

LSU RTG

Hyperplane Arrangements

Talk 1: Introduction

For this talk we will live in \mathbb{R}^n

hyperplane

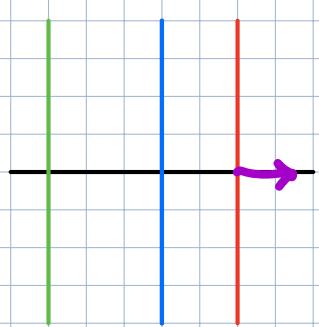
Dfn.

An affine hyperplane H is an $(n-1)$ -dimensional subspace of \mathbb{R}^n . In particular,

$$H = \{v \in \mathbb{R}^n \mid \alpha \cdot v = c\}$$

normal of H

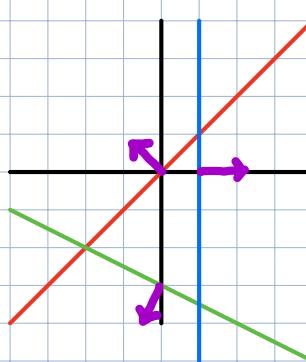
for some fixed, nonzero $\alpha \in \mathbb{R}^n$ and some $c \in \mathbb{R}$. If $c=0$, we say that H is a linear hyperplane (note that lin. h.p. are also affine).



$$H_1: (1,0) \cdot (x,y) = -3$$

$$H_2: (1,0) \cdot (x,y) = 0$$

$$H_3: (1,0) \cdot (x,y) = 2$$



$$H_1: (1,2) \cdot (x,y) = -3$$

$$H_2: (1,0) \cdot (x,y) = 1$$

$$H_3: (1,-1) \cdot (x,y) = 0$$

B1

For this talk we will live in \mathbb{R}^n

hyperplane

Dfn.

An affine hyperplane H is an $(n-1)$ -dimensional subspace of \mathbb{R}^n . In particular,

$$H = \{v \in \mathbb{R}^n \mid \alpha \cdot v = c\}$$

normal of H

nonzero

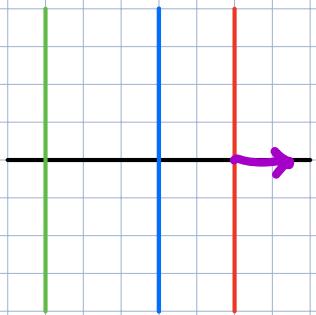
for some fixed, nonzero $\alpha \in \mathbb{R}^n$ and some $c \in \mathbb{R}$. If $c=0$, we say that H is a linear hyperplane (note that lin. h.p. are also affine).

A collection of finitely many hyperplanes $A = \{H_i\}_i$ is called a hyperplane arrangement.

The dimension of A , $\dim(A) = n$.

The rank of A , $\text{rank}(A) = \dim(\text{span}\{\alpha_1, \dots, \alpha_K\})$

If $\text{rank}(A) = \dim(A)$, A is essential.



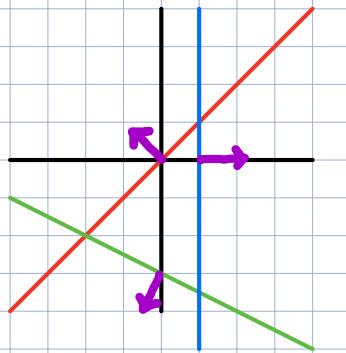
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$$H_2: (1, 0) \cdot (x, y) = 0$$

$$H_3: (1, 0) \cdot (x, y) = 2$$

$\lambda_1 \uparrow$

$$\dim(\lambda_1) = 2, \text{ rank}(\lambda_1) = 1$$



$$H_1: (1, 2) \cdot (x, y) = -3$$

$$H_2: (1, 0) \cdot (x, y) = 1$$

$$H_3: (1, 1) \cdot (x, y) = 0$$

$\lambda_2 \uparrow$

$$\dim(\lambda_2) = 2, \text{ rank}(\lambda_2) = 2,$$

For this talk we will live in \mathbb{R}^n

Dfn. hyperplane

An affine hyperplane H is an $(n-1)$ -dimensional subspace of \mathbb{R}^n . In particular,

