Designing Data-Intensive Applications

Chapter 11: Streaming Processing

Introduction

- Last chapter we talked about batch processing, which requires finite sized inputs, but in reality many data sources are unbounded and produce data continuously
- Batch processes can run daily, maybe hourly, but some use cases require more frequent updates, and that's where stream processing comes in
- This chapter introduces streaming, talks about how databases have taken a page from streaming, and then covers how stream data is processed

Transmitting Event Streams

- Polling for changes would be too expensive at high frequencies, so we have systems built to handle streaming data
- Each group of related events is called a topic or a stream

Messaging Systems

- A publish/subscribe model allows multiple producers and/or consumers
- Two questions differentiate messaging implementations:
 - What happens if consumers can't process messages as fast as they are produced: drop, buffer, or apply backpressure
 - o If nodes go down, can messages be lost: yes or no
- Direct messaging from producers to consumers is one kind of implementation
 - Options include UDP multicast, TCP or IP multicast, straight UDP, or even HTTP or RPC calls
 - Generally require application to be aware of message failure and implement any fault tolerance support

Message brokers

- Most common is a separate server that runs a message broker or message queue
- Producers and consumers connect as clients
- Queueing means delivery to consumers is asynchronous (a return message can notify the producer if it wants to wait for acknowledgement)
- Some message brokers participate in two-phase commit, but the nature of message queues where delivered messages are typically deletes is different from permanent databases that can be queried
- Multiple consumers can be:
 - Load balancing dividing the work of processing messages
 - o Fan out independent consumers each doing their own thing
- Acknowledgements and redelivery beware that the following can happen:
 - A message that wasn't acknowledged actually was processed, so redelivery to another consumer can cause it to be processed twice

 With load balancing, redelivery can cause messages to be reordered for processing

Partitioned Logs

- Log-based message brokers combine lightweight message delivery with durable log-based durable storage. This includes Apache Kafka and Amazon Kinesis.
- In this structure, producers append to a log and consumers read the log sequentially
- For scalability, the log is usually partitioned. It can also be replicated for fault tolerance.
- Load balancing consumers is usually handled by assigning consumers to partitions. If you want to balance individual messages, a traditional JMS/AMPQ style message broker is easier to implement.
- Consumer offsets instead of individual acknowledgements, the broker only needs to (periodically) track the message offset for each consumer
- Because everything is written to disk, disk space usage needs to be considered. Many systems implement a circular buffer with a fixed maximum amount of storage.
- When consumers cannot keep up with producers typically message loss is possible, although admins can be warned before this happens. A side benefit is that consumers which are shut down don't cause runaway large message queues.
- Replaying old messages is possible. The consumer offset can be intentionally set backward to an earlier value.

Databases and Streams

 A replication log is like a stream of writes, so databases can borrow from streaming to help address heterogeneous systems

Keeping Systems in Sync

- Most of the time multiple systems, such as the OLTP system and data warehouse, and they all need to be kept in sync
- Dual writes by clients can have race conditions, and ensuring both writes commit or abort is the atomic commit issue

Change Data Capture (CDC)

- Instead of treating database replication as a proprietary internal implementation detail, change data capture looks to expose all changes in an externally visible form that can be replicated by other systems
- CDC can provide changes immediately as a stream, and they can be applied by another system, such as a search index, continually
- A log-based message broker can preserve the order of messages, providing the right kind of delivery
- Basically, one database becomes the single leader, and the others act as followers
- Having snapshots to start from is a lot faster for initial synchronization than replaying every change since the beginning of time. Some CDC systems incorporate snapshots, but sometimes you have to do it manually.

