

# POMC User's Guide

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## Abstract

POMC is an implementation of the automaton construction for formulas of Precedence-Oriented Temporal Logic (POTL), and the model checking procedure thereof. This document is a reference guide to its input and output formats, and also describes at a high level its architecture and source code.

## 1 Introduction

Precedence-Oriented Temporal Logic (POTL) [6] is a novel temporal logic formalism based on the family of Operator Precedence Languages (OPL), a subclass of deterministic context-free languages. POTL is strictly more expressive than LTL and other temporal logics based on subfamilies of context-free languages, such as CaRet [3] and NWTL [2]. In particular, POTL reasons on an algebraic structure equipped with, besides the usual linear order, a binary nesting relation between word positions, which can be one-to-one, one-to-many, or many-to-one. Such a relation is more general than the one found in Nested Words [4], because the latter may only be one-to-one. POTL can be applied to the specification of several kinds of requirements on procedural programs with exceptions.

Besides some results concerning its expressiveness, we introduced an automata-based model checking procedure for POTL. This procedure consists in building an Operator Precedence Automaton (OPA), the class of pushdown automata that identifies OPL, accepting the language denoted by a given POTL formula. The size of the generated automaton is exponential in the length of the formula, which is asymptotically comparable with other linear-time temporal logic formalisms such as LTL, CaRet, and NWTL.

POMC is a tool that implements the automaton construction for POTL, and a model checking procedure for it. For the time being, only the construction for finite words has been implemented. Given a POTL formula  $\varphi$  and an input OPA modeling some system, POMC builds the OPA equivalent to  $\neg\varphi$ , computes its intersection with the input OPA, and checks the emptiness of the resulting OPA. Both the OPA construction and the intersection are done on-the-fly. POMC also supports providing input models in MiniProc, a simple procedural programming language with exceptions. MiniProc programs are automatically translated into equivalent OPA.

We used POMC to prove some interesting properties of programs which we modeled as OPA. Such experiments are contained in `pomc` files in the `opa` and `opa-more` subdirectories. Some more experiments where the model is written in MiniProc are contained in directory `miniproc`.

We show how to use POMC in Section 2. If you wish to examine the input formulas and OPA for the experiments more carefully, or to write your own, we describe the format of POMC input files in Section 3. We also demonstrate the use of the tool with

a few experiments in Section 4. Finally, Section 5 contains a high-level description of the source code.

## 2 Quick-Start Guide

POMC has been developed in the Haskell programming language, and packaged with the Haskell Tool Stack<sup>1</sup>. POMC can be built from sources by typing the following commands in a shell:

```
$ cd ~/path/to/POMC-sources
$ stack setup
$ stack build
```

Then, POMC can be executed on an input file `file.pomc` as follows:

```
$ stack exec pomc -- file.pomc
```

Directory `eval` contains several POMC input files. Such files contain POTL formulas and OPA to be checked against them. For more details on the format of POMC input files, see Section 3.

Directory `eval` also contains the Python script `mcbench.py`, which may be useful to evaluate POMC input files, as it also prints a summary of the resources used by POMC. It must be executed with a subdirectory of `~/path/to/POMC-sources` as its working directory. If invoked with no arguments, it executes POMC on all input files in the current working directory. E.g.,

```
$ cd ~/path/to/POMC-sources/eval
$ ./mcbench.py opa
```

evaluates all `*.pomc` files in directory `~/path/to/POMC-sources/eval/opa`. The script can also be invoked with POMC files as its arguments, which are then evaluated. E.g.,

```
$ cd ~/path/to/POMC-sources/eval/opa
$ ./mcbench.py 1-generic-small.pomc 2-generic-medium.pomc
```

executes POMC with input files `1-generic-small.pomc` and `2-generic-medium.pomc`. `mcbench.py` can be invoked with the following optional flags:

- `-iter <#iters>` Number of iterations of the benchmarks to be performed. The final table printed by the script contains the mean time and memory values computed on all iterations. (Default: 1)
- `-jobs <#jobs>` Number of benchmarks to be run in parallel. If you provide a value greater than 1, make sure you have enough CPU cores on your machine. (Default: 1)
- `-ms` Output time in milliseconds instead of seconds.

## 3 POMC Input/Output Format

POMC takes in input plain text files of two possible formats.

