### **OS Intro**

**Thanks to Euripides Montagne** 

- 1.- Absolute Loader
- 2.- Bootstrap Loader
- 3.- Operating System: Loading application programs
  - 3.1- The Process concept.
  - 3.2- Creating a process.
- 4.- Relocating Loaders
  - 4.1- Relocation bits
  - **4.2- Relocation maps (modification records)**

#### **Assembly program**

Label

01

<u>opcode</u>	<u>address</u>
; This is	
; a comm	ent

_	,	
02	; a comme	nt
03 start	.begin	x200
04 here	LOAD	sum
<b>)</b> 5	ADD	а
06	STORE	sum
07	LOAD	b
08	SUB	one
<b>)</b> 9	STORE	b
DA	SKIPZ	
)B	JMP	here
C	LOAD	sum
)D	HALT	
DE sum	.data	x000
OF a	.data	x005
LO b	.data	x003
L1 one	.data	x001
<b>L2</b>	.end	start

#### object code file

Program name: start Starting address text: x200 Length of text in bytes: x14 Starting address data: x20A Length of data in bytes: 8	H e a d e r
0001000000001001 0010000000001001 001100000000	Text
00000000000000000000000000000000000000	Data

**Data** 

Header

Text section

#### Assembler object code file

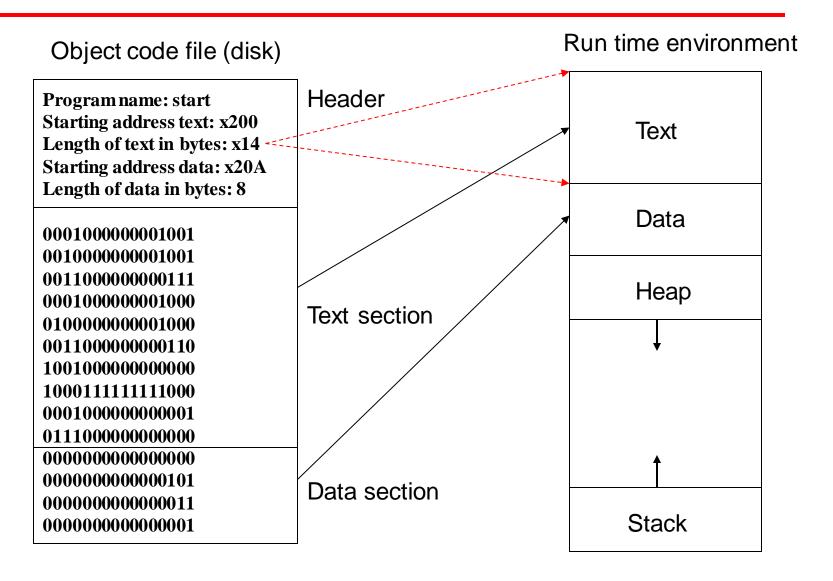
Program name: start Starting address text: x200 Length of text in bytes: x14 Starting address data: x20A Length of data in bytes: 8
Length of data in bytes. o
000400000004004
0001000000001001
001000000001001
001100000000111
0001000000001000
0100000000001000
0011000000000110
1001000000000000
1000111111111000
0001000000000001
0111000000000000
000000000000000
000000000000101
000000000000011
0000000000000001
L

#### **Absolute loader:**

The absolute loader will load the program at memory location x200:

- 1.- The header record is checked to verify that the correct program has been presented for loading.
- 2.- Each text record is read and moved to the indicate address in memory
- 3.- When the "end" record (EOF) is encountered, the loader jumps to the specified address to begin execution.

