

# Correction frames

What, how & why?

GREG RESTALL



University of  
St Andrews

NOVEMBER 2022

# COLLECTION FRAMES

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This talk reports  
joint work with  
Shawn Standefer



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but he is not responsible  
for any of the controversial  
things I say here

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# MY PLAN

1. WHAT ARE MODELS FOR?
2. THE COMPLEXITY OF ROUTLEY/MEYER FRAMES
3. THREE INTO ONE
4. BEYOND REFLEXIVITY
5. EXAMPLES
6. OPEN QUESTIONS

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MODELS

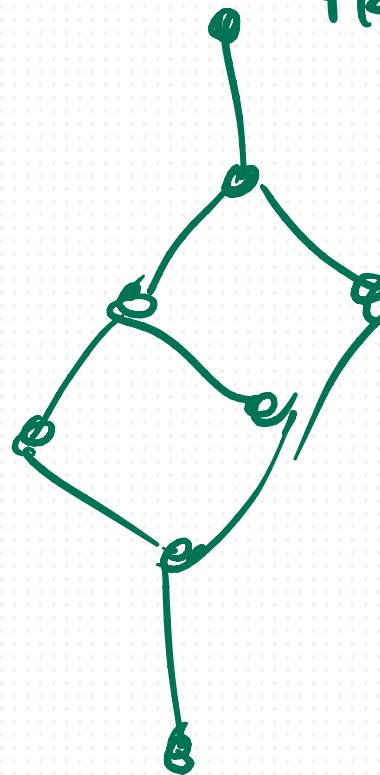
LANGUAGE  $\Rightarrow$  THINGS

LANGUAGE

MODELS

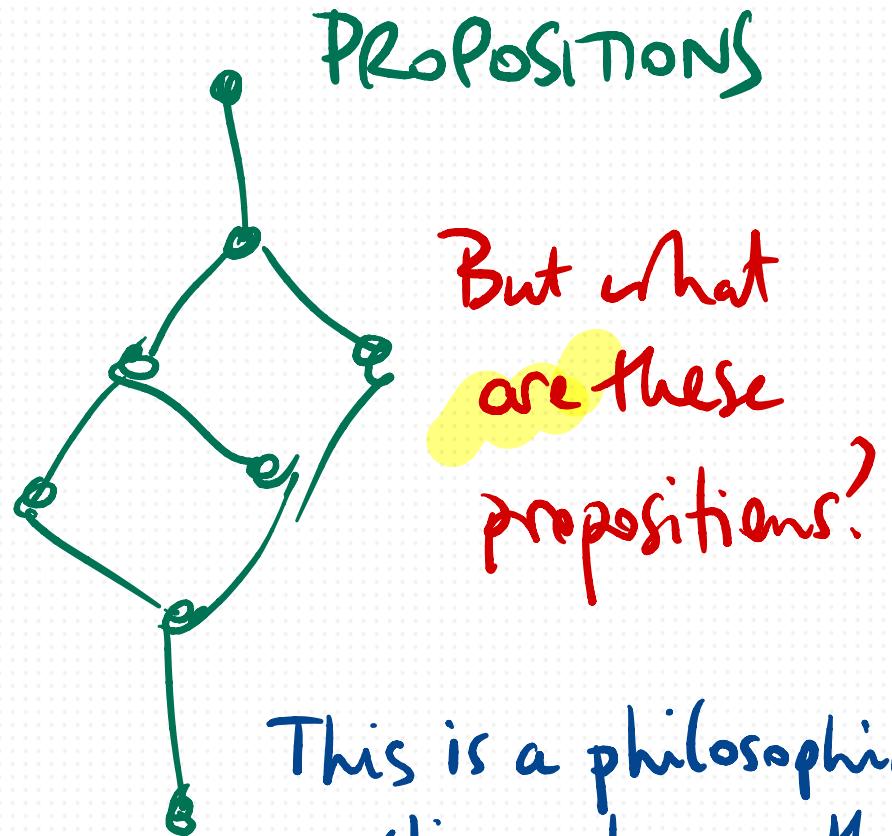


PROPOSITIONS



# MODELS

LANGUAGE  $\Rightarrow$



PROPOSITIONS

But what  
are these  
propositions?

This is a philosophical  
question & a mathematical  
one, too...

# MODELS

LANGUAGE  $\Rightarrow$

MY FOCUS TODAY

PROPOSITIONS

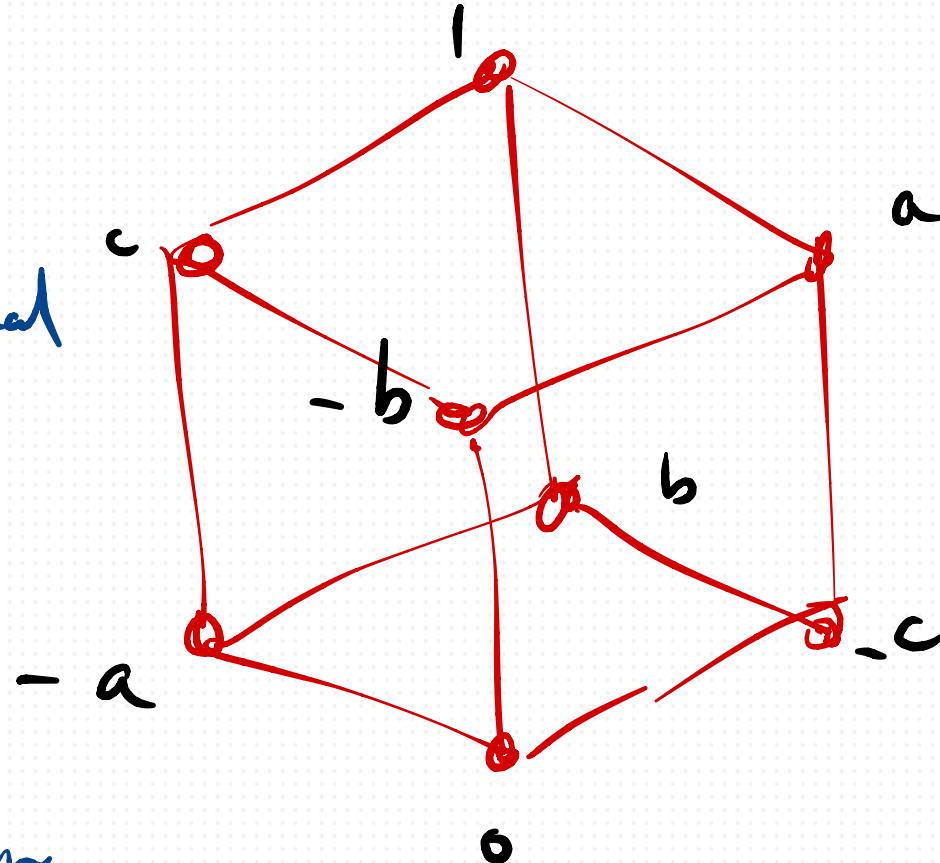
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# An EASY EXAMPLE

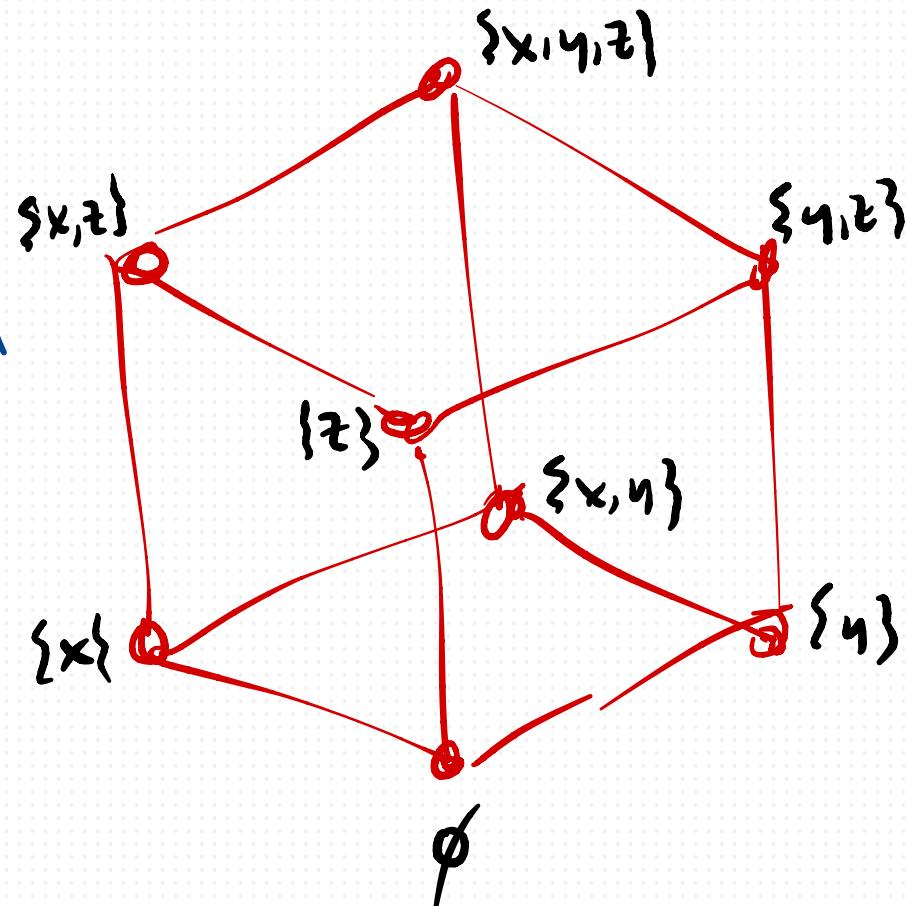
A Boolean  
Algebra  
is a propositional  
Structure  
Satisfying  
a simple  
set of  
equations on

$$\wedge \vee - = 0 1$$



# AN EASY EXAMPLE

A Boolean Algebra is a propositional structure satisfying a simple set of equations on  $\wedge \vee = 0 1$

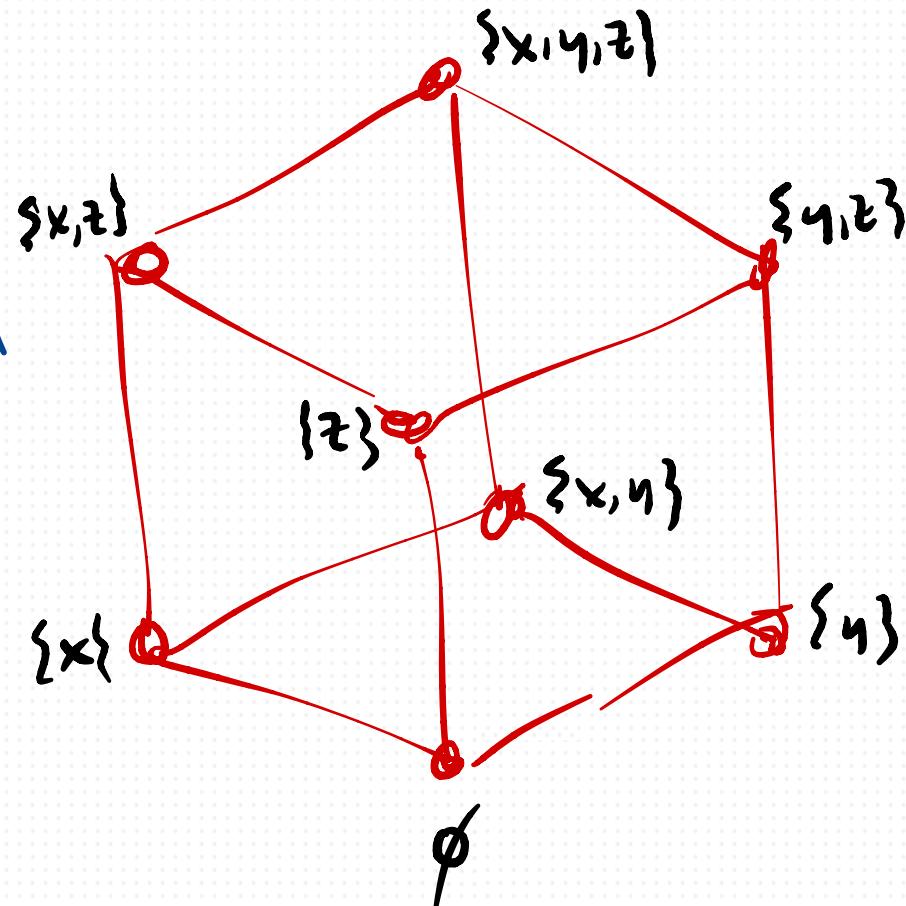


Every Boolean Algebra is isomorphic to a subalgebra of some powerset algebra on a set  $P$

$$\cap \cup - \phi P$$

# AN EASY EXAMPLE

A Boolean Algebra is a propositional structure satisfying a simple set of equations on  $\wedge \vee - \circ \perp$