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<sup>1</sup><https://www.haskellstack.org/>

|      | call | ret | han | exc |      | call | ret | han | exc | stm |
|------|------|-----|-----|-----|------|------|-----|-----|-----|-----|
| call | <    | =   | <   | >   | call | <    | =   | <   | >   | <   |
| ret  | >    | >   | >   | >   | ret  | >    | >   | >   | >   | >   |
| han  | <    | >   | <   | =   | han  | <    | >   | <   | =   | <   |
| exc  | >    | >   | >   | >   | exc  | >    | >   | >   | >   | >   |
|      |      |     |     |     | stm  | >    | >   | >   | >   | >   |

(a) OPM  $M_{\text{call}}$

(b) OPM  $M_{\text{stm}}$

Figure 1

### 3.1 Providing input models as OPA

The first input format contains a requirement specification in terms of a list of POTL formulas, and an OPA to be checked against them:

```

formulas = FORMULA [, FORMULA ...] ;
prec = SL PR SL [, SL PR SL ...] ;
opa:
  initials = STATE_SET ;
  finals = STATE_SET ;
  deltaPush = (STATE, AP_SET, STATE_SET)
               [, (STATE, AP_SET, STATE_SET) ...] ;
  deltaShift = (STATE, AP_SET, STATE_SET)
                [, (STATE, AP_SET, STATE_SET) ...] ;
  deltaPop = (STATE, STATE, STATE_SET)
              [, (STATE, STATE, STATE_SET) ...] ;

```

where STATE\_SET is either a single state, or a space-separated list of states, surrounded by parentheses. States are non-negative integer numbers (e.g. (0 1 ...)). AP\_SET is a space-separated list of atomic propositions, surrounded by parentheses (e.g. (call p1) or ("call" "p1")). In more detail:

- `prec` is followed by a comma-separated list of precedence relations between structural labels, that make up an Operator Precedence Matrix. The list is terminated by a semicolon. Precedence relations (PR) can be one of `<`, `=`, or `>`, which respectively mean `<`, `=`, and `>`. Structural labels (SL) can be any sequence of alphabetic characters.
- `formulas` is followed by a comma-separated, semicolon-terminated list of POTL formulas. The syntax of such formulas is defined later in this section.
- `opa` is followed by the explicit description of an OPA. The list of initial and final states must be given, as well as the transition relations.

Additionally, POMC input files may contain C++-style single-line comments starting with `\`, and C-style multi-line comments enclosed in `/*` and `*/`.

External files can be included with

```
include = "path/to/file.inc";
```

where the path is relative to the pomc file location.

POTL formulas can be written by using the operators in the “POMC Operator” column of Table 1, following the same syntax rules as in [5, 6].

Once POMC is executed on an input file in the format above, it checks whether the given OPA satisfies the given formulas, one by one.

Consider the example input file 1-generic-small.pomc, reported below:

```
prec = call < call, call = ret, call < han, call > exc,
      ret > call, ret > ret, ret > han, ret > exc,
      han < call, han > ret, han < han, han = exc,
      exc > call, exc > ret, exc > han, exc > exc;
```

```
formulas = G ((call And pb And (T Sd (call And pa)))
              --> (PNu exc Or XNu exc));
```

```
opa:
initials = 0;
finals = 10;
deltaPush =
  (0, (call pa), 1),
  (1, (han), 2),
  (2, (call pb), 3),
  (3, (call pc), 4),
  (4, (call pc), 4),
  (6, (call perr), 7),
  (8, (call perr), 7);
deltaShift =
  (4, (exc), 5),
  (7, (ret perr), 7),
  (9, (ret pa), 11);
deltaPop =
  (4, 2, 4),
  (4, 3, 4),
  (4, 4, 4),
  (5, 1, 6),
  (7, 6, 8),
  (7, 8, 9),
  (11, 0, 10);
```

First, OPM  $M_{\text{call}}$  from [6] (Figure 1a) is chosen.

The meaning of the formula  $G ((\text{call And pb And (T Sd (call And pa)))$   
 $\rightarrow (\text{PNu exc Or XNu exc}), \text{or } \Box((\text{call} \wedge p_B \wedge \text{Scall}(\top, p_A)) \implies \text{CallThr}(\top))$ ,  
 is explained in the paper.

