A look into the Mobile Messaging Black Box

33rd Chaos Commmunication Congress #33c3

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December 22, 2016

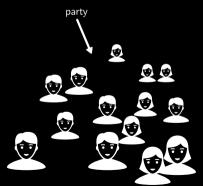
Hamburg University of Technology Security in Distributed Applications

Messaging – Identifying Our Expectations

You're at a party

- · Friend approaches you and needs to tell you something in private
- · What do you expect when you say private?
- · You enter a separate room, you trust the location
- · What does a separate room offer you?





A Private Room

You are now alone in a closed room with your Friend

- · Both of you have absolute Confidentiality that you are alone
- · Nobody can overhear your talk
- · Your exchange is completely private

We call this confidentiality



You Know Each Other

Since you're long-time friends, you're absolutely sure, whom you're talking to

- · Nobody can impersonate your friend or you, without the other noticing
- · You're talking directly, without a phone or webcam in between

We call this authenticity

In Sight of Each Other

The room you're in is small enough that you can always see each other

- · You know that the words you speak are received just as you spoke them
- · There is no way either of you hears something other than the other says

We call this integrity

It's a One-Time Talk

Suppose somebody steps into the room

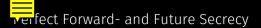
- · They could overhear your conversation
- · They would only learn the contents of this particular conversation
- · They would not learn anything about past conversations you had

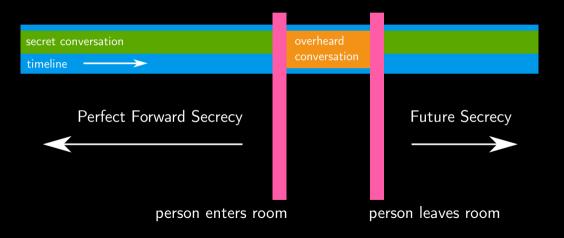
We call this forward secrecy

ightarrow After leaving they would not be able to listen to any future conversations you might have

We call this future secrecy

It's a One-Time Talk





It's a One-Time Talk Between Only You Two

There are no witnesses in the room

- · Either of you can later deny to other having made any statement
- Neither of you can prove to other that any of you have made a particular statement

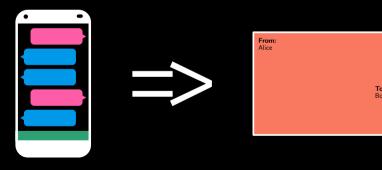
We call this deniability



Messaging – A More Technical Analogy

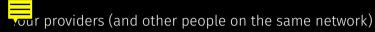
We started with a conversation analogy to identify our expectations of messaging

 \rightarrow Actually postal services are better to look at messaging from a technical point of view.



Example: Traditional Messaging

What if our party conversation had taken place via SMS?



- · would know the contents of your exchange: no confidentiality
- · could change the contents of your exchange: no integrity
- could reroute your messages and impersonate either of you: no authentication
- · would know all messages you ever exchanged: no forward Secrecy
- · would know all messages exchanged in the future: no future secrecy
- could store all messages and use them as proof of the exchange: no deniability
- ightarrow Messaging translates badly to our offline communication expectation $\overline{m{ au}}$

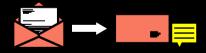
From Postcards to Letters





From Postcards to Letters





Symmetric Encryption:

 \rightarrow Encryption and decryption with the same key



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Asymmetric Encryption:

 \rightarrow Encryption and decryption with different keys



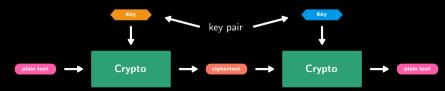
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Public-Key Cryptography – In a Nutshell



Secret Key Public Key

Identity



Secret Key

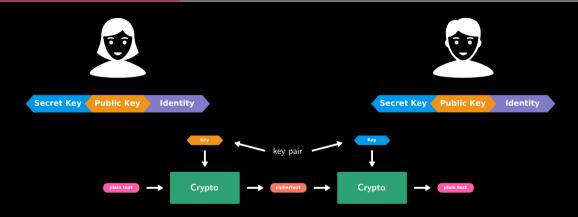
Public Key

Identity

- · Both parties publish their identities and public keys
- Any message can be encrypted with anyone's public key and only be decrypted with its corresponding secret key

