A look into the Mobile Messaging Black Box

33rd Chaos Commmunication Congress #33c3

Roland Schilling anerdingByDoing Frieder Steinmetz atwillnix

December 26, 2016

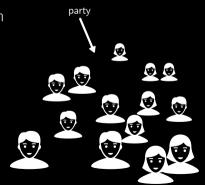
Hamburg University of Technology Security in Distributed Applications

Messaging – Identifying Our Expectations

You're at a party

- · Friend approaches you and needs to tell you something in private
- · What do you expect when you say private?
- · You enter a separate room, you trust the location
- · What does a separate room offer you?





A Private Room

You are now alone in a closed room with your Friend

- · Both of you have absolute Confidentiality that you are alone
- · Nobody can overhear your talk
- · Your exchange is completely private

We call this confidentiality



You Know Each Other

Since you're long-time friends, you're absolutely sure, whom you're talking to

- · Nobody can impersonate your friend or you, without the other noticing
- · You're talking directly, without a phone or webcam in between

We call this authenticity

In Sight of Each Other

The room you're in is small enough that you can always see each other

- · You know that the words you speak are received just as you spoke them
- · There is no way either of you hears something other than the other says

We call this integrity

It's a One-Time Talk

Suppose somebody steps into the room

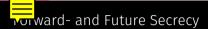
- · They could overhear your conversation
- · They would only learn the contents of this particular conversation
- · They would not learn anything about past conversations you had

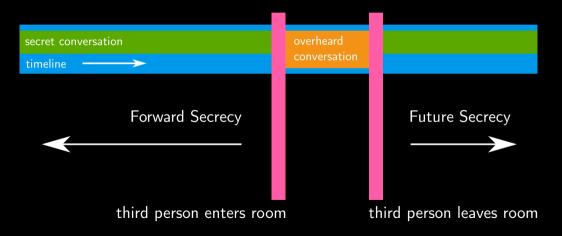
We call this forward secrecy

ightarrow After leaving they would not be able to listen to any future conversations you might have

We call this future secrecy

It's a One-Time Talk





It's a One-Time Talk Between Only You Two

There are no witnesses in the room

- · Either of you can later deny to other having made any statement
- Neither of you can prove to other that any of you have made a particular statement

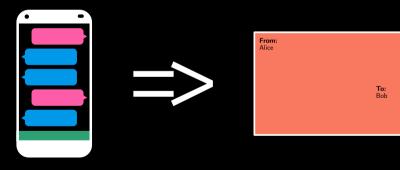
We call this deniability



Messaging – A More Technical Analogy

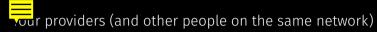
We started with a conversation analogy to identify our expectations of messaging

 \rightarrow Actually postal services are better to look at messaging from a technical point of view.



Example: Traditional Messaging

What if our party conversation had taken place via SMS?



- · would know the contents of your exchange: no confidentiality
- · could change the contents of your exchange: no integrity
- could reroute your messages and impersonate either of you: no authentication
- do not guarantee any secrecy, so we have neither forward secrecy nor future secrecy
- ightarrow We could argue having deniability though.
- ightarrow Messaging translates badly to our offline communication expectation $\overline{m{ au}}$

From Postcards to Letters





From Postcards to Letters





Symmetric Encryption:

 \rightarrow Encryption and decryption with the same key



Symmetric Encryption:

 \rightarrow Encryption and decryption with the same key



Symmetric Encryption:

ightarrow Encryption and decryption with the same key



Asymmetric Encryption:

 \rightarrow Encryption and decryption with different keys



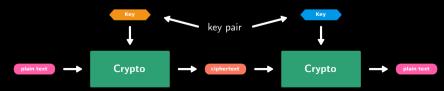
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Public-Key Cryptography – In a Nutshell



Secret Key Public Key

Identity



Secret Key

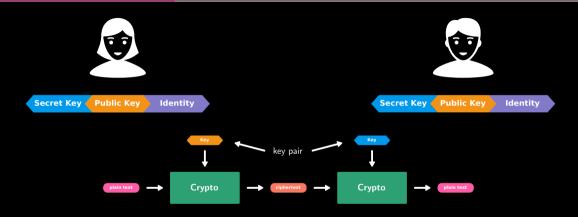
Public Key

Identity

- · Both parties publish their identities and public keys
- Any message can be encrypted with anyone's public key and only be decrypted with its corresponding secret key

