Consistency



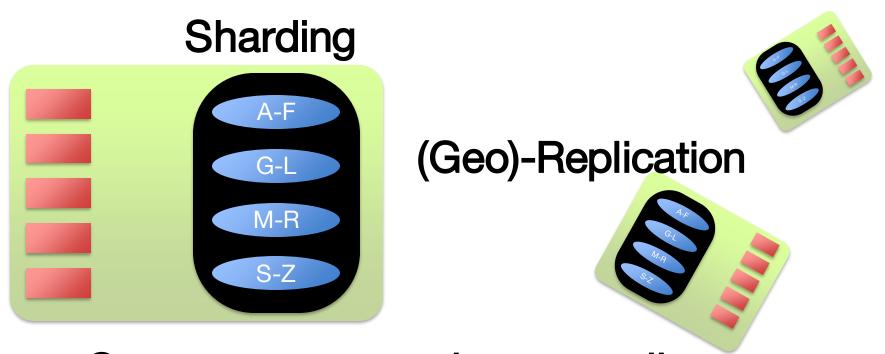
COS 316: Principles of Computer System Design

Amit Levy & Ravi Netravali

Why Do We Build Systems?

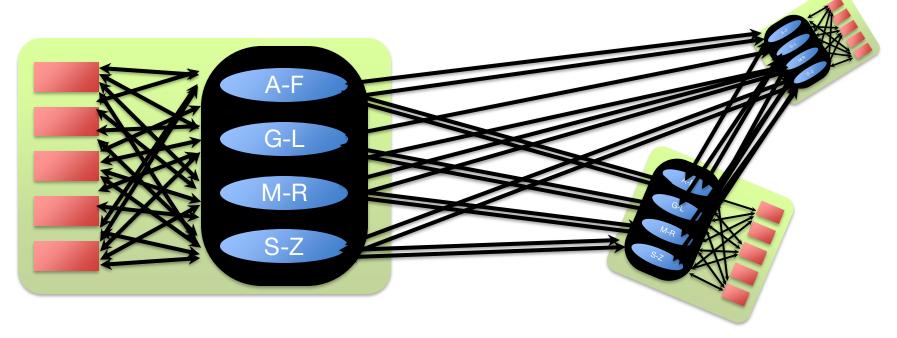
- ...
- Abstract away complexity

Distributed Systems are Highly Complex Internally



Concurrent access by many client

Distributed Systems are Highly Complex Internally Sharding, Geo-Replication, Concurrency



Distributed Systems are Highly Complex Internally Sharding, Geo-Replication, Concurrency

Consistency Models:

Control how much of this complexity is abstracted away

Consistency Models

 Contract between a (distributed) system and the applications that run on it

 A consistency model is a set of guarantees made by the distributed system

Stronger vs Weaker Consistency

Application Code

Strongly Consistent Distributed System

Application Code

Weakly Consistent Distributed System

Stronger vs Weaker Consistency

- Stronger consistency models
 - + Easier to write applications
 - System must hide many behaviors
 - Might be slow
- Fundamental tradeoffs between consistency & performance
 - (Discuss CAP, PRAM, SNOW in 418!)
- Weaker consistency models
 - Harder to write applications
 Cannot (reasonably) write some applications
 - + System needs to hide few behaviors
 - + Can be faster!

Consistency Hierarchy

Linearizability

Causal+ Consistency

Eventual Consistency

Behaves like a single machine

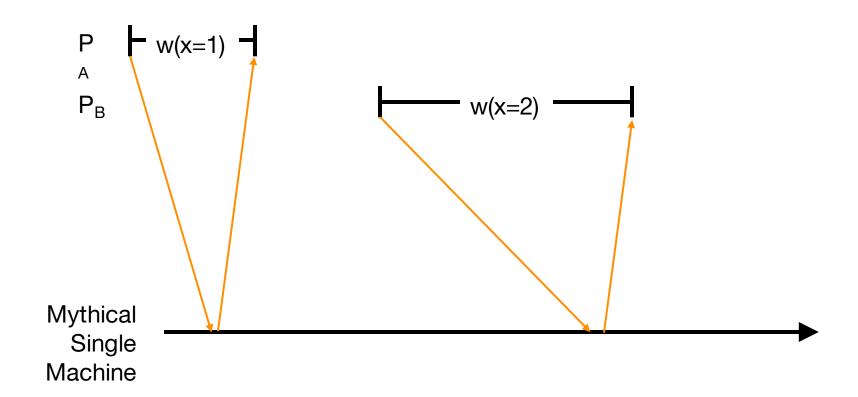
Everyone sees related operations in the same order

