

# **COP 3402 Systems Software**

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## **Assemblers**

**Thanks to Euripides Montagne**

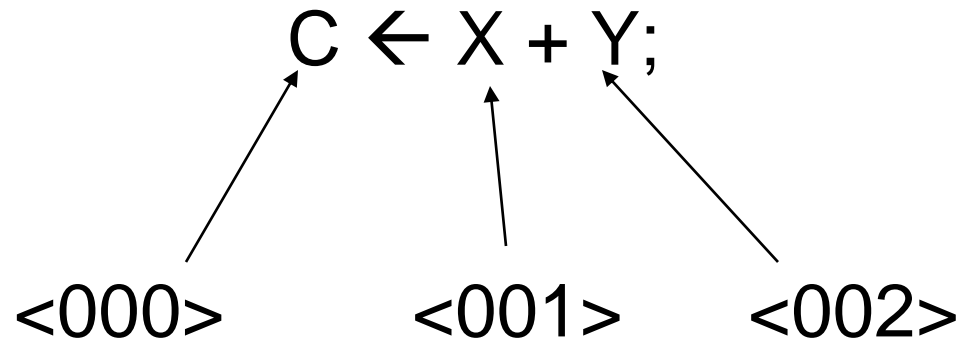
# ISA

## Instruction descriptions

opcode	mnemonic	meaning
0001	LOAD <x>	$A \leftarrow \text{Mem}[x]$
0010	ADD <x>	$A \leftarrow A + \text{Mem}[x]$
0011	STORE <x>	$\text{Mem}[x] \leftarrow A$
0100	SUB <x>	$A \leftarrow A - \text{Mem}[x]$
0101	IN <Device_#>	$A \leftarrow \text{read from Device}$
0110	OUT <Device_#>	$A \rightarrow \text{output to Device}$
0111	HALT	Stop
1000	JMP <x>	$\text{PC} \leftarrow x$
1001	SKIPZ	If Z = 1 Skip next instruction
1010	SKIPG	If G = 1 Skip next instruction
1011	SKIPN	If L = 1 Skip next instruction

# Assembly language Programming examples

Assign a memory location to each variable:



If necessary to use temporary memory locations, assign labels (names) to them.