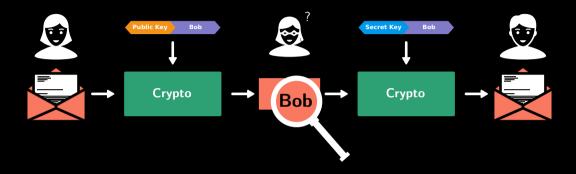


Public-Key Cryptography – In a Nutshell



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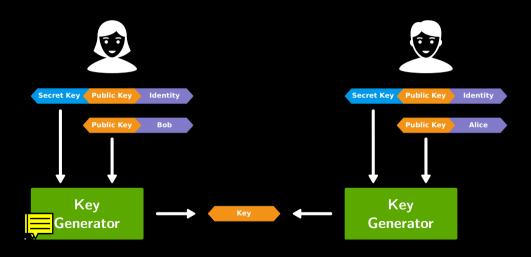
Public-Key Cryptography – In a Nutshell



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Key Establishment



Recap

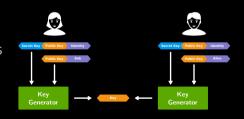
Asymmetric Encryption gives us IDs but is very expensive.



Symmetric Encryption is cheap, but a key has to be shared by all participants before communication starts.



Key Agreement allows us to create symmetric keys based on asymmetric key pairs.



But there's more...

Confidentiality





Deniability

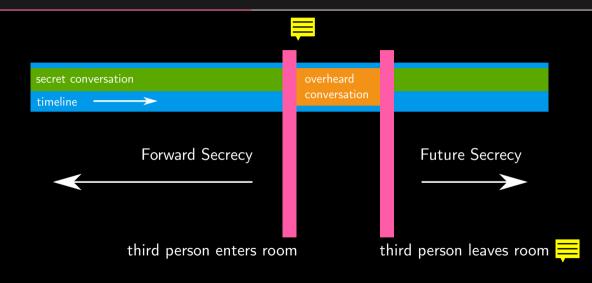
From:

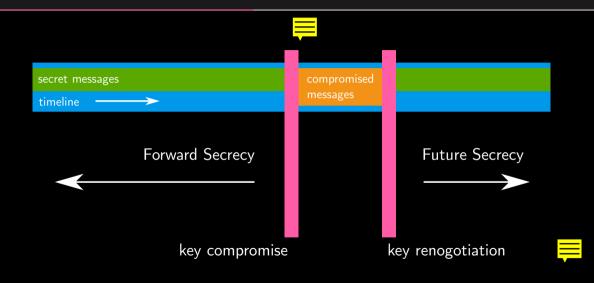
either of us

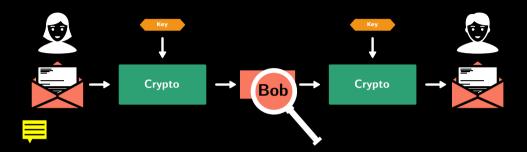


To: both of us

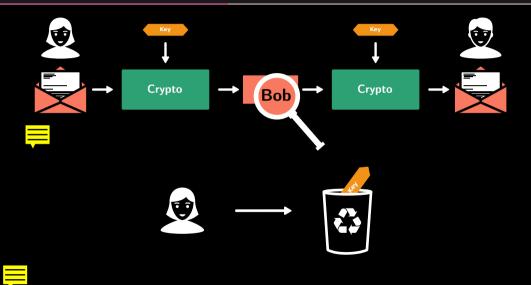












Recap

Our key agreement protocol gives us:

- Confidentiality
- Deniability
- Authenticity

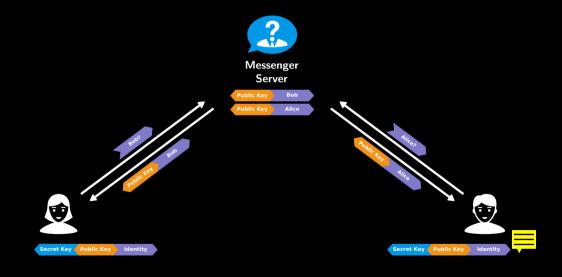
We don't have:

- Forward Secrecy
- Future Secrecy

 \rightarrow We are ignoring Integrity here, but we have that, too.

Cryptography is rarely, if ever, the solution to a security problem. Cryptography is a translation mechanism, usually converting a communications security problem into a key management problem.

-Dieter Gollmann



- · A phone number?
- An email address?
- Something else?