$$H = \{v \in \mathbb{R}^n \mid \alpha \cdot v = c\}$$

normal of H

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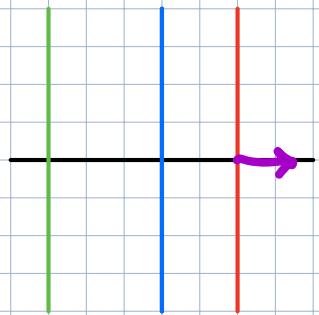
If $\text{rank}(A) = \dim(A)$, A is essential.

Let $R(A)$ be the set of connected components of

$$X = \mathbb{R}^n \setminus \bigcup_{H \in A} H$$

set of regions

$$\text{and } r(A) = |R(A)|.$$



$$H_1: (1, 0) \cdot (x, y) = -3$$

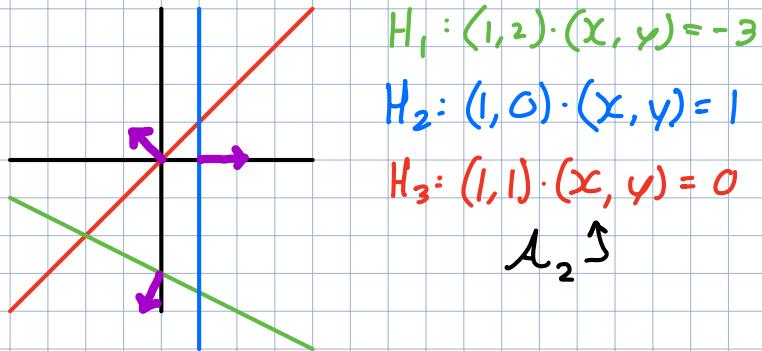
$$H_2: (1, 0) \cdot (x, y) = 0$$

$$H_3: (1, 0) \cdot (x, y) = 2$$

$\lambda_1 \uparrow$

$$\dim(\lambda_1) = 2, \text{ rank}(\lambda_1) = 1$$

$$r(\lambda_1) = 4.$$



$$H_1: (1, 2) \cdot (x, y) = -3$$

$$H_2: (1, 0) \cdot (x, y) = 1$$

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$\lambda_2 \uparrow$

$$\dim(\lambda_2) = 2, \text{ rank}(\lambda_2) = 2,$$

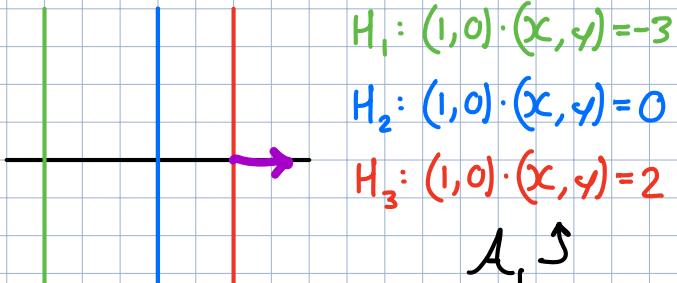
$$r(\lambda_2) = 7$$

Q: How to "make" an arrangement essential?

A: Essentialization.

- ① Let A be a h.p. arr. in \mathbb{R}^n w/ h.p.s H_1, \dots, H_K .
- ② Form $W = \text{span}\{\alpha_1, \dots, \alpha_K\}$ for α_i normal of H_i .
- ③ Let $\text{ess}(A) = \{H_1 \cap W, \dots, H_K \cap W\}$

Now $\text{ess}(A)$ is an essential hyperplane arrangement!



