NetCache: Balancing Key-Value Stores with Fast In-Network Caching

COS 316: Principles of Computer System Design

Lecture 10

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What is a Key-Value Store?

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Key-Value Store: A Very Simple Database

- Hash table
  - key: an arbitrary string (e.g., “Joe”, ”jrex.jpg”)
  - value: unstructured data (e.g., number, string, complex object)
- Primitive operations
  - put(key, value)
  - value = get(key)
  - delete(key)
- A “dictionary”
  - Array, map…

key	value
“system”	“a set of things working together as parts of a mechanism or an interconnecting network.”
“internet”	“the single worldwide computer network that interconnects other computer networks”
“database”	“a comprehensive collection of related data organized for convenient access, generally in a computer.”
“cache”	“ a temporary storage space or memory that allows fast access to data”

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Applications of Key-Value Stores

- When to use a key-value store
  - Real-time, random data access
  - Caching for frequently accessed data
  - Simple applications
- Many example applications
  - Online shopping: user id (key) and shopping cart (value)
  - User profile: user id (key) and preferences (value)
  - Movies/music: file name (key) and file contents (value)
  - Blockchain: hash value (key) and block entity (value)

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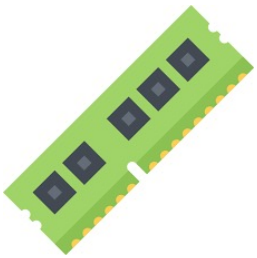
High-Speed Key-Value Store

- A need for speed
  - "web pages routinely fetch thousands of key-value pairs" [Facebook]
  - Systems with many users
- A key-value store needs resources
  - Storage: to store many key-value pairs
  - Bandwidth: to handle many operations
- In-memory key-value store
  - Avoid reads and writes to slow disk-based systems
  - But, what if you need more storage and/or bandwidth?



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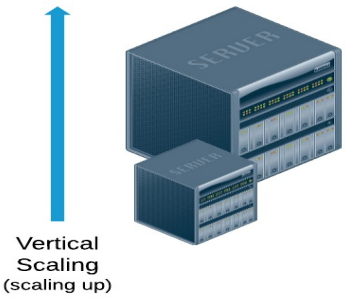
Distributed, High-Speed Key-Value Store

- Vertical scaling
  - Add more resources to a server
  - E.g., more processing, memory, or network resources
- Horizontal scaling
  - Add more servers
  - Reduce the responsibilities of each server



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Horizontal Scaling: Partitioning the Keys

- **Sharding**: breaking data up into partitions
  - Each server handles operations for a portion of the keys
  - Reduces query load and storage space on individual servers

key	value
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Server 0

Server 1

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Goals in Partitioning the Keys

- Easy to direct requests to the right server
- Even distribution of data storage
- Even distribution of request load

- E.g., range of keys
  - Alphabet: letters A-I, J-R, S-Z
  - Months: Jan-Apr, Jun–Sep, Oct-Dec

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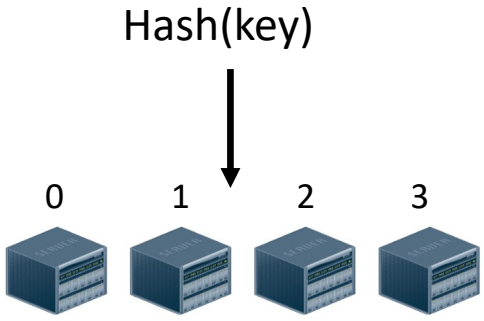
Better Design: Hash-Based Sharding

- Hashing on the key
  - Randomize and distribute requests
  - Easy to compute without coordination



- Achieves the goals?
  - Easy to direct requests to the right server?
  - Even distribution of data storage?
  - Even distribution of request load?