### Loading object code into memory



### Bootstrapping:

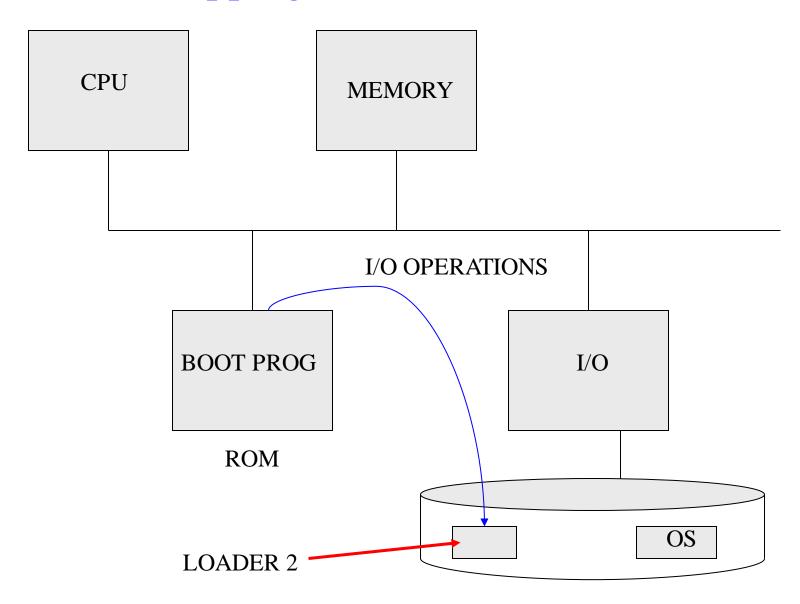
Computers execute programs stored in main memory, and initially the operating system is on the hard disk.

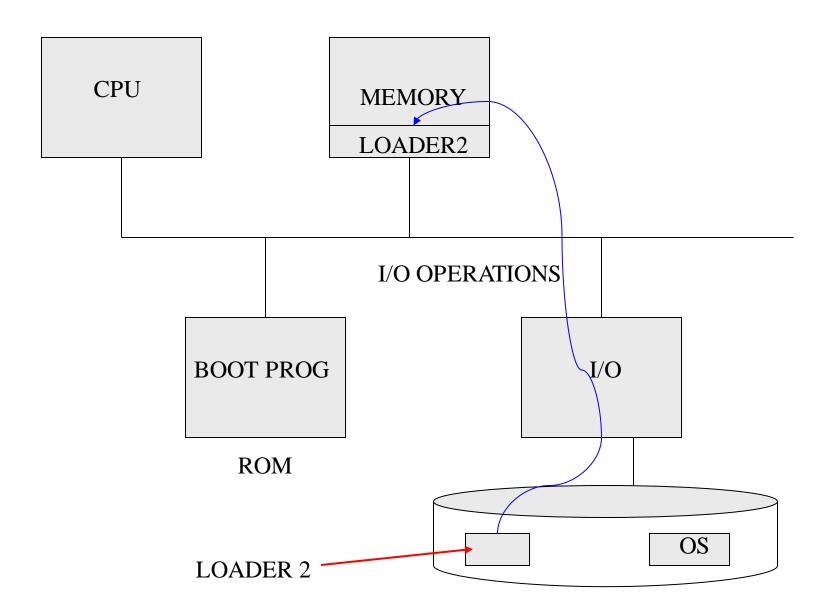
When the computer is turned on it does not have an operating system loaded in memory and the hardware alone cannot do the operations of an OS. To solve this paradox a special program called bootstrap loader is created.

■ This program does not have the full functionality of an operating system, but it is capable of loading into memory a more elaborated software (i.e. loader2) which in its turn will load the operating system.

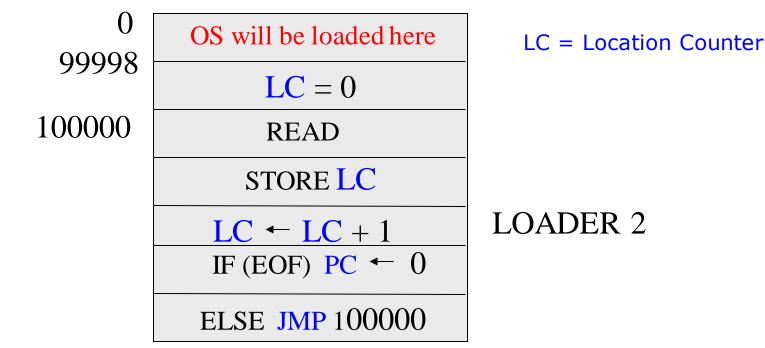
Once the OS has been loaded the loader transfers the control of the computer system to the operating system.