$$\dim(\lambda_1) = 2, \text{ rank}(\lambda) = 1$$

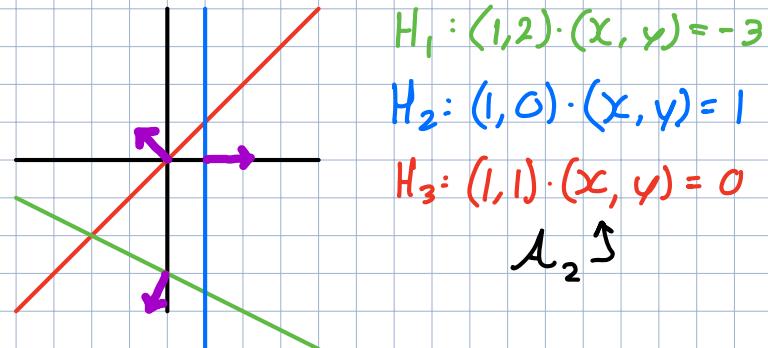
$$r(\lambda) = 4$$

$$X = x\text{-axis}$$

↓ ess

$H_1: x = 3$
 $H_2: x = 0$
 $H_3: x = 2$

$\dim(\text{ess}(\lambda_1)) = 1, \quad r(\lambda_1) = 4$
 $\text{rank}(\text{ess}(\lambda_1)) = 1,$

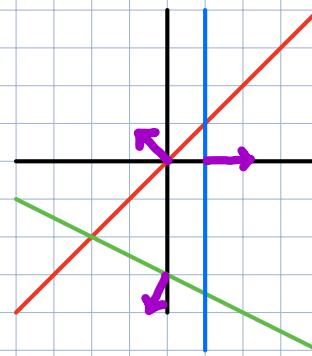


$$\dim(\lambda_2) = 2, \text{ rank}(\lambda_2) = 2,$$

$$r(\lambda_2) = 7$$

$$X = \mathbb{R}^2$$

↓ ess



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Hyperplane Arrangements

Talk 7: Generic Translations and Face
Counting

Dfn

For \mathcal{A} an arrangement of M hyperplanes given by

$$H_i = \{v \in V \mid \alpha_i \cdot v = c_i\},$$

we say c_1, \dots, c_M are generic if

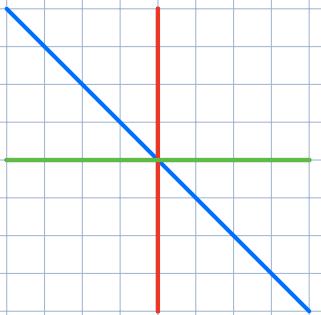
$\alpha_1, \dots, \alpha_M$ are linearly independent iff $H_{i_1} \cap \dots \cap H_{i_k} \neq \emptyset$.

This is a description of our " c_i 's, of which there are many choices for an arbitrary set of " α_i 's.

Recall that if $c_i \neq 0$, H_i is a translate of some linear hyperplane, so to fulfill this condition is to have translated some hyperplanes in a linear arrangement.

This allows us to speak of "generic arrangements"

An arrangement $\mathcal{A} = \{H_1, \dots, H_M\}$ with H_i as above is generic if c_1, \dots, c_M are generic.



not generic

↓ translation



generic

Dfn

For \mathcal{A} an arrangement of M hyperplanes given by

$$H_i = \{v \in V \mid \alpha_i \cdot v = c_i\},$$

we say c_1, \dots, c_M are generic if

$\alpha_1, \dots, \alpha_K$ are linearly independent iff $H_1 \cap \dots \cap H_K \neq \emptyset$.

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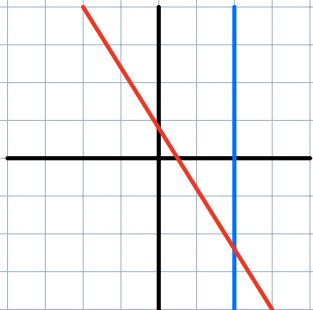
This allows us to speak of "generic arrangements"

An arrangement $\mathcal{A} = \{H_1, \dots, H_M\}$ with H_i as above is generic if c_1, \dots, c_M are generic.

Now some implications

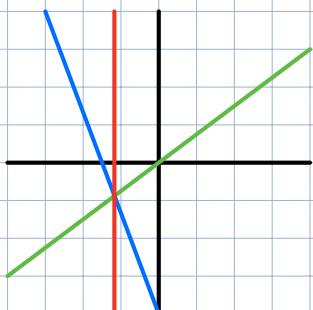
A is central and $m \leq n$

$\Rightarrow A$ is generic.