POMC will check the OPA against the formula, yielding the following output:

```
Model Checking
Formula: G (((("call" And "pb") And (T Sd ("call" And "pa"))
              --> ((PNu "exc") Or (XNu "exc"))))
Input OPA state count: 12
Result: True
Elapsed time: 14.59 s
```

| Group        | POTL Operator     | POMC Operator | Notation | Associativity |
|--------------|-------------------|---------------|----------|---------------|
| Unary        | $\neg$            | $\sim$ , Not  | Prefix   | –             |
|              | $\bigcirc^d$      | PNd           | Prefix   | –             |
|              | $\bigcirc^u$      | PNu           | Prefix   | –             |
|              | $\ominus^d$       | PBd           | Prefix   | –             |
|              | $\ominus^u$       | PBu           | Prefix   | –             |
|              | $\chi_F^d$        | XNd           | Prefix   | –             |
|              | $\chi_F^u$        | XNu           | Prefix   | –             |
|              | $\chi_P^d$        | XBd           | Prefix   | –             |
|              | $\chi_P^u$        | XBu           | Prefix   | –             |
|              | $\bigcirc_H^d$    | HNd           | Prefix   | –             |
|              | $\bigcirc_H^u$    | HNu           | Prefix   | –             |
|              | $\ominus_H^d$     | HBd           | Prefix   | –             |
|              | $\ominus_H^u$     | HBu           | Prefix   | –             |
|              | $\diamond$        | F, Eventually | Prefix   | –             |
|              | $\square$         | G, Always     | Prefix   | –             |
| POTL Binary  | $\mathcal{U}_X^d$ | Ud            | Infix    | Right         |
|              | $\mathcal{U}_X^u$ | Uu            | Infix    | Right         |
|              | $\mathcal{S}_X^d$ | Sd            | Infix    | Right         |
|              | $\mathcal{S}_X^u$ | Su            | Infix    | Right         |
|              | $\mathcal{U}_H^d$ | HUd           | Infix    | Right         |
|              | $\mathcal{U}_H^u$ | HUu           | Infix    | Right         |
|              | $\mathcal{S}_H^d$ | HSd           | Infix    | Right         |
|              | $\mathcal{S}_H^u$ | HSu           | Infix    | Right         |
| Prop. Binary | $\wedge$          | And, &&       | Infix    | Left          |
|              | $\vee$            | Or,           | Infix    | Left          |
|              | $\oplus$          | Xor           | Infix    | Left          |
|              | $\implies$        | Implies, -->  | Infix    | Right         |
|              | $\iff$            | Iff, <-->     | Infix    | Right         |

Table 1: This table contains all currently supported POTL operators, in descending order of precedence. Operators listed on the same line are synonyms. Operators in the same group have the same precedence. Note that operators are case sensitive.

Total elapsed time: 14.59 s (1.4593e1 s)

Indeed, the OPA does satisfy the formula. POMC also outputs the time taken by each acceptance check and, when a formula is rejected, a (partial) counterexample trace.

### 3.2 Providing MiniProc input models

The second kind of input files also contain POTL formulas, and a program in the *MiniProc* language to be checked against them. MiniProc is a simplified procedural programming language, where variables are all global and only take Boolean values (note that MiniProc is not Turing-complete, so any use of the word ‘program’ when referring to it is a deliberate abuse of terminology). This limitation allows POMC to translate every MiniProc program into an OPA, that is then checked against the

supplied formulas. This kind of input files have this form:

```
formulas = FORMULA [, FORMULA ...] ;
program:
PROGRAM
```

MiniProc programs have the following syntax:

```
PROGRAM := [DECLS] FUNCTION [FUNCTION ...]
DECLS := var IDENTIFIER [, IDENTIFIER ...] ;
FUNCTION := IDENTIFIER () { STMT; [STMT ...] }
STMT := IDENTIFIER := BEXPR
      | while (BEXPR) { [STMT ...] }
      | if (BEXPR) { [STMT ...] } else { [STMT ...] }
      | try { [STMT ...] } catch { [STMT ...] }
      | IDENTIFIER()
      | throw
BEXPR := BEXPR && BDISJ | BDISJ
BDISJ := BDISJ || BTERM | BTERM
BTERM := !BTERM | (BEXPR) | IDENTIFIER | true | false
```

In the definition, non-terminal symbols are uppercase, and keywords lowercase. Parts surrounded by square brackets are optional, and ellipses mean that the enclosing group can be repeated zero or more times. An IDENTIFIER is any sequence of letters, numbers, or characters '.', ':', and '\_', starting with a letter or an underscore.

The program starts with a variable declaration, which must include all variables used in the program. Then, a sequence of functions are defined, the first one being the entry-point to the program. Function bodies consist of semicolon-separated statements. Assignments, while loops and ifs have the usual semantics. The try-catch statement executes the catch block whenever an exception is thrown by any statement in the try block (or any function it calls). Exceptions are thrown by the throw statement, and they are not typed (i.e., there is no way to distinguish different kinds of exceptions). Functions can be called by prepending their name to the () token (they do not admit arguments, as all variables are global). Since all variables are Boolean, expressions can be composed with the logical and (&&), or (||) and negation (!) operators.