Every Boolean Algebra is isomorphic to a subalgebra of some powerset algebra on a set  $P$

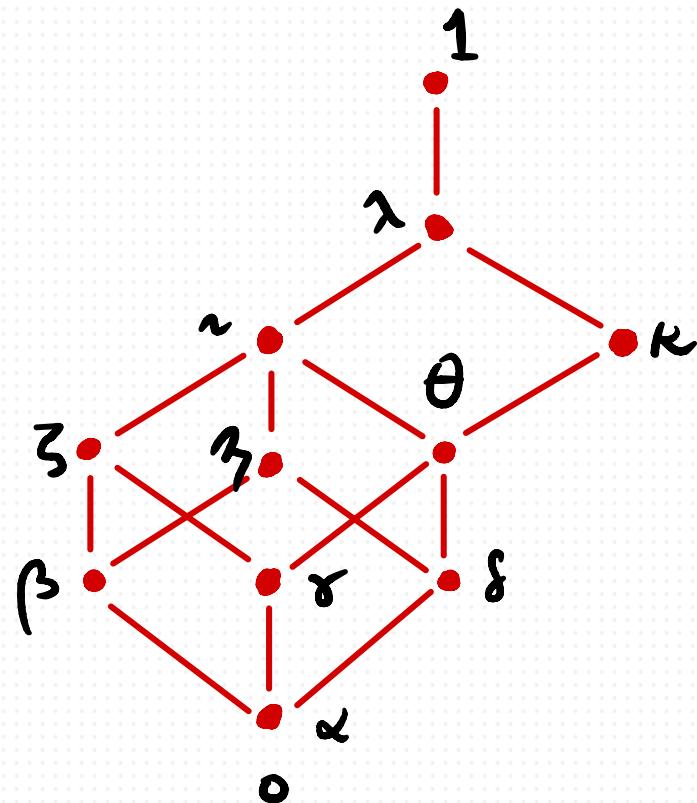
$$\wedge \vee - \circ \perp$$

$$\wedge \vee - \circ \perp$$

THIS IS PHILOSOPHICALLY & MATHEMATICALLY SUGGESTIVE

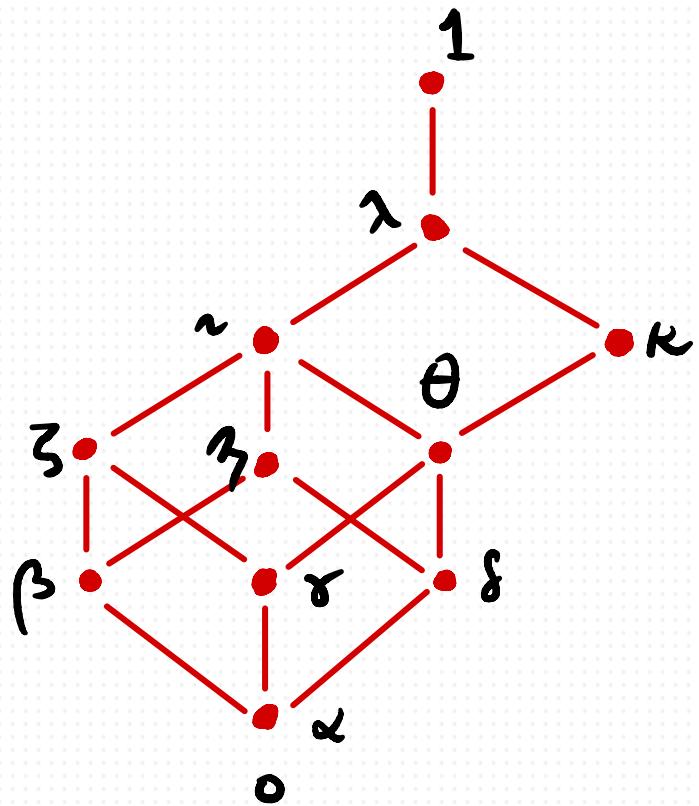
Slightly less  
A  $\wedge$  EASY EXAMPLE

A Heyting Algebra  
is a propositional  
Structure  
Satisfying  
a simple  
set of  
equations on  
 $\wedge \vee \rightarrow 0 1$



A Heyting Algebra is a propositional structure satisfying a simple set of equations on  $\wedge \vee \rightarrow 0 1$

# <sup>Slightly less</sup> A <sub>1</sub> EASY EXAMPLE



Every Heyting Algebra is isomorphic to a subalgebra of open sets on a topological space  $T$

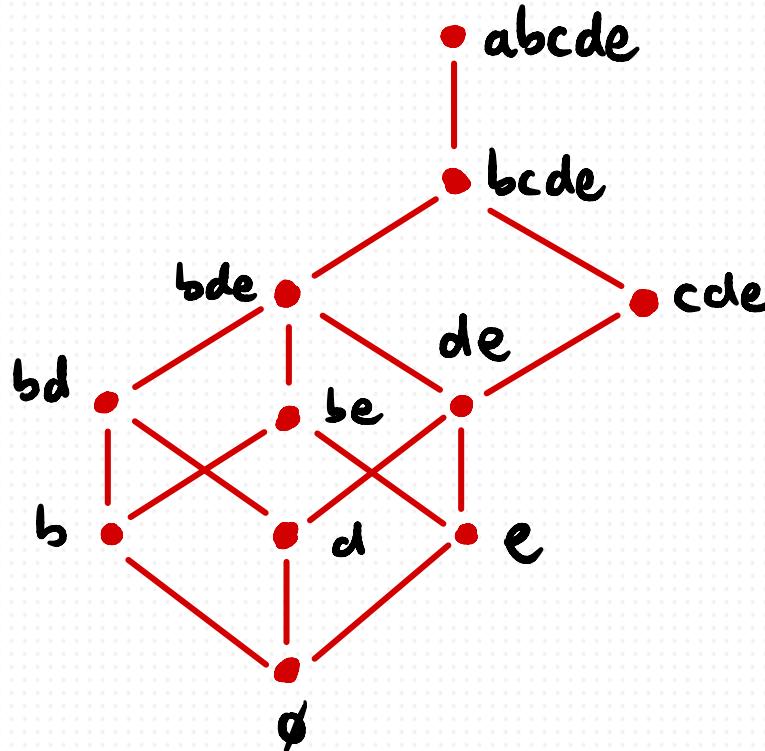
$$\cap \cup \Rightarrow \emptyset T$$

A Heyting Algebra is a propositional structure satisfying a simple set of equations on  $\wedge \vee \rightarrow 0 1$

*also*  
THIS IS PHILOSOPHICALLY & MATHEMATICALLY SUGGESTIVE

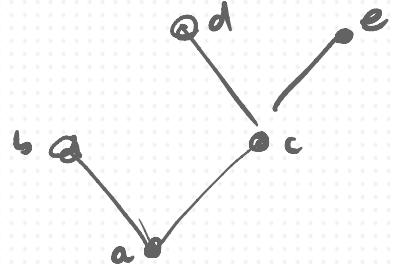
# A <sub>1</sub> EASY EXAMPLE

Slightly less



Every Heyting Algebra is *also* isomorphic to a subalgebra of upward closed sets on a partially ordered set of points P

$$\cap \cup \Rightarrow \emptyset P$$



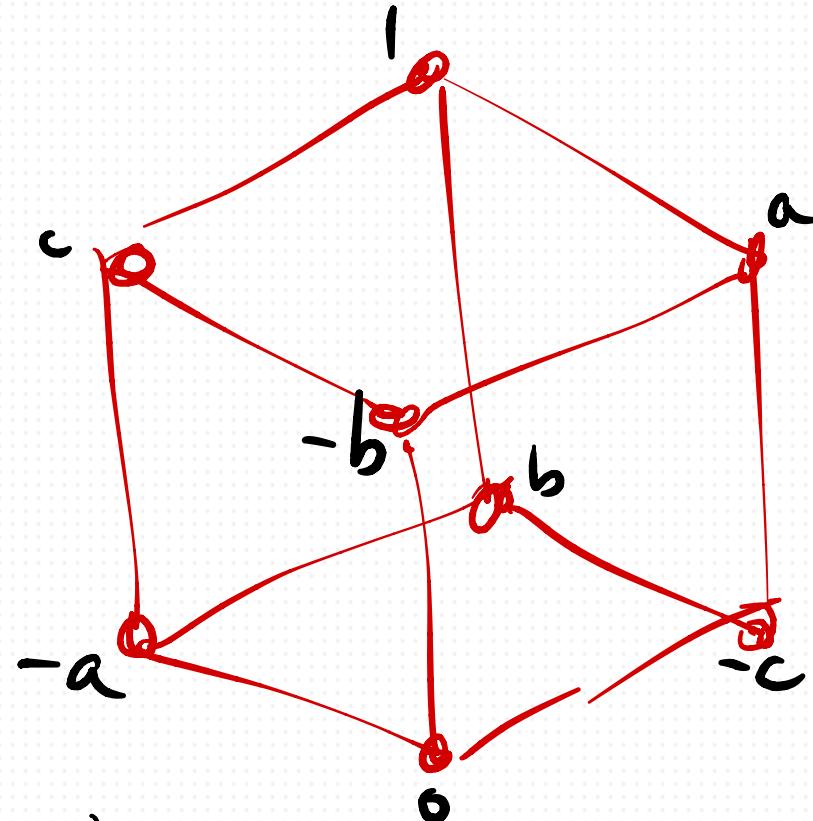
# YET ANOTHER Example

An S4 -  
 Algebra  
 is a Boolean  
 Algebra with  
 an operator  $\square$   
 satisfying