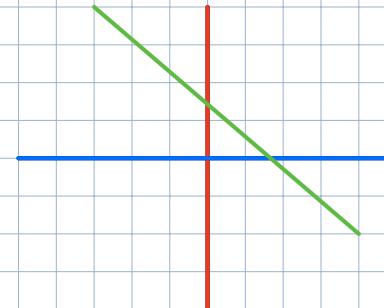
A is central and $m > n$

$\Rightarrow A$ is not generic.



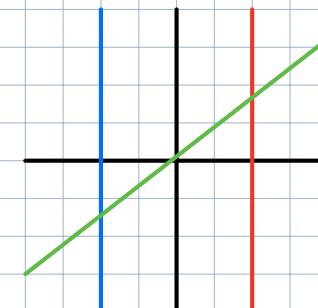
A is in gen. pos.

$\Rightarrow A$ is generic.



both

So generic is a weaker (but analogous) version of gen. pos.



not in gen.
pos.

In particular, generic $\Rightarrow \dim(H_1, n \dots, n H_K)$
is zero or $n-K$.

B2

Dh

For \mathcal{A} an arrangement of M hyperplanes given by

$$H_i = \{v \in V \mid \alpha_i \cdot v = c_i\},$$

we say c_1, \dots, c_M are generic if

$\alpha_1, \dots, \alpha_K$ are linearly independent iff $H_1 \cap \dots \cap H_K \neq \emptyset$.

This is a description of our " c_i 's, of which there are many choices for an arbitrary set of " α_i 's.

Recall that if $c_i \neq 0$, H_i is a translate of some linear hyperplane, so to fulfill this condition is to have translated some hyperplanes in a linear arrangement.

This allows us to speak of "generic arrangements"

An arrangement $\mathcal{A} = \{H_1, \dots, H_M\}$ with H_i as above is generic if c_1, \dots, c_M are generic.

Fact: If \mathcal{A} is generic, $[\hat{0}, x] \cong B_x$, the boolean algebra of rank $K = \dim(x)$.

$$\Rightarrow \chi_{\mathcal{A}}(t) = \sum_{B \text{ central}} (-1)^{|B|} t^{n-|B|}$$

b.c. $x = H_1 \cap \dots \cap H_K$ all nonparallel

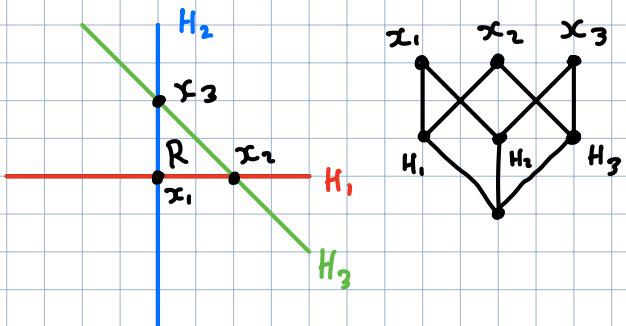
$\Rightarrow H_1, \dots, H_K$ in gen pos.

B1

Dfn.

The closure of a region $R \in \mathcal{R}(A)$, denoted \bar{R} , is the union of R and the portions of the hyperplanes forming its boundary.

A k -face of A is a set $F = x \cap \bar{R}$ for some $R \in \mathcal{R}(A)$ and $x \in L(A)$ s.t. $\dim(\text{span}(F)) = k$.



