

## Lecture 12

# Simple Calculator

# Simple Calculator

Evaluate:

- $1 + 2 * 3$
- $10 - 4 - 3$
- $2^3^3$
- $4/2/2$

We will use integer math for our examples

# What about?

1 - 2 - 4 ^ 5 \* 3 \* 6 / 7 ^ 2 ^ 2

# Postfix Notation

- **infix** notation -- Humans generally write expressions with the operators between the operands,  
as in  $2 * 3 + 4$
- **postfix** notation -- notation in which the operators are put after their operands,  
as in  $2 3 * 4 +$

This notation is preferred by computers

- With “**infix**” (the method you are used to) :  
you put operators between operands:  $a + b$
- With “**postfix**” (the method computer prefers) :  
you put operators after operands:  $a b +$

# Postfix Notation

- Why Postfix Notation is preferred by the computer:
  - It is the most efficient method for representing arithmetic expressions
  - There is never any need to use ()'s with postfix notation and there is never any ambiguity
- To evaluate an infix expression, the compiler:
  - Converts the infix expression to postfix form
  - Evaluates the postfix expression

# Application of stacks: Evaluating Postfix Expression

- **Stacks** are used by compilers to help in the process of evaluating expressions.
- Steps for evaluating postfix expressions

Make an empty stack **s**

For each token (operator \* - + / or single digit integers 0, 1, ..., 9) in the postfix expression:

if operand

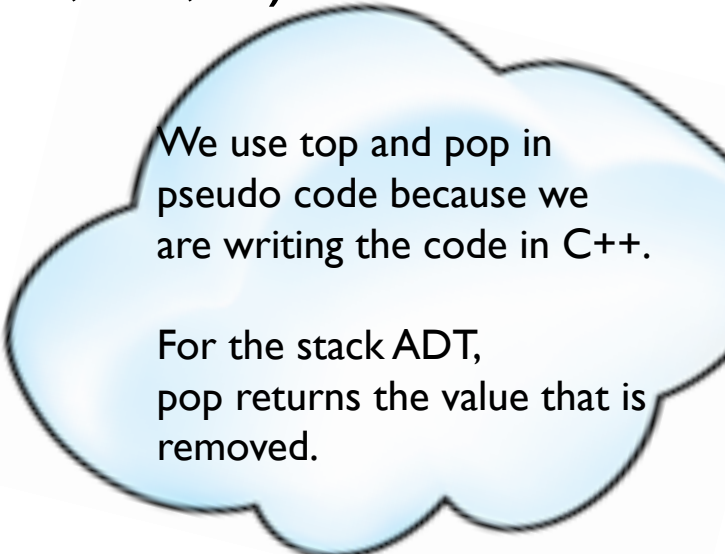
**s.push(operand);**

if operator

right = **s.top(); s.pop();**

left = **s.top(); s.pop();**

push the value of the operator applied to the left and right



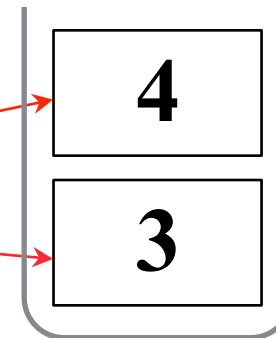
We use top and pop in pseudo code because we are writing the code in C++.

For the stack ADT, pop returns the value that is removed.

# Example:

Evaluate 3 4 - 5 3 \* -

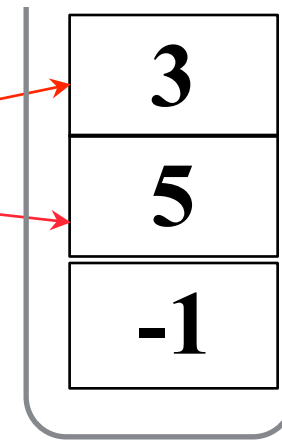
```
s.push(3);  
s.push(4);  
// found operator - so pop  
right = s.top(); s.pop();  
left = s.top(); s.pop();  
s.push(left - right);  
// 3 - 4
```



- The stack now has one value -1
- The remainder of the expression: 5 3 \* -

# Continue with 5 3 \* -

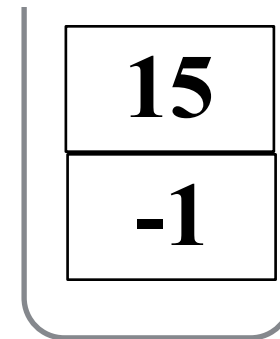
```
s.push( 5 );  
s.push( 3 );  
right = s.top(); s.pop();  
left = s.top(); s.pop();  
s.push(left*right);  
// 5 * 3
```



- The Stack has 2 values
- Only one token remains



# Continue with –



```
right = s.top(); s.pop();    // found operator -  
left  = s.top(); s.pop();  
s.push(left-right);  
      -1 - 15
```

- The expression has been processed.
- The value at the top of the stack is the value of the expression is -16
- Now evaluate 2 3 4 \* 5 \* - ???