POMC automatically translates such programs into OPA in a way similar to [6]. There are two main differences. First, assignment statements are explicitly represented as push transitions with label **stm**. We use OPM  $M_{\text{stm}}$  of Figure 1b. Second, a dummy **exc** is raised at the end of try blocks to match the **han** and pop its stack symbol. Such **exc**s do not cause the execution of the catch block, and can be recognized because they do not terminate any function, so  $\neg \ominus^u \text{ call}$  holds in them.

It is possible to declare *modules* by including a double colon (:) in function names. E.g., function A::B::C() is contained in module A::B, which is contained in A. In the OPA resulting from the program, the module names hold whenever a contained function is called or returns. This is useful for referring to multiple functions at once in POTL formulas, hence drastically reducing formula length and closure size.

An example input file is given below:

```
formulas = G ((call And pb And (call Sd (call And pa)))
--> (PNu exc Or XNu exc));
```

```

program:
var foo;

pa() {
  foo = false;
  try {
    pb();
  } catch {
    pc();
  }
}

pb() {
  if (foo) {
    throw;
  } else {}
}

pc() { }

```

POMC prints the following:

```

Model Checking
Formula: G (((("call" And "pb") And ("call" Sd ("call" And "pa"))))
--> ((PNu "exc") Or (XNu "exc")))
Input OPA state count: 28
Result: True
Elapsed time: 803.7 ms

```

Total elapsed time: 803.7 ms (8.0370e-1 s)

## 4 Some experiments

In this section we report the results of some experiments provided in the `eval` directory. The experiments were executed on a laptop with a 2.2 GHz Intel processor and 15 GiB of RAM, running Ubuntu GNU/Linux 20.04.

### 4.1 Directory opa

This directory contains a few programs modeled as OPA, on which POMC proves or disproves some interesting specifications. The resources employed by POMC on such tasks are reported in Table 2. If you wish to repeat such experiments, you may run the following commands:

```

$ cd ~/path/to/POMC-sources/eval
$ ./mcbench.py opa

```

|   | Benchmark name          | # states | Time (ms) | Memory (KiB) |         | Result |
|---|-------------------------|----------|-----------|--------------|---------|--------|
|   |                         |          |           | Total        | MC only |        |
| 1 | generic small           | 12       | 867       | 70,040       | 10,166  | True   |
| 2 | generic medium          | 24       | 673       | 70,064       | 4,043   | False  |
| 3 | generic larger          | 30       | 1,014     | 70,063       | 14,160  | True   |
| 4 | Jensen                  | 42       | 305       | 70,050       | 3,154   | True   |
| 5 | unsafe stack            | 63       | 1,493     | 109,610      | 43,177  | False  |
| 6 | safe stack              | 77       | 637       | 70,089       | 7,234   | True   |
| 7 | unsafe stack neutrality | 63       | 5,286     | 383,312      | 167,654 | True   |
| 8 | safe stack neutrality   | 77       | 840       | 70,077       | 16,773  | True   |

Table 2: Results of the evaluation.

#### Generic procedural programs. Formula

$$\Box((\text{call} \wedge p_B \wedge \text{Scall}(\top, p_A)) \implies \text{CallThr}(\top))$$

means that whenever procedure  $p_B$  is executed and at least one instance of  $p_A$  is on the stack,  $p_B$  is terminated by an exception. We checked it against three OPA representing some simple procedural programs with exceptions and recursive procedures. The formula holds on benchmarks no. 1 and 3, but not on no. 2.

**Stack Inspection.** [7] contains an example Java program for managing a bank account, which uses the security framework of the Java Development Kit to enforce user permissions. The program allows the user to check the account balance, and to withdraw money. To perform such tasks, the invoking program must have been granted permissions `CanPay` and `Debit`, respectively. We modeled such program as an OPA (bench. 4), and proved that the program enforces such security measures effectively by checking it against the formula

$$\Box(\text{call} \wedge \text{read} \implies \neg(\top \mathcal{S}_\chi^d(\text{call} \wedge \neg\text{CanPay} \wedge \neg\text{read})))$$

meaning that the account balance cannot be read if some function in the stack lacks the `CanPay` permission (a similar formula checks the `Debit` permission).

**Exception Safety.** [8] is a tutorial on how to make exception-safe generic containers in C++. It presents two implementations of a generic stack data structure, parametric on the element type  $T$ . The first one is not exception-safe: if the constructor of  $T$  throws an exception during a pop action, the topmost element is removed, but it is not returned, and it is lost. This violates the strong exception safety [1] requirement that each operation is rolled back if an exception is thrown. The second version of the data structure instead satisfies such requirement.