# Assembly language

## Programming examples

Memory

000 1245

001 1755

002 0000

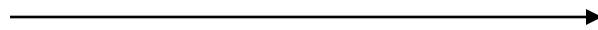
003 Load <000>

004 Add <001>

005 Store <002>

006 Halt

After execution



Memory

000 1245

001 1755

002 3000

003 Load <000>

004 Add <001>

005 Store <002>

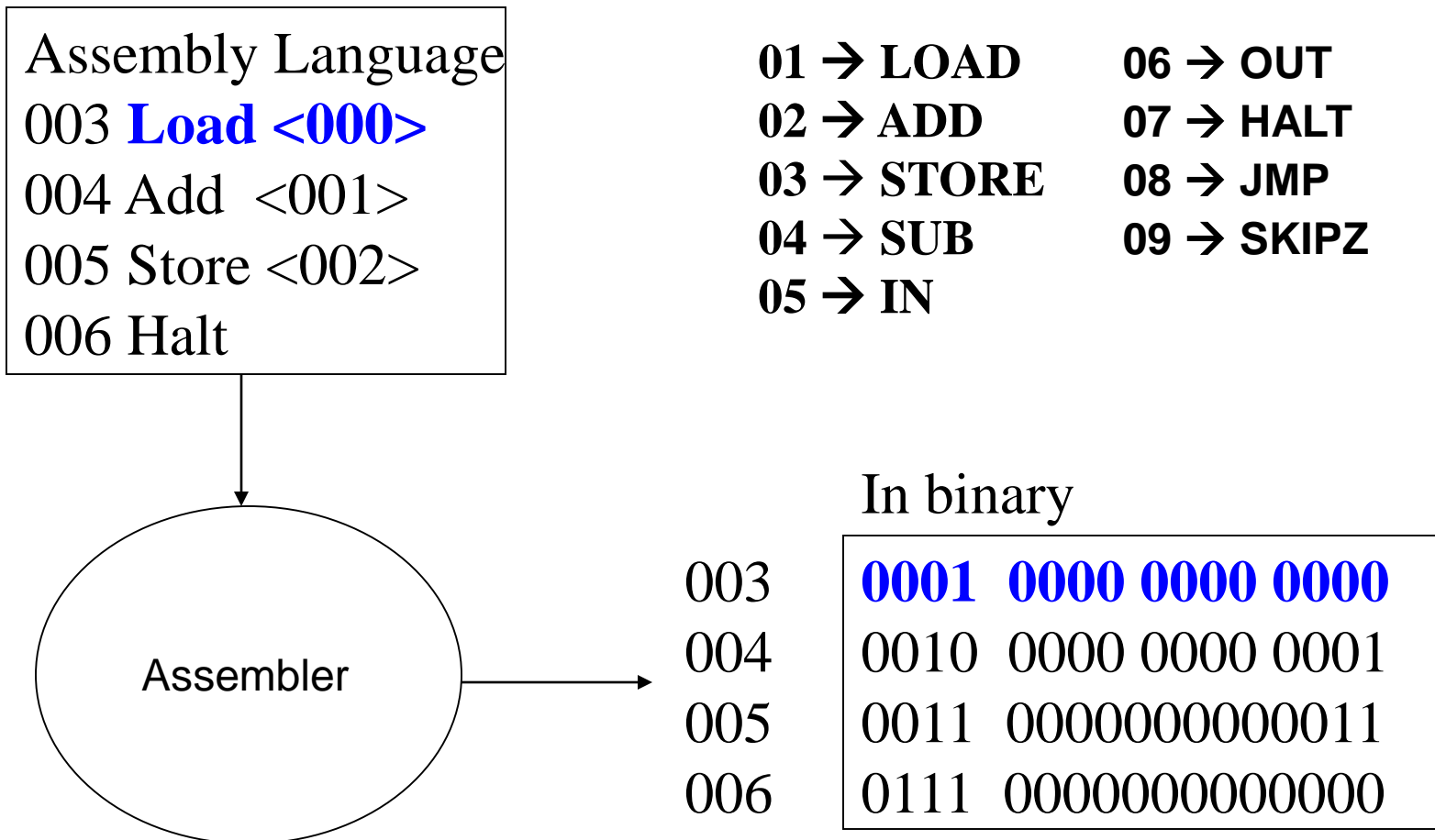
006 Halt

# One Address Architecture

- The instruction format of this one-address architecture consists of 16 bits: 4 bits to represent instructions and 12 bits for addresses :

OP	ADDRESS
0001	0000 0001 0001

# Assembler: translate Symbolic code to object code(binary)



# Assembler Directives

- The next step to improve our assembly language is the incorporation of pseudo-ops (assembler directives) to invoke a *special service from the assembler (pseudo-operations do not generate code)*

`.begin` → tells the assembler where the program starts

`.data` → to reserve a memory location.

`.end` → tells the assembler where the program ends.

**Labels** are symbolic names used to identify memory locations.

# Assembler Directives

This is an example of the usage of assembler directives

**.begin**

*“Assembly language instructions”*

**halt** *(return to OS)*

**.data** *(to reserve a memory location)*

**.end** *( tells the assembler where the program ends)*

note:

the directive **.end** can be used to indicate where the program Starts (for eample: “.end <insert label here>”



# Assembly language Programming

## Example 1

Label

opcode   address

start

.begin	
in	x005
store	a
in	x005
store	b
load	a
sub	TWO
add	b
out	x009
halt	

Text section (code)

a

b

TWO

.data	0
.data	0
.data	2
.end	start

Data section

# Assembly language Programming

## Example 2

	<u>Label</u>	<u>opcode</u>	<u>address</u>
01		; This is	
02		; a comment	
03	start	.begin	x200
04	here	LOAD	sum
05		ADD	a
06		STORE	sum
07		LOAD	b
08		SUB	one
09		STORE	b
0A		SKIPZ	
0B		JMP	here
0C		LOAD	sum
0D		HALT	
0E	sum	.data	x000
0F	a	.data	x005
10	b	.data	x003
11	one	.data	x001
12		.end	start

This program is computing  
5 x 3.