# Evaluate the Postfix expressions

1 2 3 \* +

10 4 - 3 -

2 3 3 ^ ^

4 2 / 2 /

# Running Time?

- **Linear** in the input size.

```
enum TokenType { EOL, VALUE, OPAREN, CPAREN, EXP,  
                MULT, DIV, PLUS, MINUS };
```

```
// PREC_TABLE matches order of Token enumeration
```

```
struct Precedence
```

```
{
```

```
    int inputSymbol;
```

```
    int topOfStack;
```

```
};
```

```
vector<Precedence> PREC_TABLE =
```

```
{
```

```
    { 0, -1 }, { 0, 0 },    // EOL, VALUE
```

```
    { 100, 0 }, { 0, 99 }, // OPAREN, CPAREN
```

```
    { 6, 5 },              // EXP
```

```
    { 3, 4 }, { 3, 4 },    // MULT, DIV
```

```
    { 1, 2 }, { 1, 2 }     // PLUS, MINUS
```

```
};
```

# Converting Infix Expressions to Equivalent Postfix Expressions

- An infix expression can be evaluated by first being converted into an equivalent postfix expression, and then the postfix version of the expression is evaluated
- Facts about converting from infix to postfix
  - Operands always stay in the same order with respect to one another
  - An operator will move only “to the right” with respect to the operands
  - All parentheses are removed

# Converting Infix to Postfix

–e.g.  $1 + 2 * 3 ^ 4$

–in postfix  $1 2 3 4 ^ * +$

*Note:  $^$  is a symbol in some languages for exponentiation*

–Operators are in reverse order in this example

- So we need to store them on a stack
- When an operator is encountered, pop higher order operators before pushing the lower order operator

# Create the Postfix expression

- $1 + 2 * 3 \rightarrow 1 \ 2 \ 3 \ * \ +$
- $10 - 4 - 3 \rightarrow 10 \ 4 \ - \ 3 \ -$
- $2^3^3 \rightarrow 2 \ 3 \ 3 \ \wedge \ \wedge$
- $4/2/2 \rightarrow 4 \ 2 \ / \ 2 \ /$

- Associativity

- Left associative: e.g.  $+$ ,  $-$ ,  $*$ ,  $/$

- Right associative: e.g.  $^$  (exponential)

- Left-associative: Input  $+$  is lower than stack  $+$

$$2 + 3 + 4 \longrightarrow 2 \ 3 + 4 +$$

- Right-associative: Input  $^$  is higher than stack  $^$

$$2 ^ 3 ^ 4 \longrightarrow 2 \ 3 \ 4 ^ ^$$



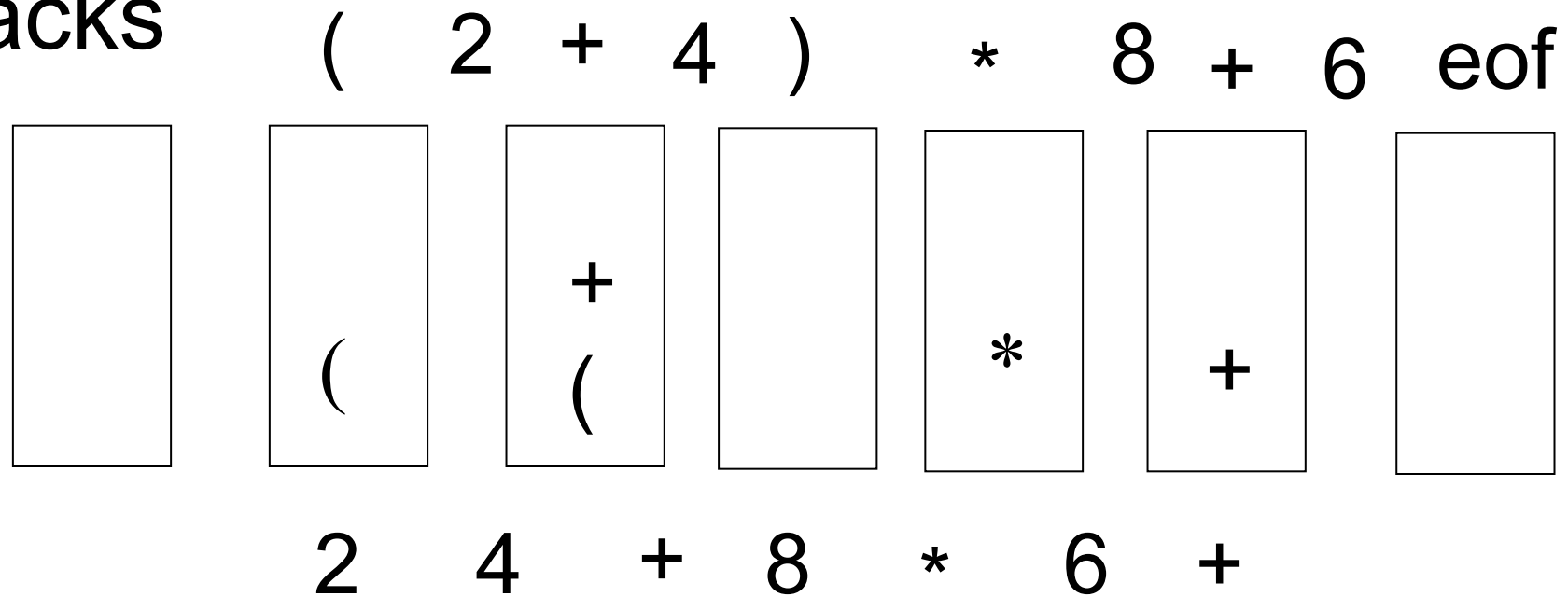
# Conversion Algorithm

- Algorithm:
  - Read infix expression as input
  - If input is operand, output the operand
  - If input is an operator +, -, \*, /, then
    - while (top of stack is an operator with greater precedence than the input operator)
      - pop** and output operator on top of stack
    - push** the input operator
  - If input is (, then **push**
  - If input is ), then **pop** and output all operators until see a ( on the stack. **Pop** the ( without output
  - If no more input then **pop** and output all operators on stack