$$\square 1 = 1$$

$$\square(x \wedge \square y) = \square(x \wedge y)$$

$$\square x \leq x; \quad \square \square x = \square x$$



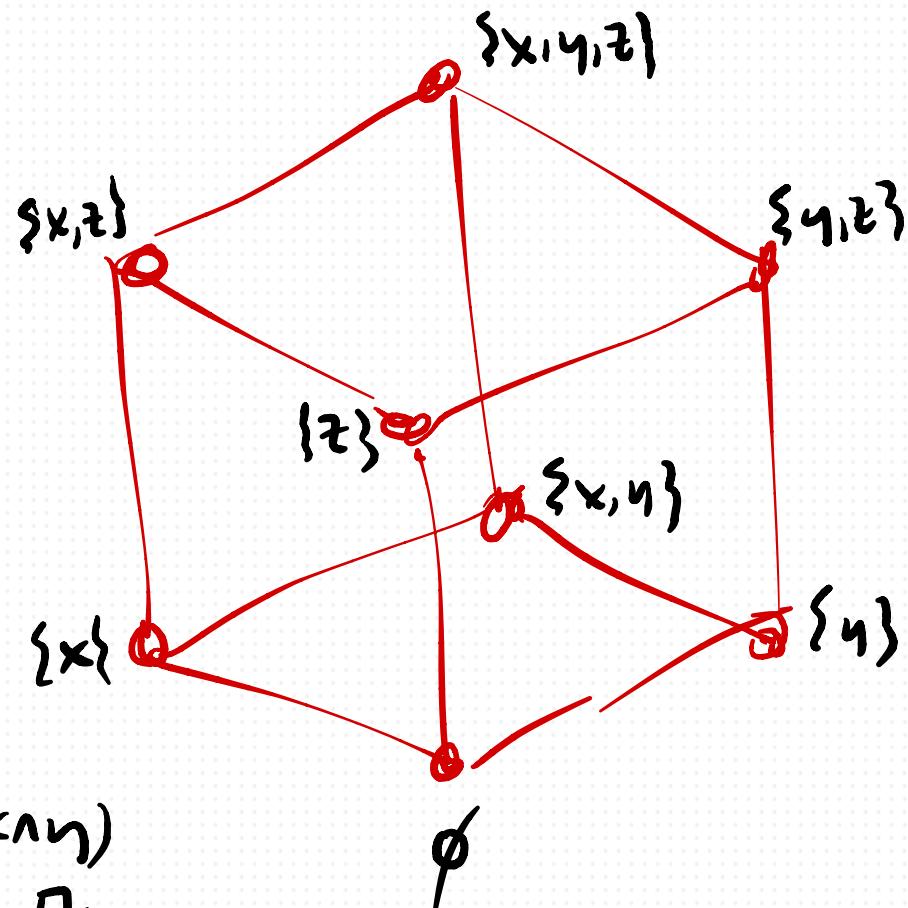
$\alpha$	$\emptyset$	$-a$	$-b$	$-c$	$b$	$c$	$a$	$1$
$\square \alpha$	$\emptyset$	$\emptyset$	$-b$	$-c$	$-c$	$-b$	$a$	$1$

An S4 -  
Algebra  
is a Boolean  
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$$\square 1 = 1$$

$$\square x \wedge \square y = \square(x \wedge y)$$

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$\alpha$	$\phi$	$x$	$y$	$z$	$x \wedge y$	$x \wedge z$	$y \wedge z$	$x \wedge y \wedge z$
$\square \alpha$	$\phi$	$\phi$	$y$	$z$	$y$	$z$	$y \wedge z$	$x \wedge y \wedge z$

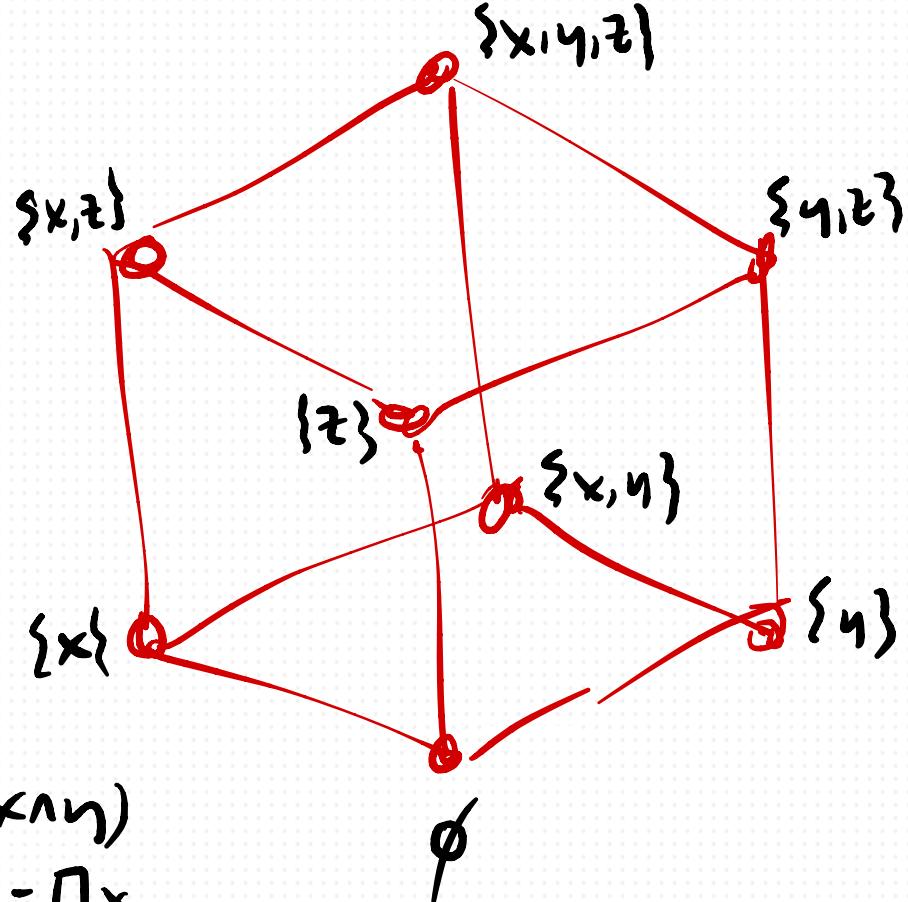
Every S4 algebra  
is isomorphic to  
a subalgebra of  
a powerset algebra  
equipped with a  
topology for which  
 $\square$  is an interior  
operator.

An S4 -  
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is a Boolean  
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$$\square 1 = 1$$

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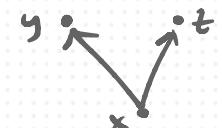
$$\square x \leq x; \quad \square \square x = \square x$$



$\alpha$	$\phi$	$x$	$y$	$z$	$x \vee y$	$x \vee z$	$y \vee z$	$x \vee y \vee z$
$\square \alpha$	$\phi$	$\phi$	$\phi$	$y$	$z$	$y$	$z$	$y \vee z$

Also  
THIS IS PHILOSOPHICALLY & MATHEMATICALLY SUGGESTIVE

Every S4 algebra  
is isomorphic  
to a subalgebra of  
a powerset algebra  
equipped with a  
topology generated  
by a partial order  
on the ground set,  
where the opens are  
up-closed sets.



# MY PLAN

1. WHAT ARE MODELS FOR?
2. THE COMPLEXITY of Routley/MAYER FRAMES
3. THREE INTO ONE
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# RELEVANT ALGEBRAS

An **RW<sup>+</sup> algebra** is a **DISTRIBUTIVE LATTICE**  $\langle L; \wedge, \vee \rangle$  equipped with an associative & commutative operator  $\circ$ , residuated by  $\rightarrow$  (so  $x \circ y \leq z$  if  $x \leq y \rightarrow z$ ) & with an identity,  $e$  (s.  $e \circ x = x$ ).

## RELEVANT ALGEBRAS

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These generalise Heyting Algebras, by distinguishing  $\circ$  &  $\wedge$ .

You'd expect to be able to generalise the representation theorem for Heyting Algebras to represent these ...

# RELEVANT FRAMES

An  $RW^+$  frame is a 4-tuple  $\langle P, \leq, R, N \rangle$  consisting of a non-empty set  $P$  equipped with

- A partial order  $\leq$  on  $P$
- A 3-place relation  $R$  on  $P$
- A subset  $N$  of  $P$

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where •  $R \downarrow \downarrow \uparrow$  •  $N \uparrow$

- $\exists x(N_x \notin Rxyz) \Leftrightarrow y \leq z$

# RELEVANT FRAMES

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Propositions are sets of points closed upwards along  $\leq$ .

where •  $R \downarrow \downarrow \uparrow$  •  $N \uparrow$

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# RELEVANT FRAMES

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Propositions are sets of points closed upwards along  $\sqsubseteq$ .

where

- $R \downarrow \downarrow \uparrow$
- $N \uparrow$
- $\exists x(N_x \notin Rxyz) \Leftrightarrow y \sqsubseteq z$

- $Rxyz \Leftrightarrow Ryxz \quad \swarrow R(yz)xv$
- $\exists z(Rxyz \nmid Rzuv) \Leftrightarrow \exists w(Ryzw \nmid Rxwv) \quad \swarrow Rx(yz)v$

# RELEVANT FRAMES

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Propositions are sets of points closed upwards along  $\sqsubseteq$ .

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- $\exists z(Rxyz \notin Rzuv) \Leftrightarrow \exists w(Ryzw \notin Rxwv) \quad \swarrow Rx(yz)v$

NOT SO

THIS IS PHILOSOPHICALLY OR MATHEMATICALLY SUGGESTIVE

# RELEVANT FRAMES

COMES WITH THE TERRITORY,  
Since every Heyting lattice  
is an  $\text{RW}^+$  algebra

- A partial order  $\leq$
- A 3-place relation  $R$  ← Used for  $\circ, \rightarrow$
- A subset  $N$  ← Used for  $e.$

where

- $R \downarrow \downarrow \uparrow$
- $N \uparrow$
- $\exists x(N_x \nmid R_{xyz}) \Leftrightarrow y \leq z$

- $R_{xyz} \Leftrightarrow R_{yxz}$
- $\exists z(R_{xyz} \nmid R_{zuv}) \Leftrightarrow \exists w(R_{yzw} \nmid R_{xwv})$

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- $\exists z(R_{xyz} \wedge R_{zuv}) \Leftrightarrow$   
 $\exists w(R_{yw} \wedge R_{xwv})$

These make  $N, \leq, R$   
play nicely together.

# RELEVANT FRAMES

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- $R_{xyz} \Leftrightarrow R_{yxz}$
- $\exists z(R_{xyz} \nmid R_{zuv}) \Leftrightarrow$   
 $\exists w(R_{yzw} \nmid R_{xwv})$

Required for commutativity  
& associativity of  $\circ.$

# UNDERSTANDING RELEVANT FRAMES

- A partial order  $\leq$
- A 3-place relation  $R$
- A subset  $N$

where •  $R \downarrow \downarrow \uparrow$    •  $N \uparrow$

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# UNDERSTANDING RELEVANT FRAMES

- A partial order  $R^2$
- A 3-place relation  $R^3$
- A subset  $R^1$

where

- $R^3 \downarrow \downarrow \uparrow$
- $R^1 \uparrow$
- $\exists x (R^1_x \notin R^3_{xyz}) \Leftrightarrow R^2_{yz}$

Hmm...

- $R_{xyz} \Leftrightarrow R_{yxz}$
- $\exists z (R_{xyz} \wedge R_{zuv}) \Leftrightarrow \exists w (R_{yzw} \wedge R_{xwv})$

# UNDERSTANDING RELEVANT FRAMES

- A partial order  $R^2$
- A 3-place relation  $R^3$
- A subset  $R^1$

We also have  $R^2 \downarrow \uparrow$ ,  
of course, since  
 $R^2$  is the partial  
order itself!

where

- $R^3 \downarrow \downarrow \uparrow$
- $R^2 \uparrow$
- $\exists x (R^2_x \notin R^3_{xyz}) \Leftrightarrow R^2_{yz}$

Hmm...

