

Introduction to Caching

COS 316: Principles of Computer System Design

Amit Levy & Wyatt Lloyd



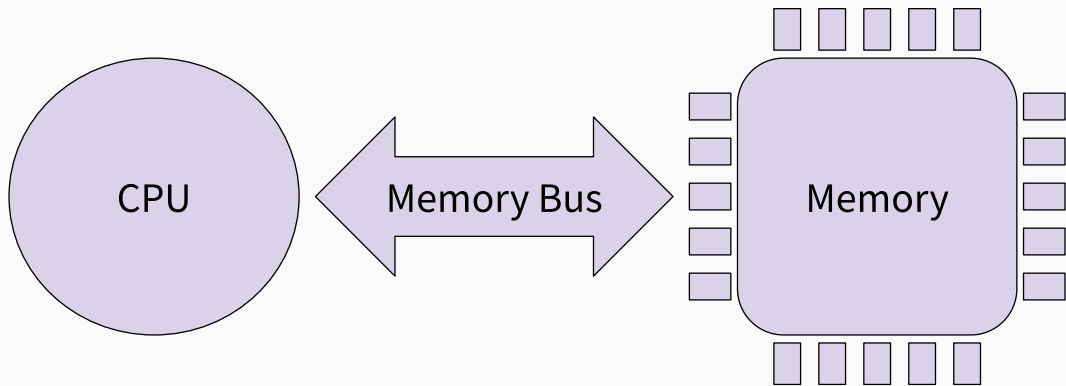


Figure 1: CPU Connected Directly to Memory

How long to run this code?

Characteristics

- CPU Instructions & Register accesses: 0.5ns (2GHz)
- Memory access: 50ns

```
int arr[1000];  
for (i = 0; i < arr.len(); i++) { ++arr[i]; }  
  
    mov    r3, #1000  
loop: ldr    r1, [r0]  
      subs r3, r3, #1  
      add   r1, r1, #1  
      str   r1, [r0], #4  
      bne   <loop>
```

How long to run this code?

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loop: ldr    r1, [r0]
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      bne  <loop>
```

CPU instruction	0.5ns
Memory access	50ns

1. $2.5\mu S$ ($2,505ns$)
2. $250\mu S$ ($250,000ns$)
3. $101.5\mu S$ ($201,505ns$)

How long to run this code?

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    mov    r3, #1000
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Solution

In each loop iteration:

- 2 instructions manipulate registers ($0.5ns$)
- 3 instructions manipulate memory ($100ns$)

$$1*0.5 + 1000*(3*0.5 + 2*50) = 101,505ns$$

Why not just make everything fast?

Type	Access Time	Typical Size	\$/MB
Registers	$< 0.5ns$	~256 bytes	\$1000
SRAM/"Cache"	$5ns$	1-4MB	\$100
DRAM/"Memory"	$50ns$	GBs	\$0.01
Solid state	$20\mu S$	TBs	\$0.0001
Magnetic Disk	$5ms$	10-100s TB	\$0.000001

- High cost of fast storage
- Physical limitations
- Not necessarily possible—e.g. accessing a web page across the world

A Solution: Caching

What is caching?

- Keep *all* data in bigger, cheaper, slower storage
- Keep *copies* of “active” data in smaller, more expensive, faster storage

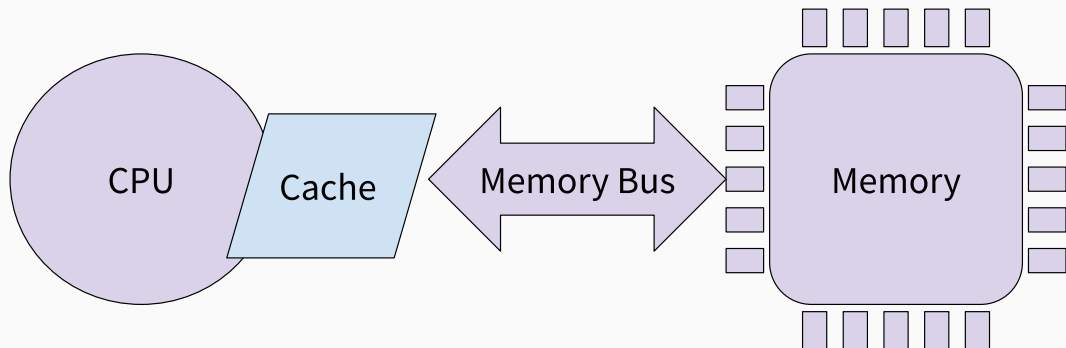


Figure 2: CPU + Cache + Memory

What do we cache?

- Data stored verbatim in slower storage
- Previous computations—recomputation is also a kind of slow storage
- Examples:
 - CPU memory hierarchy
 - File system page buffer
 - Content distribution network
 - Web application cache
 - Database cache
 - Memoization

How long to run this code?

```
    mov    r3, #1000
loop: ldr    r1, [r0]
      subs  r3, r3, #1
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      bne   <loop>
```

CPU instruction	0.5ns
CPU cache access	5ns
Memory access	50ns

1. $2.5\mu S$
2. $11.5\mu S$
3. $101.5\mu S$

How long to run this code?

```
    mov    r3, #1000
loop: ldr    r1, [r0]
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```

1. $2.5\mu S$
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CPU instruction	0.5ns
CPU cache access	5ns
Memory access	50ns

It's complicated!

We don't have enough information to answer. Yet!

- **Hit**: when a requested item was in the cache
- **Miss**: when a requested item was *not* in the cache
- **Hit ratio** and **Miss ratio**: proportion of hits and misses, respectively
- **Hit time** and **Miss time**: time to access item in cache and not in cache, respectively

When is caching effective?