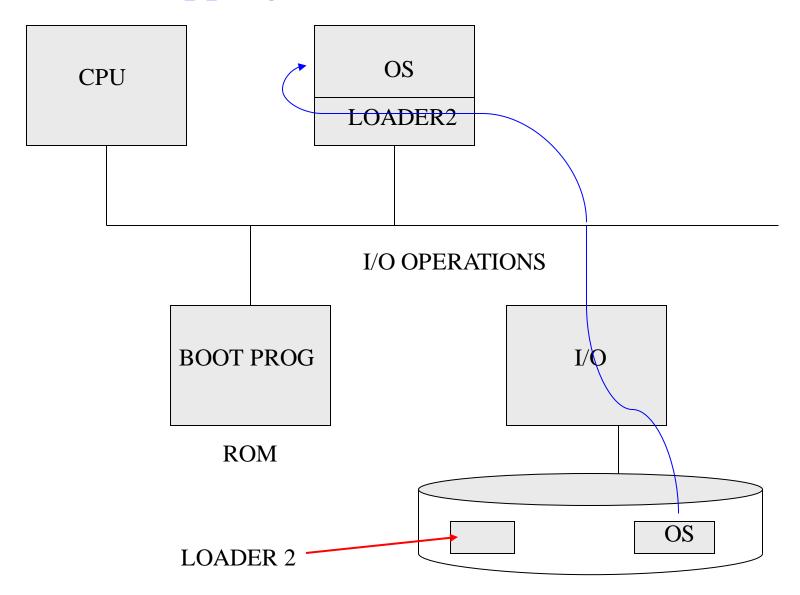
- Early programmable computers had toggle switches on the front panel to allow the operator to place the bootloader into the program store before starting the CPU.
- In modern computers the bootstrapping process begins with the CPU executing software contained in ROM at a predefined address whose elementary functionality is to search for devices eligible to participate in booting, and load a small program from a special section of a device.





■ In earlier computers data had to be hand loaded as specified before, but nowadays a small piece of software called loader helps us to avoid the manual loading.





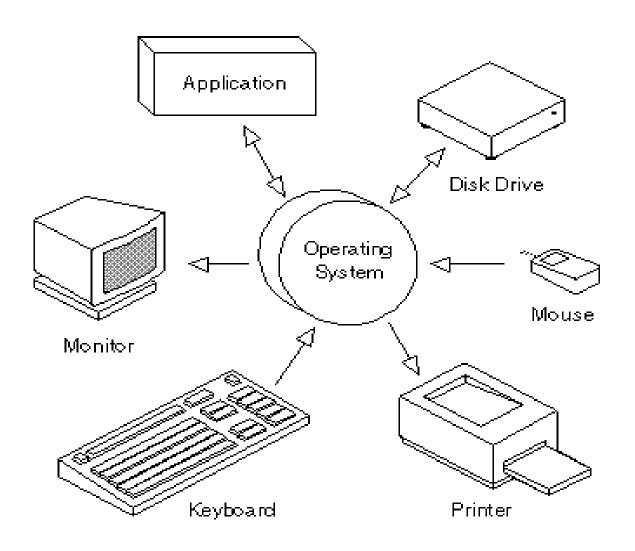
- The above diagram can be explained in the following steps.
  - 1. Check hardware
  - 2. Initiate I/O to load the loader 2 program into memory
  - 3. Loader 2 loads the OS and passes control to it

#### Conclusion

■ We have seen that once the OS has control over the system, it can create an environment for programs to run.

■ The operating system will load device drivers and other programs that are needed for the normal operation of the computer system.

## Operating system



# The Process Concept

#### **Programs and processes**

Once the operating system takes control of the computer system, an applications program (object module or ELF) can be loaded into memory to be executed.

When the program is loaded into memory a process is created.

What is a process?

#### **Process**

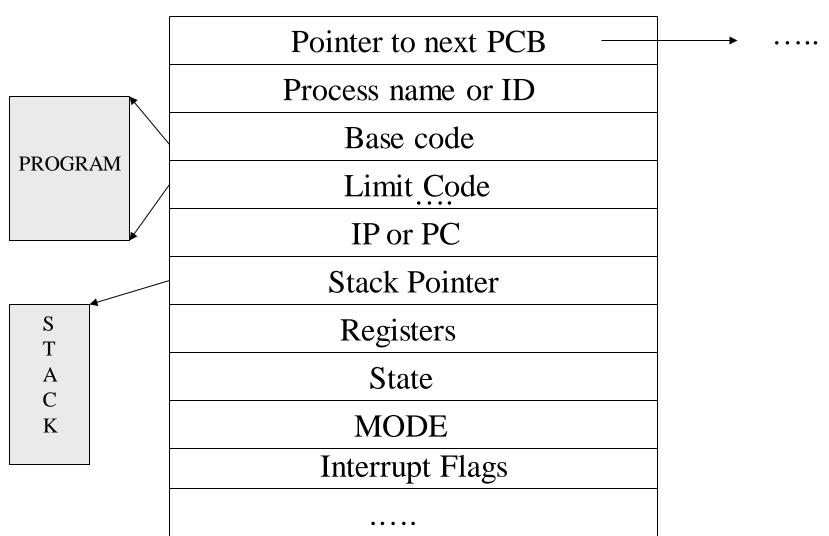
#### Definition:

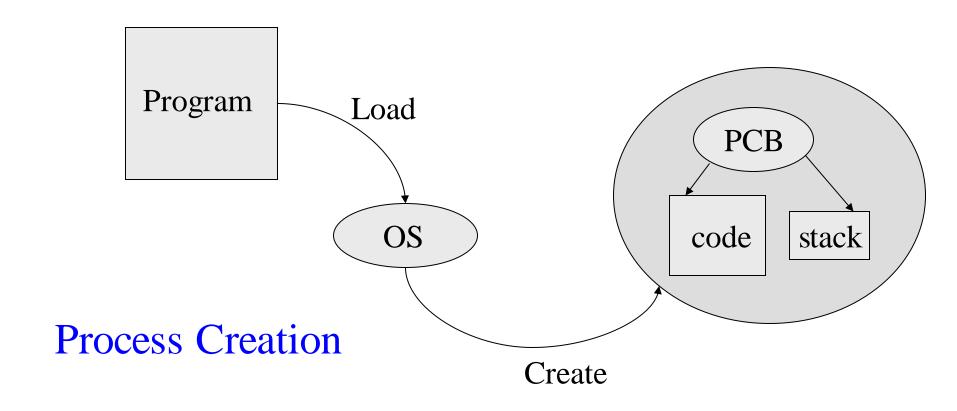
- A program in execution
- An asynchronous activity
- The "locus of control" of a procedure in execution
- It is manifested by the existence of a process control block (PCB) in the operating system.

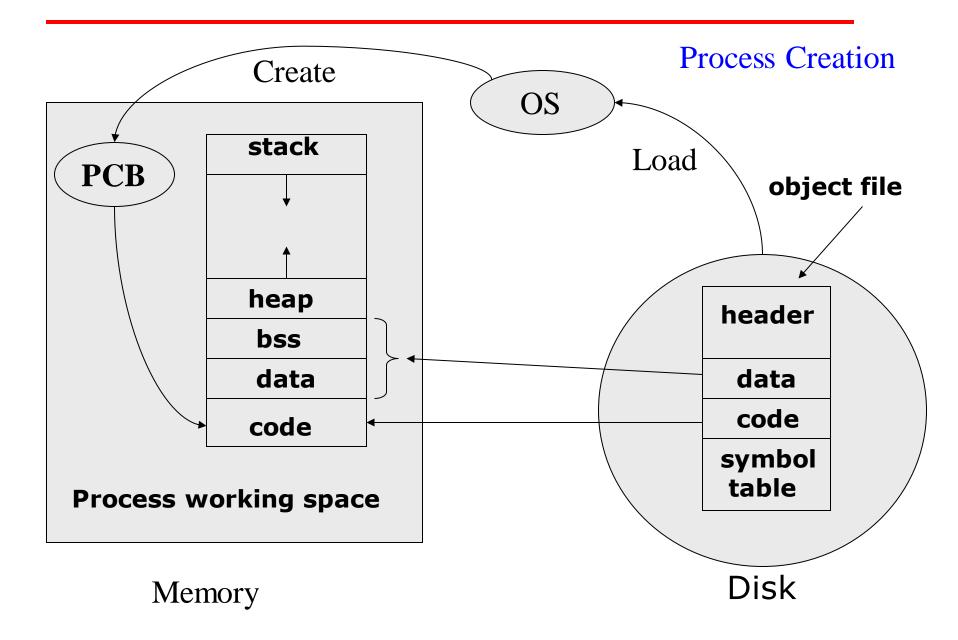
#### Process Continued...

- The activity of a process is controlled by a data structure called Process Control Block(PCB).
- A PCB is created every time a program is loaded to be executed.
- So, a process is defined by a PCB-Program couple.