# UNDERSTANDING RELEVANT FRAMES

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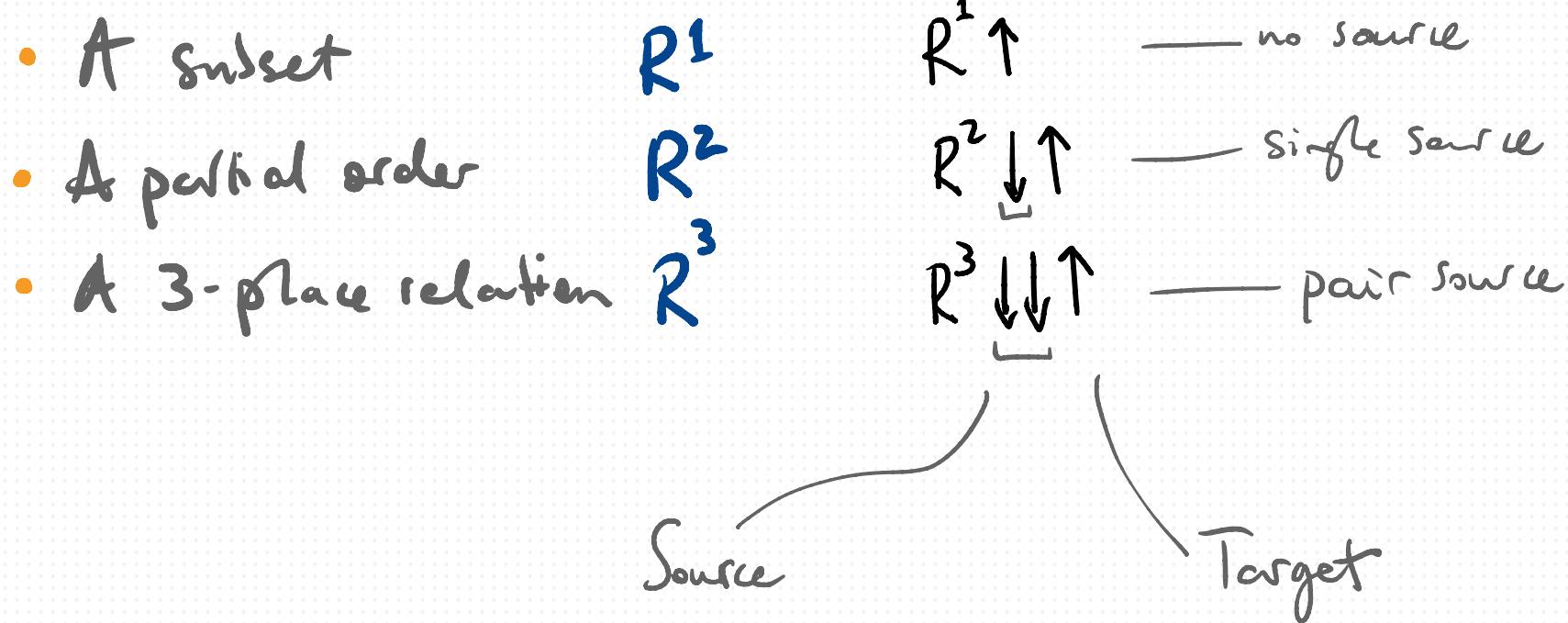
where

- $R^3 \downarrow \downarrow \uparrow$
- $R^2 \uparrow$
- $\exists x (R^2 x \notin R^3 xyz) \Leftrightarrow R^2 yz$

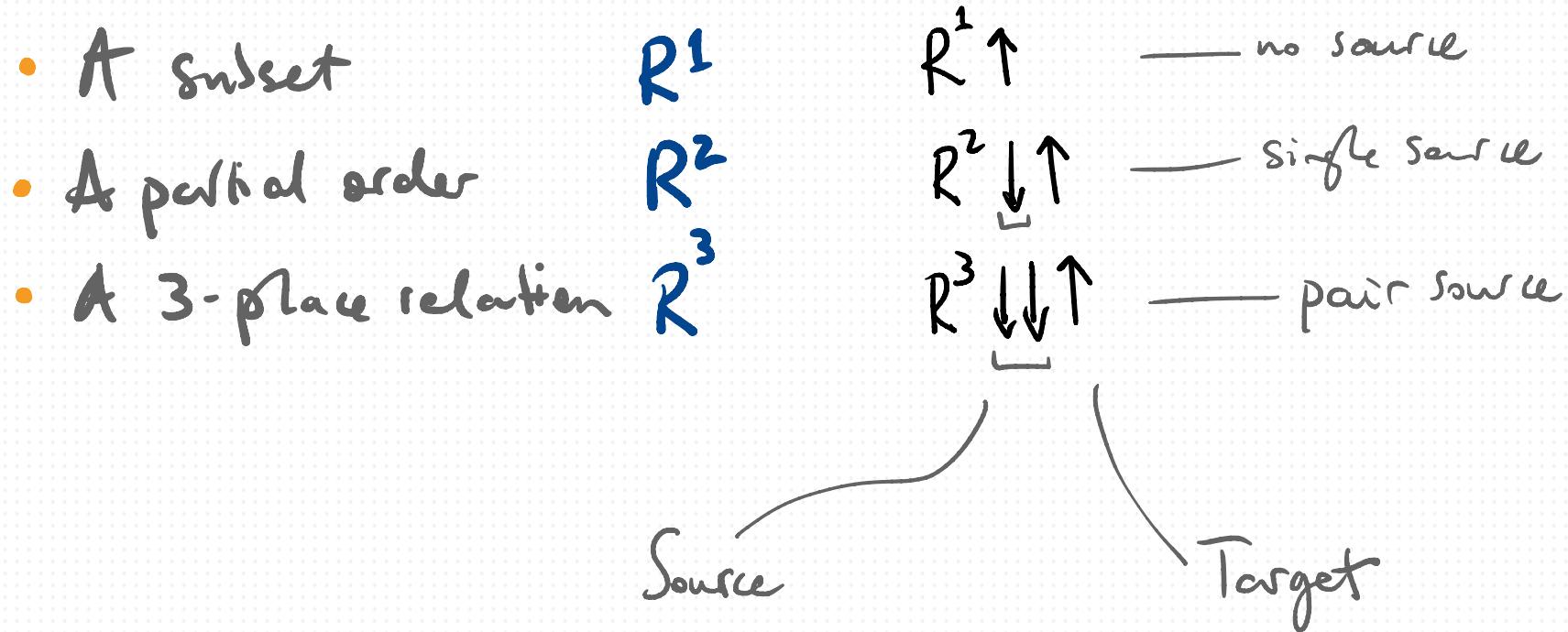
HMM...  
This generalises, too!

e.g.  $\exists x (R^2 x' x \notin R^3 xyz)$   
 $\Leftrightarrow R^3 x' yz$

- A subset  $R^1$
  - A partial order  $R^2$
  - A 3-place relation  $R^3$
- $R^1 \uparrow$  — no source
- $R^2 \downarrow \uparrow$  — single source
- $R^3 \downarrow \downarrow \uparrow$  — pair source
- 
- Source
- Target

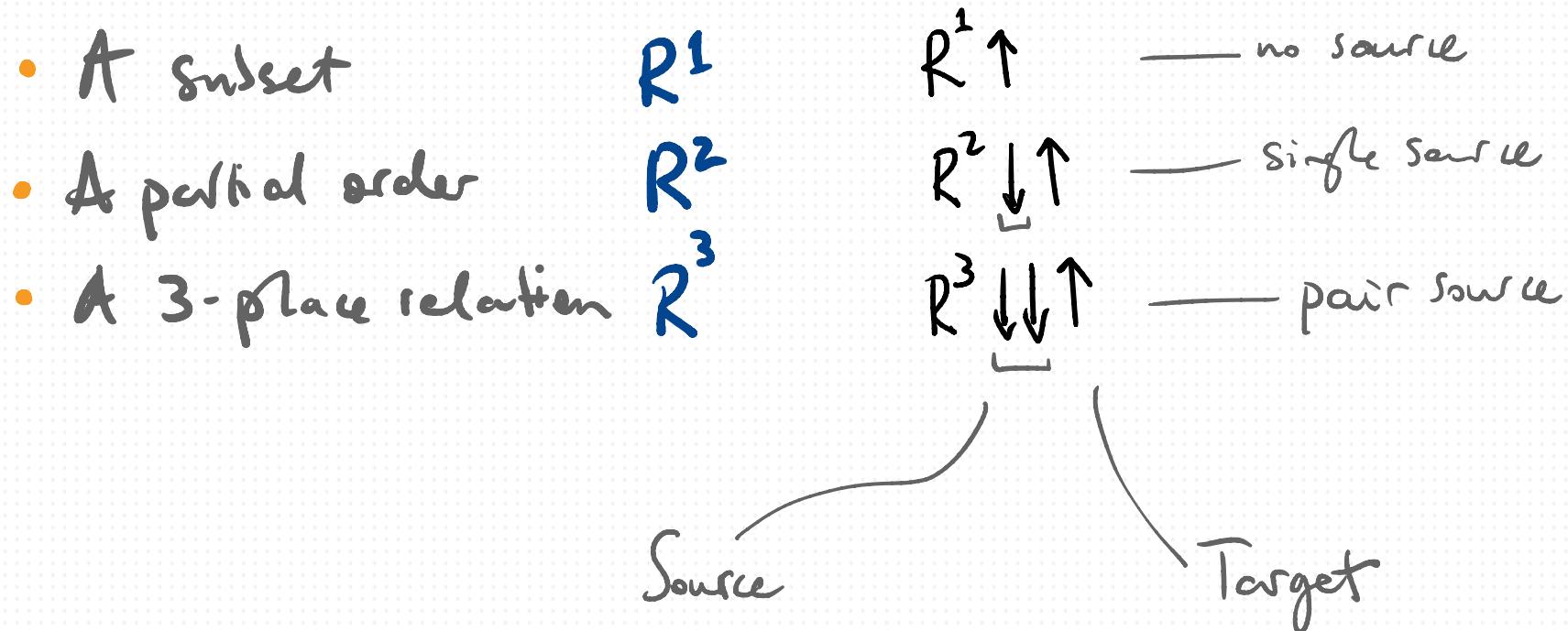


How is the source structured?



How is the source structured?

- Recall:
- $R^3_{xyz} \leftrightarrow R^3_{yxz}$  ————— the order does not matter
  - $\exists z (R^3_{xyz} \wedge R^3_{zuv}) \leftrightarrow \exists w (R^3_{yzw} \wedge R^3_{xwv})$
- └ This has something to do with associativity



How is the Source structured?

Recall:

- $R^3xyz \Leftrightarrow R^3yxz$  ————— the order does not matter
- $\exists z(R^3xyz \wedge R^3zuv) \Leftrightarrow \exists w(R^3yzw \wedge R^3xvw)$

└ This has something to do with associativity

The Source is just one thing: a multiset of points !!

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## MULTISet RELATIONS

$R$  is a multiset relation on  $P$  iff  $R \subseteq M(P) \times P$ .

That is, it relates finite multisets of points to points.

- A subset  $R^1 z$  if  $R[]z$
- A partial order  $R^2 xz$  if  $R[x]z$
- A 3-place relation  $R^3 xyz$  if  $R[x,y]z$

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- A 3-place relation  $R^3 xyz$  iff  $R[x,y]z$

The commutativity postulate  $R^3 xyz \leftrightarrow R^3 yxz$  comes for free.  
for the other frame postulates, we need some kind of condition  
on  $R$  analogous to the requirement for a partial order  
in Heyting frames.

- $R$  is reflexive iff  $[x]Rx$  for every  $x$ .

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Recall this condition:  $\exists x (R^2 x' x \wedge R^3 x u z) \Leftrightarrow R^3 x' u z$

- $R$  is reflexive iff  $[x]Rx$  for every  $x$ .

Recall this condition:  $\exists x (R^2 x' x \notin R^3 x y z) \Leftrightarrow R^3 x' y z$

i.e.:  $\exists x (R[x']x \notin R[x,y]z) \Leftrightarrow R[x',y]z$

- $R$  is reflexive iff  $[x]Rx$  for every  $x$ .

Recall this condition:  $\exists x (R^2x'x \notin R^3xuz) \Leftrightarrow R^3x'uz$

i.e.:  $\exists x (R[x']x \notin R[x,y]z) \Leftrightarrow R[x',y]z$

generalising over  $y$ :  $\exists x (R[x']x \notin R([x] \cup Y)z) \Leftrightarrow R([x'] \cup Y)z$

- $R$  is reflexive iff  $[x]Rx$  for every  $x$ .