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 - ightarrow Can identify a user. But is also considered personal information.
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 - ightarrow Dedicated IDs offer anonymous usage, but ID ownership must be verifyable.
- ightarrow Dedicated IDs are preferrable. But only if we find a way to verify ID ownership

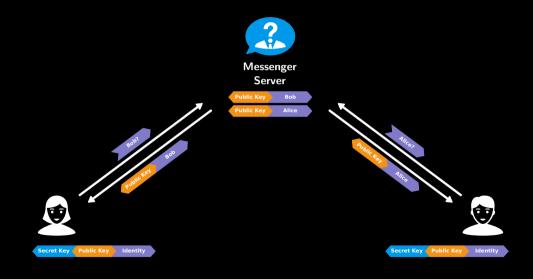
Key and ID Management

How does Alice know which is Bob's public key?

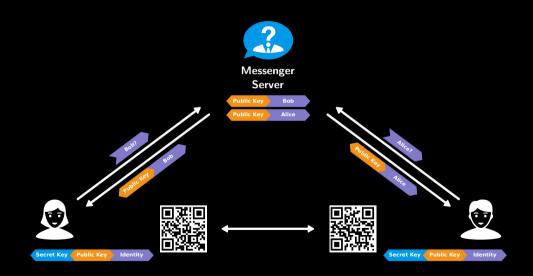




Mobile Messaging Key Management



Mobile Messaging Key Management



Authenticity

We have now solved the Authentiticy problem

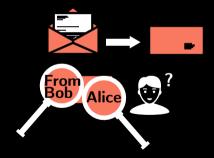
- · User can be identified by their phone number or email address
 - \rightarrow But they have dedicated IDs.
 - \rightarrow Personal verification is possible.

The remaining unsolved problem is a user changing their ID.

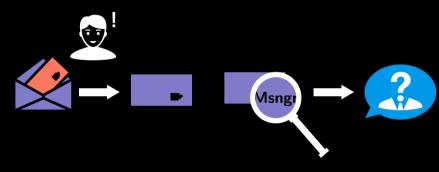
- ightarrow At this point, the problem starts anew.
- \rightarrow We will get back to that later.

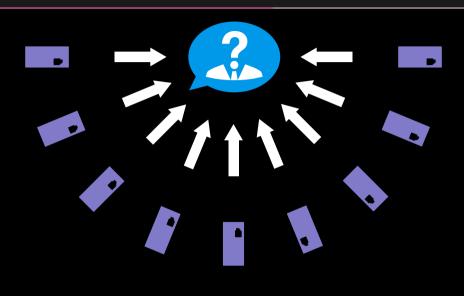
Everybody on the network can see:

- the sender of the message
- \cdot the intended receiver of the message $\overline{\ }$



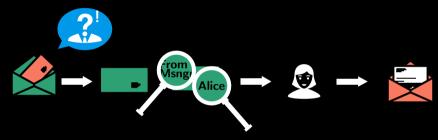
Solution: wrap encrypted message in a second layer of encryption and address it only to the message server.







The message server will remove the outer layer and add a new one, targeted at the receiver.



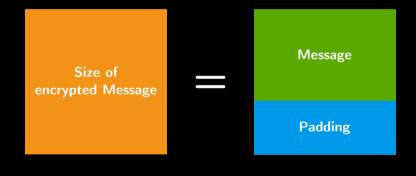
This leaves us with an encrypted end-to-end tunnel, transmitted through two transport layer encryption tunnels.



The message server still knows both communication partners!