## Running Time?

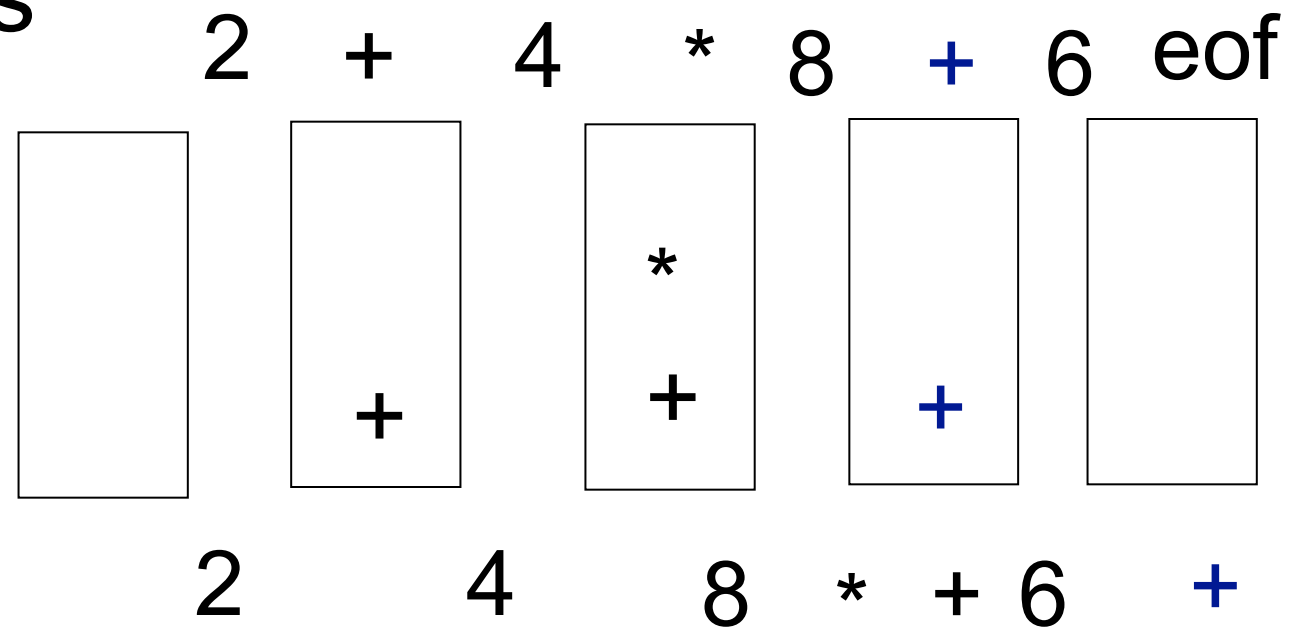
# Example: $(2 + 4) * 8 + 6$

- stacks



# Example: $2 + 4 * 8 + 6$

- stacks



# Conversion example

- Infix:  $1 - 2^3^3 - (4 + 5 * 6) * 7$
- Postfix:  $1\ 2\ 3\ 3\ \wedge\ \wedge\ -\ 4\ 5\ 6\ *\ +\ 7\ *\ -$
- -134217965

# Enumerated types

using symbols instead of numbers for constant values  
improves the readability of your code

- Simplest way to create your *own* type
  - you declare an enumerated type by using the **enum** keyword
  - you list all the values (the values are called *enumerators*) the type can hold:  

```
enum Seasons { Winter, Spring, Summer, Fall };
```

0            1            2            3
  - every enumerator is assigned an integer value, either explicitly or by default.

# A collection of named integer constants

```
#define EOL 0  
#define VALUE 1  
#define OPAREN 2
```

**enum:**

```
enum TokenType { EOL, VALUE, OPAREN, CPAREN, EXP, MULT, DIV, PLUS, MINUS};
```

**It is possible to create explicit values:**

```
enum seasons_t {spring = 10, summer = 100, fall = 50, winter = 5};  
enum months_t {January = 1, February, March, April};
```

# An alternative to if for multi-way branching if the condition being tested is equality for an integral type

```
enum Months { January = 1, February, March, April, May};  
Months month = January;
```

```
switch (month)//expression must evaluate to an integral type  
{  
    case January:  
        cout << "First month of the year!\n:";  
    case February:  
    case March:  
        cout << "It is cold this month!\n";  
        break;  
  
    case April:  
        cout << "Spring\n";  
    default:  
        cout << "One third of the year is over.\n";  
        break;  
}
```