Recall this condition:  $\exists x (R^2 x' x \notin R^3 x y z) \Leftrightarrow R^3 x' y z$

i.e.:  $\exists x (R[x']x \notin R[x,y]z) \Leftrightarrow R[x',y]z$

generalising over  $y$ :  $\exists x (R[x']x \notin R([x] \cup Y)z) \Leftrightarrow R([x'] \cup Y)z$

generalising over  $x'$ :  $\exists x (Rx \notin R([x] \cup Y)z) \Leftrightarrow R(X \cup Y)z$

- $R$  is reflexive iff  $[x]Rx$  for every  $x$ .

Recall this condition:  $\exists x (R^2 x' x \notin R^3 x y z) \Leftrightarrow R^3 x' y z$

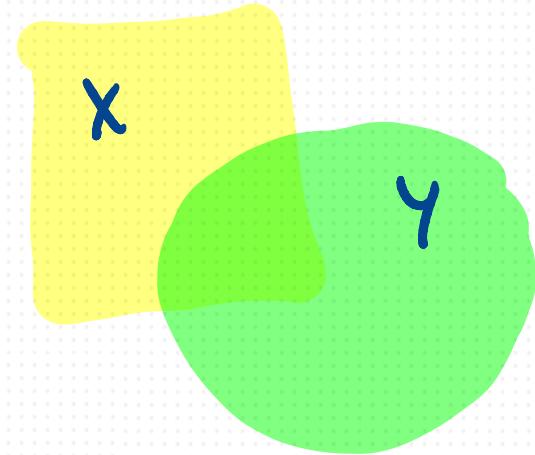
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generalising over  $y$ :  $\exists x (R[x']x \notin R([x] \cup Y)z) \Leftrightarrow R([x'] \cup Y)z$

generalising over  $x'$ :  $\exists x (R X x \notin R([x] \cup Y)z) \Leftrightarrow R(X \cup Y)z$

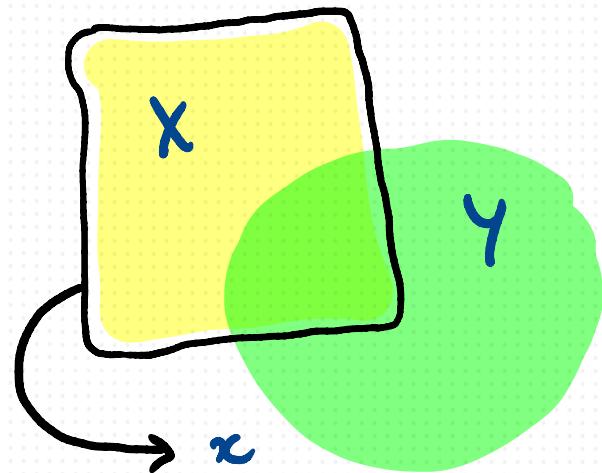
- $R$  is compositional iff  $\exists x (R X x \notin R([x] \cup Y)z) \Leftrightarrow R(X \cup Y)z$

# COMPOSITIONALITY



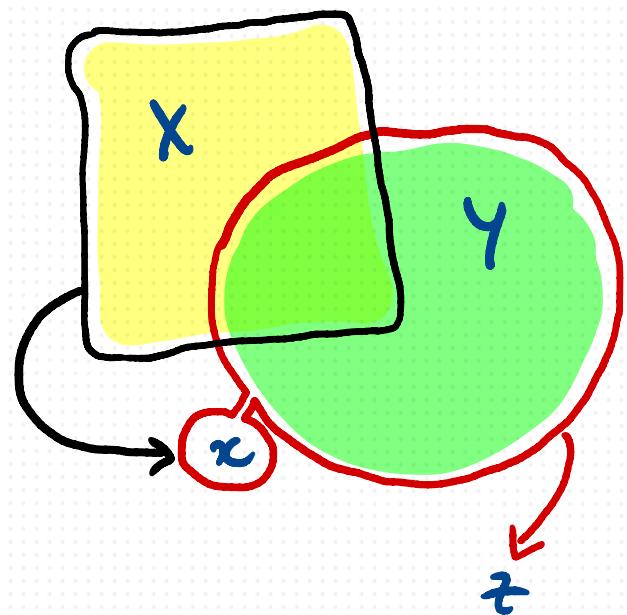
$$\exists x (R(x \in X) \wedge R(x \in Y)) \Leftrightarrow R(X \cup Y)$$

# COMPOSITIONALITY



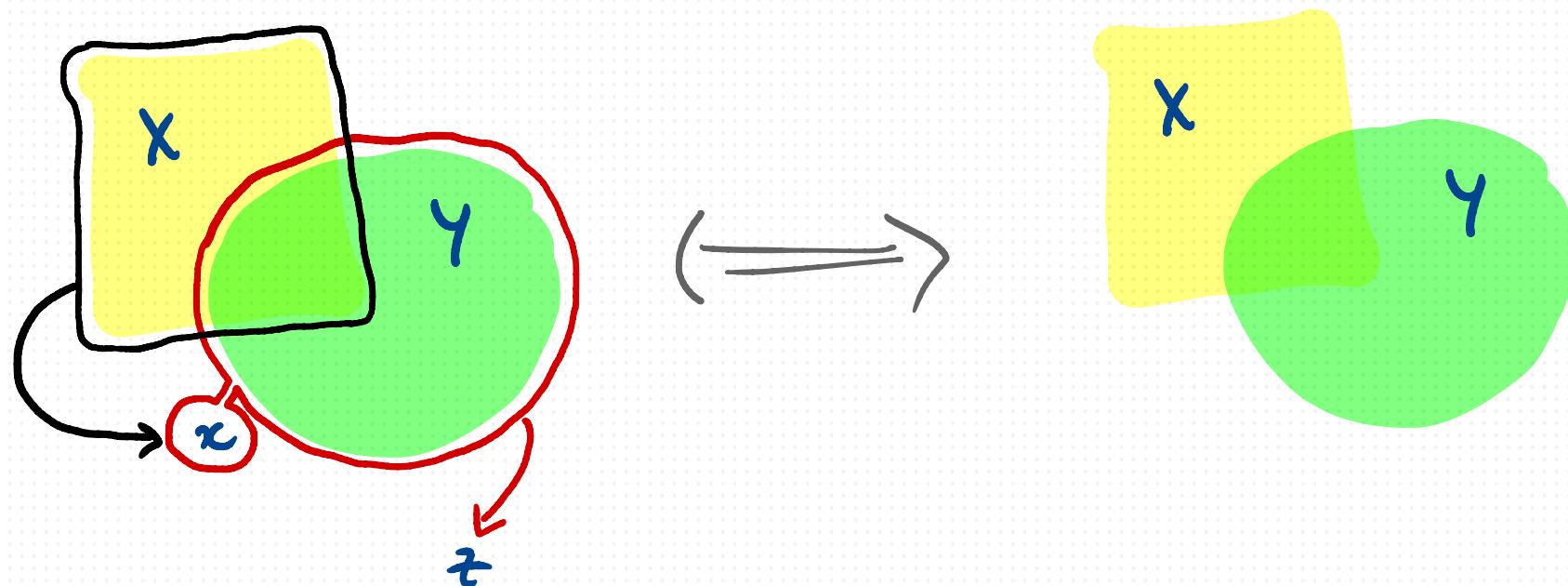
$$\exists x (R X_x \notin R([x] \cup Y) z) \Leftrightarrow R(X \cup Y)z$$

# COMPOSITIONALITY



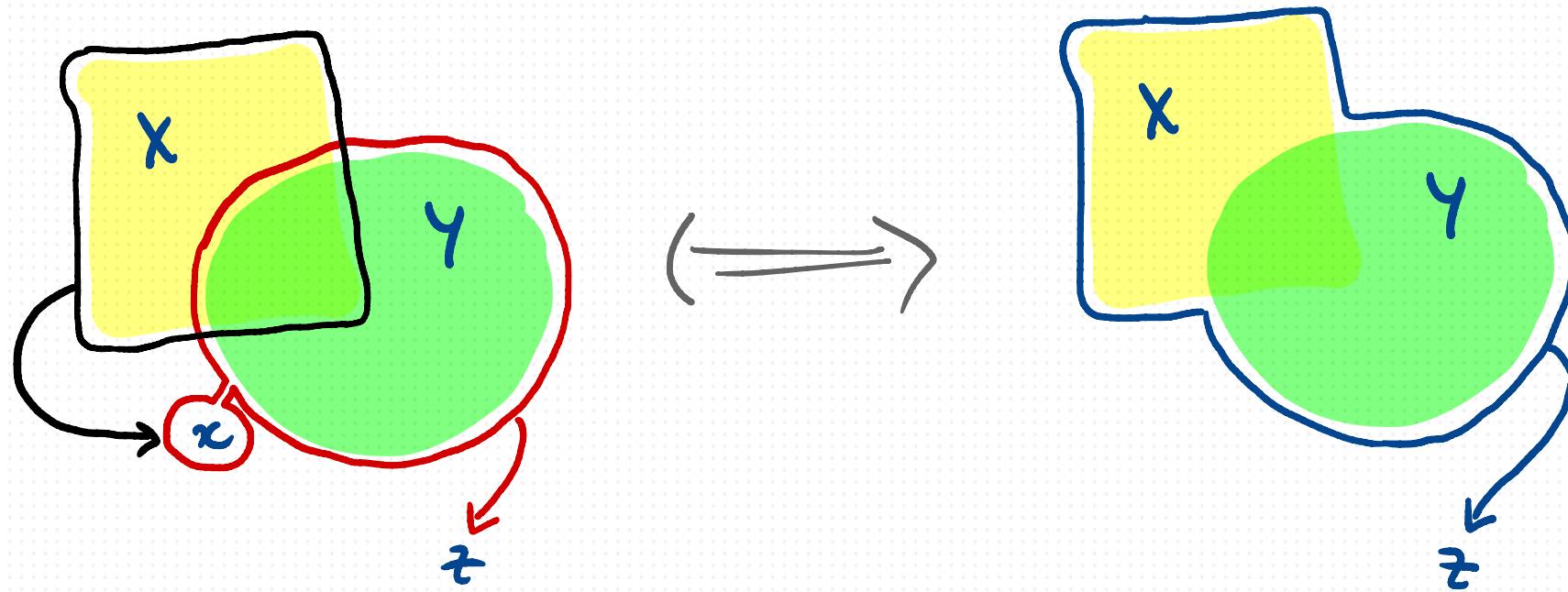
$$\exists x (R[X_x] \notin R([x] \cup Y) z) \Leftrightarrow R(X \cup Y) z$$

