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CS 161

## Homework 5

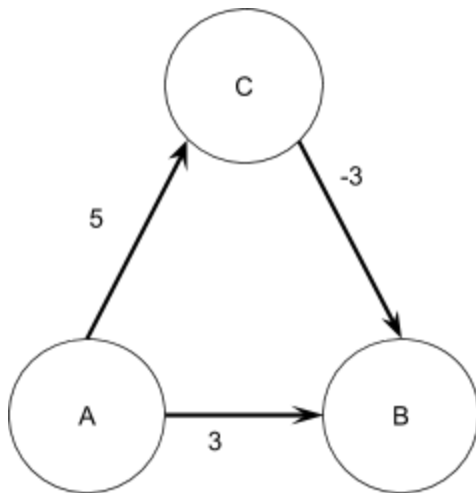
### R-13.4

LA15, LA22, LA16, LA31, LA32, LA127, LA141, LA126, LA169

### R-13.12

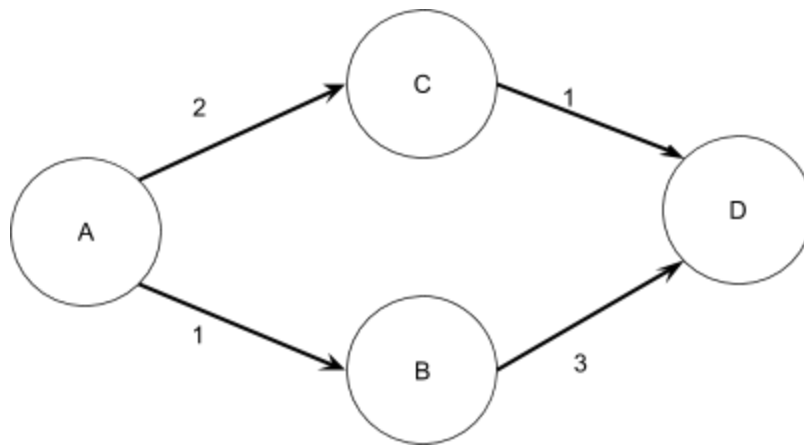
$A \rightarrow E$ ,  $E \rightarrow F$ ,  $I \rightarrow M$ ,  $M \rightarrow N$ ,  $N \rightarrow K$ ,  $K \rightarrow O$ ,  $O \rightarrow P$ ,  $P \rightarrow L$ ,  $L \rightarrow H$ ,  $H \rightarrow D$ ,  $D \rightarrow C$ ,  $C \rightarrow B$ ,  $B \rightarrow F$ ,  
 $C \rightarrow G$ ,  $G \rightarrow J$

### C-14.2



In the example above, A is the starting vertex and B is the ending vertex. Dijkstra's algorithm will compute that the shortest path from A  $\rightarrow$  B is 3. When in actuality the shortest path would be A  $\rightarrow$  C  $\rightarrow$  B which would come out to 2.

### C-14.4



In the example above A is the starting vertex and D is the goal vertex. According to the greedy strategy it will first add B to the path because it is the minimum weight edge and then will make B = start. The minimum weight edge from B to D is 3 so D is added to the path and then D = start and the greedy strategy will exit since start = goal. This means that the greedy algorithm is wrong because there is clearly a shorter path to D from A (A->C->D) however there is no method to backup and make provisions.