1) longest wit.

to find mere that; mars-length=1 current-l=1 nums[i] > nums[i-1] WC Start from i=1 compare nums[1] > nums[0] so we all +1 +0 curr. then, wring morp function, compare max to curr 4 [.e.: i=3, hums(3) <= nums[2] max is 3, cour is 3 since 4 isn't greater, We set current to 1 then => nums[4] > nums[3] 80, aurs-len + 1 => curr= 2 we take maximum of max and cur => 3 > 2 map wille 3

2) Mergo arrey
cine arrand are sorted we can start from the end
Gince arrays are sorted, we can start from the end. filling 0's of numl
(123000) [256] ve need two variables as last elements
pointer_runs = m-1 and we need variable that
pointer-numz= n-1 indicates last poor
$\frac{1}{10000000000000000000000000000000000$
to make jure we merge array in non-decreasing order we should make sure that element played in the and
is greater so.
is greater so
P $1=2$
$\begin{bmatrix} 1230007 & \begin{bmatrix} 256 \end{bmatrix} & p-n1=2 \\ position=5 \\ \end{bmatrix}$ $\begin{bmatrix} 256 & 200 \\ position=5 \\ \end{bmatrix}$
2 2 6 position= 5
L'a la gracese
pointer_nums and position $p-n2=1, posish=4$
p-n2=1, Posish-1
[123006] [256]
2 < 5 p-n2=0, position=3
C123056] [256]
con l'ave pointer?
[103356] and move pointers
1 127-1011100CaCi
2=2=> hums1[k]=humsq[j.
now we reach to the end of the neems cerrary

3) Insersell

Since we need elements that are present in both lists and one unique, we can simply use Sets. using seas remove duplicates, so elemens will be unique. after we create 2 sets:

we are going to use a new list for result

we are going to use a new list for test cand iterate over sets like; hums 1=[1,2,2,1] set $1=\{1,2\}$ nums 2=[2,2] set $2=\{2\}$

iterate through set2, E and check if set1 contains it, so that they intersect since both have 2, result = [2]