

# Path Stage 1 – Group 2 Mid Report

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## 1 Introduction

In this project, we analyze the OpenMP version of Floyd-Warshall algorithm by profiling the code using Vtune to identify bottlenecks. The Vtune profiling result is shown in section §2. To tune the code, we explore different compiler flags (section §??) and slightly vectorized the code (section §4). Finally we perform strong and weak scaling in section §5.

## 2 Profiling

### 2.1 Identify bottlenecks

To define the bottlenecks of `path.c`, we profiled our code using Intel's VTune Amplifier to understand which part of the code takes the most time to run. VTune Amplifier commands is shown in Figure 1.

```
amplxe-cl -collect advanced-hotspots ./path
amplxe-cl -R hotspots -report-output vtune-report.csv -format csv -csv-delimiter comma
```

Figure 1: VTune Amplifier Command

### 2.2 Initial Timing Result

As shown in Figure 2, the majority of time is spent inside the `square` function. We will take about how we optimize the performance of the `square` function in §4.

Function	Module	CPU Time	CPU Time:Ideal
square	path.x	45.14s	39.03s
_kmpbarrier	libiomp5.so	6.485s	0.16s
_kmpc_reduce_nowait	libiomp5.so	2.979s	0s
_kmp_fork_barrier	libiomp5.so	2.553s	0.03s
gen_graph	path.x	0.030s	0s
_intel_ssse3_memcpy	path.x	0.030s	0s
fletcher16	path.x	0.030s	0s

Figure 2: Initial Profile Result

## 3 Compiler Flags

With some experience from the previous project, we decided to play around with different compiler flags to examine whether it will help us the speedup of the code. We found that the `icc` compiler flag shown below gave us 8 times speedup to the `deinfiniteize` function.

```
OPTFLAGS=-O3 -no-prec-div -opt-prefetch -xHost -ansi-alias -ipo -restrict
```

## 4 Vectorization

As we discussed in section §2, the majority of time is spent in the `sqare` function. Looking at the vectorization report generated by the compiler default code, we see that the compiler did not vectorize the `for` loop inside the `sqare` function because it assumes dependencies in the loop.

```
LOOP BEGIN at mt19937p.c(59,5) inlined into path.c(228,14)
remark #15344: loop was not vectorized: vector dependence prevents vectorization
remark #15346: vector dependence: assumed FLOW dependence between mt line 60 and mt line
LOOP END
```

To clear the dependency, we simply added `#pragma vector aligned` in the inner loop as shown below to ensure the compiler that the `lij` is indeed aligned. Adding the `#pragma` instruction gave us a 4 times speed up to the `square` function.

```
#pragma omp parallel for shared(l, lnew) reduction (&& : done)
for (int j = 0; j<n; ++j){
for (int k = 0; k < n; ++k) {
int lkj = l[j*n+k];

#pragma vector aligned
for (int i = 0; i < n; ++i){
int lij = lnew[j*n+i];
int lik = l[k*n+i];

if (lik +lkj < lij){
lij = lik+lkj;
done=0;
}

lnew[j*n+i] = lij;
}
}
}
```

## 5 Evaluation

The result of strong and weak scaling studies of the vectorized code is shown in Figure Figure 3 and Figure Figure 4. In the strong scaling study, we examined the time spent with varies frame sizes while keeping the number of threads constant; in our case. we used 24 threads. The result shows that the time spent increases linearly as we increase the frame

size. In the weak scaling study, we tested the time spent using varies number of thread while keeping the frame size 4000. The time spent decreases in log scale as we increase the number of thread.