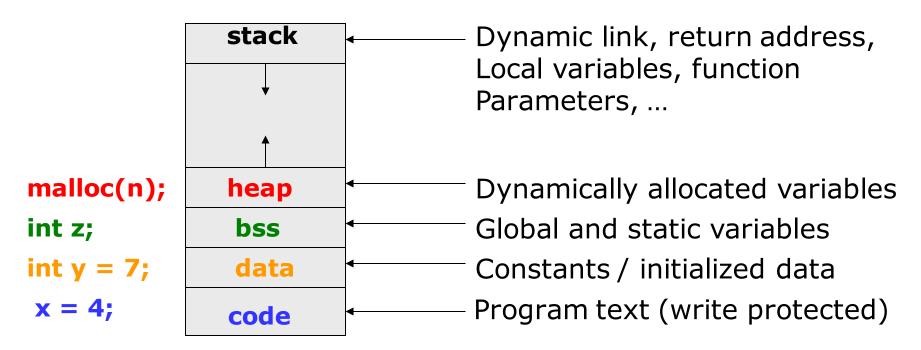
#### **Structure of the PCB**







#### Process working space (run-time environment)



Process working space or Run-time environment

#### **Process working space**

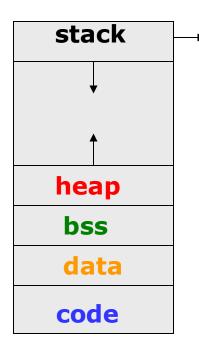
bss: means "block started by symbol" and has that name for historical reasons.

malloc(n);

int y = 7;

x = 4;

int z;



```
sum
    Local
    Local
    Local
    Local
    Local
    Local
  Parameter
  Parameter
 Dynamic Link
Return Address
```

```
void sub(float total, int part ) {
int List[5];
float sum;
...
}
```

#### **Relocating Loaders**

Absolute loaders loads a program on a specific memory location but it is often desirable to have two or more programs residing in memory sharing the different resources of a computer system.

It would be impractical to assign starting addresses to each program to plan program execution.

A loader able to load a program into memory wherever there is room for it is called a <u>relocating</u> loader.

#### **Relocation bits**

Assemblers generate code that starts at address zero but it can also emit with each line of text (code) relocation bits indicating what fields in the object code must be modified when the program is loaded in an address different from zero.

For example, if the program will be loaded at address 40, a relocation bits equal to "1" indicates what part of the instruction must be modified:

Loc#	<u>Len</u>	<u>reloc</u>	<u>text</u>		Loc#	<u>text</u>	
00	3	011	13 <mark>33 35</mark>	<b>→</b>	40	13 <mark>73 75</mark>	

	source program				before relocation				after relocation	
Lab	<u>bel</u> o	<u>pcode</u>	<u>address</u>	<u>address</u>	Loc#	<u>Len</u>	<u>reloc</u>	<u>text</u>	Loc#	<u>text</u>
00 03 06 08 10 con 12 14 16 18 20 22 25 28 30 fina 32 33 zer 34 one 35 olde 36 old 37 nev 38 limit	o Constant of Cons	opy opy ead vrite oad tore ub rpos vrite opy opy r vrite top ONST ONST PACE PACE	zero one limit old older old new limit finalL new old new comp limit	older old	00 03 06 08 10 12 14 16 18 20 22 25 28 30 32 33 34 35 36 37 38	3 2 2 2 2 2 2 2 2 2 3 3 2 2 1 1 1	011 011 01 01 01 01 01 01 011 011 01 00 0	13 33 35 13 34 36 12 38 08 36 03 35 02 36 07 37 06 38 01 30 08 37 13 36 35 13 37 36 00 10 08 38 11 00 01	40 43 46 48 50 52 54 56 58 60 62 65 68 70 72 73 74 75 76 77 78	13 73 75 13 74 76 12 78 08 76 03 75 02 76 07 77 06 78 01 70 08 77 13 76 75 13 77 76 00 50 08 78 11 00 01

#### 2.- Relocation maps (modification records)

Interleaving relocation bits with the program text makes cumbersome the process of loading the text directly into memory.

This problem can be resolve by collecting all relocation bits into a single contiguous relocation map that we will call the relocation section of the object code file (ELF).

The relocation section will be appended to the text and data sections.

The header will contain the entry point and length of the relocation section in the object module.

Program name: start
Starting address text
Length of text in bytes
Starting address data
Length of data in bytes
Starting address reloc. Sect.
Length of relocation section

Header

Text section

Data section

Relocation section