$$R = \square, \bar{R} = \triangle$$

$$R^2 \cap \bar{R} = \bar{R} \text{ a 2d face}$$

$$H_1 \cap \bar{R} = \text{---} \text{ a 1d face}$$

$$H_2 \cap \bar{R} = \dots \dots \dots$$

$$H_3 \cap \bar{R} = \curvearrowleft \dots \dots \dots$$

$$x_1 \cap \bar{R} = x_1 \text{ a 0d face}$$

$$x_2 \cap \bar{R} = x_2 \dots \dots \curvearrowright$$

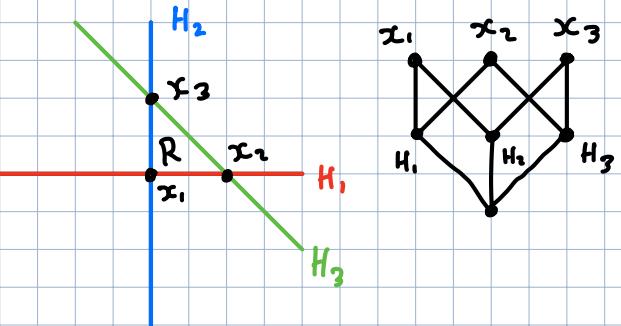
$$x_3 \cap \bar{R} = x_3 \dots \dots \curvearrowleft$$

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 Fact: Every k -face of A corresponds uniquely to a region in A^x for some $x \in L(A)$ satisfying $\dim(x) = k$.



$$R = \square, \quad \bar{R} = \triangle$$

$$R^2 \cap \bar{R} = \bar{R} \text{ a 2-face}$$

$$H_1 \cap \bar{R} = \square \text{ a 1-face}$$

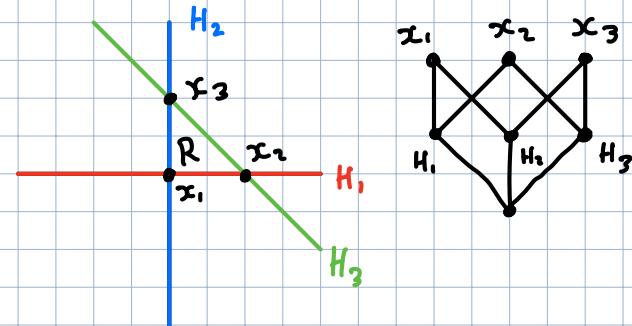
$$H_2 \cap \bar{R} = \square \dots \dots \dots$$

$$H_3 \cap \bar{R} = \square \dots \dots \dots$$

$$x_1 \cap \bar{R} = x_1 \text{ a 0-face}$$

$$x_2 \cap \bar{R} = x_2 \dots \dots \square$$

$$x_3 \cap \bar{R} = x_3 \dots \dots \square$$



Note that 2-faces are just regions in \mathbb{A}^R .

1-faces are regions in H_i

0-faces are regions in x_i

Dfn.

The closure of a region $R \in \mathcal{R}(A)$, denoted \bar{R} , is the union of R and the portions of the hyperplanes forming its boundary.

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Fact: Every k -face of A corresponds uniquely to a region in A^x for some $x \in L(A)$ satisfying $\dim(x) = k$.

Thm. 2.6.

Let $f_k(A) = \# k\text{-faces of } A$. Then

$$f_k(A) = \sum_{\dim(x)=k} \sum_{x \leq y}^{d(x) - d(y)} (-1)^{\mu(x,y)}.$$

Dfn.

The closure of a region $R \in \mathcal{R}(A)$, denoted \bar{R} , is the union of R and the portions of the hyperplanes forming its boundary.

A K -face of A is a set $F = x \cap \bar{R}$ for some $R \in \mathcal{R}(A)$ and $x \in L(A)$ s.t. $\dim(\text{span}(F)) = K$.

Fact: Every K -face of A corresponds uniquely to a region in A^x for some $x \in L(A)$ satisfying $\dim(x) = K$.

Thm. 2.6.

Let $f_K(A) = \# K\text{-faces of } A$. Then

$$f_K(A) = \sum_{\dim(x)=K} \sum_{x \leq y} (-1)^{\dim(x) - \dim(y)} \mu(x, y).$$

Pf. By our Fact,

$$\# K\text{-faces} = \sum_{\dim(x)=K} r(A^x) = \sum_{\dim(x)=K} \sum_{x \leq y} (-1)^{\dim(x) - \dim(y)} \mu(x, y). \quad \triangle$$

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Hyperplane Arrangements

Talk 12: Geometric Lattices and Matroids
Corresponding to Arrangements

Notation: I will write \bar{X} for the closure of $X \subseteq S$, i.e.

$$\bar{X} = \{y \in M \mid r_K(X \cup y) = r_K(X)\}$$

Writing $x \in M$ is equiv. to $x \in S$ for $M = (S, \mathcal{I})$.