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Distributed, In-Memory Key-Value Stores Need a Cache

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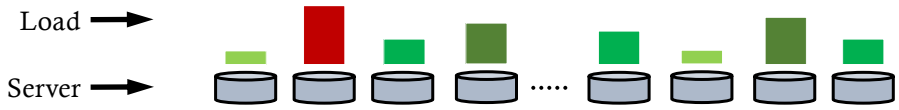
Challenge: Skewed Workload

low throughput & high tail latency



- Popular items receive far more queries than others
- E.g., "10% of items account for 60-90% of queries in the Memcached deployment at Facebook"

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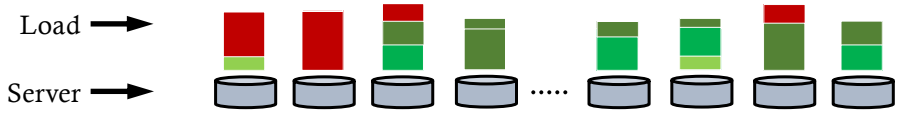
Challenge: Dynamic Workloads

low throughput & high tail latency



- The popular items change rapidly
- E.g., popular posts, limited-time offers, trending events

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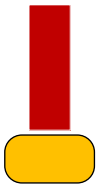
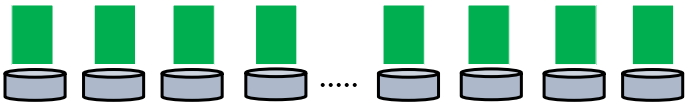
Fast, Small Cache for Better Load Balancing

Cache absorbs hottest queries



Balanced load



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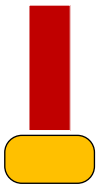
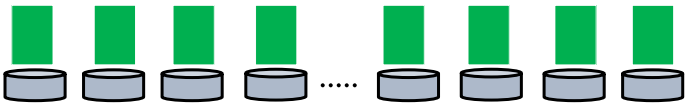
Fast, Small Cache for Better Load Balancing

Cache **O(*N* log *N*)** hottest items

E.g., 10,000 hot objects



N: # of servers



E.g., 100 backends with 100 billions items

[B. Fan et al. **SoCC'11**, X. Li et al. **NSDI'16**]

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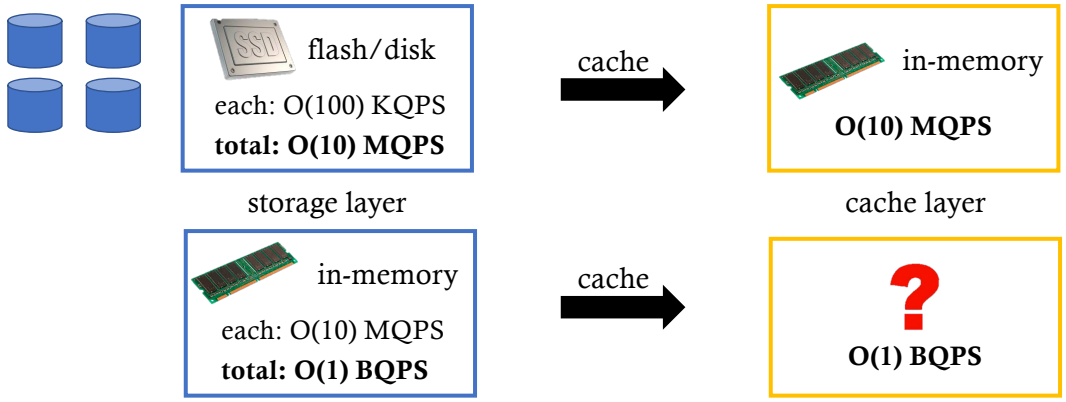
Key-Value Store is *Already* in High-Speed Memory



Cache needs to achieve the aggregate throughput of the storage layer.