While exception safety is, in general, undecidable, it is possible to prove the stronger requirement that each modification to the data structure is only committed once no more exceptions can be thrown. We modeled both versions as OPA, and checked such requirement with the following formula:

$$\Box(\text{exc} \implies \neg((\ominus^u \text{modified} \vee \chi_P^u \text{modified}) \wedge \chi_P^u (\text{Stack} :: \text{push} \vee \text{Stack} :: \text{pop})))$$

POMC successfully found a counterexample for the first implementation (5), and proved the safety of the second one (6).



| Formula  | Time (ms) | Memory (KiB) |         | Result |
|--|-----------|--------------|---------|--------|
|  |           | Tot.         | MC      |        |
| $\chi_F^d p_{Err}$   | 1.1       | 70,095       | 175     | False  |
| $\circ^d(\circ^d(\text{call} \wedge \chi_F^u \text{exc}))$   | 21.0      | 70,095       | 1,290   | False  |
| $\circ^d(\text{han} \wedge (\chi_F^d(\text{exc} \wedge \chi_P^u \text{call})))$  | 42.2      | 70,088       | 2,297   | False  |
| $\square(\text{exc} \implies \chi_P^u \text{call})$  | 10.7      | 70,099       | 839     | True   |
| $\top \mathcal{U}_\chi^d \text{exc}$   | 2.2       | 70,093       | 121     | False  |
| $\circ^d(\circ^d(\top \mathcal{U}_\chi^d \text{exc}))$   | 4.3       | 70,094       | 113     | False  |
| $\square((\text{call} \wedge p_A \wedge (\neg \text{ret } \mathcal{U}_\chi^d \text{WRx})) \implies \chi_F^u \text{exc})$           | 3,257.7   | 238,833      | 102,582 | True   |
| $\circ^d(\circ^u \text{call})$   | 0.7       | 70,094       | 139     | False  |
| $\circ^d(\circ^d(\circ^d(\circ^u \text{call})))$   | 3.4       | 70,108       | 126     | False  |
| $\chi_F^d(\circ^d(\circ^u \text{call}))$   | 1.3       | 70,096       | 137     | False  |
| $\square((\text{call} \wedge p_A \wedge \text{CallThr}(\top)) \implies \text{CallThr}(e_B))$                                       | 7,793.7   | 402,420      | 173,639 | False  |
| $\diamond(\circ_H^d p_B)$  | 2.1       | 70,097       | 114     | False  |
| $\diamond(\circ_H^d p_B)$  | 2.8       | 70,097       | 114     | False  |
| $\diamond(p_A \wedge (\text{call } \mathcal{U}_H^d p_C))$  | 594.9     | 77,806       | 29,786  | True   |
| $\diamond(p_C \wedge (\text{call } \mathcal{S}_H^d p_A))$  | 676.6     | 96,296       | 37,949  | True   |
| $\square((p_C \wedge \chi_F^u \text{exc}) \implies (\neg p_A \mathcal{S}_H^d p_B))$  | —         | —            | —       | OOM    |
| $\square(\text{call} \wedge p_B \implies \neg p_C \mathcal{U}_H^u p_{Err})$  | 198.2     | 70,088       | 10,606  | True   |
| $\diamond(\circ_H^u p_{Err})$  | 1.1       | 70,093       | 114     | False  |
| $\diamond(\circ_H^u p_{Err})$  | 1.2       | 70,089       | 114     | False  |
| $\diamond(p_A \wedge (\text{call } \mathcal{U}_H^u p_B))$  | 10.3      | 70,105       | 115     | False  |
| $\diamond(p_B \wedge (\text{call } \mathcal{S}_H^u p_A))$  | 10.8      | 70,095       | 115     | False  |
| $\square(\text{call} \implies \chi_F^d \text{ret})$  | 3.0       | 70,095       | 112     | False  |
| $\square(\text{call} \implies \neg \circ^u \text{exc})$  | 1.9       | 70,106       | 113     | False  |
| $\square(\text{call} \wedge p_A \implies \neg \text{CallThr}(\top))$   | 110.7     | 70,094       | 4,937   | False  |
| $\square(\text{exc} \implies \neg(\circ^u(\text{call} \wedge p_A) \vee \chi_P^u(\text{call} \wedge p_A)))$                         | 28.9      | 70,095       | 112     | False  |
| $\square((\text{call} \wedge p_B \wedge (\text{call } \mathcal{S}_\chi^d(\text{call} \wedge p_A))) \implies \text{CallThr}(\top))$ | 926.1     | 70,104       | 13,310  | True   |
| $\square(\text{han} \implies \chi_F^u \text{ret})$   | 17.0      | 70,079       | 1,252   | True   |
| $\top \mathcal{U}_\chi^u \text{exc}$   | 7.7       | 70,101       | 121     | True   |
| $\circ^d(\circ^d(\top \mathcal{U}_\chi^u \text{exc}))$   | 44.6      | 70,104       | 2,376   | True   |
| $\circ^d(\circ^d(\circ^d(\top \mathcal{U}_\chi^u \text{exc})))$  | 123.7     | 70,090       | 5,261   | False  |
| $\square(\text{call} \wedge p_C \implies (\top \mathcal{U}_\chi^u \text{exc} \wedge \chi_P^d \text{han}))$                         | 92.9      | 70,096       | 1,346   | False  |
| $\text{call } \mathcal{U}_\chi^d(\text{ret} \wedge p_{Err})$   | 1.8       | 70,107       | 114     | False  |
| $\chi_F^d(\text{call} \wedge ((\text{call} \vee \text{exc}) \mathcal{S}_\chi^u p_B))$  | 10.8      | 70,086       | 117     | False  |
| $\circ^d(\circ^d((\text{call} \vee \text{exc}) \mathcal{U}_\chi^u \text{ret}))$  | 5.3       | 70,094       | 114     | False  |