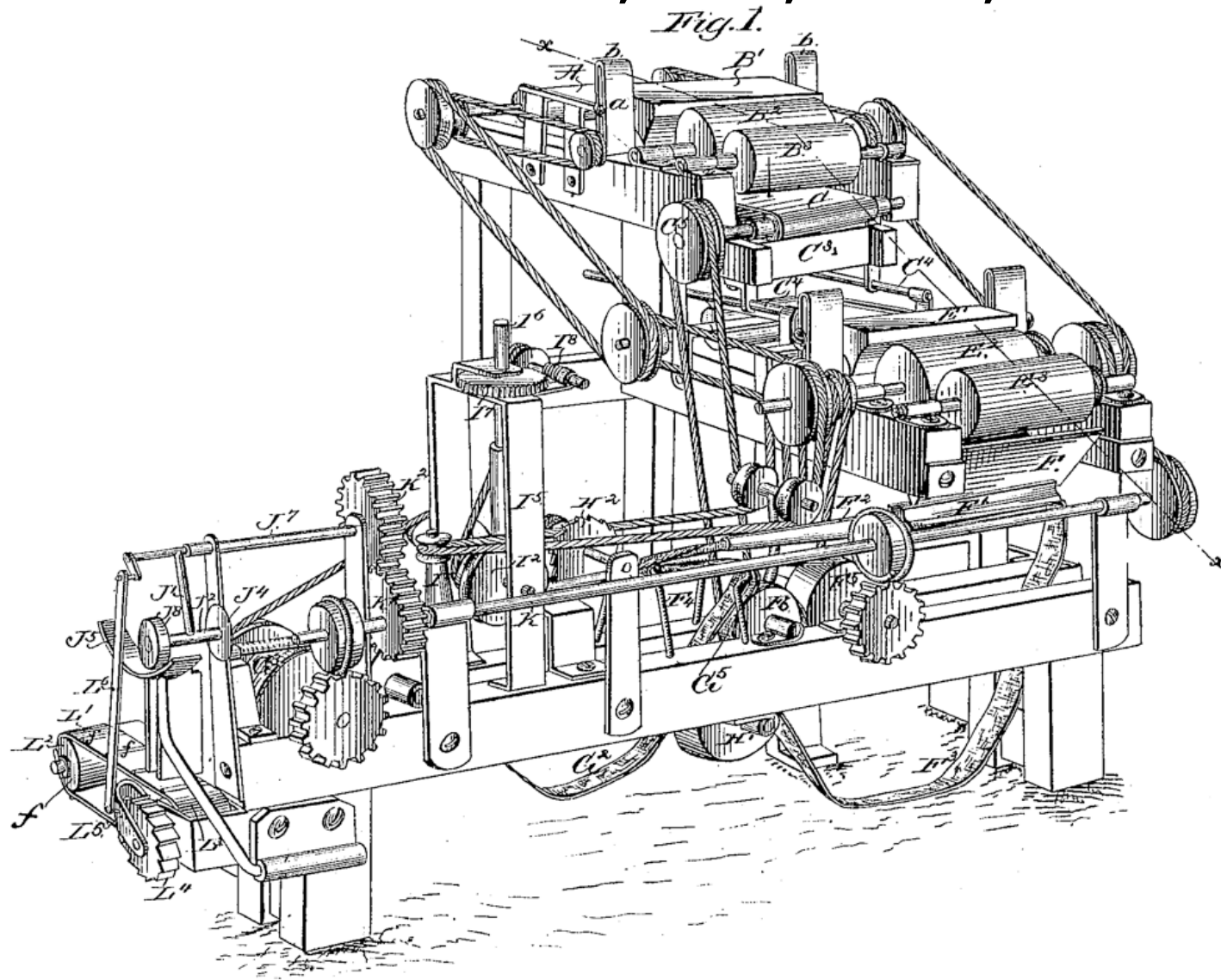
# Implementation

- Generate the tokens
  - **Lexical Analysis** - process of recognizing tokens in a stream of input.
  - **Tokenizer** - the program that does the lexical analysis of converting the input into tokens
  - **Token** - individual instance
- Evaluate the tokens
  - Template class called Evaluator
    - One stack to go from infix to postfix
    - One stack to evaluate a postfix expression