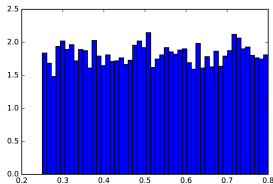
Which of these workloads could we cache effectively?

Repeated Access



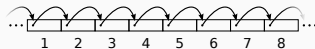
A few popular items
E.g. most social media

Random Access



No pattern to accesses
E.g. large hash tables

Sequential access



Access items in order
E.g. streaming a video

- Temporal locality: nearness in time
 - Data accessed now probably accessed recently
 - Useful data tends to continue to be useful
- Spatial locality: nearness in name
 - Data accessed now “near” previously accessed data
 - Memory addresses, files in the same directory, frames in a video...

Effective access time is a function of:

- Hit and miss ratio
- Hit and miss times

$$t_{effective} = (hit_ratio)t_{hit} + (1 - hit_ratio)t_{miss}$$

aka, Average Memory Access Time (AMAT)

- *Effective access time*
- Look-aside vs. Look-through
- Write-through vs. Write-back
- Write-allocation
- Eviction Policy

Who handles misses?

What happens when a requested item is not in the cache?

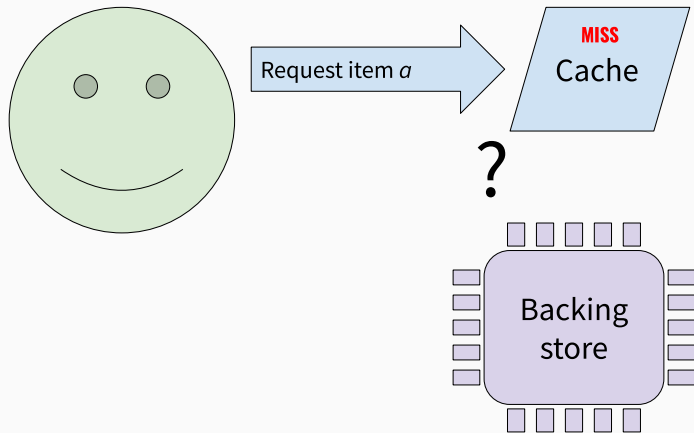


Figure 3: User requests an item not in the cache

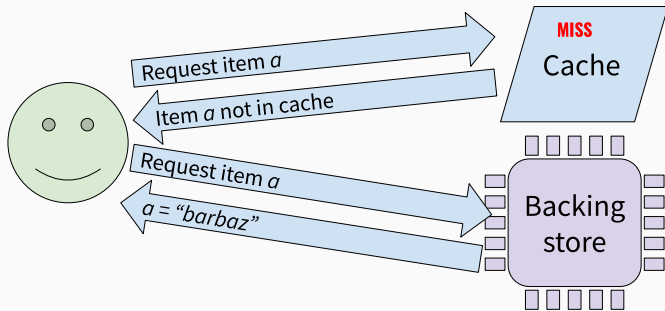


Figure 4: Look-aside Cache

- Advantages: easy to implement, flexible
- Disadvantages: application handles consistency, can be slower on misses

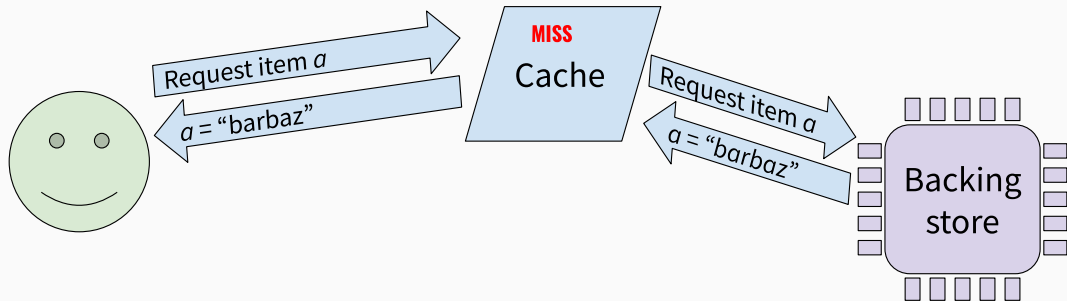


Figure 5: Look-through Cache

- Advantages: helps maintain consistency, simple to program against
- Disadvantages: harder to implement, less flexible

- Caching creates a replica/copy of the data
- When you write, the data needs to be synchronized *at some point*
 - But when?

Write to backing store on every update

- Advantages:
 - Cache and memory are always consistent
 - Eviction is cheap
 - Easy to implement
- Disadvantages:
 - Writes are at least as slow as writes to the backing store

Update only in the cache. Write “back” to the backing store only when evicting item from cache

- Advantages:
 - Writes always at cache speed
 - Multiple writes to same item combined
 - Batch writes of related items
- Disadvantages:
 - More complex to maintain consistency
 - Eviction is more expensive

When writing to items *not* currently in the cache, do we bring them into the cache?

Yes == Write-Allocate

- Advantage: Exploits temporal locality: written data likely to be accessed again soon

No == Write-No-Allocate

- Advantage: Avoids spurious evictions if data is not accessed soon

Eviction policies

Which items to we evict from the cache when we run out of space?

Many possible algorithms:

- Least Recently Used (LRU), Most Recently Used (MRU)
- Least Frequently Used (LFU)
- First-In-First-Out (FIFO), Last-In-First-Out (LIFO)
- ...

Deciding factors include:

- Workload
- Performance

Challenges in Caching

- Speed: making the cache itself fast
- Cache Coherence: dealing with out-of-sync caches
- Performance: maximizing hit ratio
- Security: avoiding information leakage through the cache

Remainder of this Section

- Caching in the CPU Memory Hierarchy
- CDN Caching
- From the research: Learning Relaxed Belady
- Next assignment: in-memory Web application cache