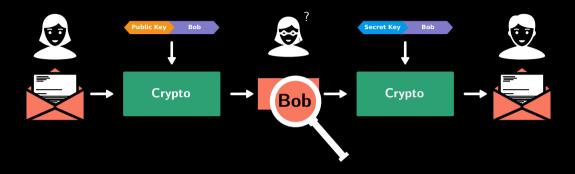


Public-Key Cryptography – In a Nutshell



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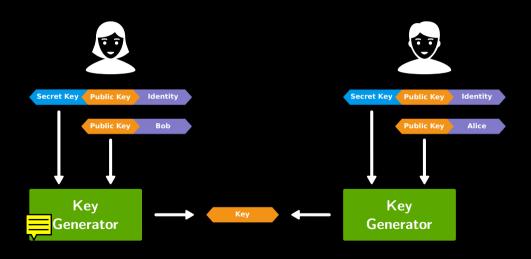
Public-Key Cryptography – In a Nutshell



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Authenticated Enryption



Recap

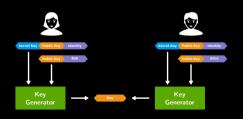
Asymmetric Encryption gives us IDs but is very expensive.



Symmetric Encryption is cheap, but a key has to be shared by all participants before communication starts.

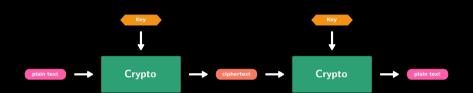


Authenticated Encryption allows us to create symmetric keys based on asymmetric key pairs.



But there's more...

Confidentiality





Deniability

From:

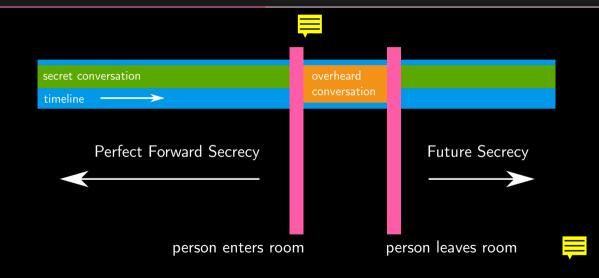
either of us

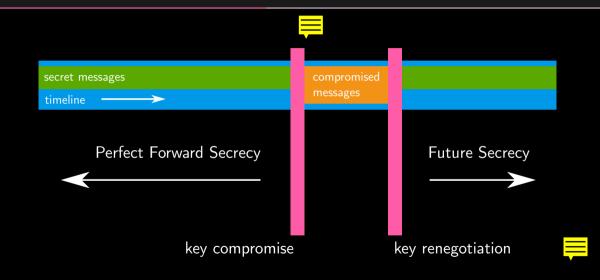


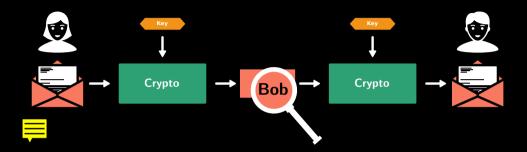
To:

both of us

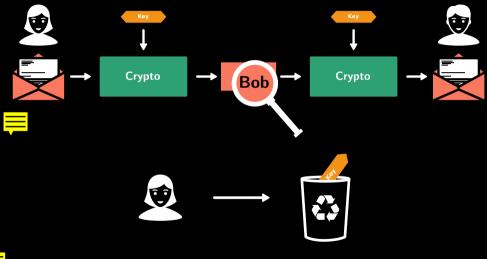














Recap

Authenticated Encryption gives us:

