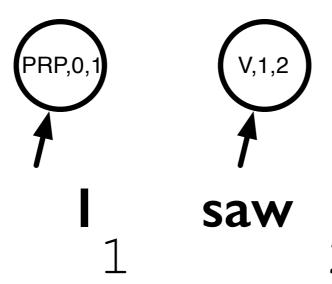
V → saw

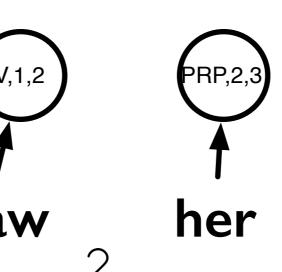
NN → duck

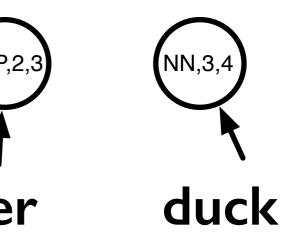
V → duck

 $PRP \rightarrow I$ 











 $VP \rightarrow V NP$ 

 $VP \rightarrow V SBAR$ 

SBAR → PRP V

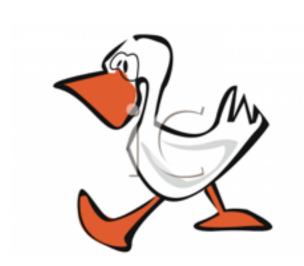
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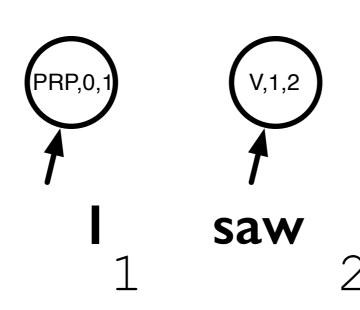
V → saw

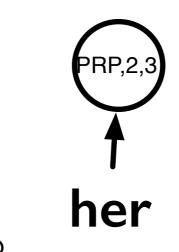
NN → duck

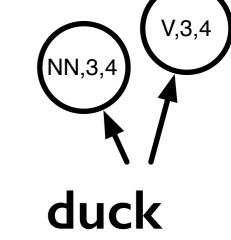
V → duck

 $PRP \rightarrow I$ 













 $VP \rightarrow V NP$ 

 $VP \rightarrow V SBAR$ 

SBAR → PRP V

### NP → PRP NN

V → saw

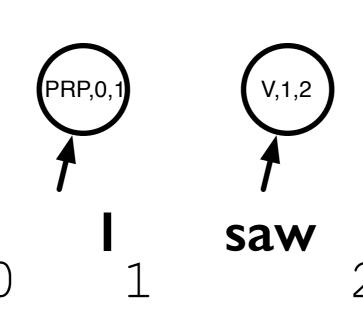
NN → duck

V → duck

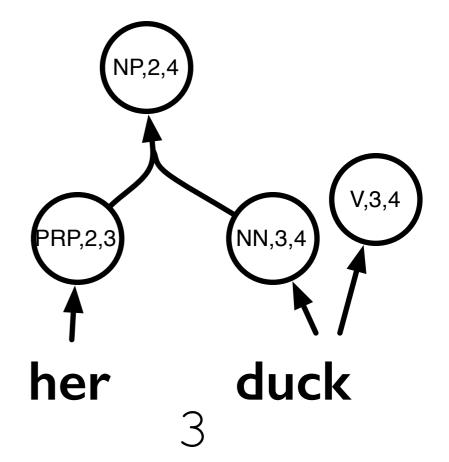
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PRP → her