# COMPOSITIONALITY



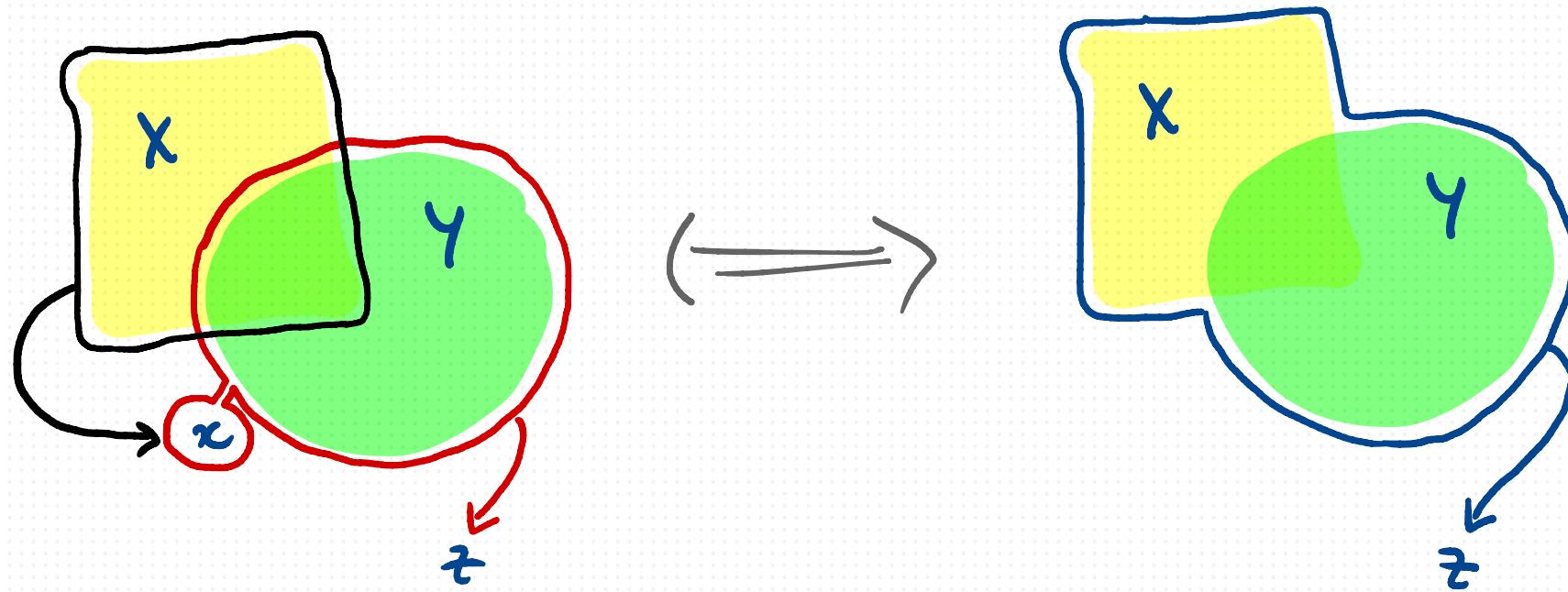
$$\exists x (R[X_x \notin R([x] \cup Y)]z) \Leftrightarrow R(X \cup Y)z$$

# COMPOSITIONALITY



$$\exists x (R[X_x] \notin R([x] \cup Y)z) \Leftrightarrow R(X \cup Y)z$$

# COMPOSITIONALITY



$$\exists x (R[X_x \notin R([x] \cup Y)]z) \Leftrightarrow R(X \cup Y)z$$

When  $|X| = m, |Y| = m$ , call this **Compositionality<sup>m</sup>**

fun FACT #1 : This is conceptually way simpler

### TERNARY FRAME LANGUAGE

- $N_x$
- $x \in q$
- $Rxyz$

### MULTISET FRAME LANGUAGE

- $R[]x$
- $R[x]y$
- $R(x,y)z$

fun FACT #1 : This is conceptually way simpler

### TERNARY FRAME LANGUAGE

- $N_x$
- $x \in y$
- $Rxyz$
- $\subseteq$  Reflexive

### MULTISET FRAME LANGUAGE

- $R[]x$
- $R[x]y$
- $R(x,y)z$
- $R$  reflexive

fun FACT #1 : This is conceptually way simpler

### TERNARY FRAME LANGUAGE

- $N_x$
- $x \in y$
- $Rxyz$
- $\subseteq$  Reflexive
- $\subseteq$  Transitive

### MULTISSET FRAME LANGUAGE

- $R[]x$
- $R[x]y$
- $R(x,y)z$
- $R$  reflexive
- $R$  compositional

fun FACT #1 : This is conceptually way simpler

### TERNARY FRAME LANGUAGE

- $N_x$
- $x \in y$
- $Rxyz$
- $\subseteq$  Reflexive
- $\subseteq$  Transitive
- $N^\uparrow$

### MULTISSET FRAME LANGUAGE

- $R[]x$
- $R[x]y$
- $R(x,y)z$
- R reflexive
- R compositional<sup>1</sup>
- R compositional<sup>2</sup>

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- $\subseteq$  Transitive
- $N^\dagger$
- $R \downarrow \dagger \uparrow$

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- R compositional<sup>3</sup>

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- $R(x,y)z$
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- $R$  compositional<sup>1</sup>
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- $R$  compositional<sup>4</sup>

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- $R$  compositional<sup>4</sup>
- $R$  compositional<sup>5</sup>
- $-$

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- $Rxyz \Leftrightarrow Ryxz$
- $\exists z (Rxyz \wedge Rzuv) \Leftrightarrow \exists w (Ryzw \wedge Rxwv)$

### MULTISSET FRAME LANGUAGE

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- $R$  compositional<sup>1</sup>
- $R$  compositional<sup>0</sup>
- $R$  compositional<sup>1</sup>
- $R$  compositional<sup>0</sup>
- $R$  compositional<sup>1</sup>
- $R$  compositional<sup>0</sup>
- $-$
- $R$  compositional<sup>2</sup>

## Fun Facts #2 & #3

★ Every RW+ ternary frame  $\langle P; R, E, N \rangle$  uniquely determines a multiset frame  $\langle P; R \rangle$ , where  $R$  is reflexive & compositional.

- Define •  $[ ]Rz := Nz$     •  $[x,y]Rz := Rxz \wedge Ryz$
- $[x]Rz := x \subseteq z$     •  $([x] \cup Y)Rz := \exists y (YRz \wedge Rxzy)$

## Fun Facts #2 & #3

★ Every RW+ ternary frame  $\langle P; R, \sqsubseteq, N \rangle$  uniquely determines a multiset frame  $\langle P; R \rangle$ , where  $R$  is reflexive & compositional.

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- $[x]Rz := x \sqsubseteq z$     •  $([x] \cup Y)Rz := \exists y (YRz \wedge Rxz)$

★ Every multiset frame  $\langle P; R \rangle$ , where  $R$  is reflexive & compositional, determines an equivalent ternary frame  $\langle P; R, \sqsubseteq, N \rangle$ .

(Just reverse the process!)

# NOTICE SOMETHING INTERESTING?

## TERNARY FRAME LANGUAGE

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- $x \sqsubseteq y$
- $Rxyz$
- $\sqsubseteq$  Reflexive
- $\sqsubseteq$  Transitive
- $N^\dagger$
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- $\exists x (N_x \wedge Rxyz) \Leftrightarrow y \sqsubseteq z$
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 $\exists w (Ryzw \wedge Rxwv)$

## MULTISSET FRAME LANGUAGE

- $R[]x$
  - $R[x]y$
  - $R(x,y)z$
  - $R$  reflexive ???
  - $R$  compositional ✓
  - —
  - $R$  compositional ✓
- Why is this here? ↗

# MY PLAN

1. WHAT ARE MODELS FOR?
2. THE COMPLEXITY of ROUTER/MAYER FRAMES
3. THREE INTO ONE
4. BEYOND REFLEXIVITY
5. EXAMPLES
6. OPEN QUESTIONS

REFLEXIVITY of  $R$  is not NEEDED for  $RW^+$ !

A multiset frame  $\langle P; R \rangle$  for  $RW^+$  is a non-empty set  $P$  equipped with a compositional multiset relation  $R$  on  $P$ .

That's it !

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A multiset frame  $\langle P; R \rangle$  for  $RW^+$  is a non-empty set  $P$  equipped with a compositional multiset relation  $R$  on  $P$ .

That's it !

This is much more mathematically suggestive, at least!

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## EXAMPLE #1 Multiplication on $\omega$

$[a_1, \dots, a_n] R b$  iff  $a_1 \times \dots \times a_n = b$

$[] R b$  iff  $1 = b$

## EXAMPLE #1 Multiplication on $\omega$

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$[] R b$  iff  $1 = b$

Compositionality:

$$\begin{aligned} & \exists x (\Pi X = x \wedge \Pi((x) \cup Y) = z) \\ \Leftrightarrow & \quad \Pi(\Pi X \cup Y) = z \\ \Leftrightarrow & \quad \overline{\Pi}(X \cup Y) = z \end{aligned}$$

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Compositionality:  $\exists x (\Pi X = x \wedge \Pi((x) \cup Y) = z)$   
 $\Leftrightarrow \Pi(\Pi X \cup Y) = z$   
 $\Leftrightarrow \Pi(X \cup Y) = z$

Ordering:  $[x] R y$  iff  $x = y$

## EXAMPLE #2 Multiplication on $\omega$ with DIVISION

$[a_1, \dots, a_n] R b$  iff  $a_1 \times \dots \times a_n \mid b$

$[] R b$  iff  $1 \mid b$

$x \mid y$  iff  
 $(\exists n > 0) nx = y$

## EXAMPLE #2    MULTIPLICATION on $\omega$ with DIVISION

$$[a_1, \dots, a_n] R b \quad \text{iff} \quad a_1 \times \dots \times a_n \mid b$$

$$[ ] R b \quad \text{iff} \quad 1 \mid b$$

Compositionality:  $\exists x (\Pi X | x \notin \Pi([x] \cup Y)) | z$

$$\Leftrightarrow \exists n > 0 (\Pi([n \Pi X] \cup Y)) | z$$

$$\Leftrightarrow \overline{\Pi}(X \cup Y) | z$$

$x | y$  iff  
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$$\Leftrightarrow \exists n > 0 (\Pi([n \Pi X] \cup Y)) | z)$$

$$\Leftrightarrow \Pi(X \cup Y) | z$$

Ordering:  $[x] R y \quad \text{iff} \quad x | y$

$x | y \text{ iff}$   
 $(\exists n > 0) nx = y$

## EXAMPLE #3    MULTIPLICATION on $\omega$ with DIVISION — backwards!

$[a_1, \dots, a_n] R b \text{ iff } b \mid a_1 \times \dots \times a_n$

$[] R b \text{ iff } b \mid 1, \text{ i.e.}$

Compositionality:  $\exists x (\Pi X | x \notin \Pi([x] \cup Y)) | z$

$\Leftrightarrow \exists n > 0 (\Pi([n \Pi X] \cup Y)) | z$

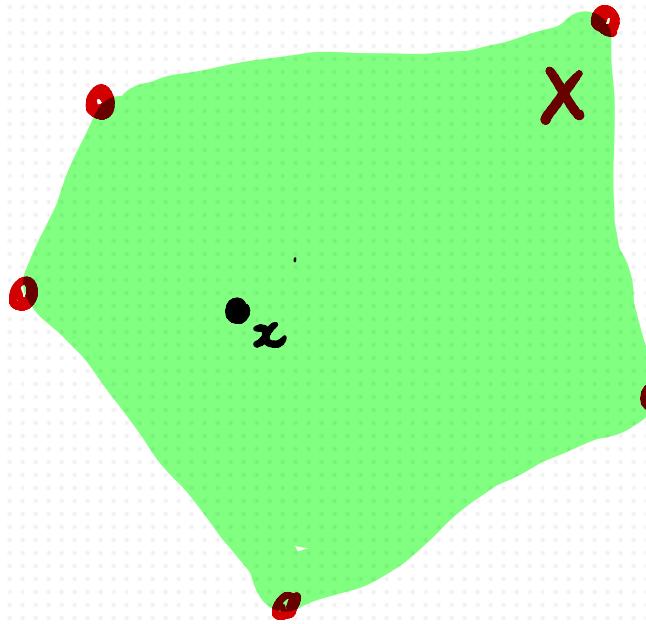
$\Leftrightarrow \Pi(X \cup Y) | z$

Ordering:  $[x] R y \text{ iff } y | x$

$x | y \text{ iff}$   
 $(\exists n > 0) nx = y$

## EXAMPLE #4? — CONVEX SHAPES IN $\mathbb{R}^2$

$xRy$  if  $y$  is in the shape bounded by the points  $X$ .