- Confidentiality
- Deniability
- Authenticity

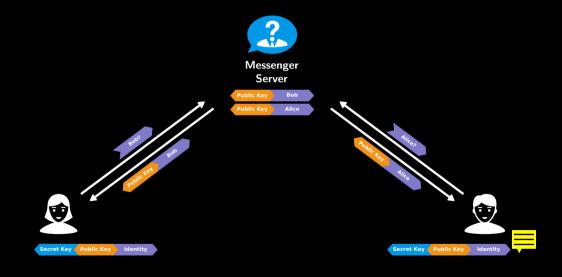
We don't have:

- Perfect Forward Secrecy
- Future Secrecy

 \rightarrow We are ignoring Integrity here, but we have that, too.

Cryptography is rarely, if ever, the solution to a security problem. Cryptography is a translation mechanism, usually converting a communications security problem into a key management problem.

-Dieter Gollmann



- A phone number?
- · An email address?
- Something else?

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 - ightarrow Can identify a user. But is also considered personal information.
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- · Something else?
 - ightarrow Dedicated IDs offer anonymous usage, but ID ownership must be verifyable.
- ightarrow Dedicated IDs are preferrable. But only if we find a way to verify ID ownership

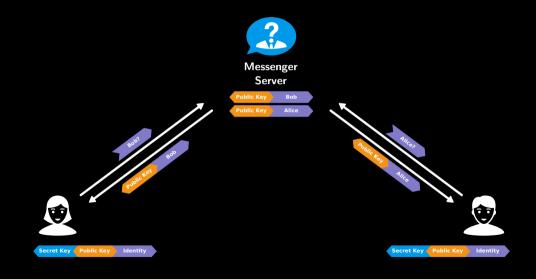
Key and ID Management

How does Alice know which is Bob's public key?

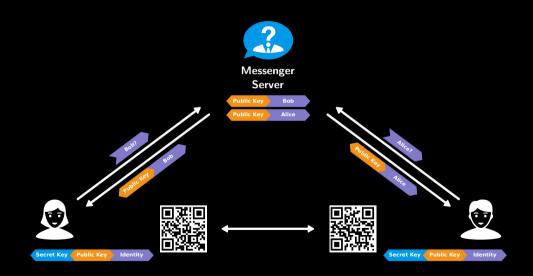




Mobile Messaging Key Management



Mobile Messaging Key Management

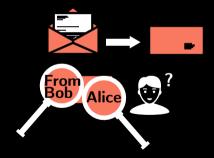


Authenticity

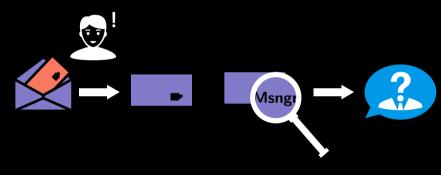
- How to connect a key to a person?
 - → Key signing (PGP)
 - → Certificates (trusted third party)
 - → (Messenger rvice-based directory (based on phone numbers or email addresses)
- How to deal with changing keys?
 - warnings are annoying
 - Threema's traffic light system encourages authentication but doesn't deal with changing keys (other than new identities for known phone numbers with yellow dots)

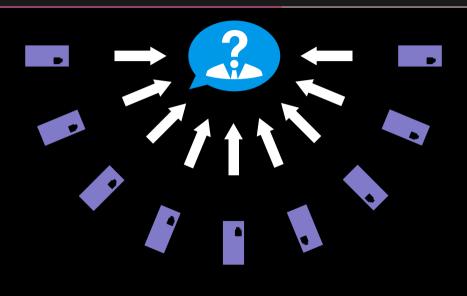
Everybody on the network can see:

- the sender of the message
- \cdot the intended receiver of the message $\overline{\ }$



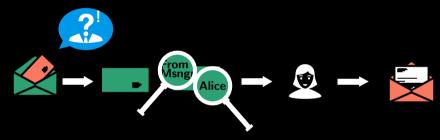
Solution: wrap encrypted message in a second layer of encryption and address it only to the message server.