 $VP \rightarrow V NP$ 

 $VP \rightarrow V SBAR$ 

SBAR → PRP V

 $NP \rightarrow PRP NN$ 

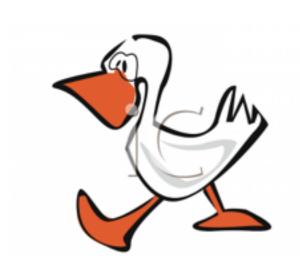
V → saw

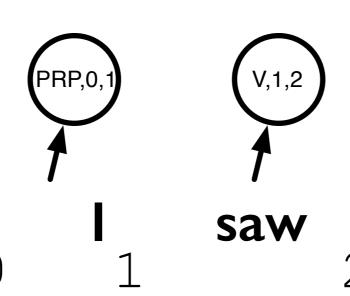
NN → duck

V → duck

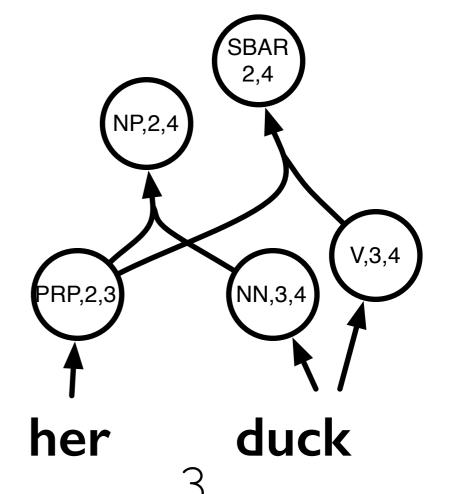
 $PRP \rightarrow I$ 

PRP → her









 $VP \rightarrow V NP$ 

 $VP \rightarrow V SBAR$ 

SBAR → PRP V

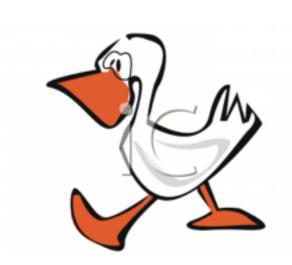
 $NP \rightarrow PRP NN$ 

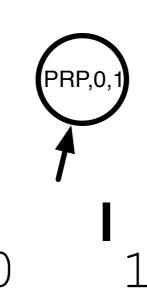
V → saw

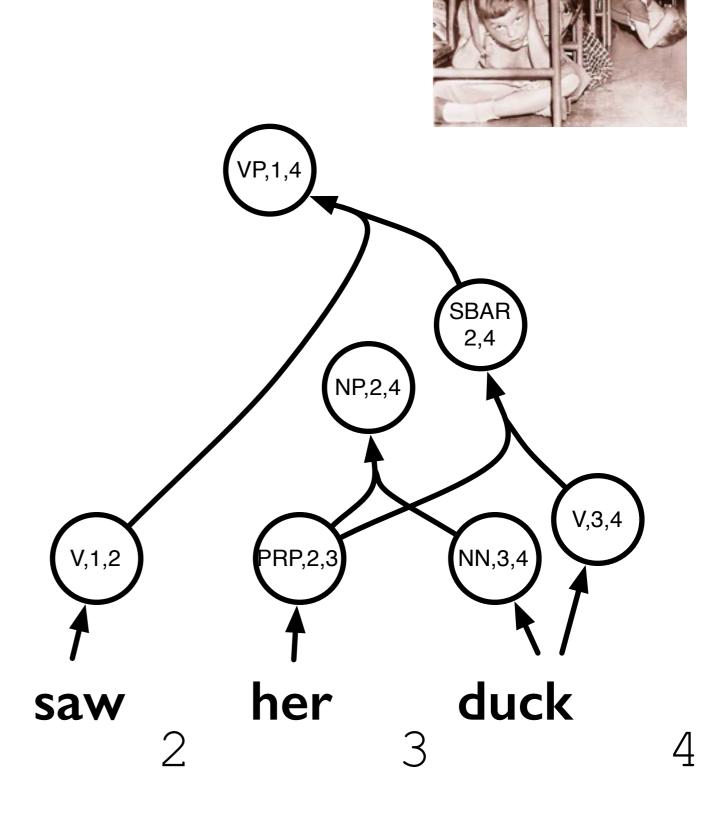
NN → duck

V → duck

 $PRP \rightarrow I$ 







#### $VP \rightarrow V NP$

 $VP \rightarrow V SBAR$ 

SBAR → PRP V

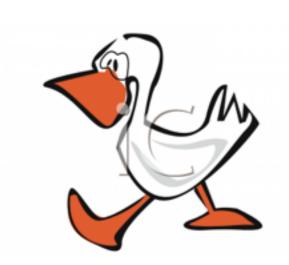
 $NP \rightarrow PRP NN$ 

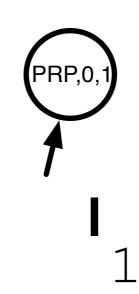
V → saw

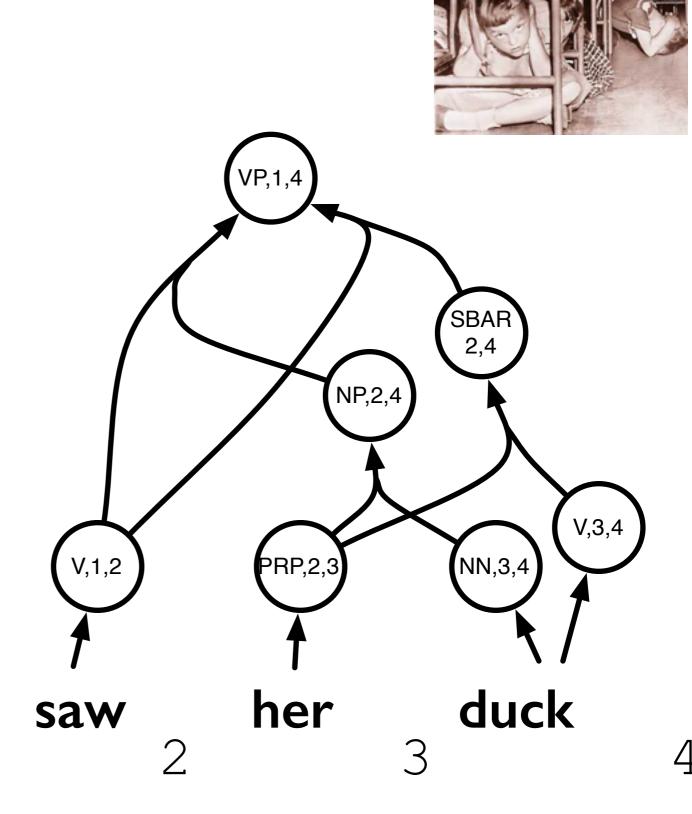
NN → duck

V → duck

 $PRP \rightarrow I$ 







### S → PRP VP

 $VP \rightarrow V NP$ 

 $VP \rightarrow V SBAR$ 

SBAR → PRP V

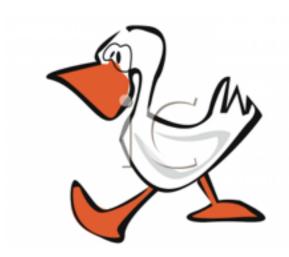
 $NP \rightarrow PRP NN$ 

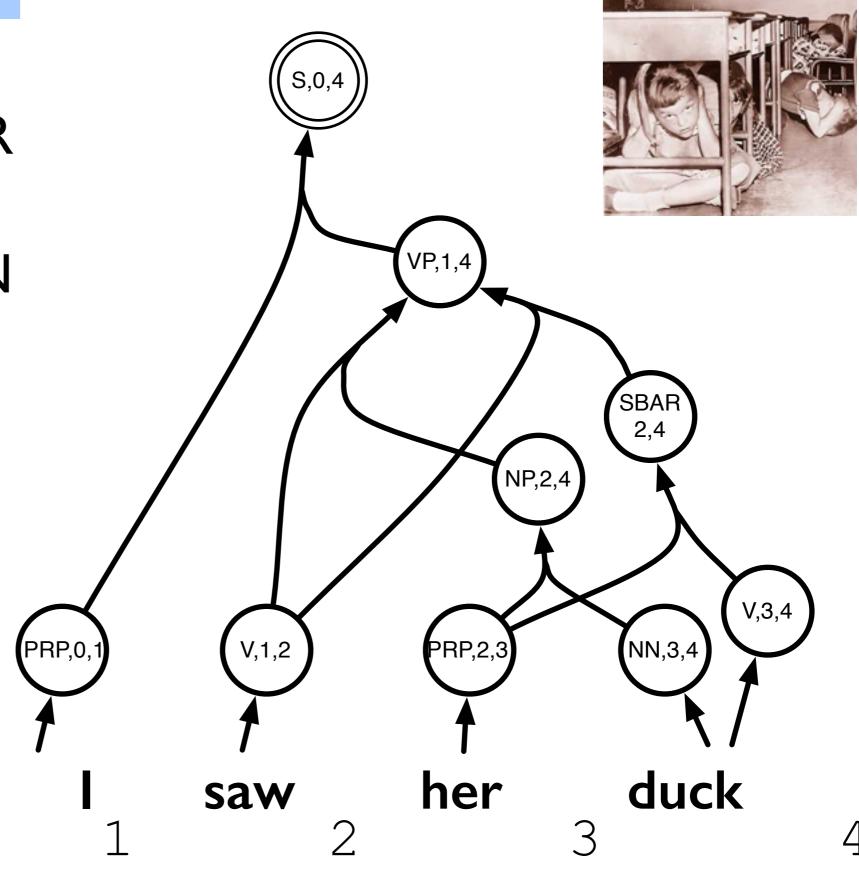
V → saw

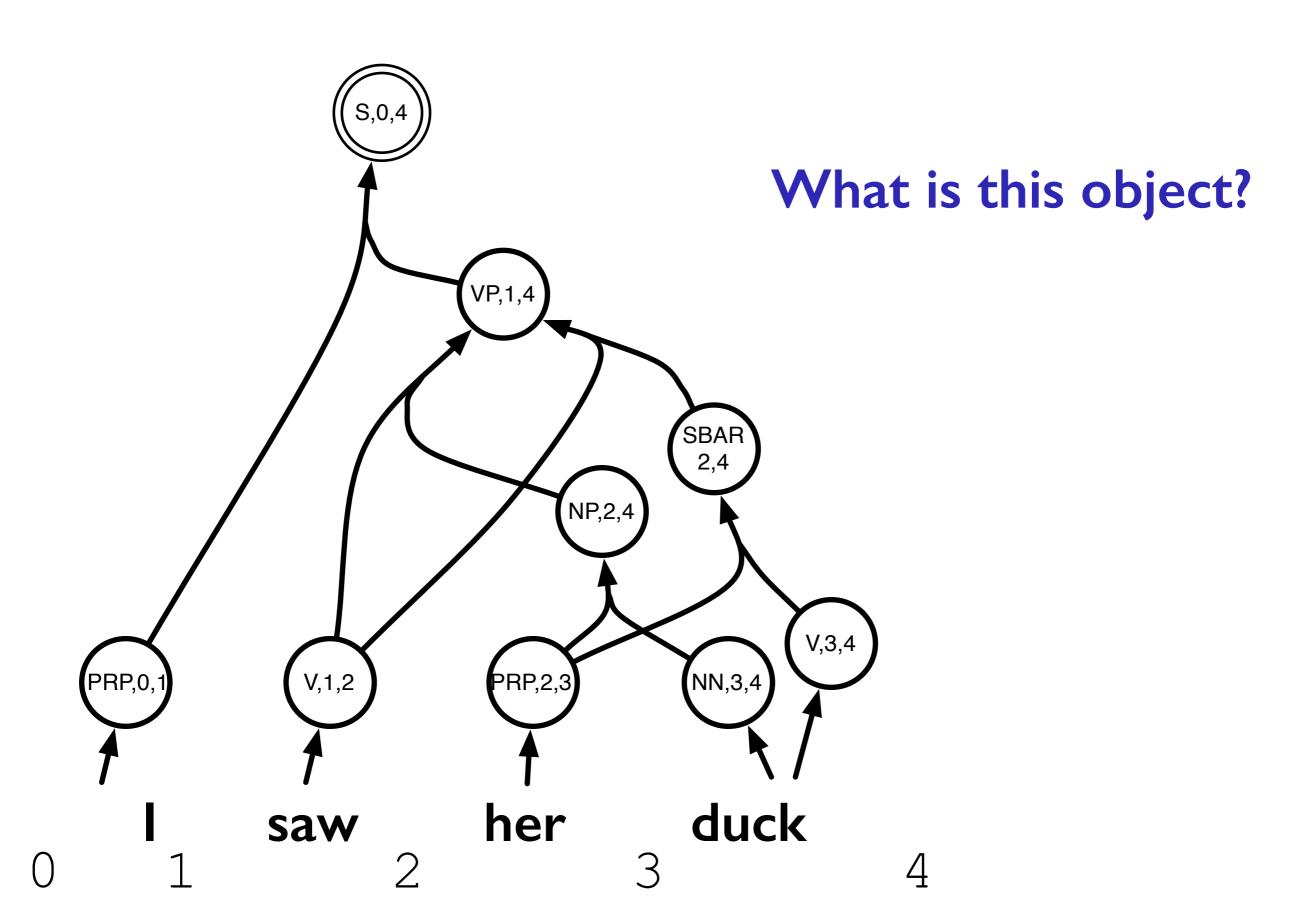
NN → duck

V → duck

 $PRP \rightarrow I$ 







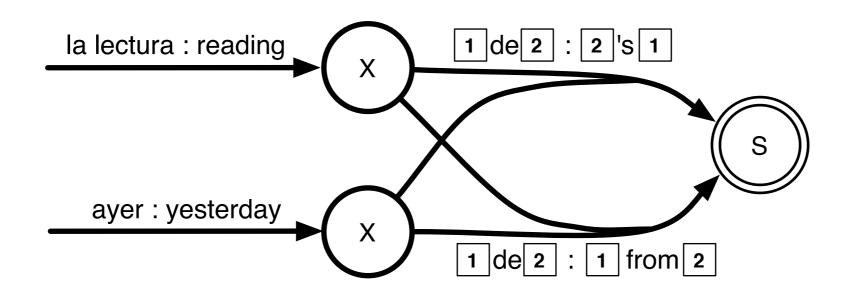
# Semantics of hypergraphs

- Generalization of directed graphs
- Special node designated the "goal"
- Every edge has a single head and 0 or more tails (the arity of the edge is the number of tails)
- Node labels correspond to LHS's of CFG rules
- A derivation is the generalization of the graph concept of path to hypergraphs
- Weights multiply along edges in the derivation, and add at nodes (cf. semiring parsing)

# Edge labels

- Edge labels may be a mix of terminals and substitution sites (non-terminals)