Recall. A matroid $M = (S, \mathcal{I})$ is simple if it is loopless and has no parallel points.

No rank zero singletons No rank one 2-subsets

$$r_K(x) > 0 \quad \forall x \in M \quad r_K(\{x, y\}) \geq 1 \quad \forall \{x, y\} \in 2^S$$

Note. Setting $\widehat{M} = \{\bar{x} \mid x \in M, x \notin \emptyset\}$, we may construct a matroid \widehat{M} s.t.

$$\{\bar{x}_1, \dots, \bar{x}_n\} \in \mathcal{I}(\widehat{M}) \text{ iff } \{x_1, \dots, x_n\} \in \mathcal{I}(M).$$

So \widehat{M} is simple and $L(M) \cong L(\widehat{M})$

This means that, as far as lattices go, we only care about simple Matroids.

Q: Can we have \bar{X} contain enough ind. vs s.t. $r_K(\bar{X}) > r_K(X)$?

A: No. Supp. B a basis of X . Then $\forall x \in X \setminus X$ we have

$$r_K(B \cup x) = r_K(X \cup x) = r_K(X) = |B| = r_K(B) \leq r_K(B \cup x)$$

So now $r_K(B \cup x) = r_K(B) \Rightarrow B \cup x$ is dependent $\Rightarrow B$ a basis for \bar{X} .

$$M \text{ given by } \begin{bmatrix} 1 & 2 & 0 & 0 & 1 \\ 0 & 0 & 1 & 0 & 1 \\ 0 & 0 & 0 & 1 & 1 \end{bmatrix}$$

v_1, v_2, v_3, v_4, v_5

$$\mathcal{X}(M) = \{\emptyset, v_1, v_2, v_3, v_4, v_5, v_{13}, v_{14}, v_{15}, v_{23}, v_{24}, v_{25}, v_{34}, v_{35}, v_{45}, v_{134}, v_{135}, v_{145}, v_{234}, v_{235}, v_{345}\}$$

$$\varphi: v_1 \sim v_2 \mapsto v_1, v_3 \mapsto v_3, v_4 \mapsto v_4, v_5 \mapsto v_5$$

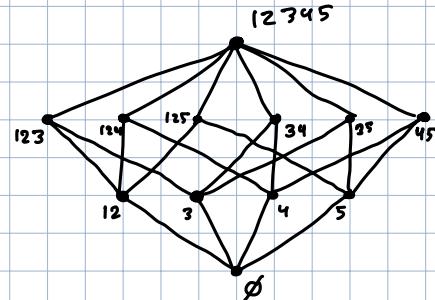
$$\Rightarrow \widehat{M} = \{v_1, v_3, v_4, v_5\} \text{ with}$$

$$\mathcal{X}(\widehat{M}) = \{\emptyset, v_1, v_3, v_4, v_5, v_{13}, v_{14}, v_{15}, v_{34}, v_{35}, v_{45}, v_{134}, v_{135}, v_{145}, v_{245}\}$$

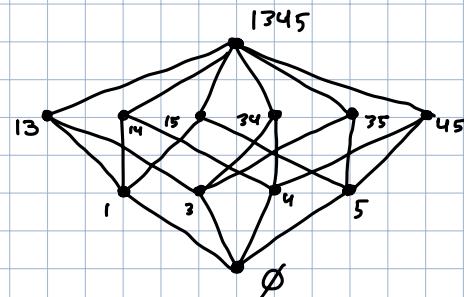
Note that $\widehat{M} \neq M$!

"I know that if $\text{rk}(\mathcal{X}_{13}) = 1$, x and y are always in flats together, so just move them one guy"

$L(M)$



$L(\widehat{M})$



Notation: I will write \bar{X} for the closure of $X \subseteq S$, i.e.

$$\bar{X} := \{ y \in M \mid r_K(X \cup y) = r_K(X) \}$$

Recall. A matroid $M = (S, \mathcal{I})$ is simple if it is loopless and has no parallel points.