The message server will remove the outer layer and add a new one, targeted at the receiver.



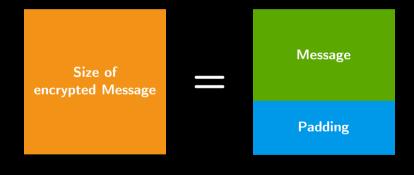
This leaves us with an encrypted end-to-end tunnel, transmitted through two transport layer encryption tunnels.



The message server still knows both communication partners!

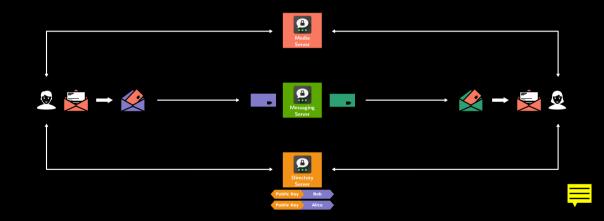
We can obfuscate the size of a message with padding

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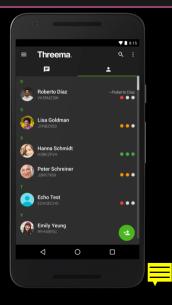




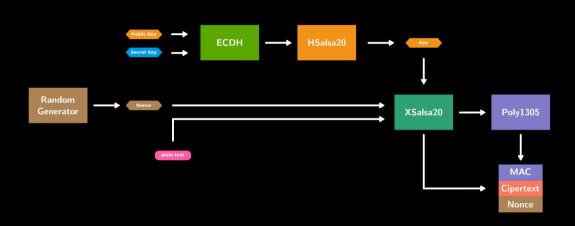
Threema's Architecture



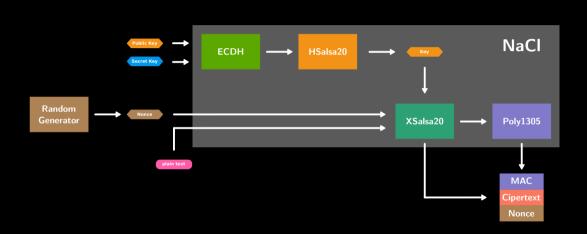
Threema Fingerprints

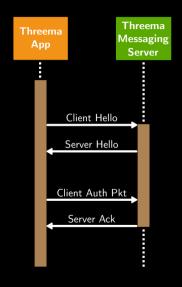


NaCl and Threema

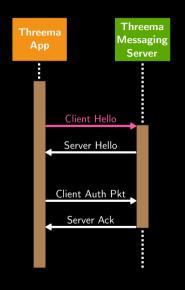


NaCl and Threema





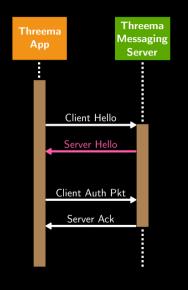
Exchange a set of ephemeral keys and verify each others long term identity keys.



Client Hello Packet



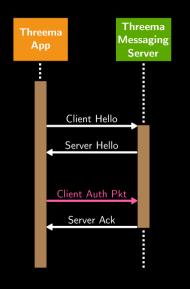
- · Client generates a ephemeral key pair
- · Client generates random nonce prefix



Server Hello Packet



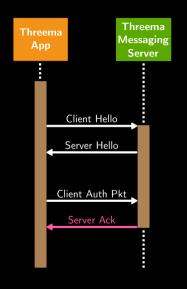
- · Server generates ephemeral key pair
- Server generates random nonce
- Ciphertext encrypted with Server Nonce, Client
 Ephemeral Key and Server Long-Term Key



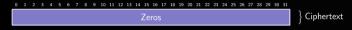
Client Authentication Packet



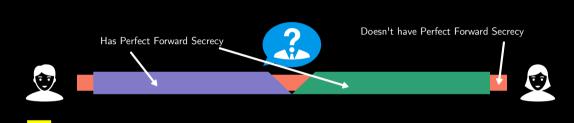
- Outer Encryption with ephemeral Keys
- Ciphertext links clients ephemeral key pair to it's long term key pair