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NetCache

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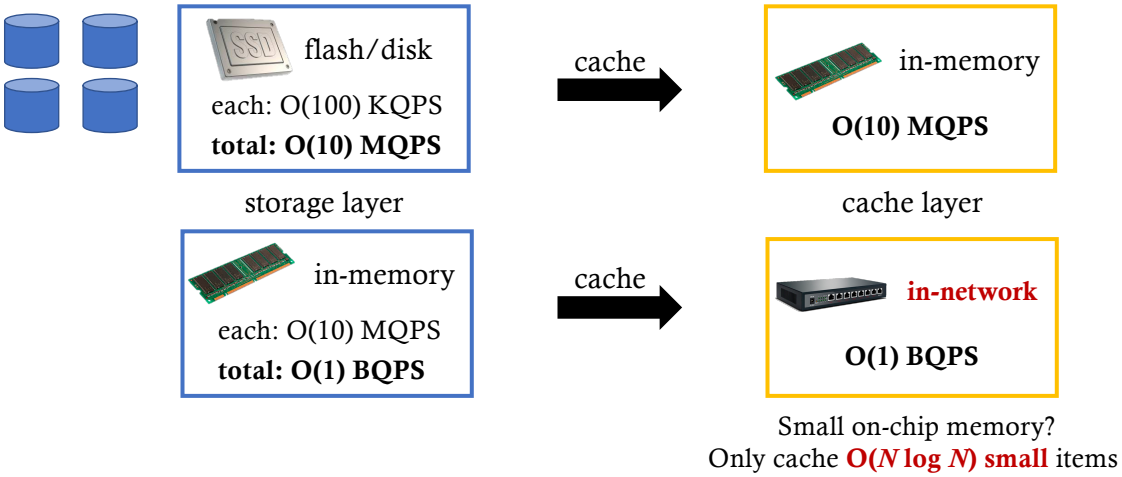
Cache Directly in the Network Devices!



Small on-chip memory?
Only cache **O(N log N) small** items

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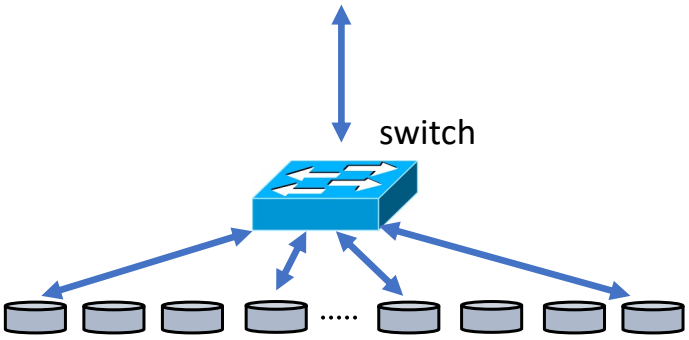
In-Network Load-Balancing Cache



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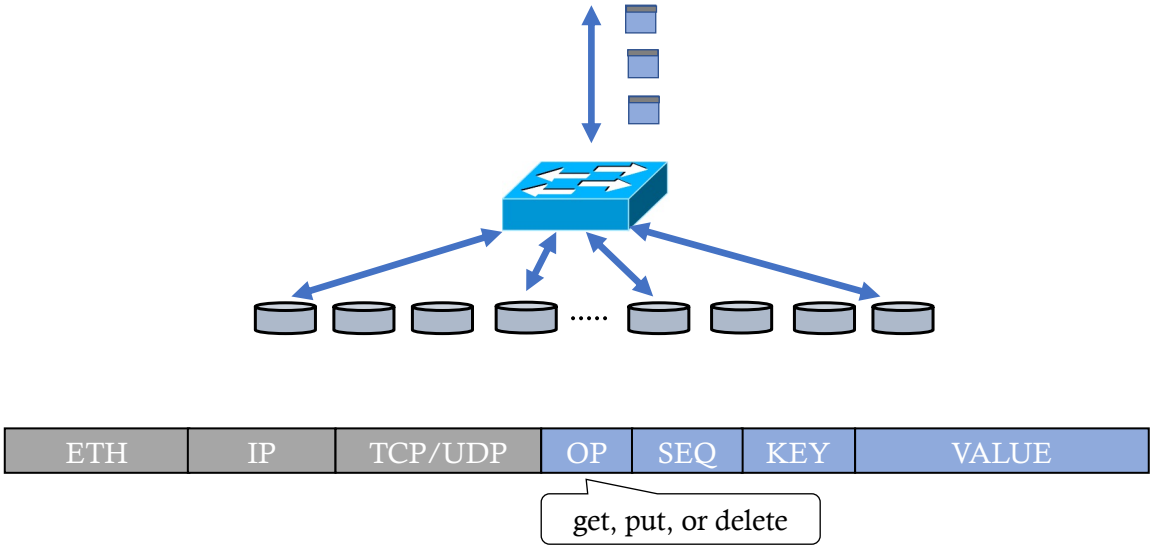
Individual Packets Carry the Messages