No rank zero singletons No rank one 2-subsets

$$r_K(x) > 0 \quad \forall x \in M \quad r_K(\{x, y\}) > 1 \quad \forall \{x, y\} \in 2^S$$

Note. Setting $\hat{M} = \{\bar{x} \mid x \in M, r_K(x) > 0\}$, we may construct an isomorphism

$$\{\bar{x}_1, \dots, \bar{x}_n\} \in \mathcal{I}(\hat{M}) \text{ iff } \{x_1, \dots, x_n\} \in \mathcal{I}(M).$$

So \hat{M} is simple and $L(M) \cong L(\hat{M})$

This means that, as far as lattices go, we only care about simple matroids.

Fact. M being simple $\Rightarrow L(M)$ uniquely determines it.

How? Identify $(x_1, \dots, x_n) \in \mathcal{I} \iff r_K(x_1 \vee \dots \vee x_n) = K$ in $L(M)$.

Simple \Rightarrow rank in $L(M)$ works exactly like rank in M .

Prop. 3.6.

Maybe clarify this?

If A is central, then the linear Matroid whose ground set is given by the set of normals in A is simple and satisfies

$$L(M) \cong L(A).$$

p.f. (sketch) For simplicity note that A central \Rightarrow no parallel hyperplanes, and no loops b.c. nonzero normals.

For lattice isomorphism, we have a bijection given by

$$X = \bigcap_{H \in \mathcal{B}} H \longmapsto \left\{ \alpha_H \mid \underbrace{\alpha_H \in \text{span} \left\{ \alpha_H \mid H \in \mathcal{B} \right\}}_{\text{normal is in span of normals of } \mathcal{B}} \right\} = \overline{\mathcal{B}}$$

and you may check that this preserves order.

M given by

$$\begin{bmatrix} 1 & 0 & 0 & 1 \\ 0 & 1 & 0 & 1 \\ 0 & 0 & 1 & 1 \end{bmatrix}$$

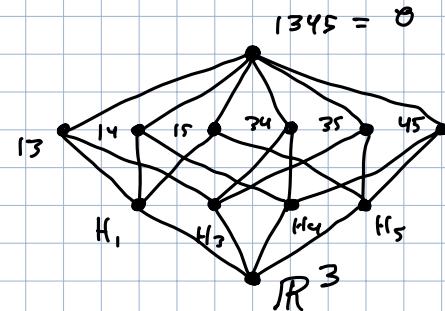
$v_1 \ v_2 \ v_3 \ v_4 \ v_5$

$$A = \left\{ H_i \mid H_i = \left\{ v \in \mathbb{R}^3 \mid v_i \cdot v = 0 \right\} \right\}$$

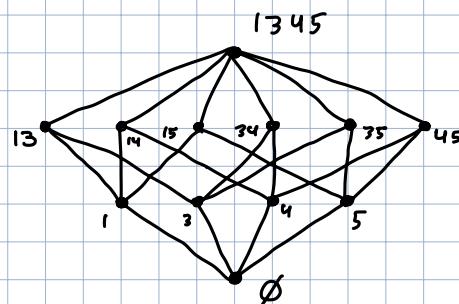
$$\Rightarrow M = \{v_1, v_3, v_4, v_5\} \text{ with}$$

$$\chi(M) = \{\emptyset, v_1, v_3, v_4, v_5, v_{13}, v_{14}, v_{15}, v_{34}, v_{35}, v_{45}, v_{134}, v_{135}, v_{145}, v_{245}\}$$

$L(A)$



$L(M)$



Prop. 3.6.

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For lattice isomorphism, we have a bijection given by

$$X = \bigcap_{H \in B} H \longmapsto \left\{ \alpha_H \mid \alpha_H \in \text{span} \left\{ \alpha_H \mid H \in B \right\} \right\}$$

normal is in span
of normals of B

and you may check that this preserves order.

