

EEL 4837

Programming for Electrical Engineers II

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Matrices

Readings:

- Cormen 28
- Weiss 1.7 (using STL)

Matrices & Sparse Matrices

In mathematics, a matrix contains *m* rows and *n* columns of elements. We write *m*×*n* to designate a matrix with *m* rows and *n* columns

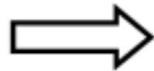
	col 0	col 1	col 2		col 0	col 1	col 2	col 3	col 4	col 5
row 0	-27	3	4	row 0	15	0	0	22	0	-15
row 1	6	82	-2	row 1	0	11	3	0	0	0
row 2	109	-64	11	row 2	0	0	0	-6	0	0
row 3	12	8	9	row 3	0	0	0	0	0	0
row 4	48	27	47	row 4	91	0	0	0	0	0
				row 5	0	0	28	0	0	0
5*3				6*6						
(a)	15/15			(b)	8/36					

A Separate Sparse Matrix Representation

- The standard representation of a matrix is a two dimensional array defined as

$a[\text{MAX_ROWS}][\text{MAX_COLS}]$

- We can locate quickly any element by writing $a[i][j]$
- Sparse matrix wastes space**
 - We must consider alternate forms of representation.
 - Our representation of sparse matrices should store only nonzero elements.
 - Each element is characterized by **<row, col, value>**.

$$\begin{bmatrix} 0 & 0 & 3 & 0 & 4 \\ 0 & 0 & 5 & 7 & 0 \\ 0 & 0 & 0 & 0 & 0 \\ 0 & 2 & 6 & 0 & 0 \end{bmatrix}$$


Row	0	0	1	1	3	3
Column	2	4	2	3	1	2
Value	3	4	5	7	2	6

```
int m[2][3] =  
{  
    {2, 7, 1},  
    {5, 6, 0}  
};
```

```
// print 7  
cout << m[0][1];  
  
// get 2nd row  
m[1];
```

Matrix Algorithms

1. Given a matrix A, write a program to compute the transpose of A.
2. Given two $m \times n$ matrices A and B, write a program to compute the sum $C = A + B$
3. Given an $m \times n$ matrix A and an $n \times p$ matrix B, write a program to compute their product $C = A \times B$
4. Given an $n \times n$ square matrix A, write a program to compute determinant of A.
5. Given an $n \times n$ square matrix A (that is not singular), write a program to compute the inverse A^{-1}

```
for (int i = 0; i < R1; i++)
    for (int j = 0; j < C1; j++) {
        transpose[j][i] = a[i][j];
    }

for (int i = 0; i < R1; i++) {
    for (int j = 0; j < C1; j++) {
        sum[i][j] = a[i][j] + b[i][j];
    }
}

for (int i = 0; i < R1; i++) {
    for (int j = 0; j < C2; j++) {
        prod[i][j] = 0;
        for (int k = 0; k < R2; k++) {
            prod[i][j] += a[i][k] * b[k][j];
        }
    }
}
```

Exercises: Sparse Matrix Algorithms

1. Given a matrix A , write a program to compute the transpose of A .
2. Given two $m \times n$ matrices A and B , write a program to compute the sum $C = A + B$
3. Given an $m \times n$ matrix A and an $n \times p$ matrix B , write a program to compute their product $C = A \times B$

Implement the above when A and B are sparse matrices.

4. Given a matrix A in regular representation, write a program to create a sparse matrix representation of A .
5. Given a matrix A in sparse matrix representation, write a program to generate a regular representation of A .

Exercises: Sparse Matrix Algorithms

1. Given a matrix A, write a program to compute the transpose of A.

Matrix 1: (4x4)

Row	Column	Value
-----	--------	-------

1	2	10
---	---	----

1	4	12
---	---	----

3	3	5
---	---	---

4	1	15
---	---	----

4	2	12
---	---	----

Result of transpose on the first matrix: (4x4)

Row	Column	Value
-----	--------	-------

1	4	15
---	---	----

2	1	10
---	---	----

2	4	12
---	---	----

3	3	5
---	---	---

4	1	12
---	---	----

- We traverse through the matrix element by element
- At each time we replace the pair $\langle \text{row}, \text{column} \rangle$ by the pair $\langle \text{column}, \text{row} \rangle$
 - This will mean that the resulting entries will not be sorted
- Then sort the entries using the pair $\langle \text{row}, \text{column} \rangle$ as key

Exercises: Sparse Matrix Algorithms

2. Given two $m \times n$ matrices A and B, write a program to compute the sum $C = A + B$

Matrix 1: (4x4)

Row Column Value

1	2	10
1	4	12
3	3	5
4	1	15
4	2	12

Matrix 2: (4X4)

Row Column Value

1	3	8
2	4	23
3	3	9
4	1	20
4	2	25

Result of Addition: (4x4)

Row Column Value

1	2	10
1	3	8
1	4	12
2	4	23
3	3	14
4	1	35
4	2	37

- We traverse through the two matrices element by element
- At each time we copy the smaller <row, column> entry to the new array
- If the same entry exists in both arrays, we put the sum of the values
- When one matrix is done, we copy the remaining entries of the other

Solving Linear Equations with Matrices

- We view a linear equation as the matrix form $AX = b$ where A is $n \times n$ matrix and X and b are $n \times 1$ matrices
- Here we will consider square full-rank/invertible matrices
- We can derive X by simply computing the product $A^{-1}b$, but that will be expensive

Linear Equation system

$$\begin{array}{rrrrrrcl} & & 2x_2 & - & 3x_3 & = & 4 \\ -2x_1 & + & x_2 & + & 2x_3 & = & -6 \\ 2x_1 & & & + & x_3 & = & 0 \end{array}$$

1. Transposing equations doesn't change their solution.
2. Scaling an equation doesn't change its solution.
3. If a set of numbers satisfies two equations, then it also satisfies the equation which is one plus a scalar multiple of the other.