# ASSEMBLER

## Pass 1

	<u>Label</u>	<u>opcode</u>	<u>address</u>	
01		; This is		
02		; a comment		
03	start	.begin	x200	
04	here	LOAD	sum	x200
05		ADD	a	x201
06		STORE	sum	x202
07		LOAD	b	x203
08		SUB	one	x204
09		STORE	b	x205
0A		SKIPZ		x206
0B		JMP	here	x207
0C		LOAD	sum	x208
0D		HALT		x209
0E	sum	.data	x000	x20A
0F	a	.data	x005	x20B
10	b	.data	x003	x20C
11	one	.data	x001	x20D
12		.end	start	x20E

Symbol Table

here	x200
sum	x20A
a	x20B
b	x20C
one	x20D

symbol    address

In pass one the assembler examines the program line by line in order to built the symbol table.

There is an entry in the symbol table for each label found in the program.

# Opcode and Symbol Tables

Opcode table

opcode	mnemonic
0001	LOAD
0010	ADD
0011	STOR
0100	SUB
0101	IN
0110	OUT
0111	HALT
1000	JMP
1001	SKIPZ
1010	SKIPG
1011	SKIPN

Symbol Table

here	x200
sum	x20A
a	x20B
b	x20C
one	x20D

symbol address

Using the symbol table and the opcode table the assembler translates the program to object code.

As the program can be loaded anywhere in memory PC-relative addressing is used to resolve the symbols.

For instance, the offset between **LOAD sum** and the declaration of **sum** is 9, because when **LOAD sum** is fetched for execution, the pc is pointing to the instruction **ADD a**. ( pc + offset = 9)

# ASSEMBLER

## Pass 2

	<u>Label</u>	<u>opcode</u>	<u>address</u>
01		; This is	
02		; a comment	
03	start	.begin	x200
04	here	LOAD	sum
05		ADD	a
06		STORE	sum
07		LOAD	b
08		SUB	one
09		STORE	b
0A		SKIPZ	
0B		JMP	here
0C		LOAD	sum
0D		HALT	
0E	sum	.data	x000
0F	a	.data	x005
10	b	.data	x003
11	one	.data	x001
12		.end	start

		<u>Object code</u>
	x200	000100000000 <b>1001</b> (9 is the offset)
PC →	x201	0010000000001001
	x202	001100000000 <b>0111</b> (7 is the offset)
	x203	0001000000001000
	x204	0100000000001000
	x205	0011000000000110
	x206	
	x207	
	x208	All addresses are
	x209	pc-relative addresses.
	x20A	(PC + offset)
	x20B	
	x20C	
	x20D	Recall: PC is always pointing to
	x20E	the next instruction to be fetch.

offset

# Assembly language Programming object code

	<u>Label</u>	<u>opcode</u>	<u>address</u>
01		; This is	
02		; a comment	
03	start	.begin	x200
04	here	LOAD	sum
05		ADD	a
06		STORE	sum
07		LOAD	b
08		SUB	one
09		STORE	b
0A		SKIPZ	
0B		JMP	here
0C		LOAD	sum
0D		HALT	
0E	sum	.data	x000
0F	a	.data	x005
10	b	.data	x003
11	one	.data	x001
12		.end	start

		<u>Object code</u>
	x200	0001000000001001 (9 is the offset)
PC →	x201	0010000000001001
	x202	0011000000000111 (7 is the offset)
	x203	0001000000001000
	x204	0100000000001000
	x205	0011000000000110
	x206	1001000000000000
	x207	1000111111111000 (-7)
	x208	0001000000000001
	x209	0111000000000000
	x20A	0000000000000000
	x20B	0000000000000101
	x20C	0000000000000011
	x20D	0000000000000001
	x20E	

offset

One's complement

# ASSEMBLER

## object code

The object code file has several sections:

Header section: Size of code, name source file, size of data

Text section (code): Object code

Data section: Data (in binary)

Relocation information section: Addresses to be fixed up by the linker

Symbol table section: Global symbols in the program, Imported symbols

Debugging section: Source file and line number information, description of data structures.