- Log compaction is beneficial for keeping down the size of logs, especially if you're going to persist everything
- Newer DBs are supporting CDC functionality, instead of it being a bolt-on afterthought

Event Sourcing

- Event sourcing is similar to CDC in that all changes are captured, but it is involved application design, and events are at a higher level
- The event store can only be appended -- updates or deletes are discouraged or forbidden
- Event sourcing records logical, immutable actions. It usually won't record entire records, like CDC does, so full history may be needed to get the complete current state of a record, and log compaction may not be possible
- Note that all events start as commands, but once all validation is completed, it is now an
 event and is immutable

State, Streams, and Immutability

- The events that change data in a database are an immutable history. You can think of the contents of the database as a cache of the latest values in the log.
- Advantages of immutable events
 - The book has an example that accounting uses and append-only ledger.
 Immutability can also help debugging and provide a richer history for analytics.
 - Deriving several views from the same event log having an explicit process for translating event log entries to the database makes logic explicit, facilitating multiple views. Separating how data is written and read can offer a lot of flexibility.
 - Concurrency control one big negative with CDC and event sourcing is that event log consumers are usually updated asynchronously, so a read after a write may be stale. On the flip side, good self-contained event design may eliminate the need for multi-object transactions.
- Limitations of immutability
 - Workloads with lots of updates and deletes may be hard. Fragmentation and how well compaction and garbage collection perform may be critical.
 - Note also that certain circumstances require deletion, such as privacy rules around someone closing their account

Processing Streams

- Input streams can be processed to create new output streams
- Mapping and filtering work similarly to batch processing, as does partitioning and parallelization

Uses of Stream Processing

- Monitoring and alerting credit card fraud detection, stock market price changes, etc.
- Complex event processing usually a high level declarative language is used to define criteria. Unlike databases, the query persists over time, and the data comes and goes.

- Stream analytics usually about metrics, not events, e.g. number of comments per minute.
 - Note that probabilistic algorithms are often used because they are much cheaper/faster, but are not required
- Maintaining materialized views can keep another copy of data, as has been discussed
- Search on streams while CEP often looks at combinations of events, sometimes you
 just want complex search on individual items, such as monitoring news articles for topics
 of interest
- Message passing and RPC streams provide fault tolerance in a way that RPC doesn't.
 Can have a stream for calls and another stream for returns, and these can scale to multiple nodes.

Reasoning About Time

- Since message delays are unbounded, usually differentiate event time from time message was received/processed
- Knowing when you're ready if you want to process all messages using a window between 1:00 and 1:01, how do you know when you have them all?
 - You can set a timeout and drop stragglers that arrive late (and you can monitor how often you get stragglers)
 - You can output a correction, voiding prior output
- Whose clock are you using, anyway?
 - Because of all the issues with clocks, some systems track three times:
 - The time when the event occurred, per the producer's clock
 - The time when the message was sent, per the producer's clock
 - The time when the message was received, per the broker's clock
 - With fairly short network delays, you can estimate how far apart the producer's clock is by calculating the difference between the last two times
- Types of windows
 - Tumbling window adjacent time slots
 - Hopping window overlapping time slots
 - Sliding window fixed time slot starting with the oldest event in buffer
 - Session window e.g. per user events until a period of inactivity

Stream Joins

- Stream-stream join (window join) typically state is maintained (e.g. a hash index) to join
 events from two different streams, such as searches and clicks mentioned in the book
- Stream-table join (stream enrichment) typically a local copy of the table is kept using change data capture, such as for a user table
- Table-table join (materialized view maintenance) the concept of streams updating a
 materialized view of a join between two tables. Changes to either stream need to inform
 potentially multiple rows from the other table in the join output.
- Time dependence of joins if ordering between two streams is undetermined, joins can be nondeterministic. Using history on a slowly changing dimension can solve determinism, but means you can't do log compaction.

Fault Tolerance

- Microbatching and checkpointing write state to durable storage periodically
- Atomic commit revisited can't easily support full XA across heterogeneous systems, but can provide some support for atomic commit internally between streams
- Idempotence events that are the same if repeated (or smart enough not to repeat even when you try) can be retried without risk
- Rebuilding state after a failure because state is required, must periodically persist it to local or remote storage for fault-tolerance

Summary

- Streams can be transported through direct messaging, message brokers, or event logs
 - AMQP/JMS-style message broker consumers acknowledge individual messages once processed, then they are deleted
 - Log-based message broker consumers read log partitions, and brokers maintain messages on disk
- Streams are a powerful way of integrating all kinds of systems
- Time windowing strategies and three kinds of joins are used in stream processing