- In translation hypergraphs, edges are labeled in both the source and target languages
- The number of substitution sites must be equal to the arity of the edge and must be the same in both languages
- The two languages may have different orders of the substitution sites
- There is no restriction on the number of terminal symbols

# Edge labels



{ la lectura de ayer : yesterday 's reading }, la lectura de ayer : reading from yesterday }

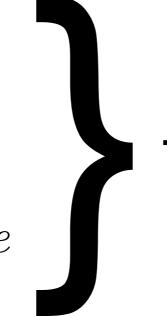
# A Lingua Franca for MT

- Translation hypergraphs are a lingua franca for translation search spaces
  - Note that FST lattices are a special case
- Decoding problem: how do I build a translation hypergraph?
  - For SCFG-translation: just parse

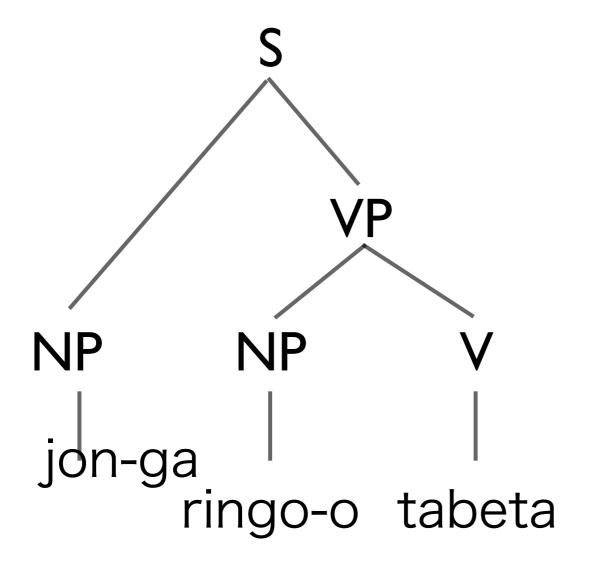
# Tree-to-string Translation

- How do we generate a hypergraph for a tree-tostring translation model?
  - Simple linear-time (given a fixed translation model) top-down matching algorithm
    - Recursively cover "uncovered" sites in tree
  - Each node in the input tree becomes a node in the translation forest
  - For details, Huang et al. (AMTA, 2006) and Huang et al. (EMNLP, 2010)

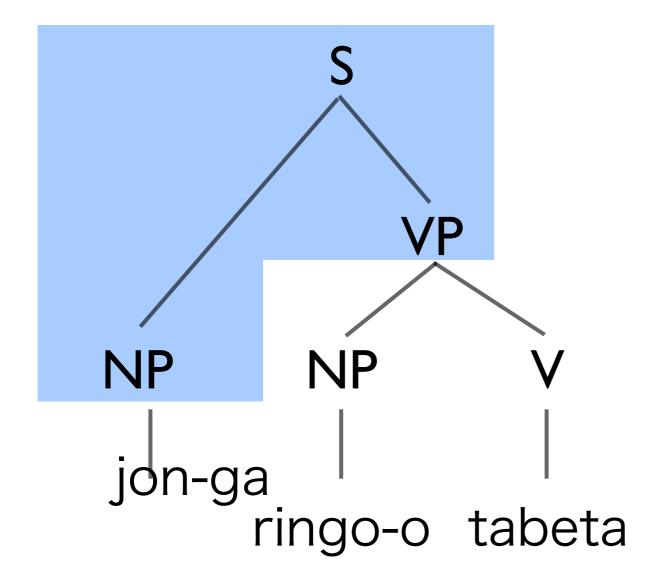
 $S(x_1:NP \ x_2:VP) \rightarrow x_1 \ x_2$   $VP(x_1:NP \ x_2:V) \rightarrow x_2 \ x_1$   $tabeta \rightarrow ate$   $ringo-o \rightarrow an \ apple$   $jon-ga \rightarrow John$ 

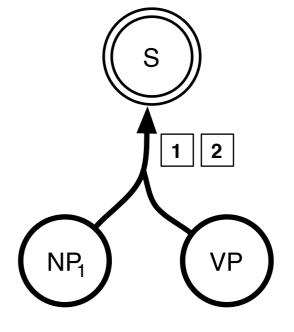


Tree-to-string grammar



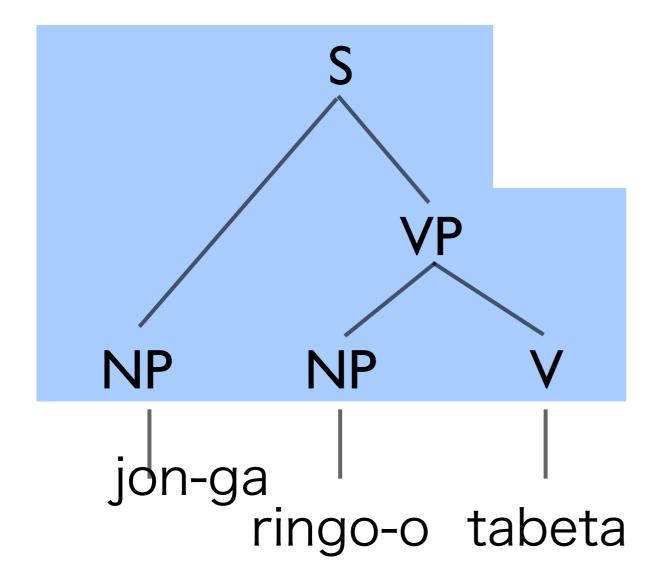
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 $VP(x_1:NP \ x_2:V) \rightarrow x_2 \ x_1$ 
 $tabeta \rightarrow ate$ 
 $ringo-o \rightarrow an \ apple$ 
 $jon-ga \rightarrow John$ 





