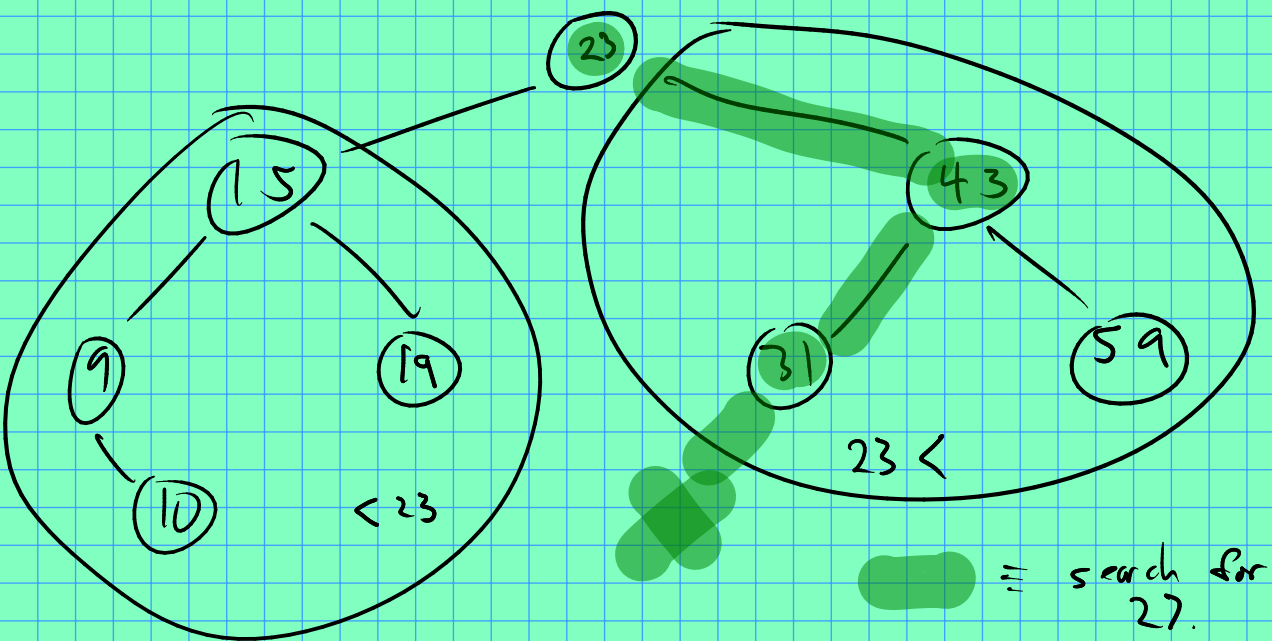


# Sets & Maps

Implemented as binary search trees:



(Quick note on implementation: use pointers!)



Advantages over arrays/vectors:

Say tree has  $n$  values.

Time for a search is  $\log_2 n$

$$1 = n/2^k$$

solve for  $k$ :

$$2^k = n \quad \text{so } k = \log_2 n.$$

$n/2^k \approx \#$  nodes left in subtree  
after  $k$ -steps (step  $\equiv$  comparison)

$\log_2 n$  is much faster than searching in  
a vector ( $n$  steps)

## Disadvantages vs vectors:

Random access (setting the  $i^{\text{th}}$  element)  
takes  $\approx i$  steps in a set, but  
 $\approx 1$  for a vector (use  $VC[i]$ )

## Features they share:

easy to add new elements at will:

$v.\text{push\_back}(x) \approx s.\text{insert}(x)$

Notes: they are modelling mathematical sets,  
so they don't store duplicates.

```
s.insert(23);  
cout << s.size(); // prints 10  
s.insert(23);  
cout << s.size(); // still 10.
```

Aside: boolean functions as sets:

Say universe is  $U$ .

how does each function  $f: U \rightarrow \{0,1\}$   
define (uniquely) a subset  $S \subseteq U$ ?

$$\{S \subseteq U\} \longleftrightarrow \{f_S: U \rightarrow \{0,1\}\}$$

The bijection is as follows:

For a set  $S \subseteq U$ , define its 'characteristic

function"  $f_S: U \rightarrow \{0,1\}$  by the rule

$$f_S(x) = \begin{cases} 1 & \text{if } x \in S \\ 0 & \text{else} \end{cases}.$$

For a function  $f: U \rightarrow \{0,1\}$  could recover  
the set as  $S_f = f^{-1}(\{1\}) \subseteq U$

$$\begin{array}{c} \nearrow \\ \{x \in U \mid f(x) \in \{1\}\} \\ \text{(i.e., } f(x) = 1) \end{array}$$

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$$f^{-1}: \mathcal{P}(Y) \rightarrow \mathcal{P}(X)$$

$$\text{for } f: X \rightarrow Y.$$