It may be worth noting that int sets are int subsets of normals, and what is here are flats.

DEFN.

A finite lattice L satisfying

$$rk(x) + rk(y) \geq rk(x \vee y) + rk(x \wedge y)$$

is said to be semimodular. A finite lattice L is atomic if every $x \in L$ is the join of some set of atoms.

A finite lattice that is both semimodular and atomic is geometric.

Thm. 3.8.

For L a finite lattice the following are equivalent,

- 1) L is geometric
- 2) $L \cong L(M)$ for some (unique, simple) matroid M .

p.f. (sketch) \Rightarrow Let A be the set of atoms and

$$\mathcal{I} = \{ T \subseteq A \mid r_K(V_T) = |T| \}.$$

Semimodularity of L yields inclusion property b.c. for any $T \subseteq A$ and $x \in A$,

$$r_K((V_T) \vee x) \leq r_K(V_T) + 1. \quad \text{i.e.: if } T \cup x \in \mathcal{I} \Rightarrow r_K((V_T) \vee x) = |T| + 1 \\ \Rightarrow r_K(V_T) = |T| \text{ b.c. } (V_T) \vee x \text{ covers } V_T \text{ by P3.7.}$$

For our property concerning sizes of maximal ind^t sets, an argument by contradiction follows b.c. L is atomic and so

\Leftarrow) Atomic because each flat is the join of some singletons (atoms).

Semimodularity follows for flats straightforwardly and

Now we have that results for simple matroids and geometric lattices can be applied to our study of central arrangements.

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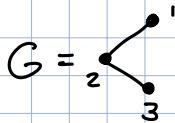
Hyperplane Arrangements

Capstone Part IV.:

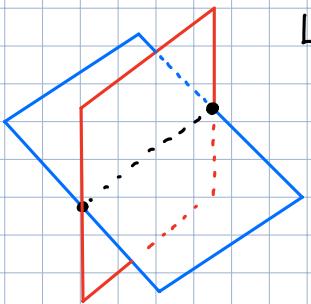
Graphic Arrangements

Graphic Arrangements: For G a graph with $V = \{1, \dots, n\}$ we define

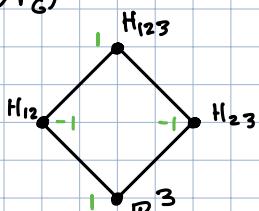
$$A_G = \{H_{ij} \mid ij \in E\}, \quad H_{ij} = \{(v_1, \dots, v_n) \in \mathbb{R}^n \mid v_i = v_j\}.$$



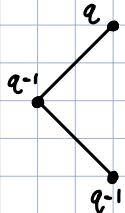
$$\begin{aligned} H_{12} : x_1 &= x_2 \\ H_{23} : x_2 &= x_3 \end{aligned}$$



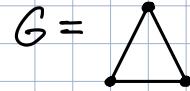
$L(A_G)$:



$$\chi_{A_G}(t) = t^3 - 2t^2 + t$$

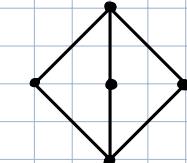


$$\begin{aligned} \chi_G(q) &= q(q-1)^2 \\ &= q^3 - 2q^2 + q \end{aligned}$$



$$A_G \approx \begin{array}{c} \text{blue line} \\[-1ex] \text{red line} \\[-1ex] \text{green line} \end{array} = A_2$$

$L(A_2)$



Now recall:

$$\chi_{A_2}(t) = t^2 - 3t + 2$$

but we claim that

A_2 is just A_G projected onto a plane

$$\Rightarrow \chi_{A_G}(t) = t^3 - 3t^2 + 2t$$

$$\chi_G(t) = t^3 - 3t^2 + 2t.$$

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$$\mathcal{A}_G = \{H_{ij} \mid ij \in E\}, \quad H_{ij} = \{(v_1, \dots, v_n) \in \mathbb{R}^n \mid v_i = v_j\}.$$

Punchline: $\chi_G(t) = \chi_{\mathcal{A}_G}(t)$

for any graph G .