Table 3: Results of the additional experiments on OPA “generic larger”.

Additionally, we proved that both implementations are *exception neutral* (7, 8), i.e. they do not block exceptions thrown by the underlying types.

## 4.2 Directory opa-more

This directory contains more experiments devised with the purpose of testing all POTL operators, also in order to find the most critical cases. In fact, the complexity of POTL model checking is exponential in the length of the formula. This is of course unsurprising, since it subsumes logics such as LTL and NCTL, whose model checking is also exponential. Actually, model checking is feasible for many specifications useful in practice. There are, however, some cases in which the exponentiality of the construction becomes evident.

In Table 3 we show the results of model checking numerous POTL formulas on one of the OPA representing generic procedural programs. Some of them are checked

| Benchmark name | # states | Time (s) | Memory (KiB) |           | Result |
|----------------|----------|----------|--------------|-----------|--------|
|                |          |          | Total        | MC only   |        |
| doubleHan      | 22       | 52.96    | 2,091,256    | 869,661   | False  |
| jensen         | 1236     | 1.97     | 73,712       | 17,339    | True   |
| simpleExc      | 19       | 65.42    | 3,278,876    | 1,353,000 | False  |
| simpleExcNoHan | 12       | 37.72    | 1,510,524    | 656,422   | False  |
| simpleIfElse   | 28       | 27.62    | 942,280      | 383,231   | False  |
| simpleIfThen   | 28       | 30.67    | 1,046,584    | 415,648   | False  |
| simpleWhile    | 16       | 0.09     | 73,768       | 3,251     | True   |
| stackSafe      | 340      | 31.51    | 653,616      | 265,363   | True   |
| stackUnsafe    | 162      | 16.48    | 532,736      | 224,573   | False  |

Table 4: Results of the evaluation of miniproc files.

very quickly, while others require a long execution time and a very large amount of memory. POMC runs out of memory on one of such formulas. We were able to run it in 367 seconds on a server with a 2.0 GHz 16-core AMD CPU and 500 GB of RAM. If you wish to repeat such experiments, you may run the following commands:

```
$ cd ~/path/to/POMC-sources/eval
$ ./mcbench.py opa-more/generic-larger
```

Of course, a machine with an appropriate amount of RAM is needed.

### 4.3 Directory miniproc

This directory contains a few verification tasks in which the model has been expressed as a MiniProc program. Each file in this directory contains multiple formulas.

`jensen.pomc`, `stackUnsafe.pomc` and `stackSafe.pomc` contain the same tasks as those with the same name described in Section 4.1. This time, however, models are expressed as MiniProc programs, and the resulting OPA contain many more states.

Other files contain simpler programs, checked against all formulas from Table 3.

Table 4 reports the results of such experiments. When more than one formula is checked in a single file, the reported result is True only if all formulas are verified, False if at least one of them is not.

## 5 Source Code

The source code of POMC is contained in the `src/Pomc` directory. We describe the contents of each file below.

**Check.hs** This file contains the data structures and functions that implement the translation of POTL formulas into OPA. The `check` and `fastcheck` functions build the OPA and check for string acceptance. `makeOpa` returns a thunk containing an un-evaluated OPA, which is built on-the-fly while the calling context evaluates the transition functions.