$$S(x_1:NP \ x_2:VP) \rightarrow x_1 \ x_2$$

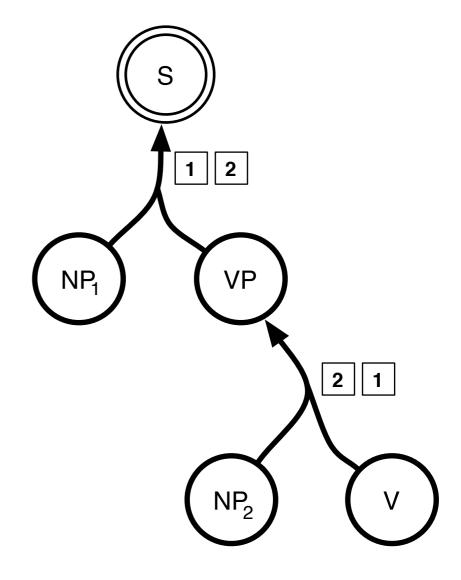
$$VP(x_1:NP \ x_2:V) \rightarrow x_2 \ x_1$$
 
$$tabeta \rightarrow ate$$
 
$$ringo-o \rightarrow an \ apple$$
 
$$jon-ga \rightarrow John$$

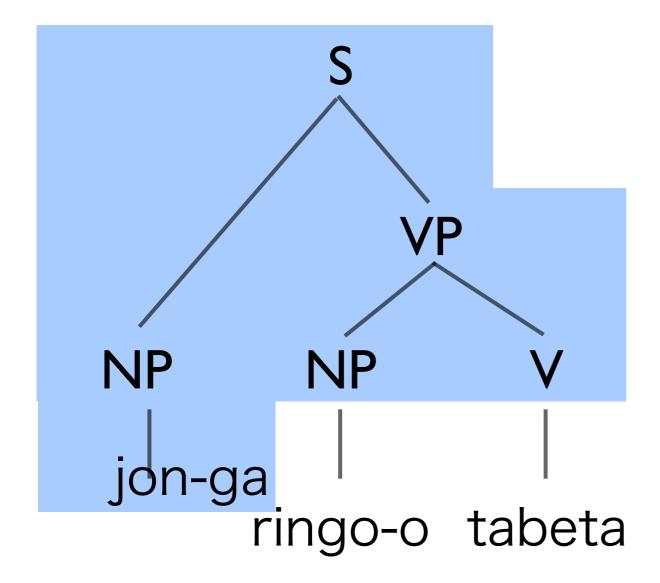


$$S(x_1:NP \ x_2:VP) \rightarrow x_1 \ x_2$$

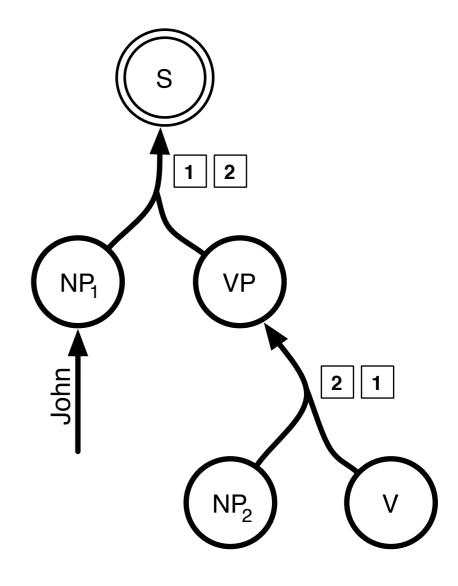
$$VP(x_1:NP \ x_2:V) \rightarrow x_2 \ x_1$$

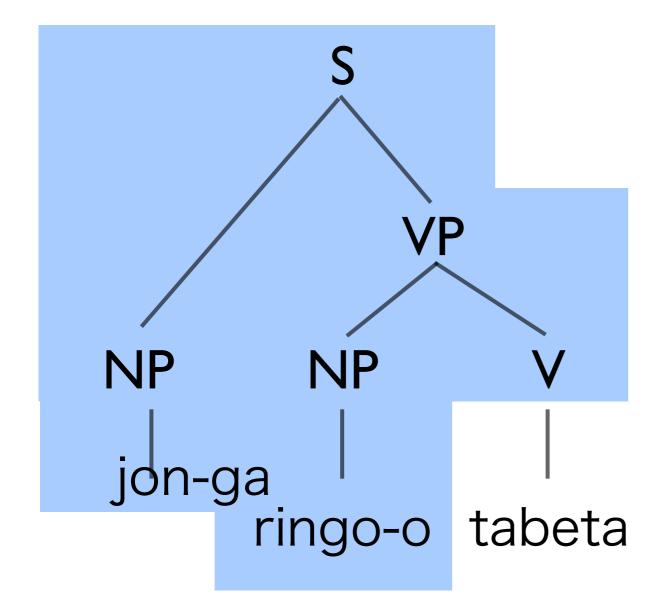
 $tabeta \rightarrow ate$   $ringo-o \rightarrow an \ apple$   $jon-ga \rightarrow John$ 





$$S(x_1:NP \ x_2:VP) \rightarrow x_1 \ x_2$$
 $VP(x_1:NP \ x_2:V) \rightarrow x_2 \ x_1$ 
 $tabeta \rightarrow ate$ 
 $ringo-o \rightarrow an \ apple$ 
 $jon-ga \rightarrow John$ 

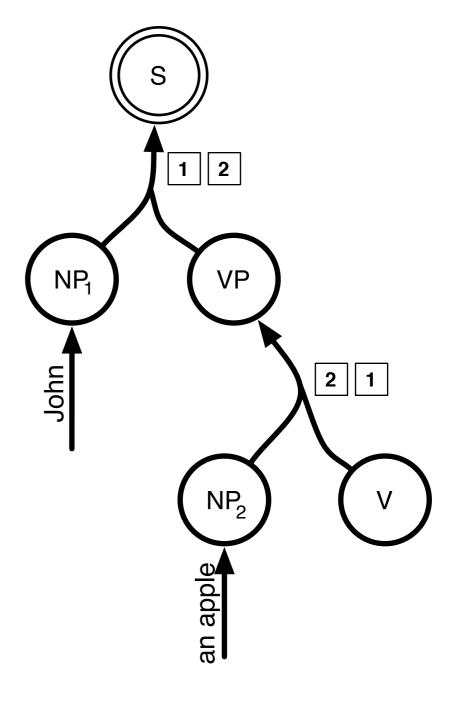


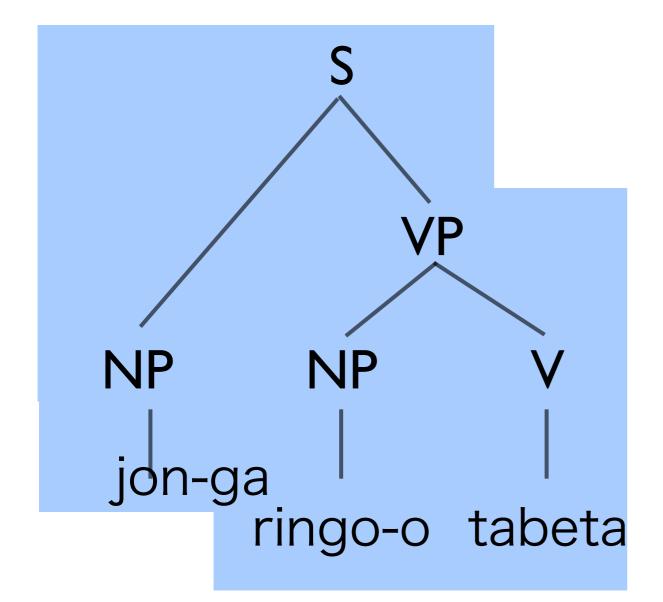


$$S(x_1:NP \ x_2:VP) \rightarrow x_1 \ x_2$$
 $VP(x_1:NP \ x_2:V) \rightarrow x_2 \ x_1$ 
 $tabeta \rightarrow ate$ 