We can obfuscate the size of a message with padding

We can obfuscate the size of a message with padding



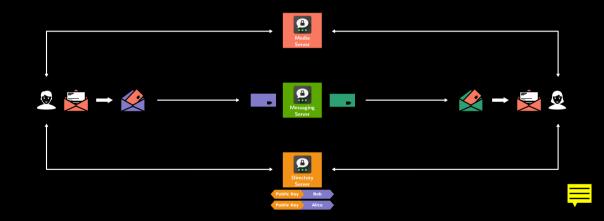


Threema

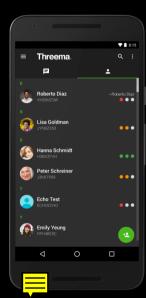




Threema's Architecture



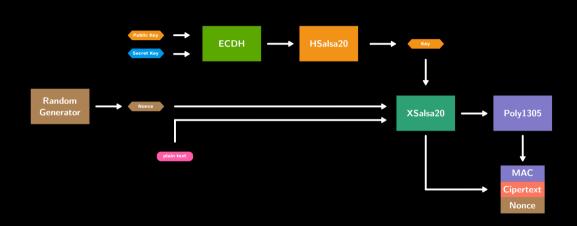
Threema Fingerprints



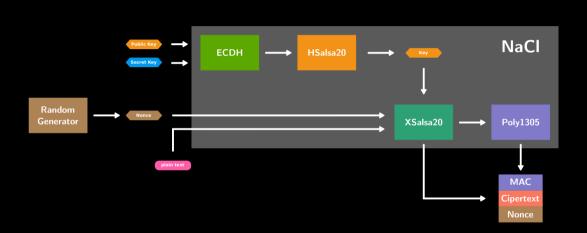
Threema offers dedicated IDs

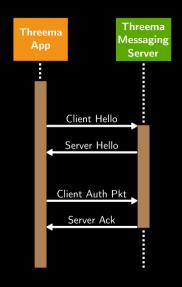
- · Users may provider their phone number and email.
- If provided, phone number and email are used for identification with the directory server.
- If no additional data is provided, IDs can only be exchanged manually.
- In either case, manual verification using QR codes is encouraged.
- The app permanently tracks the verification status of each peer ID.

NaCl and Threema

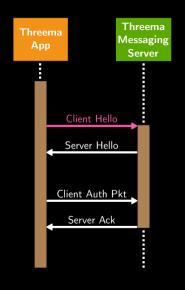


NaCl and Threema





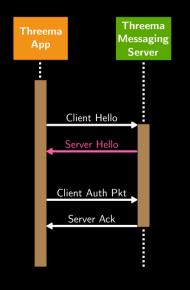
Exchange a set of ephemeral keys and verify each others long term identity keys.



Client Hello Packet



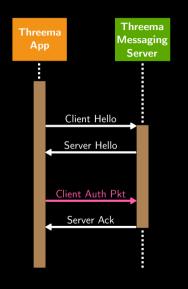
- · Client generates a ephemeral key pair
- · Client generates random nonce prefix



Server Hello Packet



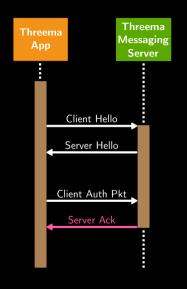
- · Server generates ephemeral key pair
- · Server generates random nonce
- Ciphertext encrypted with Server Nonce, Client
 Ephemeral Key and Server Long-Term Key



Client Authentication Packet



- Outer Encryption with ephemeral Keys
- Ciphertext links clients ephemeral key pair to it's long term key pair



Server Acknowledgement Packet

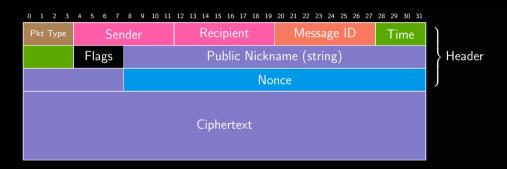


- Server comfirms everything worked fine by encrypting something with both ephemeral keys
- We have established a forward secure channel between app and messaging server.

A 2-Layer Tunnel

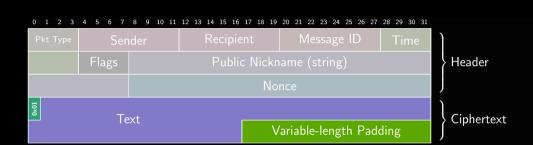


Threema Packet Format

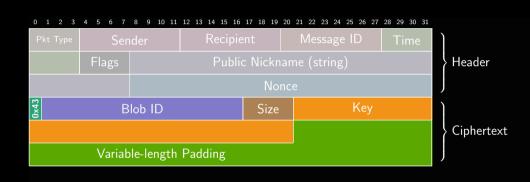


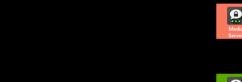


Threema Text Messages



Threema Image Messages



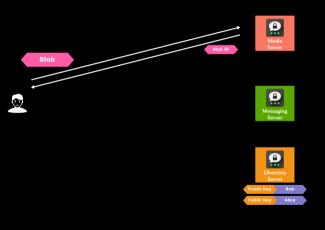




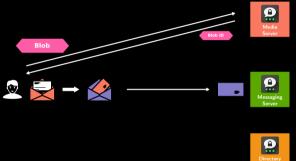






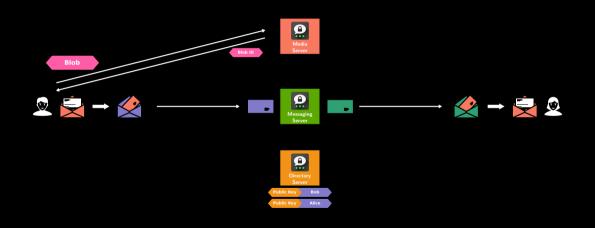


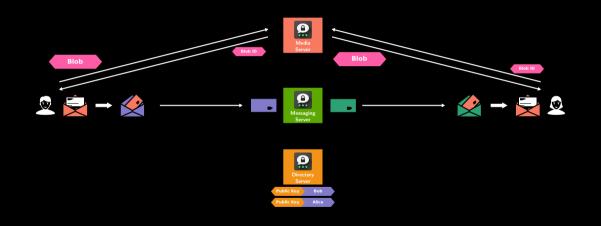














Basic messaging functionality achieved.