# Tokenizing $(2 + 4) * 8 + 6$

```
enum TokenType { EOL, VALUE, OPAREN, CPAREN, EXP,  
MULT, DIV, PLUS, MINUS };
```



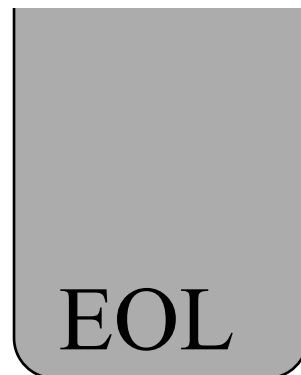
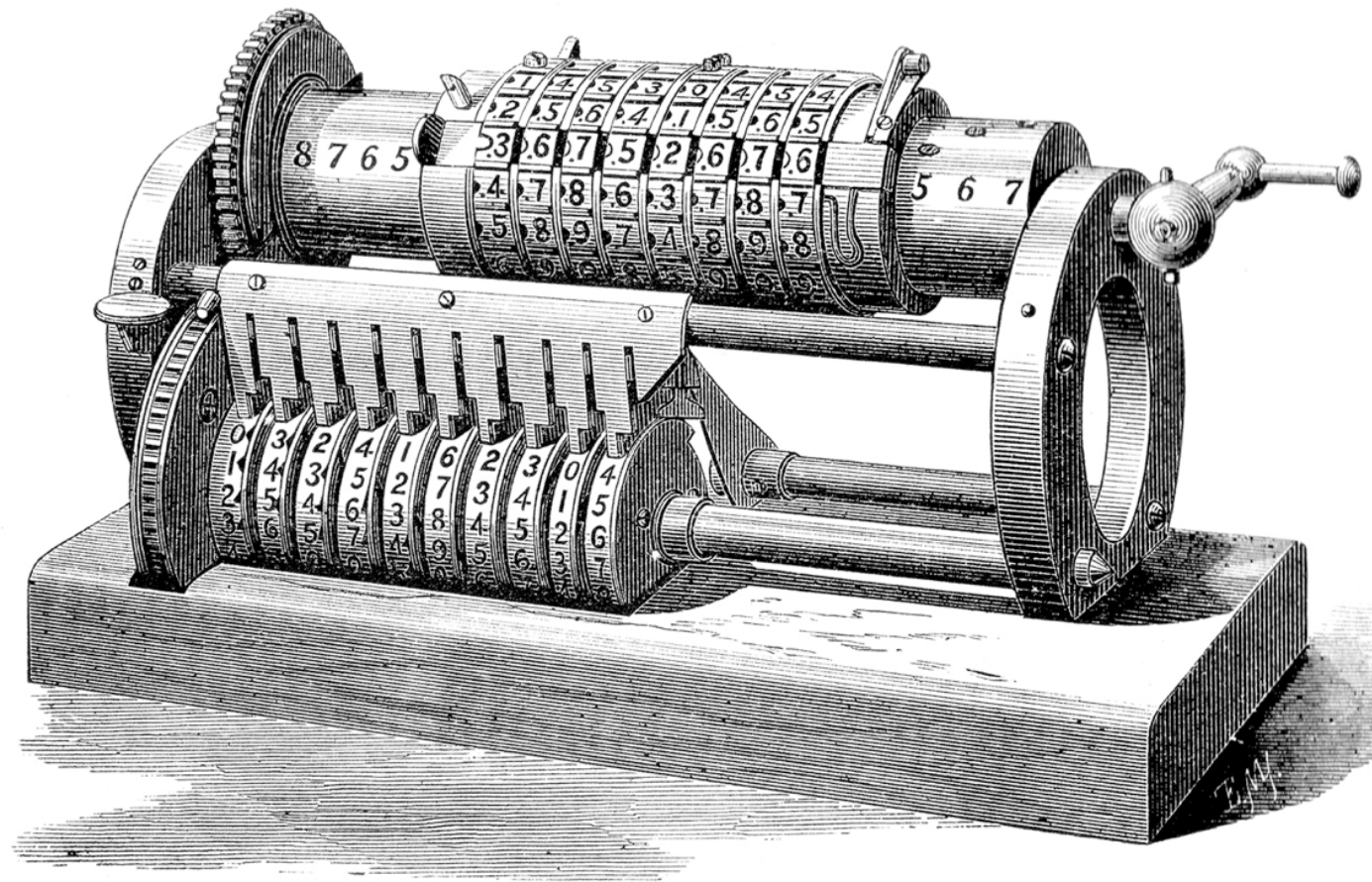
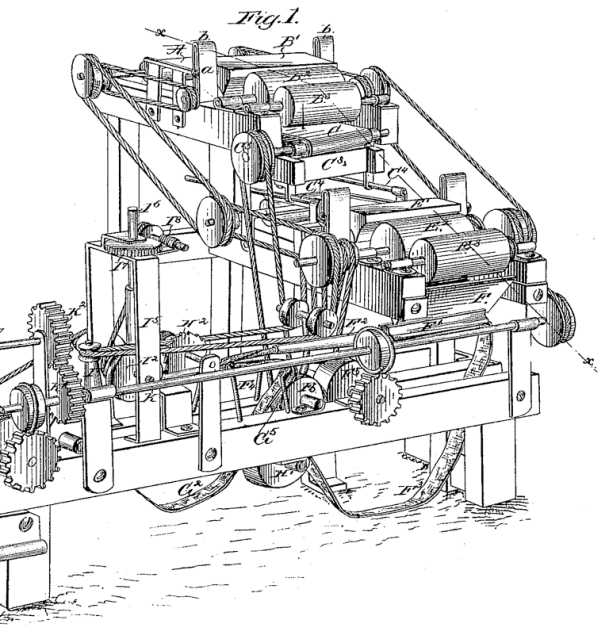
$)4 + 2($

```
enum TokenType { EOL,VALUE, OPAREN, CPAREN, EXP, MULT, DIV, PLUS, MINUS };
```