$$ringo-o \rightarrow an \ apple$$

jon- $ga \rightarrow John$ 



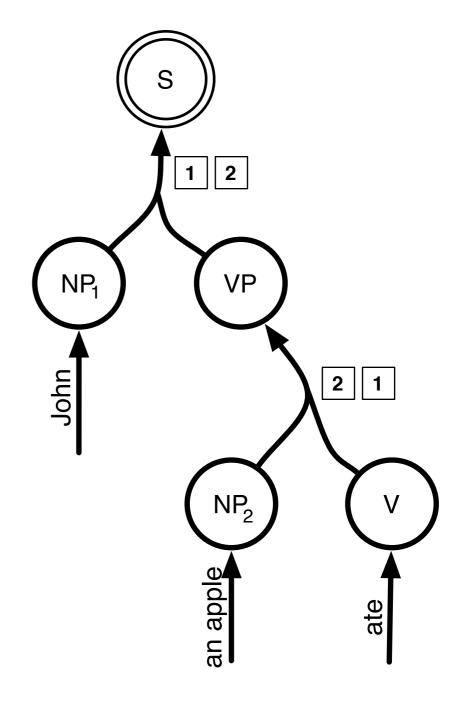


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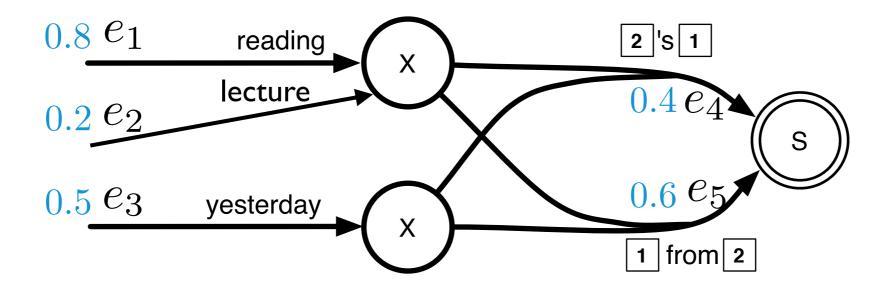
#### $tabeta \rightarrow ate$

 $ringo-o \rightarrow an \ apple$   $jon-ga \rightarrow John$ 



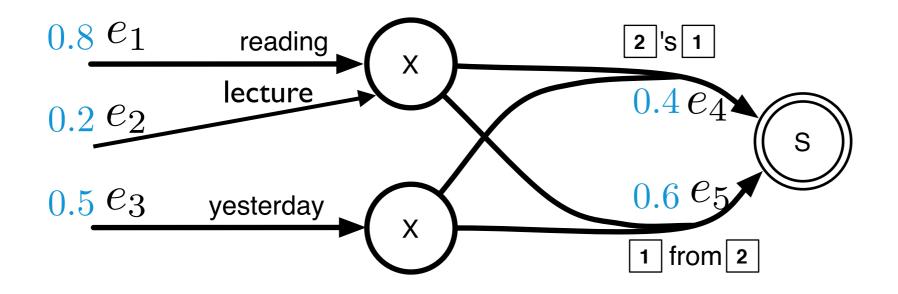
# Working With Hypergraphs

#### Derivations

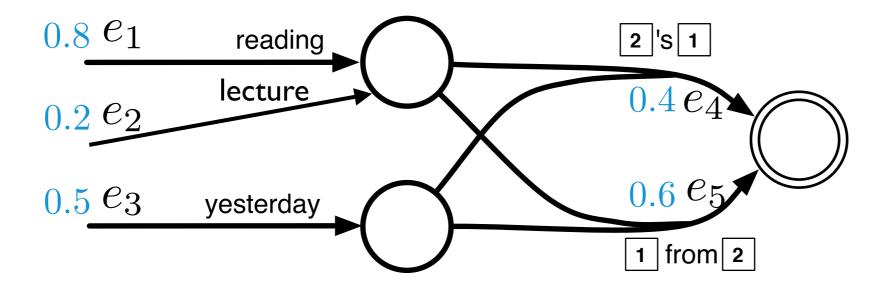


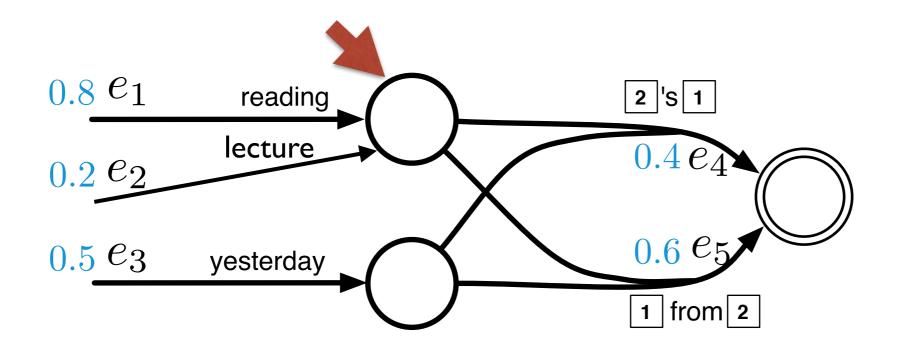
$$egin{aligned} oldsymbol{d}_1 &= e_4 e_1 e_3 & y(oldsymbol{d}_1) &= yesterday ext{'s reading} \ oldsymbol{d}_2 &= e_5 e_1 e_3 & y(oldsymbol{d}_2) &= reading ext{ from yesterday} \ oldsymbol{d}_3 &= e_4 e_2 e_3 & y(oldsymbol{d}_3) &= yesterday ext{'s lecture} \ oldsymbol{d}_4 &= e_5 e_2 e_3 & y(oldsymbol{d}_4) &= lecture ext{ from yesterday} \end{aligned}$$

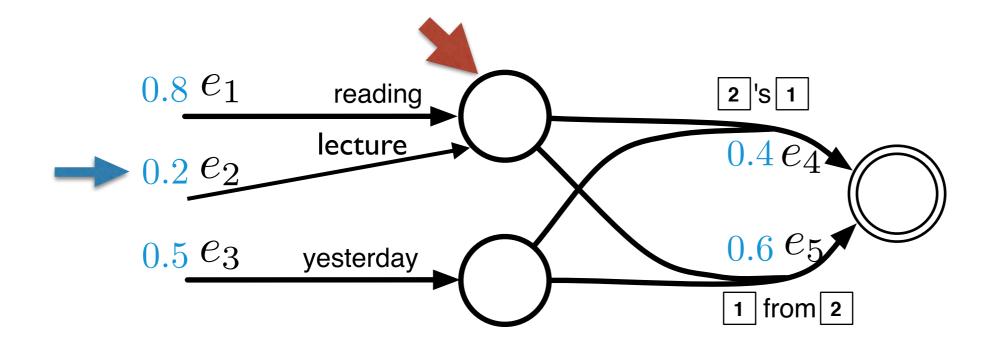
#### Derivations

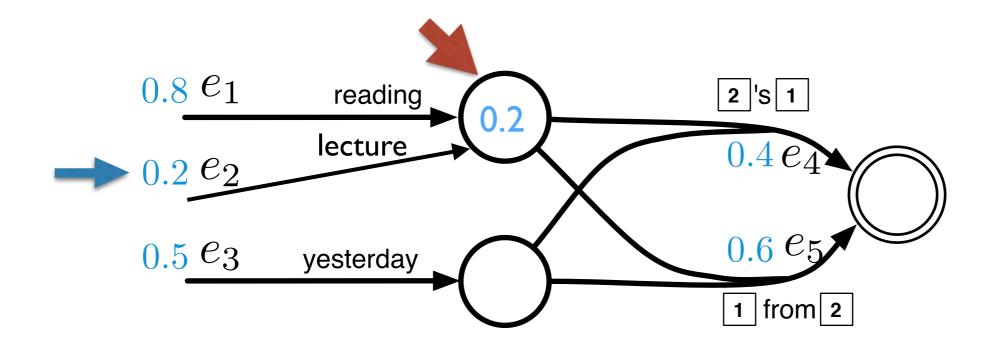


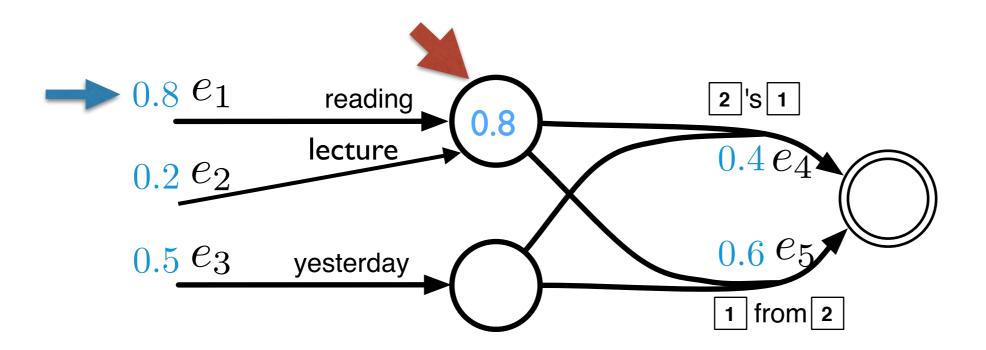
$$egin{aligned} m{d}_1 &= e_4 e_1 e_3 & w[m{d}_1] &= 0.4 \cdot 0.8 \cdot 0.5 = 0.16 \\ m{d}_2 &= e_5 e_1 e_3 & w[m{d}_2] &= 0.6 \cdot 0.8 \cdot 0.5 = 0.24 \\ m{d}_3 &= e_4 e_2 e_3 & w[m{d}_3] &= 0.4 \cdot 0.2 \cdot 0.5 = 0.04 \\ m{d}_4 &= e_5 e_2 e_3 & w[m{d}_3] &= 0.6 \cdot 0.2 \cdot 0.5 = 0.06 \end{aligned}$$