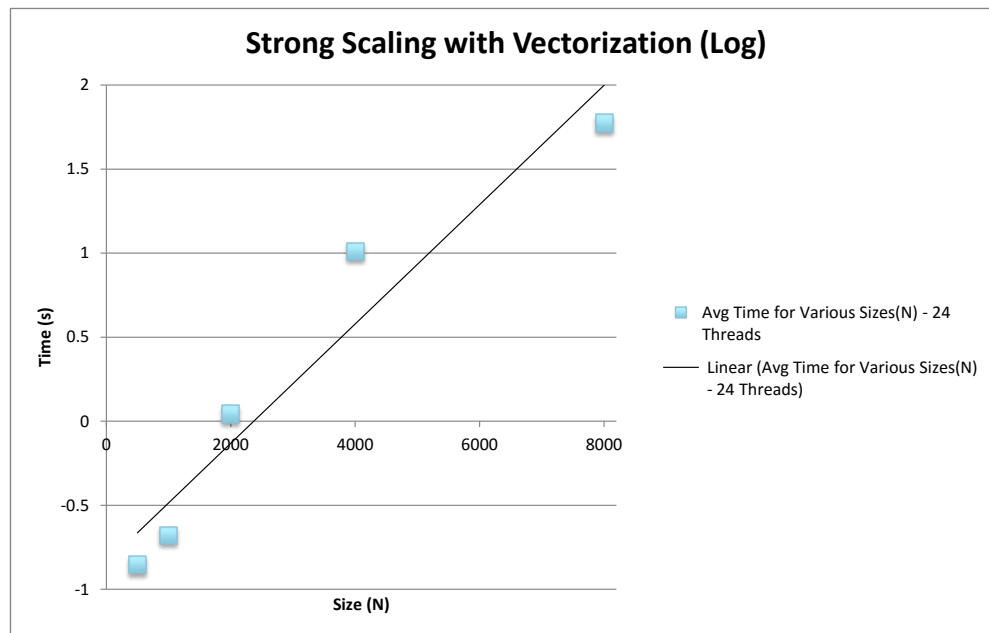


Figure 3: Strong scaling study with 24 threads and varies number of frames

## 6 Future Work

- **Blocking.** We plan on performing some kind of blocking to improve the performance of the code.
- **Using MPI.** We might try to implement MPI and compare the performance with OpenMP.

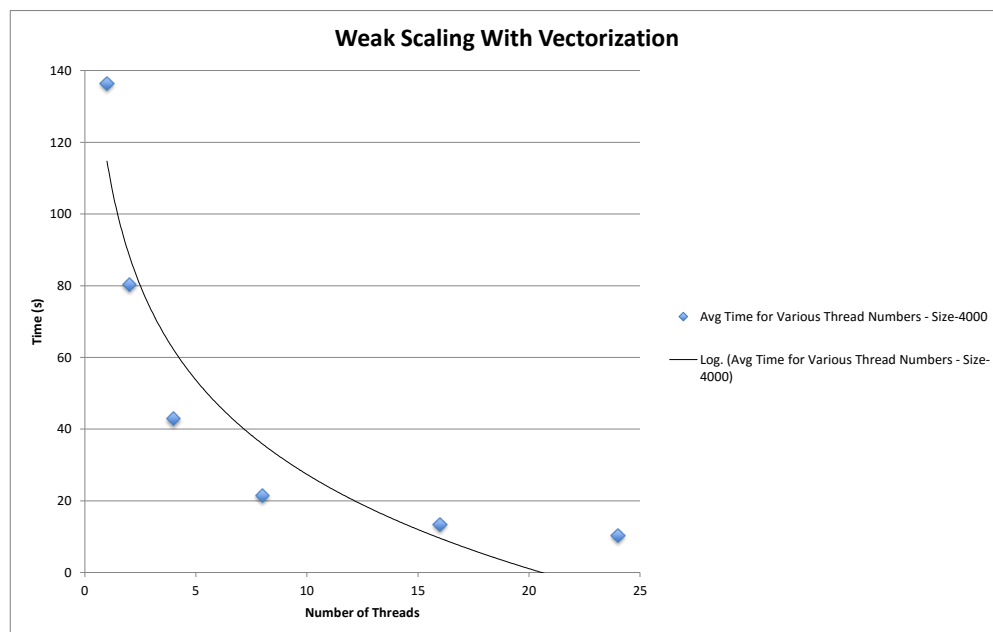


Figure 4: Weak scaling study with frame size 4000 and varies number of threads