**Encoding.hs** contains a data structure that represents a set of POTL formulas as a bit vector. It is used to encode OPA states in a memory-efficient form in `Check.hs`.

**MiniProc.hs** contains code that translates MiniProc programs into OPA.

**MiniProcParse.hs** contains a parser for MiniProc programs.

**ModelChecker.hs** contains the model checking launcher functions, and a data structure to represent the input OPA to be checked explicitly. It calls `makeOpa` to translate the negation of the specification into an equivalent OPA, creates a thunk representing an un-evaluated intersection of the two OPA, and then uses the reachability algorithm from `Satisfiability.hs` to determine emptiness.

**Opa.hs** contains an implementation of OPA, which is used to test string acceptance.

**OpaGen.hs** contains a simple automated OPA generator (still experimental).

**Parse.hs** contains a parser for POMC input files.

**Potl.hs** defines the datatype for POTL formulas.

**Prec.hs** defines the data type for precedence relations.

**Prop.hs** defines the data type for atomic propositions.

**PropConv.hs** contains some functions useful to change the representation of atomic propositions from strings to unsigned integers. This is used by other parts of the program to achieve better performances, as strings are represented as lists of `char` in Haskell, which is quite inefficient.

**Satisfiability.hs** contains the reachability algorithms used in the model checker to decide OPA emptiness. They can also be used to decide satisfiability of a formula.

**Util.hs** contains various functions used in other parts of the code.

The `test` directory contains regression tests based on the HUnit provider of the Tasty<sup>2</sup> framework. They can be run with

```
$ stack test
```

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## A From MiniProc to OPA

A MiniProc program can be converted to an equivalent OPA or  $\omega$ OPBA. This is done in two stages: first, we build an *extended* OPA whose transitions are labeled with Boolean expressions and assignments; then, we convert such OPA to a normal one, ready for model checking.

### A.1 Extended OPA

Given a MiniProc program  $P$  and the set  $I_P$  of identifiers in  $P$ , we call  $L_P = BExp_P \cup Ass_P$  the set of labels on  $P$ , where  $BExp_P$  and  $Ass_P$  are resp. the sets of Boolean expressions and assignments on  $I_P$ . We build the extended OPA

$$\mathcal{A}_P^E = (\Sigma_P, M_{\text{call}}, Q_P^E, \{q_0\}, \{q_f\}, \delta_P^E)$$

with  $\Sigma_P = \Sigma_{\text{call}} \cup L_P$ .  $Q_P$  and  $\delta_P^E$  are built inductively on the program structure. For each statement  $s$  in  $P$ , we define the set of entry state/label pairs  $En_s \subseteq Q_P \times L_P$ . Each entry state is labeled with an element form either  $BExp_P$  or  $Ass_P$ , but not both.

**Functions** For each function  $f$  in  $P$  we define a set of entry states  $En_f = En_s$ , where  $s$  is the first statement in the function's body; we also add transitions and states  $q_f^l \xrightarrow{\text{ret } f} q_f^r$ , to which we link the last statement in  $f$ , and  $q_f^t \xrightarrow{\text{exc}} q_f^e$ , which implements throw statements.

**Function Call** For a call  $s$  to function  $f$ , we add  $q_s \xrightarrow{\text{call } f} q$  for all  $(q, l) \in En_f$ , and  $q_f^t \xrightarrow{q_s} q_f^r$ , where  $f'$  is the function containing  $s$ . Let  $s'$  be the successor of  $s$ : we add  $q_f^r \xrightarrow{q_s} q$  for all  $(q, l) \in En_{s'}$ .

**Assignments** For each assignment  $s$  we add  $q_s \xrightarrow{\text{stm } s} q_s$ , and set  $En_s = Ex_s = \{(q_s, \top)\}$ . Let  $s'$  be the successor of  $s$ : we add  $q_s \xrightarrow{(q_s, l)} q$  for all  $(q, l) \in En_{s'}$ .

**If-then-else** For each statement  $s$  of the form **if**  $b_s$  **then**  $\{s_1; \dots; s_n\}$  **else**  $\{s_{n+1}; \dots; s_m\}$  we have  $En_s = \{(q, b_s \wedge l) \mid (q, l) \in En_{s_1}\} \cup \{(q, \neg b_s \wedge l) \mid (q, l) \in En_{s_{n+1}}\}$ .