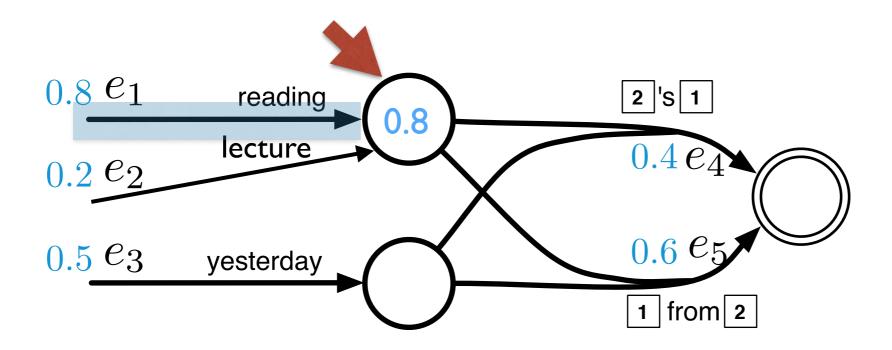


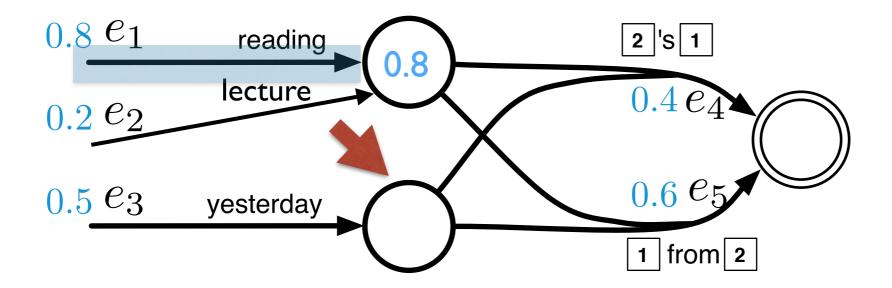


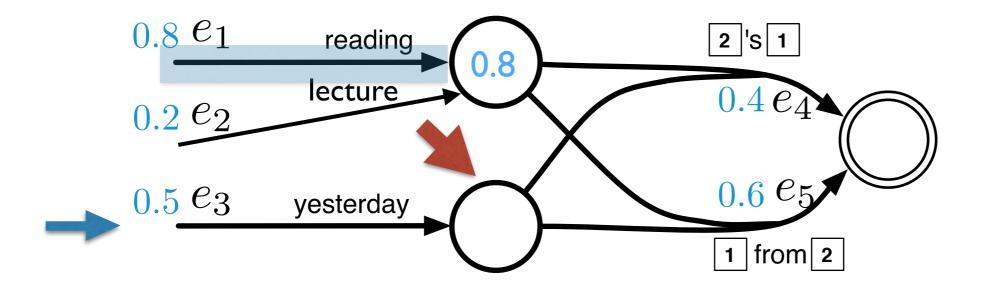


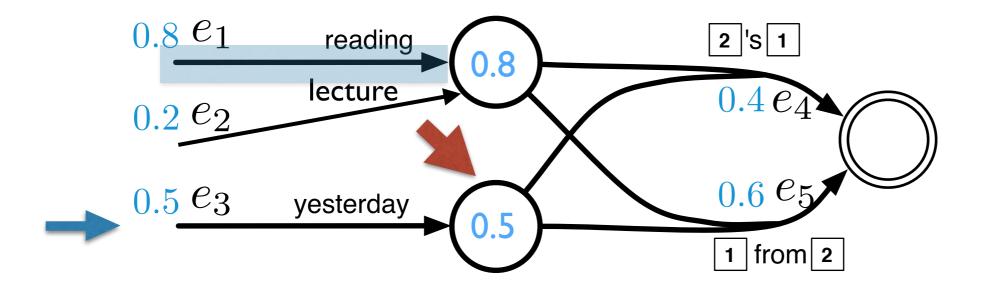


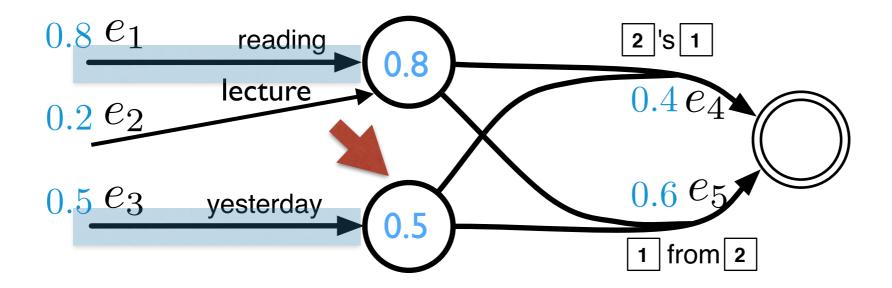


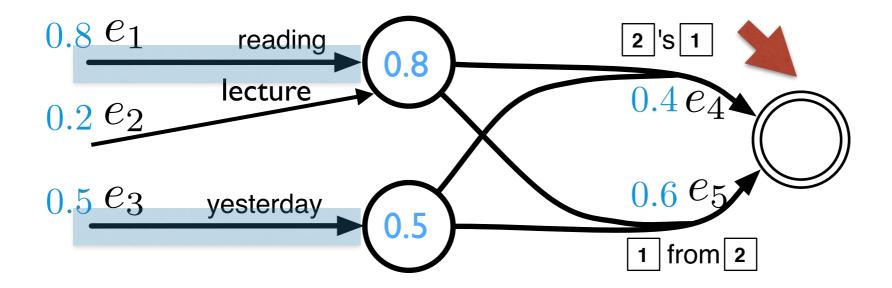


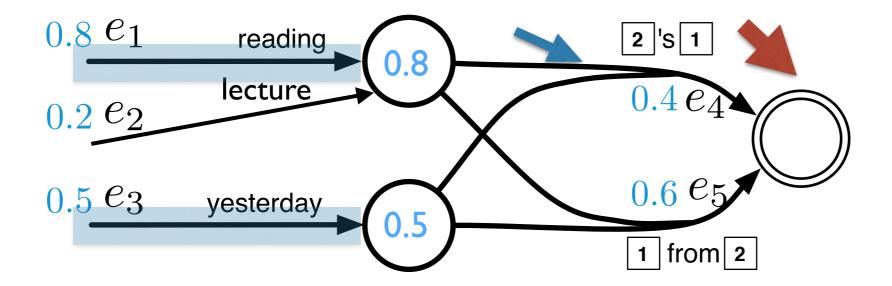


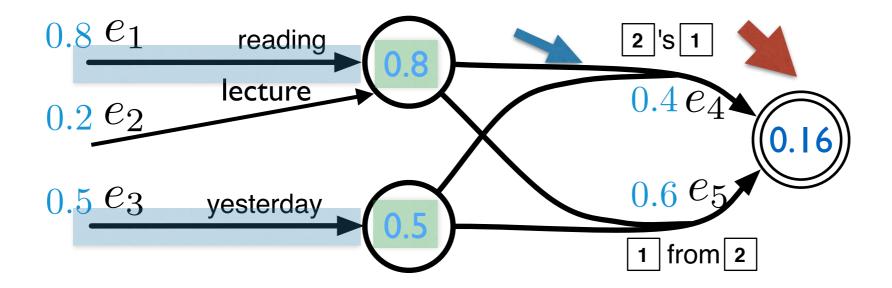




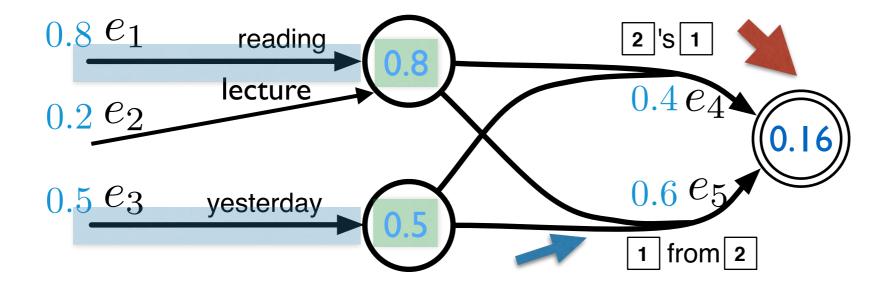


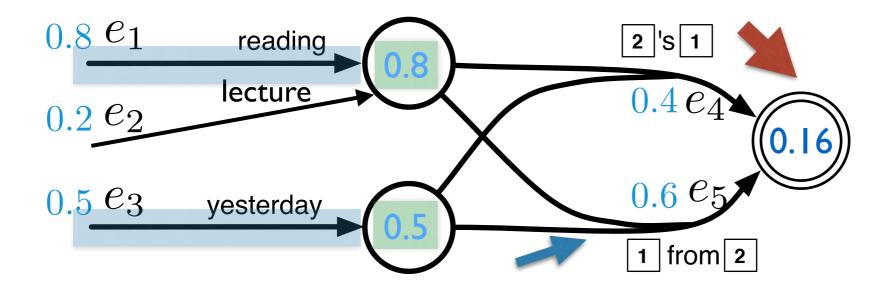




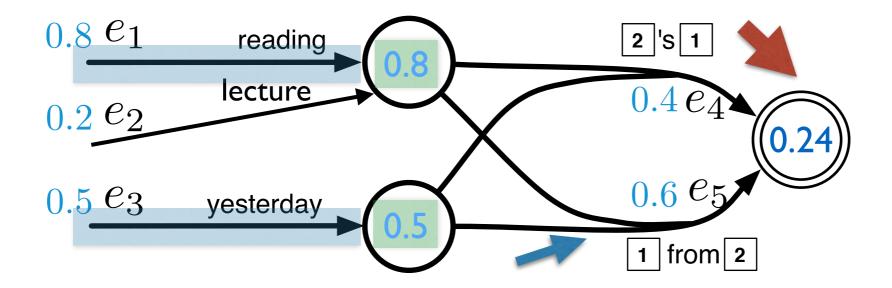


 $0.8 \times 0.5 \times 0.4 = 0.16$ 

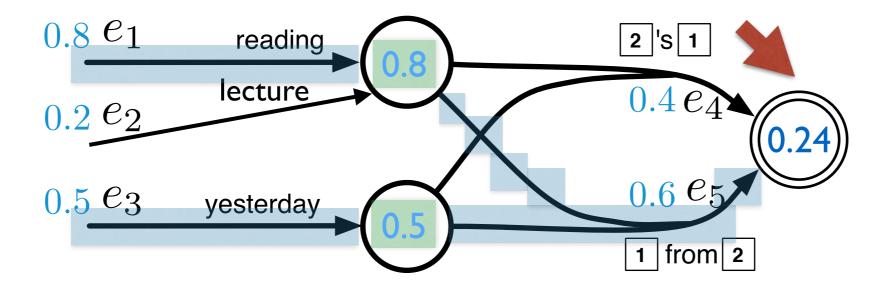


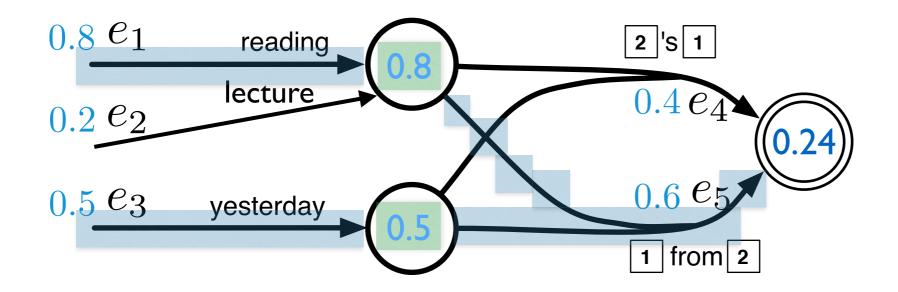


 $0.8 \times 0.5 \times 0.6 = 0.24$ 



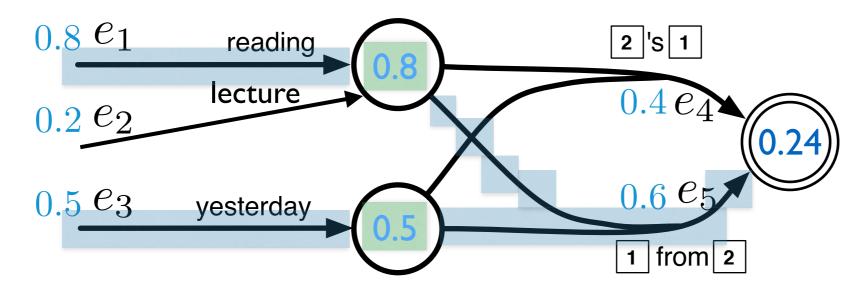
 $0.8 \times 0.5 \times 0.6 = 0.24$ 





Best yield: reading from yesterday

Best path: 0.24



Best yield: reading from yesterday Best path: 0.24