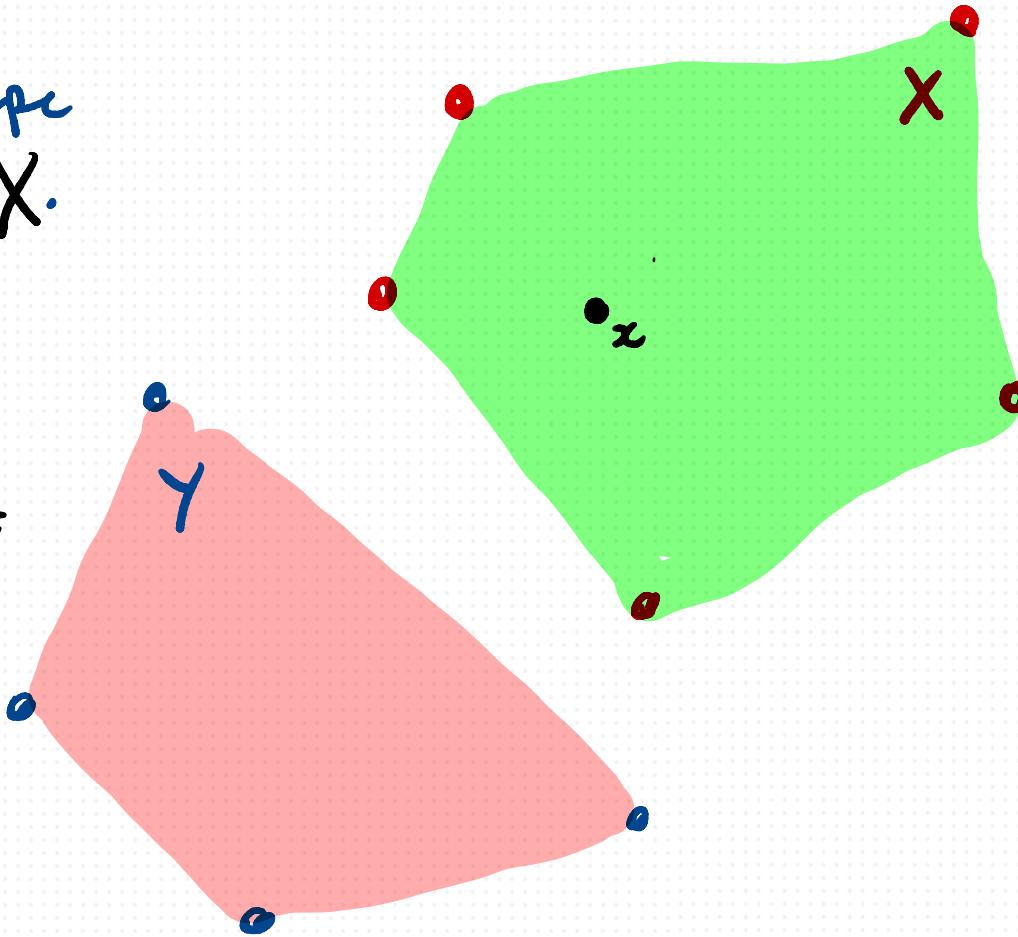


## EXAMPLE #4? — CONVEX SHAPES IN $\mathbb{R}^2$

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If  $X R x \notin [x] \cup Y R z$

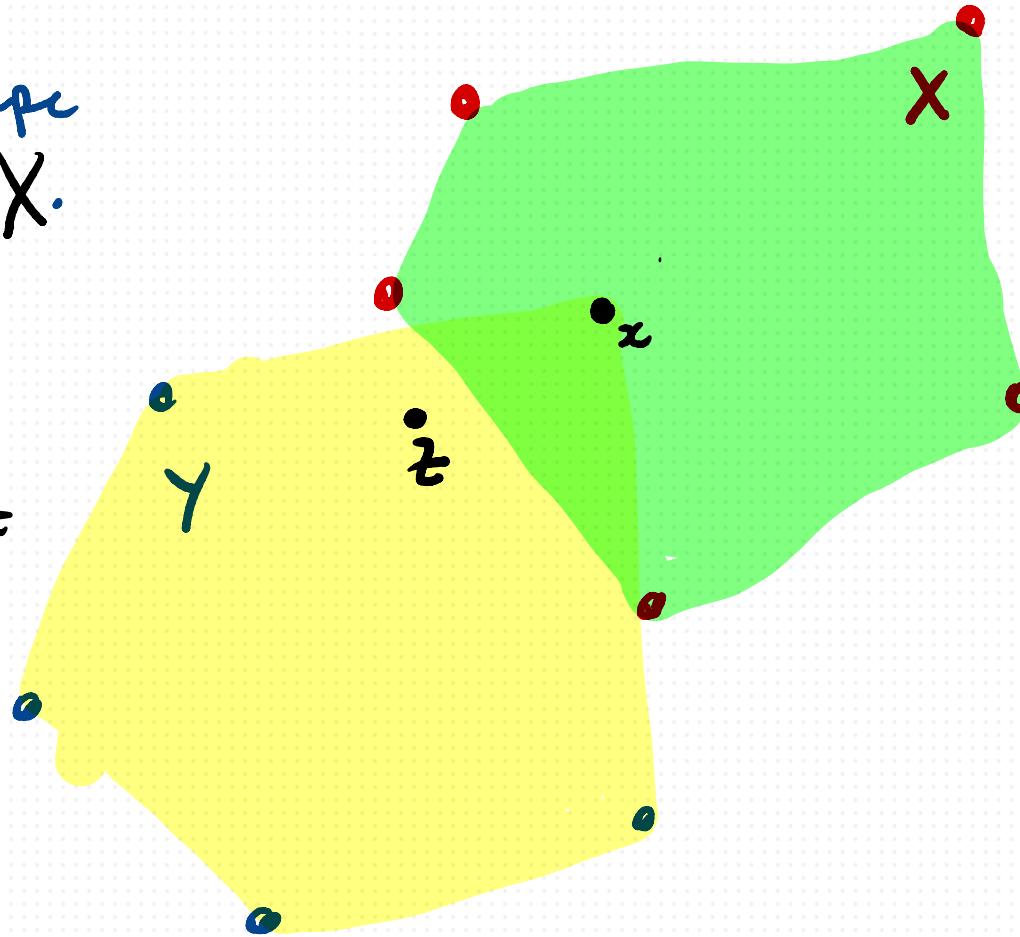
then  $(X \cup Y) R z$



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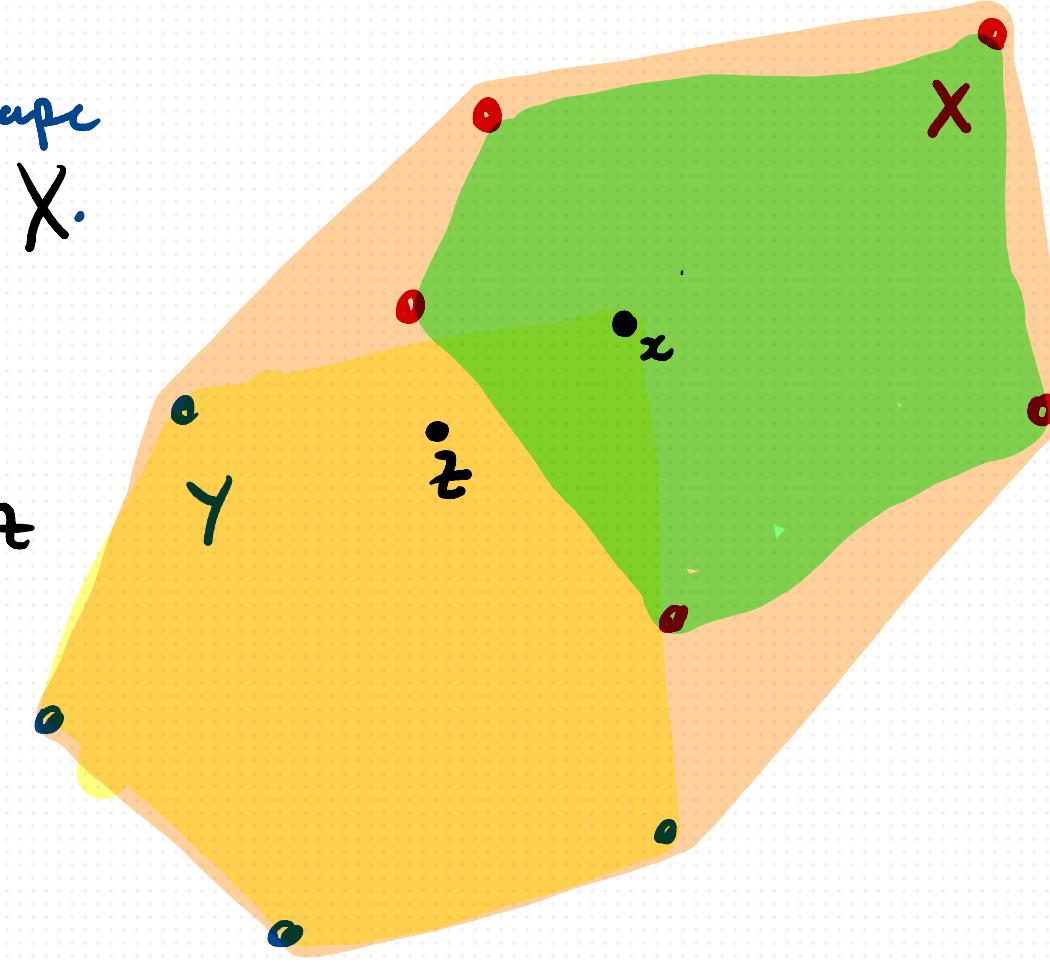
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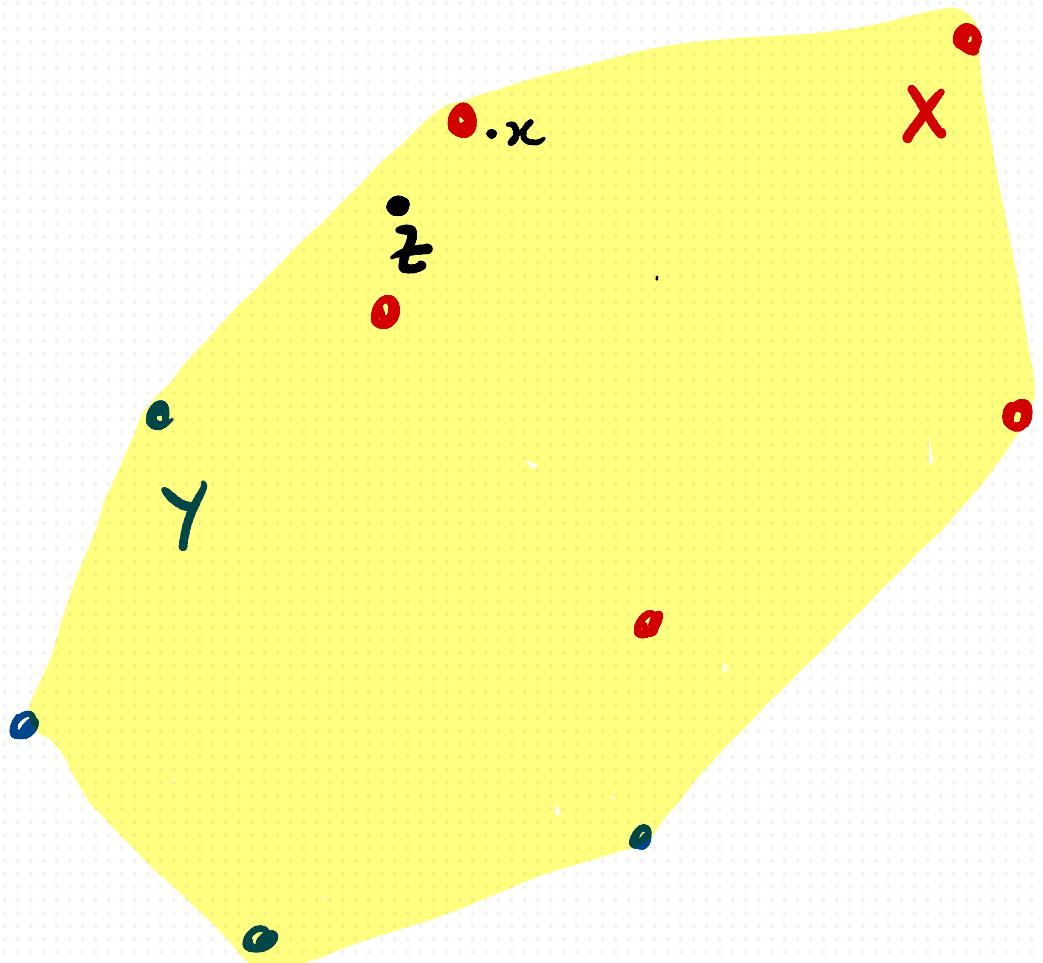


## EXAMPLE #4? — CONVEX SHAPES IN $\mathbb{R}^2$

$xRy$  if  $y$  is in the shape bounded by the points  $X$ .

