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## CS 171 Homework 5

### Problem 8.8

No,  $\neg \text{Spouse}(\text{George}, \text{Laura})$  does not follow from the facts  $\text{Jim} \neq \text{George}$  and  $\text{Spouse}(\text{Jim}, \text{Laura})$ . The additional axiom needed to prove it would be:  $\text{Spouse}(x, y)$  and  $(z \neq x) \Rightarrow \neg \text{Spouse}(z, y)$ . If we use  $\text{Spouse}$  as a unary function symbol instead of a binary predicate then no additional axiom is needed.

### Problem 8.10

- a.  $\text{Occupation}(\text{Emily}, \text{Surgeon}) \vee \text{Occupation}(\text{Emily}, \text{Lawyer})$
- b.  $\exists o \neq \text{actor} \wedge \text{Occupation}(\text{Joe}, \text{Actor}) \wedge \text{Occupation}(\text{Joe}, o)$
- c.  $\forall p \text{Occupation}(p, \text{Surgeon}) \rightarrow \text{Occupation}(p, \text{Doctor})$
- d.  $\neg \exists p \text{Occupation}(p, \text{Lawyer}) \wedge \text{Customer}(\text{Joe}, p)$
- e.  $\exists p \text{Occupation}(p, \text{Lawyer}) \wedge \text{Boss}(p, \text{Emily})$
- f.  $\exists p1 \text{Occupation}(p1, \text{Lawyer}) \wedge \forall p2 \text{Customer}(p2, p1) \rightarrow \text{Occupation}(p2, \text{Doctor})$
- g.  $\forall p1 \text{Occupation}(p1, \text{Surgeon}) \rightarrow \exists p2 \text{Occupation}(p2, \text{Lawyer}) \wedge \text{Customer}(p1, p2)$

### Problem 8.28

- b.  $\neg \text{Wrote}(\text{Gershwin}, \text{Eleanor Rigby})$
- d.  $\exists s \text{Wrote}(\text{Joe}, s)$
- g.  $\neg [\exists s \text{Wrote}(\text{Gershwin}, s) \wedge \exists p \text{Sings}(p, s, \text{Revolver})]$
- h.  $\forall s \text{Wrote}(\text{Gershwin}, s) \rightarrow \exists p, a \text{Sings}(p, s, a)$
- k.  $\forall a [\exists s \text{Sings}(\text{McCartney}, s, a)] \rightarrow \exists d \text{CopyOf}(d, a) \wedge \text{Owns}(\text{Joe}, d)$

### Problem 9.9

- 1.  $0 \leq 3$ .

2.  $7 \leq 9$ .
3.  $\forall x x \leq x$ .
4.  $\forall x x \leq x + 0$ .
5.  $\forall x x + 0 \leq x$ .
6.  $\forall x, y x + y \leq y + x$ .
7.  $\forall w, x, y, z w \leq y \wedge x \leq z \Rightarrow w + x \leq y + z$ .
8.  $\forall x, y, z x \leq y \wedge y \leq z \Rightarrow x \leq z$

a. Start with  $7 \leq 3 + 9$  using Rule 8 you're left with  $7 + 0 \leq 3 + 9$

Using Rule 6 gives  $9 + 3 \leq 3 + 9$

Using Rule 7 gives  $7 + 0 \leq 9 + 3$

Using Rule 1 gives  $0 \leq 3$

Using Rule 2 gives  $7 \leq 9$

b. Using Rule 7, and  $w = 0, x = 7, y = 3, z = 9$  gives  $0 + 7 \leq 3 + 9$

Rule 6 gives  $9 + 3 \leq 3 + 9$

Rule 4 gives  $7 \leq 7 + 0$

Rule 8 and  $x = 0 + 7, y = 7 + 0, z = 3 + 9$  gives  $7 + 0 \leq 3 + 9$

Strip away the 0 with Rules 5 and 8 and you're left with  $7 \leq 3 + 9$

### Problem 9.20

- a.  $\exists p \forall q \text{ person}(p) \wedge \text{person}(q) \wedge ((\neg S(q, q) \leftrightarrow S(p, q)))$
- b.  $\text{person}(p) \wedge \text{person}(q) \wedge (S(q, q) \vee S(p, q)) \wedge (\neg S(p, q) \vee \neg S(q, q))$
- c. The clausal form resolves to empty clause meaning the logic is not satisfiable because it is false.

### Problem 13.8

	<i>toothache</i>		$\neg$ <i>toothache</i>	
	<i>catch</i>	$\neg$ <i>catch</i>	<i>catch</i>	$\neg$ <i>catch</i>
<i>cavity</i>	0.108	0.012	0.072	0.008
$\neg$ <i>cavity</i>	0.016	0.064	0.144	0.576

**Figure 13.3** A full joint distribution for the *Toothache*, *Cavity*, *Catch* world.

a.  $P(\text{Toothache}) = 0.108 + 0.012 + 0.016 + 0.064 = 0.2$

b.  $P(\text{Cavity}) = 0.108 + 0.012 + 0.072 + 0.008 = 0.2$

c.  $P(\text{Toothache} \mid \text{Cavity})$

$$= \frac{P(\text{toothache} \wedge \text{cavity})}{p(\text{cavity})} = \frac{0.108 + 0.012}{0.2} = 0.6$$

d.  $P(\text{Cavity} \mid \text{toothache} \vee \text{catch})$

$$= \frac{P(\text{cavity} \wedge (\text{toothache} \vee \text{catch}))}{P(\text{toothache} \vee \text{catch})} = \frac{0.108 + 0.012 + 0.072}{0.108 + 0.012 + 0.072 + 0.016 + 0.064 + 0.144} = 0.46$$