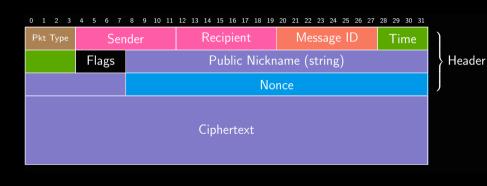
Server Acknowledgement Packet



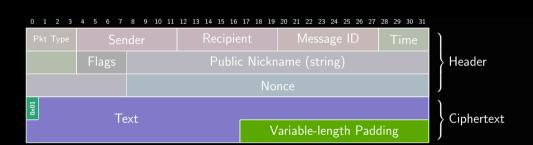
- Server comfirms everything worked fine by encrypting something with both ephemeral keys
- We have established a forward secure channel between app and messaging server.



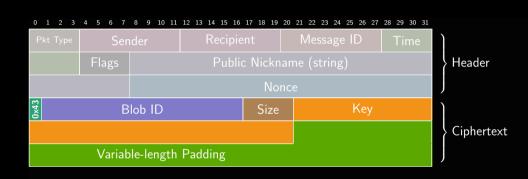
Threema Packet Format

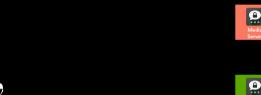


Threema Text Messages



Threema Image Messages



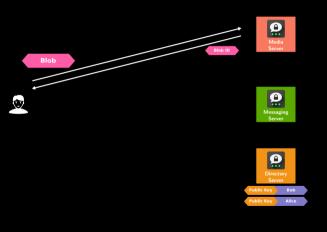




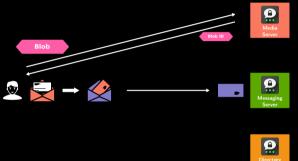






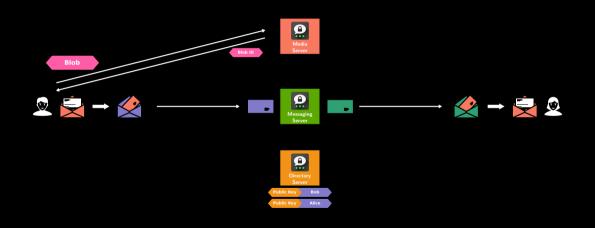


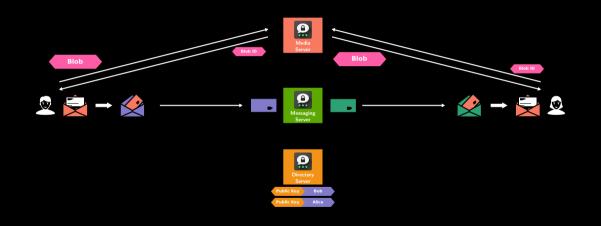








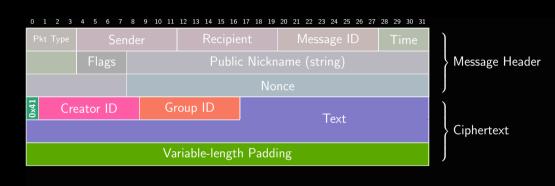






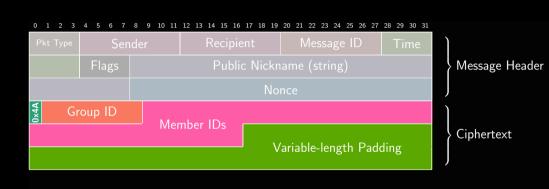
Basic messaging functionality achieved.