Augmented Matrix

$$\left(\begin{array}{ccc|c} 0 & 2 & -3 & 4 \\ -2 & 1 & 2 & -6 \\ 2 & 0 & 1 & 0 \end{array} \right)$$

1. Transposing (switching) rows in an augmented matrix does not change the solution.
2. Scaling any row in an augmented matrix does not change the solution.
3. Adding to any row in an augmented matrix any multiple of any other row in the matrix does not change the solution.

Gaussian Elimination Step 1: Echelon Matrix

Transpose the first and third equations:

$$\begin{pmatrix} 0 & 2 & -3 & 4 \\ -2 & 1 & 2 & -6 \\ 2 & 0 & 1 & 0 \end{pmatrix} \begin{array}{c} \leftarrow \\ \leftarrow \end{array} \rightsquigarrow \begin{pmatrix} 2 & 0 & 1 & 0 \\ -2 & 1 & 2 & -6 \\ 0 & 2 & -3 & 4 \end{pmatrix}$$

Now we can add the first row to the second and get another zero in that column.

$$\begin{pmatrix} 2 & 0 & 1 & 0 \\ -2 & 1 & 2 & -6 \\ 0 & 2 & -3 & 4 \end{pmatrix} \begin{array}{c} \leftarrow \\ \leftarrow + \end{array} \rightsquigarrow \begin{pmatrix} 2 & 0 & 1 & 0 \\ 0 & 1 & 3 & -6 \\ 0 & 2 & -3 & 4 \end{pmatrix}$$

We add (-2) times the second row to the third row.

$$\begin{pmatrix} 2 & 0 & 1 & 0 \\ 0 & 1 & 3 & -6 \\ 0 & 2 & -3 & 4 \end{pmatrix} \begin{array}{c} \leftarrow \\ \leftarrow -2 \\ \leftarrow + \end{array} \rightsquigarrow \begin{pmatrix} 2 & 0 & 1 & 0 \\ 0 & 1 & 3 & -6 \\ 0 & 0 & -9 & 16 \end{pmatrix}$$

The matrix is now in "row echelon" form and we can solve it using back substitution

Gaussian Elimination Step 2: back substitution

Starting with the last matrix above, we scale the last row by $-\frac{1}{9}$:

$$\left(\begin{array}{cccc} 2 & 0 & 1 & 0 \\ 0 & 1 & 3 & -6 \\ 0 & 0 & -9 & 16 \end{array}\right) \mid -\frac{1}{9} \rightsquigarrow \left(\begin{array}{cccc} 2 & 0 & 1 & 0 \\ 0 & 1 & 3 & -6 \\ 0 & 0 & 1 & -\frac{16}{9} \end{array}\right)$$

Now we can zero out the third column above that bottom entry, by adding (-3) times the third row to the second row, then adding (-1) times the third row to the first row.

$$\left(\begin{array}{cccc} 2 & 0 & 1 & 0 \\ 0 & 1 & 3 & -6 \\ 0 & 0 & 1 & -\frac{16}{9} \end{array}\right) \begin{array}{l} \xleftarrow{+} \\ \xleftarrow{+} \\ \xleftarrow{-3} \end{array} \rightsquigarrow \left(\begin{array}{cccc} 2 & 0 & 0 & \frac{16}{9} \\ 0 & 1 & 0 & -\frac{6}{9} \\ 0 & 0 & 1 & -\frac{16}{9} \end{array}\right)$$

The top row can be scaled by $\frac{1}{2}$, and we finally have

$$\left(\begin{array}{cccc} 2 & 0 & 0 & \frac{16}{9} \\ 0 & 1 & 0 & -\frac{6}{9} \\ 0 & 0 & 1 & -\frac{16}{9} \end{array}\right) \mid \frac{1}{2} \rightsquigarrow \left(\begin{array}{cccc} 1 & 0 & 0 & \frac{8}{9} \\ 0 & 1 & 0 & -\frac{6}{9} \\ 0 & 0 & 1 & -\frac{16}{9} \end{array}\right)$$

The matrix is now in "reduced row echelon" form

We can also implement back substitution implicitly without actually creating a full reduced row echelon matrix

Example: Back Substitution Implementation

```
double x[N]; // An array to store solution

// Start calculating from last equation up to the first
for (int i = N-1; i >= 0; i--) {
    // Start with the RHS of the equation
    x[i] = mat[i][N];

    // Initialize column j to i+1 since matrix is upper triangular
    for (int j=i+1; j<N; j++)
    {
        /* Subtract all the lhs values except the coefficient of the variable
           whose value is being calculated */
        x[i] -= mat[i][j]*x[j];
    }

    // Divide the RHS by the coefficient of the unknown being calculated
    x[i] = x[i]/mat[i][i];
}
```

A diagram illustrating the back substitution process on a 3x4 matrix. The matrix is shown as $\begin{pmatrix} 2 & 0 & 1 & 0 \\ 0 & 1 & 3 & -6 \\ 0 & 0 & 1 & -\frac{16}{9} \end{pmatrix}$. A blue 'i' with an upward arrow indicates the current row index, pointing to the first row. A blue 'j' with an upward arrow indicates the current column index, pointing to the third column. A horizontal arrow points from the 'i' to the matrix, and another horizontal arrow points from the 'j' to the matrix.

What is the time complexity?