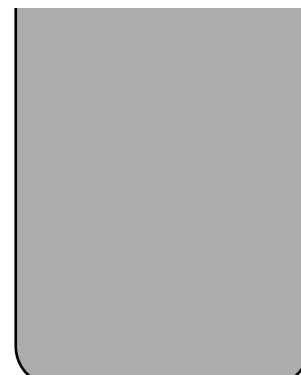
Tokenizer

# Simple Calculator

$(2 + 4) * 8 + 6$



opStack



postFixStack

```
// PREC_TABLE
struct Precedence
{
    int inputSymbol;
    int topOfStack;
}
vector<Precedence> PREC_TABLE =
{
    { 0, -1 }, // EOL = 0
    { 0, 0 }, // VALUE = 1
    { 100, 0 }, // OPAREN = 2
    { 0, 99 }, // CPAREN = 3
    { 6, 5 }, // EXP = 4
    { 3, 4 }, // MULT = 5
    { 3, 4 }, // DIV = 6
    { 1, 2 }, // PLUS = 7
    { 1, 2 }, // MINUS = 8
};
```

$$2 + 4$$



```
enum TokenType { EOL,VALUE, OPAREN, CPAREN, EXP, MULT, DIV, PLUS, MINUS };
```

```
template <class NumericType>  
class Token  
{  
public:
```

notice the default values for  
the constructor!

```
Token( TokenType tt = EOL, const NumericType & nt = 0 )  
: theType( tt ), theValue( nt ) { }
```

```
TokenType getType( ) const{ return theType; }  
const NumericType & getValue( ) const{ return theValue; }
```

```
private:  
    TokenType theType;  
    NumericType theValue;  
};
```

VALUE,2

PLUS,0

VALUE,4

```
template <class NumericType>
class Tokenizer
{
public:
    Tokenizer( istream & is ): in( is ) { }

    Token<NumericType> getToken( );
```

Just one simple method!

```
private:
    istream & in;
};
```

where we get the input from



( 12 + 4 )

// Find the next token, skipping blanks, and return it.

// Print error message if input is unrecognized.

```
template <class NumericType>
```

```
Token<NumericType> Tokenizer<NumericType>::getToken( )
```

```
{
```

```
    char ch;
```

```
    NumericType theValue;
```

OPAREN,0

VALUE,12

PLUS,0

VALUE,4

CPAREN,0

// Skip blanks

```
while( in.get( ch ) && ch == ' ' );
```

```
if( in.good( ) && ch != '\n' && ch != '\0' ){
```

```
    switch( ch ){
```

```
        case '^': return EXP;
```

```
        case '/': return DIV;
```

```
        case '*': return MULT;
```

```
        case '(': return OPAREN;
```

```
        case ')': return CPAREN;
```

```
        case '+': return PLUS;
```

```
        case '-': return MINUS;
```

```
        default:
```

```
            in.putback( ch );
```

```
            if( !( in >> theValue ) ){
```

```
                cerr << "Parse error" << endl;
```

```
                return EOL;
```

```
            }
```

```
            return Token<NumericType>( VALUE, theValue );
```

```
    }
```

```
    return EOL;
```

```
}
```

# Simple Calculator

```
int main()
{
    string str;

    while( getline( cin, str ) )
    {
        Evaluator<int> e( str );
        cout << e.getValue( ) << endl;
    }
    return 0;
}
```

```
enum TokenType { EOL,VALUE, OPAREN, CPAREN, EXP, MULT, DIV, PLUS, MINUS };
```

```
template <class NumericType>
```

```
class Evaluator
```

```
{
```

```
public:
```

```
    Evaluator( const string & s ) : str( s )  
        { opStack.push( EOL ); }
```

```
    // The only publicly visible routine
```

```
    NumericType getValue( );    // Do the evaluation
```

```
private:
```

```
    stack<TokenType> opStack;    // Operator stack for conversion
```

```
    stack<NumericType> postfixStack; // Stack for postfix machine
```

```
    istream str;                // String stream
```

```
    // Internal routines
```

```
    NumericType getTop( );
```

```
    void binaryOp( TokenType topOp );
```

```
    void processToken( const Token<NumericType> & lastToken );
```

```
    // Get top of postfix stack
```

```
    // Process an operator
```

```
    // Handle LastToken
```

```
};
```

```
// Public routine that performs the evaluation.
// Examines the postfix machine to see if a single result
// is left and if so, returns it; otherwise prints error.
```

```
template <class NumericType>
NumericType Evaluator<NumericType>::getValue(
{
    Tokenizer<NumericType> tok( str );
    Token<NumericType> lastToken;

    do {
        lastToken = tok.getToken( );
        processToken( lastToken );
    } while( lastToken.getType( ) != EOL );

    if( postFixStack.empty( ) )
    {
        cerr << "Missing operand!" << endl;
        return 0;
    }

    NumericType theResult = postFixStack.top( );
    postFixStack.pop( );
    if( !postFixStack.empty( ) )
        cerr << "Warning: missing operators!" << endl;

    return theResult;
}
```

( 12 + 4 )

0,OPAREN	12,VALUE	0,PLUS	4,VALUE	0,CPAREN
----------	----------	--------	---------	----------

} Process each token till  
eof/eoln

} Check if operand is missing

} Output result  
check if operator  
is missing



// top and pop the postfix machine stack; return the result.

