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## CS 132 Homework 5

### 1. Random Access Protocols

#### A. Analysis of Slotted Aloha: Problem 11

- i.  $[1 - (p(1-p)^3)]^4 \cdot p(1-p)^3$
- ii.  $p(1-p)^3 \cdot 4$
- iii.  $[1 - 4p(1-p)^3]^2 \cdot 4p(1-p)^3$
- iv.  $4p(1-p)^3$

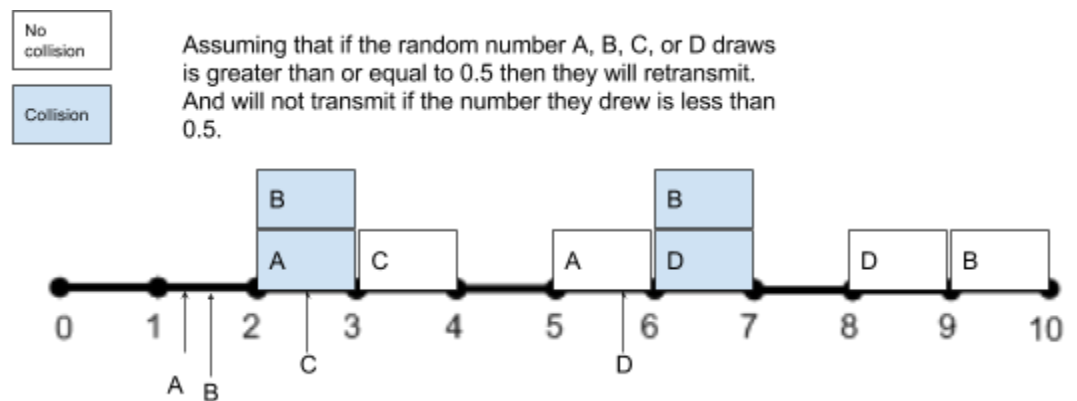
#### B. Comparing Protocols

- i. No these transmissions could not have been generated by Pure Aloha because the data would transmit as soon as it arrives.
- ii. Yes these transmission could have been generated by Slotted Aloha because M1 and M2 arrive and transmit in the same time slot and cause a collision. M3 then sends successfully. M1 has a better probability than M2 so it attempts to retransmit first and does so successfully. Then right after M2 and M4 collide with each other. Since m4 had a better probability it will successfully transmit. Then finally M2 will transmit and all four have been delivered.

- iii. No these transmissions could not have been generated by CSMA because as soon as  $M_1$  arrives it would detect a clear line and start transmitting.
- iv. No these transmissions could not have been generated by CSMA/CD for the same reason as it could not have been generated by CSMA,  $M_1$  would send the data when it arrives because the line was free.
- v. No these transmissions could not have been generated by an algorithm with exp. back off because  $M_1$  and  $M_2$  are colliding and would have to start transmitting data at either time = 12 or time = 16.

### C. Slotted Aloha in Action

i.

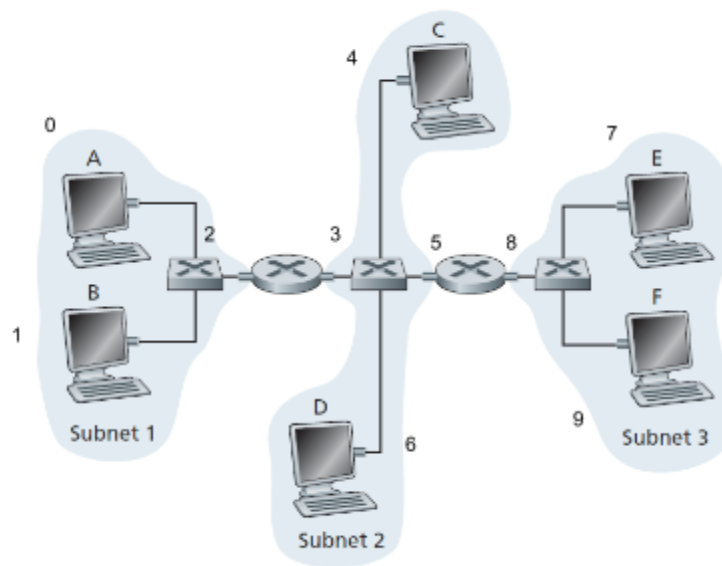


ii.

1. In order to maximize efficiency I would retransmit if the probability was less than or equal to 0.37 and wouldn't if it was greater than 0.37.
2. In order to maximize fairness I would keep the probability at greater than or equal to 0.5 would retransmit and would not transmit if less than 0.5.

## 2. LANs: Addresses and Switches

### A. MAC and IP addresses: Problem 21



- i. Source MAC: 00-00-00-00-00-00  
Destination MAC: 00-00-00-00-00-02  
Source IP: 1.1.1.0  
Destination IP: 1.1.1.9
- ii. Source MAC: 00-00-00-00-00-03

Destination MAC: 00-00-00-00-00-05

Source IP: 1.1.1.0

Destination IP: 1.1.1.9

iii. Source MAC: 00-00-00-00-00-08

Destination MAC: 00-00-00-00-00-09

Source IP: 1.1.1.0

Destination IP: 1.1.1.9

#### B. Learning Switches: Problem 26

i. B sends a frame to E

1. Switch table adds the MAC address of B
2. Link(s) on which the transmitted frame will be forwarded:  
A,C,D,E,F
3. Switch table is empty and does not know the MAC address of E so it sends packets to every link.

ii. E replies with a frame to B

1. Switch table adds the MAC address of E
2. Link(s) on which the transmitted frame will be forwarded: B
3. Switch table already knows MAC address of B so it just updates the MAC address of E.

iii. A sends a frame to B

1. Switch table adds the MAC address of A
2. Link(s) on which the transmitted frame will be forwarded: B

3. Switch table already knows the MAC address of B so it just updates the MAC address of A.

iv. B replies with a frame to A

1. Switch table doesn't change from before
2. Link(s) on which the transmitted frame will be forwarded: A
3. Switch table already knows the MAC addresses for A and B so it doesn't need to update.