

Library for reasoning on nominal algorithms in Coq

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1 Introduction

This library forms a basis for reasoning on nominal algorithms in the proof assistant Coq [4]. The source files are available as a Coq contribution (see <http://coq.inria.fr>).

As proposed by Zen [1, 2] we interpret probabilistic programs as measure transformers; the main contribution of this paper is to view this interpretation as a modal transformation on functional names. Using this

2.2 Specification of U

- Constants: 0 and 1
- Constructor: $U(n) = \frac{1}{n}$

It follows that

Hypothesis $U_{q_d}(u) \leq n \implies y \in U_{\neg \neg} \implies y \implies y$.

Hypothesis $U_{\neg t}(a) \implies y \in U(\neg y)$

3.2. Properties of

■ `mma U_ = u t_ri t ; y (× y)`.

■ `mma U_ = u t_ ft ; y (× y y)`.

Hint `R s v U_ = u t_ri t U_ = u t_ ft`.

■ `mma Uinv_ = r ri t ; y U (([] y)) y ([])`.

Hint `R s v Uinv_ = r`

3. Definition of X^n

Let point $U = (U) \in \text{nat} \times \{structure\}$ in $U =$
 $at \text{ } n \text{ } wit \text{ } O \text{ } (S) \times U \text{ } nd.$

3.9 Definition and properties of $X \& Y$

Assuming that the operation which is defined with `min` and `mult` in 0 and 1 see `M` in `an` & `M` in `en`

Hint $R_s v_{si} a_{in} rR$

■ $mmaR$ — —

■ `mma` $n_{su} \vdash n_s q U su$.

Hint $R_s v \vdash n_{su}$.

■ `mma` $u \vdash q_{us} \vdash u$ (function

aU .y t

$a \quad F330.522. \quad 3 \quad 0.5.3Tf35.0TdTJF3.3Tf \quad 0.5 \quad .5dDay \quad Uy3uL \quad q_t \quad at_ri \quad t$
 $(U \quad (U \quad u \quad t \quad ? \quad ?y) \quad (U \quad u \quad t \quad ? \quad ? \quad)) \quad a \quad y \quad U \quad u \quad t \quad - \quad - \quad at_ft$
 $(U \quad (U \quad u \quad t \quad ? \quad ? \quad) \quad (U \quad u \quad t \quad ?y \quad ? \quad)) \quad a \quad y \quad U \quad u \quad t \quad - \quad - \quad at_ri \quad t$
 $(U \quad (U \quad u \quad t \quad ? \quad ? \quad) \quad (U \quad u \quad t \quad ? \quad ?y)) \quad s \quad t \quad id \quad -r \quad writ \quad (U \quad u \quad t \quad -sy \quad y);$
 $a \quad y \quad U \quad u \quad t \quad - \quad - \quad at_ri \quad t$
 $(U \quad (U \quad u \quad t \quad ? \quad ?y) \quad (U \quad u \quad t \quad ? \quad ? \quad)) \quad s \quad t \quad id \quad -r \quad writ \quad (U \quad u \quad t \quad -sy \quad y);$
 $a \quad y \quad U \quad u \quad t$
 $- \quad q- \quad at_ft$
 $(U \quad q \quad (U \quad u \quad t \quad ? \quad ? \quad) \quad (U \quad u \quad t \quad ?y \quad ? \quad)) \quad a \quad y \quad U \quad u \quad t \quad - \quad q- \quad at_ri \quad t$
 $(U \quad q \quad (U \quad u \quad t \quad ? \quad ? \quad) \quad (U \quad u \quad t \quad ? \quad U \quad u \quad t \quad ? \quad ? \quad (U \quad q \quad (U \quad u \quad t \quad ? \quad ?y) \quad (U \quad u \quad t \quad ? \quad ? \quad)) \quad s \quad t \quad id$
 $- \quad - \quad at_ri \quad t$
 $(U \quad q \quad (U \quad u \quad t \quad ? \quad ?y) \quad (U \quad u \quad t \quad ? \quad ? \quad)) \quad a \quad y \quad U \quad u \quad t$

3.2 Expected properties of measure

Definition $n \leq m$ (Ty) (\vdash) Rr \vdash =

4. Operations on distributions

Definition 4.1. $\text{distr} (\text{Type}) (\text{distr}) := f (u \ f) (u \ 2 \ f)$.

Definition 4.2. $\text{q_distr} (\text{Type}) (\text{distr}) := f (u \ f) == (u \ 2 \ f)$.

Lemma 4.3.

Figure 1: Composition of Distribution

Required Report

■ mm
■ mma
End *Fi* *ints.*

■ mm
■ mma

End *Fi* *ints.*

Phiro . : An

7.1 Definition of connectivity

$p \models [e](q)$ is defined as $p \models \mu(\varphi)$

Definition $k(Ty)(U)$

D e f i n i t i o n $kfun \left(\begin{array}{c} \vdash Ty \\ \vdash U \end{array} \right) (\vdash U)$
 $\quad \equiv \quad \vdash \left(k \left(\begin{array}{c} \vdash Ty \\ \vdash U \end{array} \right) q \right).$

7.2 Identifications

$$(Ty)_{q,q'} = (U)_{q,q'} \text{ (distr)} (q, U)_{q,q'}$$

7.3.3 Rule for conditional

$$p_1 \vdash [e_1]($$

$$\sin \alpha = \frac{n_{\text{ist}}}{\text{add}} \quad \sin \alpha = \frac{n_{\text{ist}}}{n_{\text{ist}}}$$

Definition $d(n_{\text{ist}}) =$

$$\text{at } \text{wit}(\sin \alpha) = a \quad (\text{add } a) = a \quad \text{nd.}$$

■ point $a = (\text{wit}(\sin \alpha) \{ \text{strut} \}) \in n_{\text{ist}} =$

$$\text{at } \text{wit}(\sin \alpha) = \text{add } a \quad (\text{add } a) = \text{add } a \quad (\text{add } a) = \text{nd.}$$

■ point $r_{v_a} = (\text{wit}(\sin \alpha) \{ \text{strut} \}) \in n_{\text{ist}} =$

$$\text{at } \text{wit}(\sin \alpha) = \text{add } a \quad (\text{add } a) = r_{v_a} \quad (\text{add } a) = \text{nd.}$$

Definition $5.3Tf3.550T \quad Tjf3.3Tf2.0Tn_{\text{ist}}Tjf \quad 5.3Tf2 \quad 0 \quad 0Td333Tjf3.3Tf \quad .3 \quad 0Td \quad Tj$

$$\text{at } \text{wit}(\sin \alpha)$$

3.2.2 Assumption

We need extra premises on p_1 , p_2 and *choice*.

- p_1 and p_2 terminate with probability 1
- Q alive

Hint $R[s, v] = \dots$

■ `mma.i`