# ASSEMBLER

## object code file for the example

Program name: start  
Starting address text: x200  
Length of text in bytes: x14  
Starting address data: x20A  
Length of data in bytes: 8

Header

0001000000001001  
0010000000001001  
0011000000000111  
0001000000001000  
0100000000001000  
0011000000000110  
1001000000000000  
1000111111111000  
0001000000000001  
0111000000000000

Text section

0000000000000000  
0000000000000101  
0000000000000011  
0000000000000001

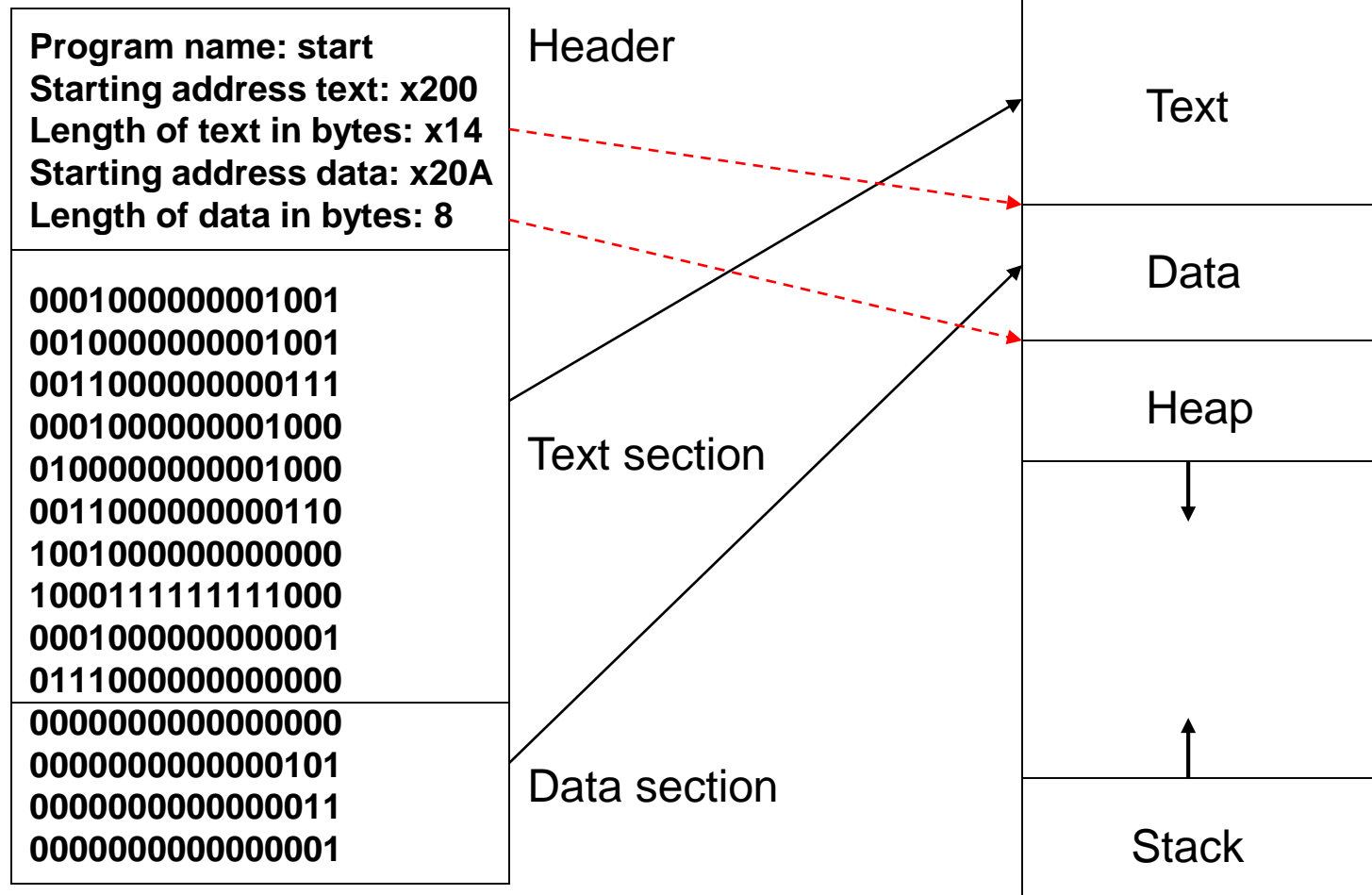
Data section



# Loading Object code in Memory

Object code file (disk)

Run time environment



# UNIX a.out format

a.out object code format

a.out header
text section
data section
symbol table information
relocation Information

a.out header

**magic number**  
text segment size  
initialized data size(data)  
uninitialized data size(**bss**)  
symbol table size  
**entry point**  
text relocation size  
data relocation size

**a.out stands for "assembler output".**

**magic number indicates type of executable file.**

**bss is an acronym for block storage start.**

**entry point: starting address of the program**