### Problem 13.13

$$P(\text{Test A positive} \mid \text{virus is present}) = 0.95$$

$$P(\text{Test A positive} \mid \text{virus is not present}) = 0.1$$

$$P(\text{Test B positive} \mid \text{virus is present}) = 0.9$$

$$P(\text{Test B positive} \mid \text{virus is not present}) = 0.05$$

$$P(\text{Virus is present}) = 0.01$$

$$P(\text{Virus is absent}) = 0.99$$

$$P(\text{Virus is present} \mid \text{Test A positive})$$

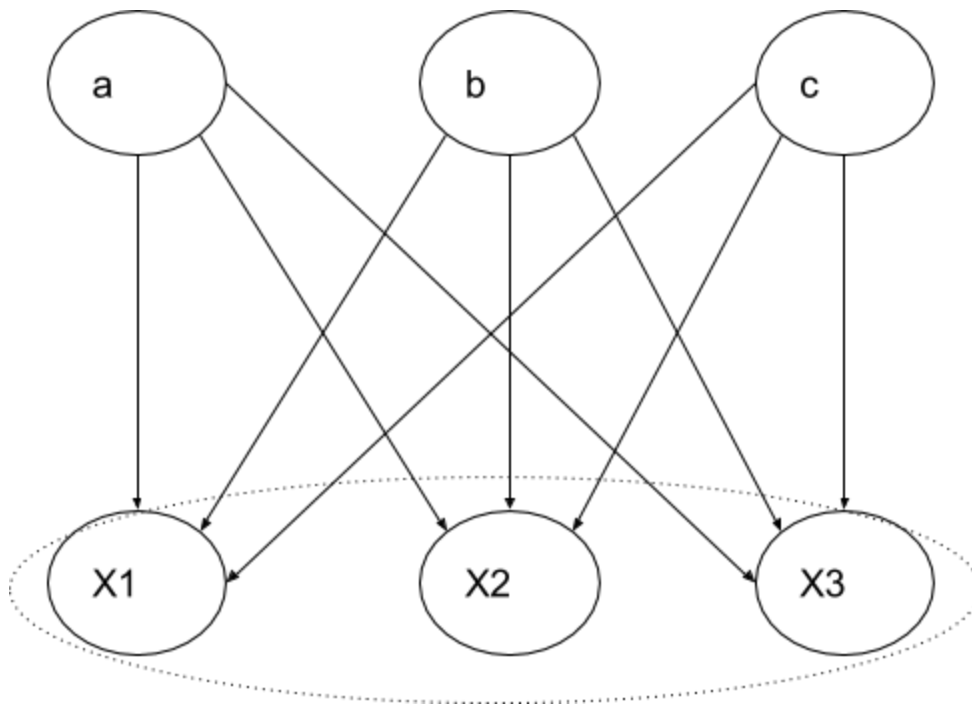
$$= \frac{0.95 * 0.01}{0.95 * 0.01 + 0.1 * 0.99} = 0.088$$

$$P(\text{Virus is present} \mid \text{Test B positive})$$

$$= \frac{0.9 * 0.01}{0.9 * 0.01 + 0.05 * 0.99} = 0.15$$

Test B is more indicative of someone really carrying the virus.

### Problem 14.1



a.

b.  $P(a \mid H, H, T) = 0.2 * 0.2 * 0.8 * 0.333 = 0.0107$

$$P(b \mid H, H, T) = 0.6 * 0.6 * 0.4 * 0.333 = 0.048$$

$$P(c \mid H, H, T) = 0.8 * 0.8 * 0.2 * 0.333 = 0.0427$$

Therefore coin b is the mostly like to be drawn from the bag.

### Problem 14.15

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function ELIMINATION-ASK( $X, e, bn$ ) returns a distribution over  $X$ 
inputs:  $X$ , the query variable
           $e$ , observed values for variables  $E$ 
           $bn$ , a Bayesian network specifying joint distribution  $P(X_1, \dots, X_n)$ 

 $factors \leftarrow []$ 
for each  $var$  in ORDER( $bn.VARS$ ) do
     $factors \leftarrow [MAKE-FACTOR(var, e) | factors]$ 
    if  $var$  is a hidden variable then  $factors \leftarrow SUM-OUT(var, factors)$ 
return NORMALIZE(POINTWISE-PRODUCT( $factors$ ))

```

**Figure 14.11** The variable elimination algorithm for inference in Bayesian networks.

- a.  $P(\text{Burglary} \mid \text{JohnCalls} = \text{true}, \text{MaryCalls} = \text{true}) = \langle 0.284, 0.716 \rangle$
- b. The number of arithmetic operations was  $2 * 11 + 3 = 25$ . While using the enumeration algorithm there are  $2 * (16 + 3) + 3$  operations giving a total of 41.
- c. The enumeration algorithm would take  $O(n * 2^{n-2})$  time. Whereas variable elimination would have a running time of  $O(n)$ .
- d. If one were to ignore the the direction of the edges in the polytree network, then we would be left with an undirected graph  $G$  with  $V$  being our variable elimination ordering. When a node is eliminated from the graph, it has no children left in the ordering  $V$ . Therefore the only remaining neighbor is the parent of the node. Thus the factor corresponding to that node will be a function of the parent alone and can be computed in linear time.