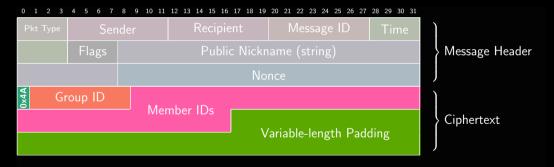
Group Messages





Group Management Messages

Group creation message



Group Management Messages

Group rename message



Implementation of Addon Features

Captions in Image Messages

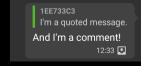




Implementation of Addon Features

Quoted Messages

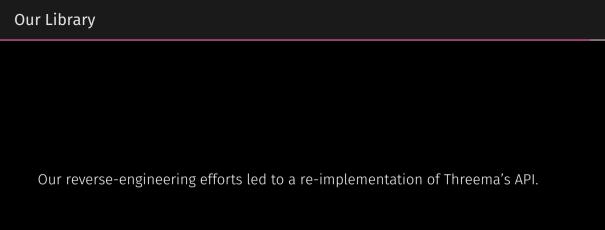




The Devil's in the Detail

Sammlung kleinerer Dinge, die uns aufgefallen sind

- Media messages could be StageFright attack vectors
- The protocol implementation looks sound to us but the message design prevents feature upgrades on the protocol (not text-protocol) level
- Retention policies for uploaded Blobs



Thank You!

Roland Schilling

x schilling@tuhh.de

∭ aNerdingByDoing

Frieder Steinmetz

├── frieder.steinmetz@tuhh.de

⋰ atwillnix

Beamer Theme: Metropolis by Matthias volgelsang

Color Theme: Owl by Ross Chirchley

Icons: The BIG collection by Sergey Demushkin Foundation Icon Fonts 3 by ZURB

NaCl slide was adapted from a figure in Threema's Cryptography Whitepaper

Threema Screenshots taken from the Threema press package

Thanks to Jan Ahrens and Philipp Berger – their work has made ours somewhat easier Thanks to Maximilian Köstler for his initial work on Threema

Message Packet (Threema Protocol Layer)



• Only the MSB of Flags is used

Message Packet on the Wire



Text Message

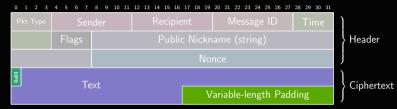
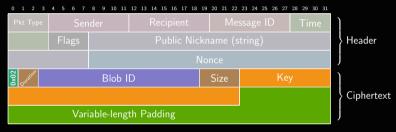


Image Message

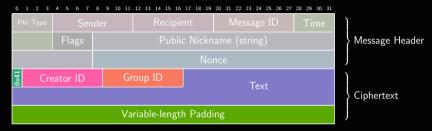


- Blob is symmetrically encrypted using Key and uploaded to asset server.
- Image captions are stored inside the image's EXIF data. These data leak upon creating such an image while the "save media to gallery" option is enabled.

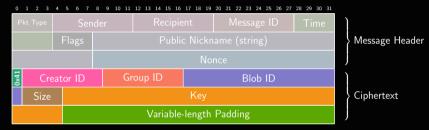
Audio Message



Group Message Packet



Group Image Message



Group Picture Update



Create/Update Group (members)



Acknowledgement Packet to Server

0	1	2	3	4	5	0	7	8	9	10	11	12	13	14	15	16	17	18	19	20	21	22	23	24	25	26	27	28	29	30	31
Len	gth	Pkt Type				Sender								Message ID																	
							+			L	.en	gt	h			\rightarrow															

Client-Server Handshake





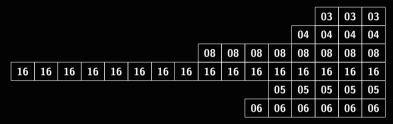
Client Authentication Packet



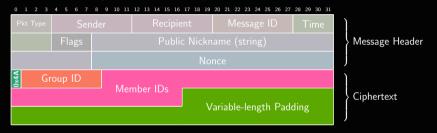
Server Acknowledgement



PKCS7 Padding



Group Management Message - Add Users



Group Management Message - Rename Group



Quoted Text Message