// If the stack is empty, print an error message.

```
template <class NumericType>
```

```
NumericType Evaluator<NumericType>::getTop( )
```

```
{
```

```
    if( postfixStack.empty( ) )
```

```
    {
```

```
        cerr << "Missing operand" << endl;
```

```
        return 0;
```

```
    }
```

```
    NumericType tmp = postfixStack.top( );
```

```
    postfixStack.pop( );
```

```
    return tmp;
```

```
}
```

```
// Process an operator by taking two items off the postfix
// stack, applying the operator, and pushing the result.
// Print error if missing closing parenthesis or division by 0.
```

```
template <class NumericType>
```

```
void Evaluator<NumericType>::binaryOp( TokenType topOp ){
```

```
    if( topOp == OPAREN ){
        cerr << "Unbalanced parentheses" << endl;
        opStack.pop( );
        return;
    }
```

} Error check

```
    NumericType rhs = getTop( );
    NumericType lhs = getTop( );
```

} top/pop operands

```
    if( topOp == EXP )
        postFixStack.push( pow( lhs, rhs ) );
    else if( topOp == PLUS )
        postFixStack.push( lhs + rhs );
    else if( topOp == MINUS )
        postFixStack.push( lhs - rhs );
    else if( topOp == MULT )
        postFixStack.push( lhs * rhs );
    else if( topOp == DIV )
        if( rhs != 0 )
            postFixStack.push( lhs / rhs );
        else {
            cerr << "Division by zero" << endl;
            postFixStack.push( lhs );
        }
    }
```

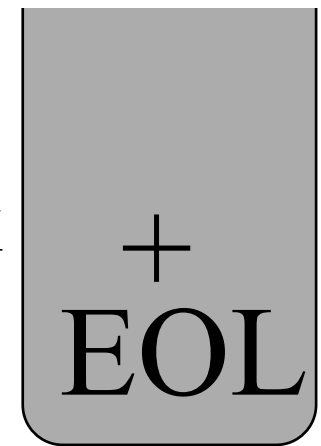
} perform  
binary  
operation;  
put result on  
postFixStack

```
    opStack.pop( );
```

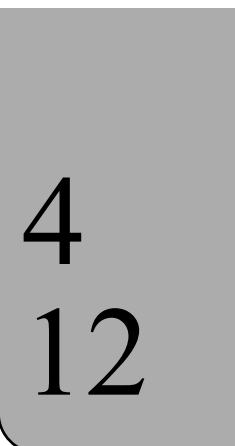
} pop operator stack

```
}
```

12+4-3  
↑

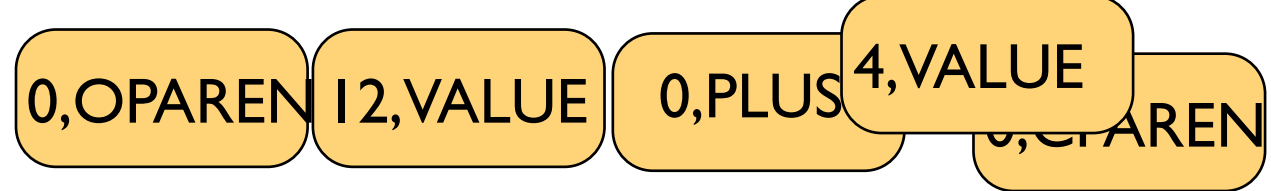


opStack



postFixStack

// After token is read, use operator precedence parsing  
 // algorithm to process it; missing opening parentheses  
 // are detected here.



```
template <class NumericType>
void Evaluator<NumericType>::processToken(const Token<NumericType> & lastToken)
```

```
{
    TokenType topOp;
    TokenType lastType = lastToken.getType( );
```

```
    switch( lastType ){
```

```
        case VALUE:
```

```
            postFixStack.push( lastToken.getValue( ) );
            return;
```

} Put operand on  
postFixStack

```
        case CPAREN:
```

```
            while( ( topOp = opStack.top( ) ) != OPAREN && topOp != EOL )
                binaryOp( topOp );
            if( topOp == OPAREN )
                opStack.pop( ); // Get rid of opening parentheses
            else
                cerr << "Missing open parenthesis" << endl;
            break;
```

} pop opStack  
and eval till “(“  
is found

```
        default: // General operator case
```

```
            while( PREC_TABLE[ lastType ].inputSymbol <=
                    PREC_TABLE[ topOp = opStack.top( ) ].topOfStack )
                binaryOp( topOp );
            if( lastType != EOL )
                opStack.push( lastType );
            break;
```

} pop operators  
with less  
precedence  
and eval

```
// PREC_TABLE
struct Precedence
{
    int inputSymbol;
    int topOfStack;
};
vector<Precedence> =
{
    { 0, -1 }, // EOL = 0
    { 0, 0 }, // VALUE = 1
    { 100, 0 }, // OPAREN = 2
    { 0, 99 }, // CPAREN = 3
    { 6, 5 }, // EXP = 4
    { 3, 4 }, // MULT = 5
    { 3, 4 }, // DIV = 6
    { 1, 2 }, // PLUS = 7
    { 1, 2 } // MINUS = 8
};
```