$$egin{aligned} m{d}_1 &= e_4 e_1 e_3 & w[m{d}_1] &= 0.4 \cdot 0.8 \cdot 0.5 = 0.16 \\ m{d}_2 &= e_5 e_1 e_3 & w[m{d}_2] &= 0.6 \cdot 0.8 \cdot 0.5 = 0.24 \\ m{d}_3 &= e_4 e_2 e_3 & w[m{d}_3] &= 0.4 \cdot 0.2 \cdot 0.5 = 0.04 \\ m{d}_4 &= e_5 e_2 e_3 & w[m{d}_3] &= 0.6 \cdot 0.2 \cdot 0.5 = 0.06 \end{aligned}$$

# Other Algorithms

- Given a weighted hypergraph
- In the Viterbi (Inside) algorithm, there are two operations
  - Multiplication (extend path)
  - Maximization (chose between paths)
- Semirings generalize these to compute other quantities

# Semirings

semiring	$\mathbb{K}$	$\oplus$	$\otimes$	$\overline{0}$	$\overline{1}$	notes
Boolean	{0,1}	V	$\wedge$	0	1	idempotent
count	$\mathbb{N}_0 \cup \{\infty\}$	+	×	0	1	
probability	$\mathbb{R}_+ \cup \{\infty\}$	+	×	0	1	
tropical	$\mathbb{R} \cup \{-\infty,\infty\}$	max	+	-∞	0	idempotent
log	$\mathbb{R} \cup \{-\infty,\infty\}$	$\oplus_{log}$	+	_∞	0	

# Inside Algorithm

$$\alpha(q_{goal}) = \bigoplus_{\mathbf{d} \in \mathcal{G}} \bigotimes_{e \in \mathbf{d}} w(e)$$