Anything goes

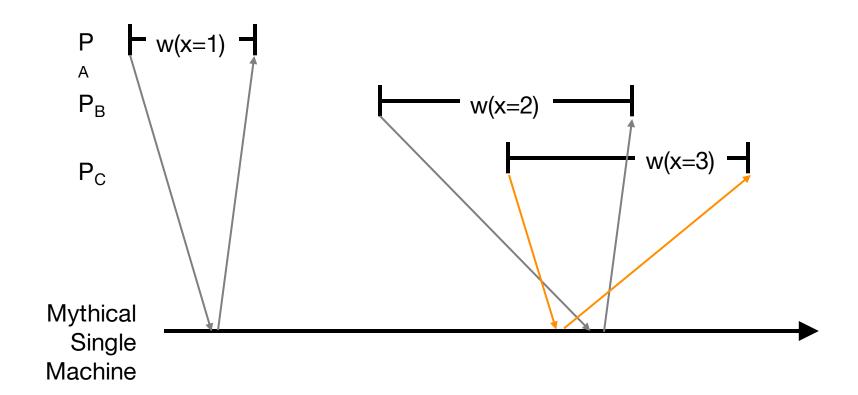
Linearizability == "Appears to be a Single Processor"

- External client submitting requests and getting responses from the system can't tell this is not a single processor!
- Consistent with some total order over all operations
 - As though all requests processed one by one in some order
 - Such that...
- Order preserves the real-time ordering between operations
 - If operation A completes before operation B begins, then A is ordered before B in real-time
 - If neither A nor B completes before the other begins, then there is no real-time order
 - (But there must be some total order)

Real-Time Ordering Examples



Real-Time Ordering Examples

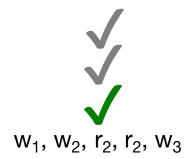


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P | + w(x=1) + + w(x=2) + + w(x=3) +
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\vdash w(x=3) \dashv
P<sub>D</sub>
\vdash r(x)=2 \dashv \vdash r(x)=3 \dashv
P<sub>D</sub>
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P \vdash w(x=1) \dashv
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P<sub>D</sub>
\vdash r(x)=2 \dashv \vdash r(x)=2 \dashv
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        r(x)=2 r(x)=2
P_{D}
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P
$$\vdash$$
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P \vdash w(x=1) \dashv \vdash w(x=2) \dashv \vdash w(x=3) \dashv
                   P_D
                                     - w(x=6) -
 P_{E}
              r(x)=2 r(x)=5 r(x)=6 r(x)=5
 P_{G}
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P_D
                 - w(x=6) -
P_{E}
      r(x)=4 r(x)=2 r(x)=3 r(x)=6
P_{H}
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 W_1 , W_4 , r_4 , W_2 , r_2 , W_3 , r_3 , W_5 , W_6 , r_6

```
P \vdash w(x=1) \dashv P \vdash w(x=2) \dashv P \vdash r(x)=1 \dashv X
```

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How to Provide Linearizability?

- 1. Use a single machine 😊
- 1. Use "state-machine replication" on top of a consensus protocol like Paxos
 - Distributed system appears to be single machine that does not fail!!
 - Covered extensively in 418
- 2. ...

Consistency Hierarchy

Linearizability

Causal+ Consistency

Eventual Consistency

Behaves like a single machine

Everyone sees related operations in the same order

Anything goes

Causal+ Consistency Informally

1.

• Potential causality: event a could have a causal effect on event b.

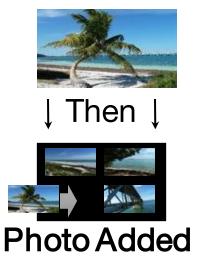
2.

- Think: is there a path of information from a to b?
 - a and b done by the same entity (e.g., me)
 - a is a write and b is a read of that write
 - + transitivity

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Causal+ Sufficient









Deletion retained

Causal+ Not Sufficient

(Need Linearizability)

- Need a total order of operations
 - e.g., Alice's bank account ≥ 0

- Need a real-time ordering of operations
 - e.g., Alice changes her password, Eve cannot login with old password

Consistency Hierarchy

Causal+ Consistency

Eventual Consistency

Behaves like a single machine

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Anything goes

Eventual Consistency

- Anything goes for now...
 - (If updates stop, eventually all copies of the data are the same)
- But, eventually consistent systems often try to provide consistency and often do
 - e.g., Facebook's TAO system provided linearizable results 99.994% of the time [Lu et al. SOSP '15]
- "Good enough" sometimes
 - e.g., 99 vs 100 likes

Consistency Model Summary

- Consistency model specifies strength of abstraction
 - Linearizability 2 Causal+ 2 Eventual
 - Stronger hides more, but has worse performance
- When building an application, what do you need?
 - Select system(s) with necessary consistency
 - Always safe to pick stronger
- When building a system, what are your guarantees?
 - Must design system such that they always hold
 - Must confront fundamental tradeoffs with performance
 - What is more important?