Group Messages

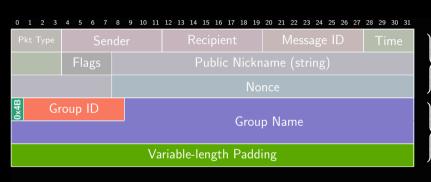




Group Messages



Group Messages



Message Header

Ciphertext



The Devil's in the Detail

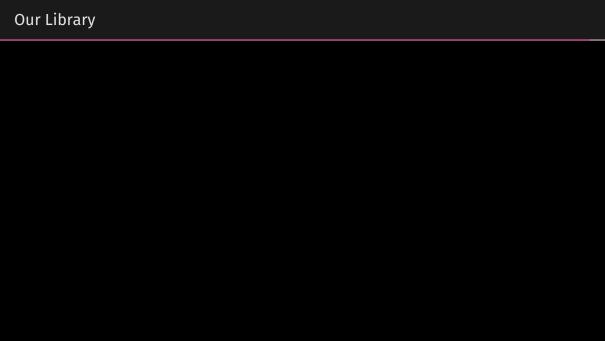
Sammlung kleinerer Dinge, die uns aufgefallen sind

- Media messages could be StageFright attach vectors
- The protocol implementation looks sound to us but the message design prevents feature upgrades on the protocol (not text-protocol) level

Reverse-Engineering – What to look for?

- Test for common pitfalls in implementation
 - Handling of TLS
 - · Handling of keys and nonces
 - NaCl implementation errors
 - · Uncommon data leaks
 - Bugs
 - · ...?
- Find out how protocol is designed
 - 1. Understand handshakes
 - 2. Understand protocol
 - 3. decipher messages

Positive side-note: Threema had released a security white paper early on



Thank You!

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Beamer Theme: Metropolis by Matthias volgelsang

Color Theme: Owl by Ross Chirchley

Icons: The BIG collection by Sergey Demushkin

Foundation Icon Fonts 3 by ZURB

NaCl slide was adapted from a figure in Threema's Cryptography Whitepaper

Thanks to Jan Ahrens and Philipp Berger – their work has made ours somewhat easier Thanks to Maximilian Köstler for his initial work on Threema

Message Packet (Threema Protocol Layer)



• Only the MSB of Flags is used

Message Packet on the Wire



Text Message

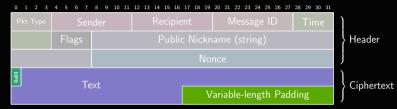
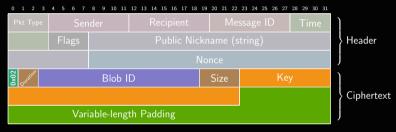


Image Message

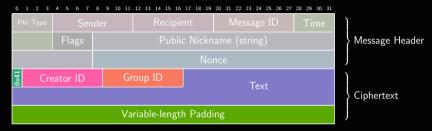


- Blob is symmetrically encrypted using Key and uploaded to asset server.
- Image captions are stored inside the image's EXIF data. These data leak upon creating such an image while the "save media to gallery" option is enabled.

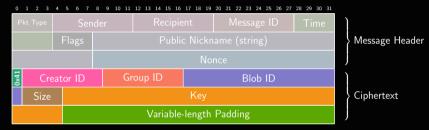
Audio Message



Group Message Packet



Group Image Message



Group Picture Update



Create/Update Group (members)



Acknowledgement Packet to Server

0 1	2	3	4	5	6	7	8	9	10	11	12	13	14	15	16	17	18	19	20	21	22	23	24	25	26	27	28	29	30	31
Length	Pk		Гур	e	Sender								Message ID																	
						+			L	.en	gt	h			\rightarrow															

Client-Server Handshake





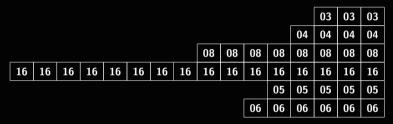
Client Authentication Packet



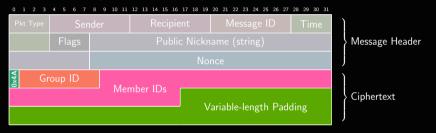
Server Acknowledgement



PKCS7 Padding



Group Management Message - Add Users



Group Management Message - Rename Group