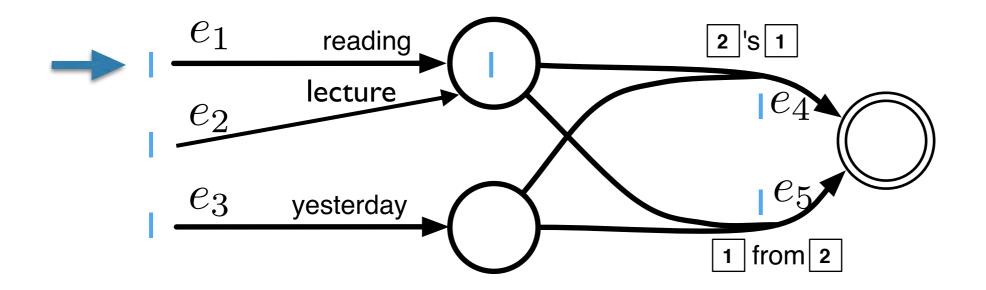
```
1: function Inside(G, K)
                                                     \triangleright G is an acyclic hypergraph and K is a semiring
        for q in topological order in G do
           if B(q) = \emptyset then
 3:
               \alpha(q) \leftarrow \overline{1}
                                                           > assume states with no in-edges are axioms
 4:
           else
 5:
               \alpha(q) \leftarrow \overline{0}
 6:
              for all e \in B(q) do

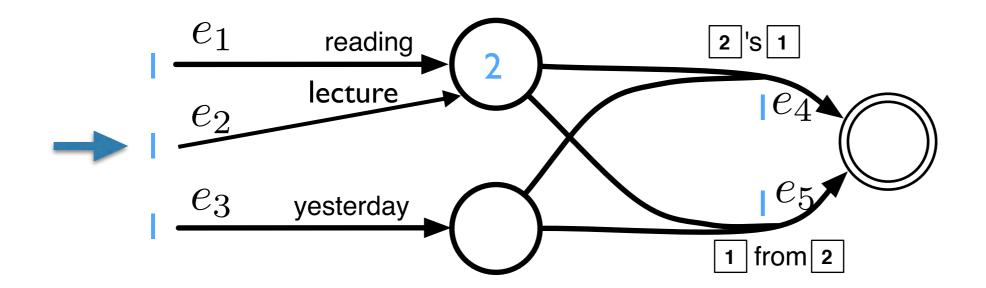
    all in-coming edges to node q

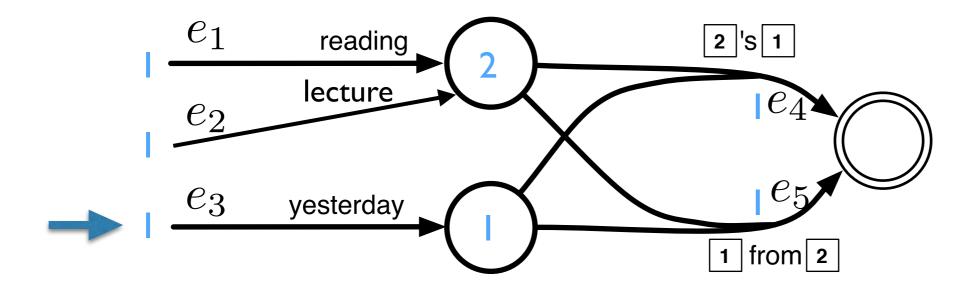
             k \leftarrow w(e)
                  for all r \in \mathbf{t}(e) do

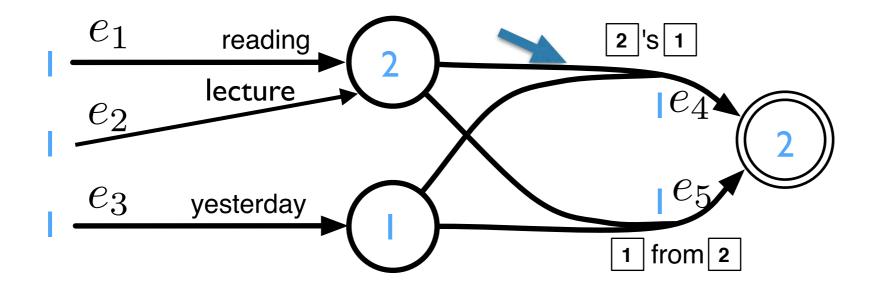
⇒ all tail (previous) nodes of edge e

                     k \leftarrow k \otimes \alpha(r)
10:
                  \alpha(q) \leftarrow \alpha(q) \oplus k
11:
12:
        return \alpha
```

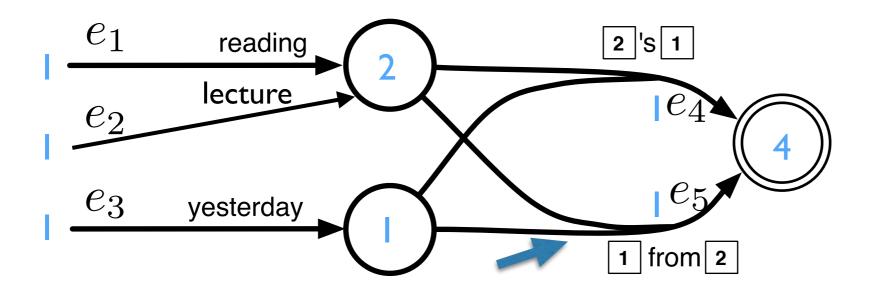




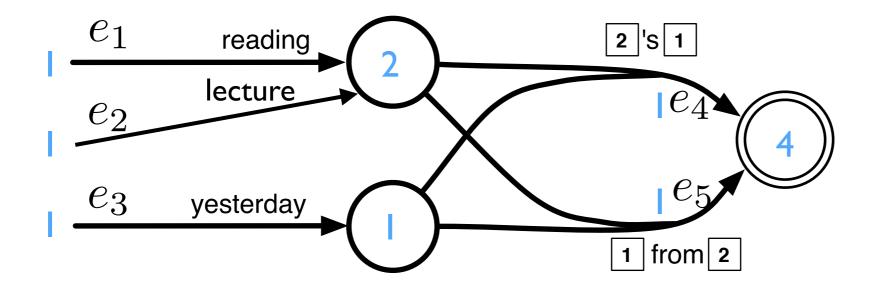




$$2 \times | \times | = 2$$



$$2 \times | \times | = 2$$



#### Inside-Outside

```
1: function Outside(G, K, \alpha)
                                                                             \triangleright \alpha is the result of INSIDE(G, K)
        for all q \in \mathcal{G} do
 2:
            \beta(q) \leftarrow \overline{0}
 3:
        \beta(q_{goal}) = \overline{1}
        for q in reverse topological order in G do
 5:
            for all e \in B(q) do
                                                                               \triangleright all in-coming edges to node q
 6:
               for all r \in \mathbf{t}(e) do
                                                                          \triangleright all tail (previous) nodes of edge e
                   k \leftarrow w(e) \otimes \beta(q)
 8:
                   for all s \in \mathbf{t}(e) do
                                                                 \triangleright all tail (previous) nodes of edge e, again
                      if r \neq s then
10:
                          k \leftarrow k \otimes \alpha(s)
                                                                                        ▷ incorporate inside score
11:
                   \beta(r) \leftarrow \beta(r) \oplus k
12:
        return B
13:
 1: function InsideOutside(G, K)
                                                                                       \alpha \leftarrow \text{INSIDE}(G, K)
 2:
        \beta \leftarrow \text{OUTSIDE}(G, K, \alpha)
        for edge e in G do
 4:
            \gamma(e) \leftarrow w(e) \otimes \beta(n(e))

    bedge weight and outside score of edge's head node

 5:
            for all q \in \mathbf{t}(e) do
 6:
               \gamma(e) \leftarrow \gamma(e) \otimes \alpha(q)

    inside score of tail nodes

                                                                                \triangleright \gamma(e) is the edge marginal of e
 8:
        return γ
```

#### Inside-Outside

- Compute lots of interesting quantities
  - The score of the best path through each edge
  - The total number of derivations that contain an edge
  - The total score of all derivations going through an edge

# Inference algorithms

- Viterbi O(|E| + |V|)
  - Find the maximum weighted derivation
  - Requires a partial ordering of weights
- Inside outside O(|E| + |V|)
  - Compute the marginal (sum) weight of all derivations passing through each edge/node
- k-best derivations  $O(|E| + |D_{max}| k \log k)$ 
  - Enumerate the k-best derivations in the hypergraph
  - See IWPT paper by Huang and Chiang (2005)

# Things to keep in mind

Bound on the number of edges (SCFG):

$$|E| \in O(n^3|G|^3)$$

Bound on the number of nodes:

$$|V| \in O(n^2|G|)$$

#### Next time

#### What about the LM?