If  $xRx \notin [x] \cup YRz$   
then  $(x \cup y)Rz$

If  $(x \cup y)Rz$  then  
 $\exists_x xRx \notin (x \cup y)Rz$

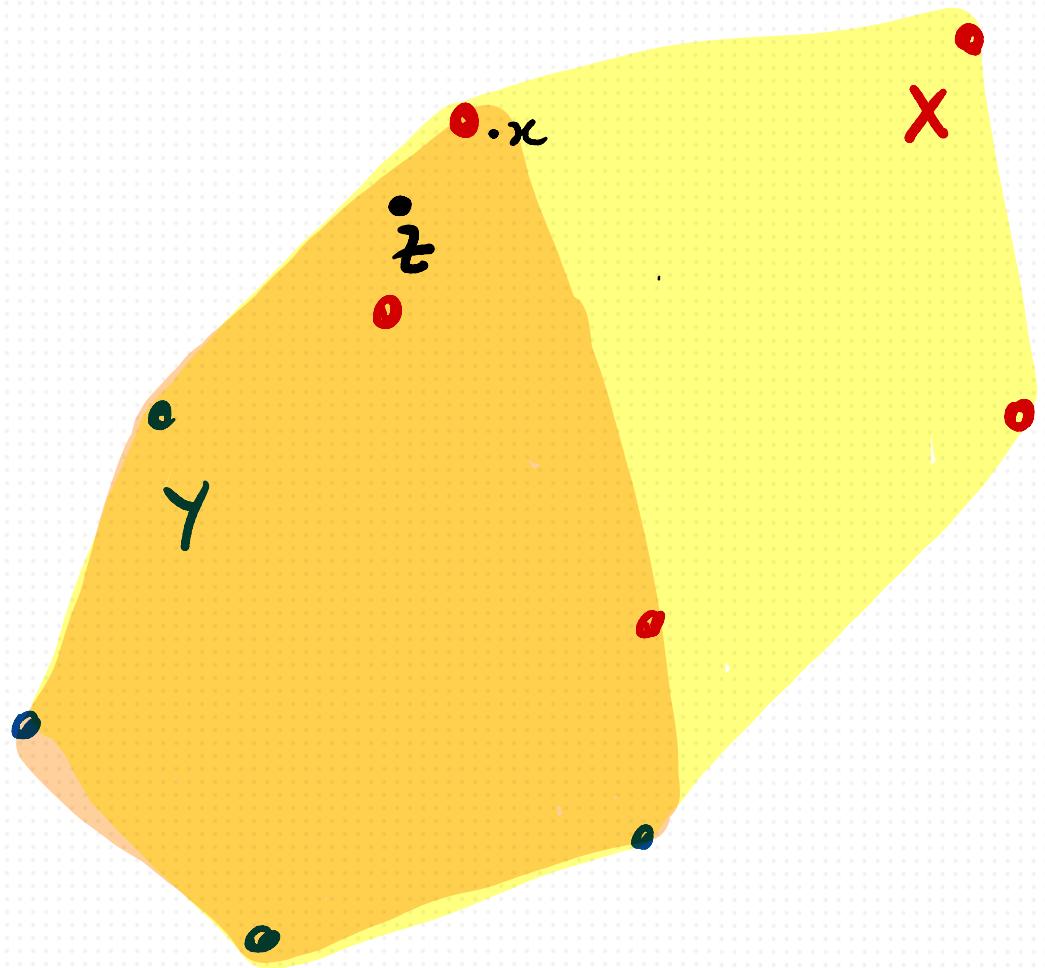


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If  $x R_x \notin [x] \cup Y R z$

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If  $(x \cup y) R z$  then

$\exists_x x R_x \notin (x \cup y) R z$

if  $X$  is nonempty!

If  $X$  is empty, there's no  $x$  where  $x R_x$ !

membership of ConvexShape is a composition relation  
on inhabited multisets (or sets) on  $\mathbb{R}^2$ .

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if  $X$  is nonempty!

If  $X$  is empty, there's no  $x$  where  $x R_x$ !

MEMBERSHIP in a convex shape is a compositional relation on inhabited multisets (or sets) of points in  $\mathbb{R}^2$

$R$ , a relation on  $M^+(P) \times P$  is compositional iff  
inhabited multisets

$$\forall X \in M^+(P) \forall Y \in M(P) \forall z (\exists x X R_x \notin (x) \cup Y) R_z \leftrightarrow (X \cup Y) R_z]$$

$R$ , a relation on  $M^+(P) \times P$  is compositional iff  
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- Interior of Convex shapes is a compositional relation on inhabited multisets (or sets) in  $\mathbb{R}^2$

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inhabited multisets

$$\forall X \in M^+(P) \forall Y \in M(P) \forall z (\exists x X R x \notin (x \cup Y) R z) \leftrightarrow (X \cup Y) R z]$$

- Interior of Convex shapes is a compositional relation on inhabited multisets (or sets) in  $\mathbb{R}^2$ .
- Multiset (or Set) membership is also compositional on inhabited sets.

# INHABITED MULTISet RELATIONS

$N_x$

$[]R_x$

$x \subseteq y$

$(x)R_y$

$R_{xyz}$

$(x,y)R_z$

## INHABITED MULTISet RELATIONS

$\exists x$	$(\lambda)R_x$
$x \in y$	$(x)R_y$
$R_{xyz}$	$(x,y)R_z$

These model the logic without  $\top$ .

They represent validity of sequents,  
but not theoremhood, as there is  
nowhere for theorems to hold.

WHAT IF COLLECTIONS ARE SETS & NOT MULTISETS?

$R_{xyz}$  iff  $\{x,y\} R z$

This means  $R_{xxy}$  iff  $\{x\} R y$  iff  $x \subseteq y$

WHAT IF COLLECTIONS ARE SETS & NOT MULTISETS?

$$Rxyz \text{ iff } \{x,y\} R z$$

This means  $Rxxz$  iff  $\{x\} R z$  iff  $x \subseteq z$

This delivers  $p \rightarrow (p \rightarrow q) \vdash p \rightarrow q$ .

Set frames model  $R^+$  (at least).

WORRY

$$\{x, x\} R y \Leftarrow \{x\} R y \Leftarrow x \Sigma y$$

Looks a lot like  $p \circ p \vdash p$ ,

which delivers  $p \vdash p \rightarrow p$ .

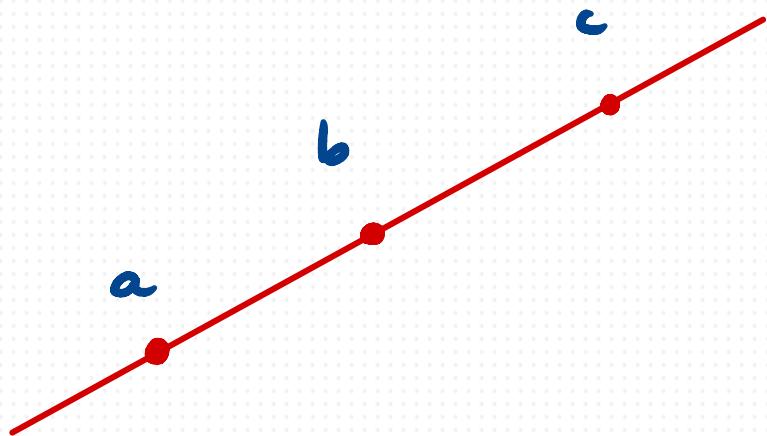
the MINCLE rule,

which doesn't hold in  $R^+$ .

## MINGLE CAN FAIL IN SET FRAMES

Take the  $\mathbb{R}^2$  frame with  
 $X R y$  iff  $y$  is in the  
convex shape bound by  $X$ .

Let  $a, b, c$  be three distinct  
points on a line, with  $b$   
between  $a$  and  $c$ .



The relation  $R$  is compositional only  
on inhabited subsets of  $\mathbb{R}^2$ .  
You could add a point at infinity,  
 $\infty$  & extend  $R$  to  $\mathbb{R}^\times$  on  $\mathbb{R}^2 \cup \{\infty\}$  by

setting  $X R_z^\times$  if  $\begin{cases} z = \infty & \text{if } X = \{\gamma\} \text{ or } \{\delta\} \\ (X \backslash \{\alpha\}) R_z & \text{otherwise.} \end{cases}$

$R^\times$  is compositional on all finite sets.

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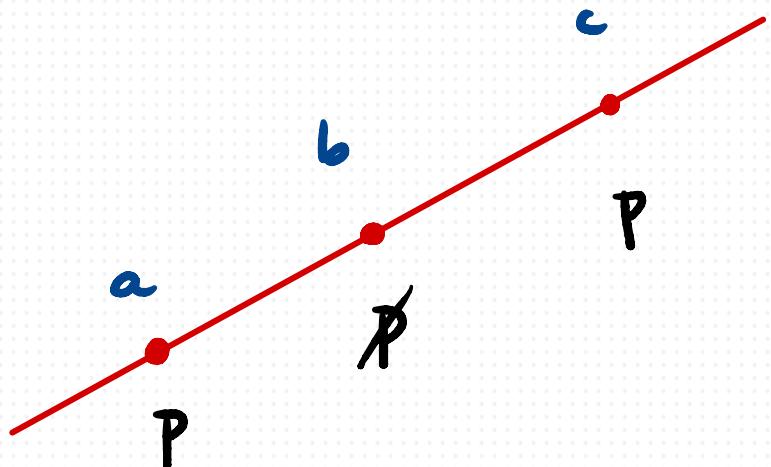
Let  $a, b, c$  be three distinct  
points on a line, with  $b$   
between  $a$  and  $c$ .

Let  $[p] = \{a, c\}$

$a \Vdash p \rightarrow p$ , since  $c \Vdash p$

$\notin \{a, c\} R b \notin b \Vdash p$ .

So  $p \not\Vdash p \rightarrow p$ .



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CONJECTURE: SET FRAMES ARE COMPLETE  
FOR  $R^+$ .

Fact: Some  $R^+$  multi-set frames ( $RW^+$  frames where  $[x,x]R x$ )  
do not satisfy  $[x,x]R y \Leftrightarrow [x]R y$ .

Tiny Case:

$R^2$	0	a	b
0	0	a	b
a	a	ab	0ab
b	b	0ab	ab

But - can any such frame be embedded into some  $R^+$  set frame?

Thank You /  
.

<https://consequently.org/>