# COMP 312 Assignment 2

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## 1 Python

#### 1.1 Dice Problem

#### 1.1.1 Part A

```
import random
random.seed(123)
def roll_die(n):
        """ Rolls_n_die_and_returns_the_results_as_array"""
        rolls = [];
        for i in range (0, n):
                rolls.append(random.randint(1, 6))
        return rolls;
if __name__ == '__main__':
        n = 1000000
        events = 0
        for i in range (0, n):
                o = False
                for j in range (0, 24):
                         roll = roll_die(2)
                         if roll[0] = 6 and roll[1] = 6:
                                 o = True
                if o:
                         events += 1
        p = float (events)/float (n)
        print p
```

The output of this program is 0.491206 so yes the probability of the event occurring is less than  $\frac{1}{2}$ . The value of n was very large so we can be confident in this finding.

#### 1.1.2 Part B

```
import random
random.seed(123)
```

```
def roll_die(n):
        """ Rolls _n_die _and_returns _the _results _as _array """
        rolls = [];
        for i in range (0, n):
                 rolls.append(random.randint(1, 6))
        return rolls;
if = name_{-} = ' = main_{-}':
        n = 1000000
        min_num_throws = 24
        max_num_throws = 26
        events = \{\}
        for i in range (0, n):
                 for j in range(min_num_throws, max_num_throws + 1):
                          o = False
                          events.setdefault(j, 0)
                          for k in range (0, j):
                                  roll = roll_die(2)
                                  if roll [0] = 6 and roll [1] = 6:
                                           o = True
                          if o:
                                  events[j] += 1
        for i in range(min_num_throws, max_num_throws+1):
                 events [i] = float (events [i]) / float (n)
        print events
```

The modified code calculates the probability of the event occurring with a different number of throws. Over the range 24 to 26 it outputs the the following  $\{24:0.491443,25:0.506207,26:0.518883\}$ 

From this we see that it requires 25 throws for the probability to be  $> \frac{1}{2}$  but as close to  $\frac{1}{2}$  as possible

#### 1.2 Optimized Table Look-up

 $\mathbf{import} \hspace{0.2cm} \mathrm{random}$ 

```
\mathbf{def} tablelookup(y,p):
        "" Sample_from_y[i]_with_probabilities_p[i]""
        u = random.random()
        sumP = 0.0
        b = 0
         for i in range(len(p)):
                 sumP += p[i]
                 b += 1
                 if u < sumP:
                          return { 'val ': y[i], 'count ': b}
def run_lookup(y, p):
        """Runs_a_lookup_for_given_probabilities"""
        m = 1000000
        b = 0.0
         valuetotal = 0.0
         for k in range(m):
                 d = tablelookup(y,p)
                 b += d['count']
                 valuetotal += d['val']
         print valuetotal/m
        print b/m
        print "\n"
random. seed (123)
y1 = [0, 1, 2, 3, 4, 5]
p1 = \begin{bmatrix} 1.0/1024, & 15.0/1024, & 90.0/1024, & 270.0/1024, & 405.0/1024, & 243.0/1024 \end{bmatrix}
y2 = [5,4,3,2,1,0]
p2 = [243.0/1024, 405.0/1024, 270.0/1024, 90.0/1024, 15.0/1024, 1.0/1024]
y3 = [4,3,5,2,1,0]
p3 = [405.0/1024, 270.0/1024, 243.0/1024, 90.0/1024, 15.0/1024, 1.0/1024]
print "Part_I"
run_lookup(y1, p1)
print "Part_II"
```

```
run_lookup(y2, p2)
print "Part_III"
run_lookup(y3, p3)
The output from running this code is
```

Part I Part II Part III 3.750354 3.750114 3.750704 4.750354 2.249886 2.064436

As we can see from this all results are mostly the same. We also see that option 2 and 3 take about half as many steps as option 1. However option 3 is just slightly faster

### 2 Queues

#### 2.1 Problem 3

We know that the service times are distributed exponentially. And we assume that all the servers have equal average service times, let this time be m. Once one of the other services has finished you can begin being served (you are now 'racing' against the remaining 6 people). As the exponential distribution has the property of being memory-less all the services have the same probability of taking time m. This means that  $\frac{6}{7}$  probability of finishing before atleast one of the other customers being served.

#### 2.2 Problem 4

 $state \leftarrow number of active pizza shops$ 

For this problem I used python to compute the steady state vector. The following code is what I used to do this.

```
coeffs = {}
# Compute the coefficients
l = 1
u = 1

for i in range(1, 32 + 1):
    price = max(0, 16 - 0.5 * (i-1))

l = l * price
u = u * (i * 1.0/(10 + price))

coeffs[i] = float(l)/float(u)
```

```
# Compute pi 0
pi0 = 0
s = 1
for i in range (1, 32 + 1):
        s = s + coeffs[i]
pi0 = 1/s
pies = \{\}
pies[0] = pi0
# Compute the rest of the pis
for i in range (1, 32 + 1):
        pies[i] = coeffs[i] * pi0
# Ensure that the sum of pies is 1
piSum = 0
for i in range (32 + 1):
        piSum += pies[i]
if not piSum == 1.0:
        raise Exception("Sum_of_pies_should_be_1")
\#\!\!/\!\!/ Solve~Assignment~Problems~\#\!\!/\!\!/
# Compute average number of pizza restrants
avg = 0
for i in range (32 + 1):
        avg += i * pies[i]
print "Average_Num_Pizza_Shops:_" + str(avg)
# Compute fraction of time more than 20 restrants
timeM20 = 0.0
for i in range (20 + 1, 32 + 1):
        timeM20 += pies[i]
print "Fraction_of_time_with_more_than_20_restrants:_" + str(timeM20)
```

To get the average state of the system we take the following sum

$$\sum_{n=0}^{32} n\pi_n$$

for which we get the following result

 $Avg\ number\ of\ pizza\ shops:\ 27.6082698882$ 

To get the fraction of time spent in state i we know this value is  $\pi_i$  so if we want to find the fraction of time with 20 or more restaurants this is

$$\sum_{n=20}^{32} \pi_n = 0.999741852504 \tag{1}$$

 $Fraction\ of\ time\ with\ >\ 20\ restrants:\ 0.999741852504$