# Key Ideas

- Evaluator class
  - Evaluates **infix** expression by “converting” to **postfix** expression and evaluating
  - Does creation and **evaluation** of **postfix** expression in one step (evaluates as it creates)
- Precedence table
  - Establishes **precedence** of operators
  - Establishes whether operators are **left** or **right** associative

# Expression Trees

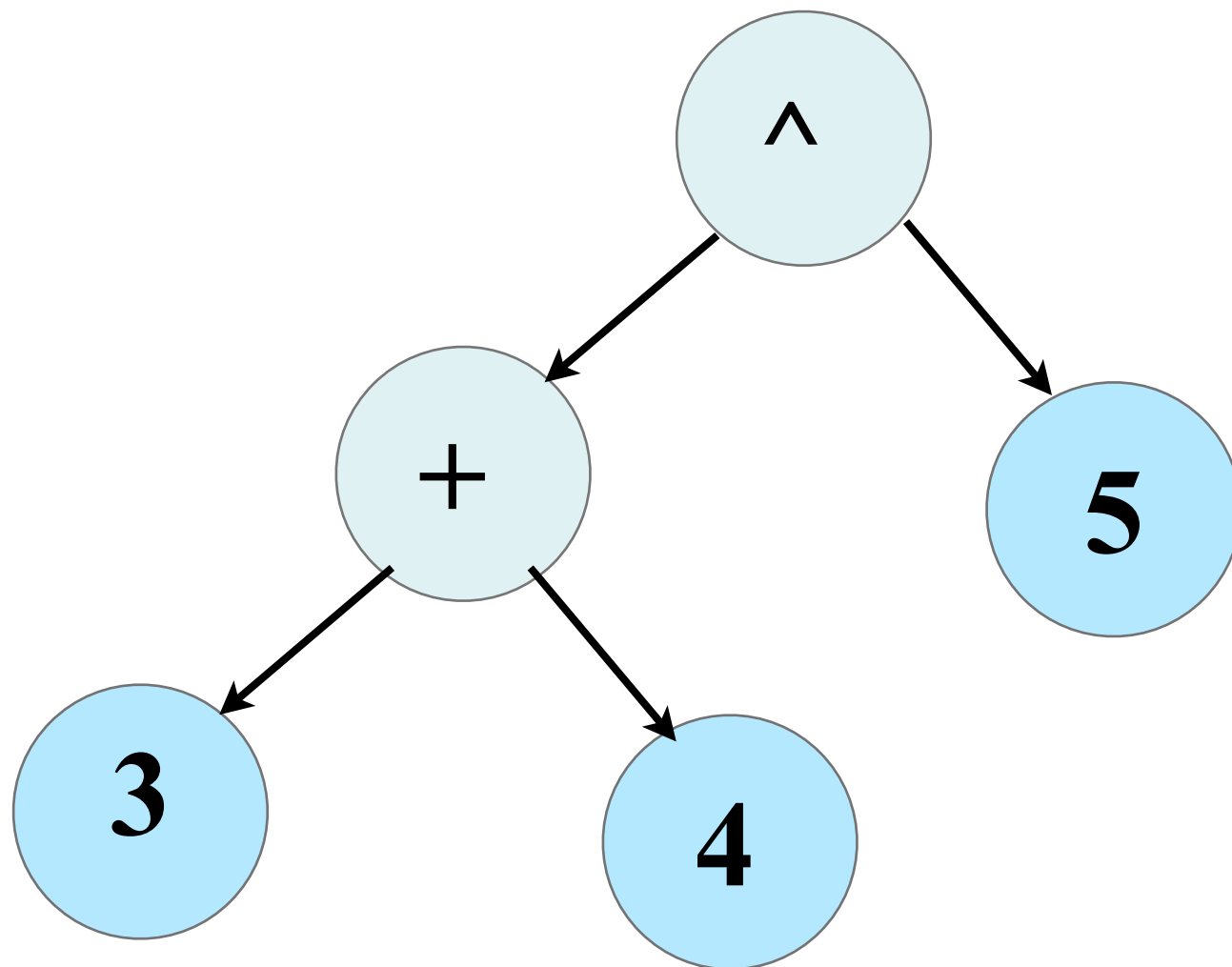
- leaves are operands
- other nodes are operators
- binary operators implies binary tree
- a node would have only one child if unary (e.g. -)
- evaluate by applying the operator at the root and recursively evaluating the left and right subtrees

# Running Time?

- Every step involves a single push (and if it is an operator, two pops occur before the push.)
- E.g. 1 2 - 4 5 ^ 3 \* 6 \* 7 2 2 ^ ^ / -
- Has 9 operands and 8 operators, thus 17 steps and 17 pushes.
- **Linear** in the input size.

# Expression Tree

$$(3 + 4)^5$$



# Expression Trees

- Leaves contain operands (e.g., constants or variable names)
- Non-leaf nodes contain operators (e.g.,  $^$ ,  $*$ ,  $/$ ,  $+$ ,  $-$ )
- Nodes can have 1, 2, 3, or more children



# Expression Tree

$$3 + 4^5$$