**While** For a statement  $s$  of the form  $\text{while } b_s \{ s_1; \dots; s_n \}$  we set  $En_s = \{(q, b_s \wedge l) \mid (q, l) \in En_{s_1}\} \cup \{(q, \neg b_s \wedge l) \mid (q, l) \in En_{s_{n+1}}\}$ , where  $s_{n+1}$  is the successor of  $s$ . Also, both  $s_{n+1}$  and  $s$  itself are considered as successors of  $s_n$ , and their entry sets are merged.

**Throw** For a throw statement  $s$  in a function  $f$  we just set  $En_s = \{(q_f^t, \top)\}$ .

**Try-Catch** For a statement  $s$  in function  $f$  of the form  $\text{try } \{ s_1; \dots; s_n \} \text{ catch } \{ s_{n+1}; \dots; s_m \}$ , we add a new state  $q_s$  and set  $En_s = \{(q_s, \top)\}$ , and a push transition  $q_s \xrightarrow{\text{han}}^l q$  for each  $(q, l) \in En_{s_1}$  that installs the handler. We first deal with the case when an exception is caught. We add pop transitions  $q_f^e \xrightarrow{q_s}^l q$  for each  $(q, l) \in En_{s_{n+1}}$  that pop the handler when an exception is thrown in the try block, and pass the execution flow to the catch block. Then, statement  $s_m$  is linked to the entry states of  $s'$ , the first statement after  $s$  (how this is done depends on what kind of statement  $s_m$  is). For the case when no exception is thrown, we add a shift transition that simulates a dummy throw statement  $t$  after  $s_n$ , to uninstall the handler. When lowering  $s_n$ , we consider  $t$  as its next statement, add states  $q_t$  and  $q'_t$ , and set  $En_t = \{(q_t, \top)\}$ . Then we add  $q_t \xrightarrow{\text{exc dummy}} q'_t$ , and  $q'_t \xrightarrow{q_s}^l q$  for all  $(q, l) \in En_{s'}$ , which pop the handler and continue the execution with the first statement after  $s$ .

Finally, if  $f_0$  is the first function listed in the MiniProc program, we add transitions  $q_0 \xrightarrow{\text{call } f_0}^l q$  for all  $(q, l) \in En_{f_0}$ , and  $q_{f_0}^r \xrightarrow{q_0} q_f$ .

## A.2 From extended OPA to OPA

We expand states of  $\mathcal{A}_P^E$  with all possible variable valuations, to obtain OPA

$$\mathcal{A}_P = (\Sigma_{\text{call}} \times I_P, M_{\text{call}}, Q_P, \{q_0\} \times \{0, 1\}^{|I_P|}, \{q_f\} \times \{0, 1\}^{|I_P|}, \delta_P),$$

where  $Q_P \subseteq Q_P^E \times \{0, 1\}^{|I_P|}$ . Each state is a pair  $(q, v)$  with  $q \in Q_P^E$  and  $v$  is a bitvector representing a possible valuation of variables that hold in  $q$ . By  $v \models l$  we mean that the variable valuation  $v \in \{0, 1\}^{|I_P|}$  satisfies Boolean expression  $l \in BExp_P$ ; if  $l = (x := e) \in Ass_P$  with  $x \in I_P$  and  $e \in BExp_P$  we mean  $v \models x \iff e$ . By  $\text{vars}(v)$  we denote the set of variables satisfied by  $v \in \{0, 1\}^{|I_P|}$ . We define  $Q_P := \bigcup_{i \in \mathbb{N}} Q_P^i$  inductively through the following equations:

$$\begin{aligned} Q_P^0 &:= \{q_0\} \times \{0, 1\}^{|I_P|} \\ Q_P^{n+1} &:= \{(q, v) \mid q' \in Q_P^n, (q', a, l, q) \in \delta_P^E, v \models l\} \end{aligned}$$

This is implemented through a depth-first visit of  $\mathcal{A}_P^E$ , from which we derive

$$\begin{aligned} \delta_P &:= \{q \xrightarrow{a \text{ vars}(v)} q' \mid q \xrightarrow{a, l} q' \in \delta_P^E, q, q' \in Q_P\} \\ &\cup \{q \xrightarrow{a \text{ vars}(v)} q' \mid q \xrightarrow{a, l} q' \in \delta_P^E, q, q' \in Q_P\} \\ &\cup \{q \xrightarrow{p} q' \mid q \xrightarrow{p, l} q' \in \delta_P^E, q, q', p \in Q_P\} \end{aligned}$$

Note that  $\mathcal{A}_P$  has size exponential in  $|I_P|$  in the worst case, but not in general, since only reachable variable assignments are considered.