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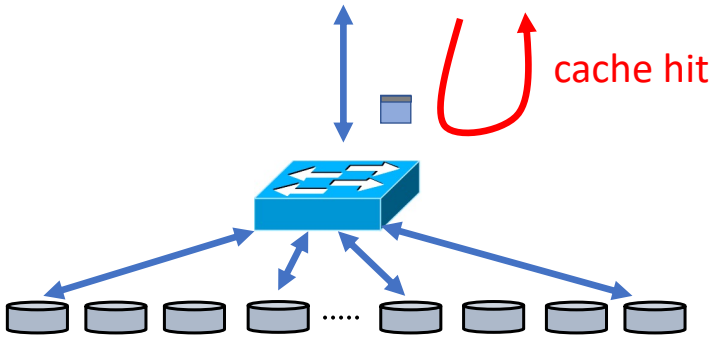
Cache Hit and Cache Miss



- Cache hit: switch responds directly
- Cache miss: switch forwards to the server, which responds

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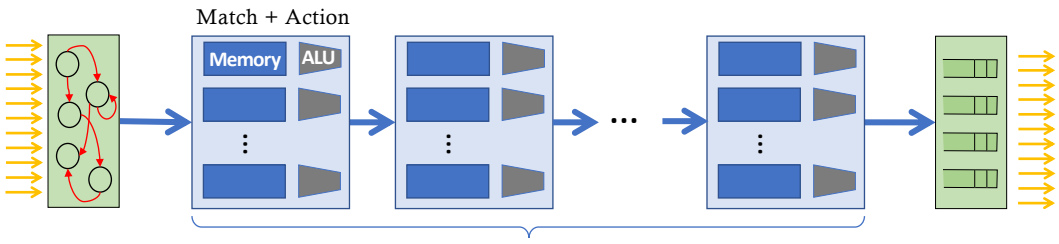
Programmable Packet-Processing Hardware

- Modern switches are increasingly programmable
  - Parsing of packet headers
  - Match-action processing on header fields
  - Registers for storing data across packets

www.p4.org



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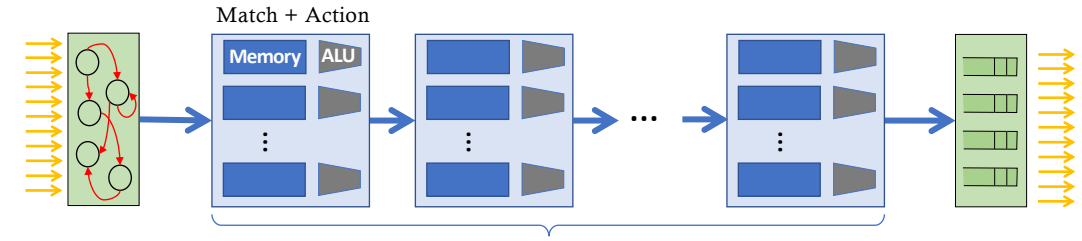
NetCache Uses Programmable Switches

- Programmable parser
  - Parse custom key-value fields in the packet
- Programmable match-action pipeline
  - Read and update key-value data
  - Maintain query statistics for cache updates



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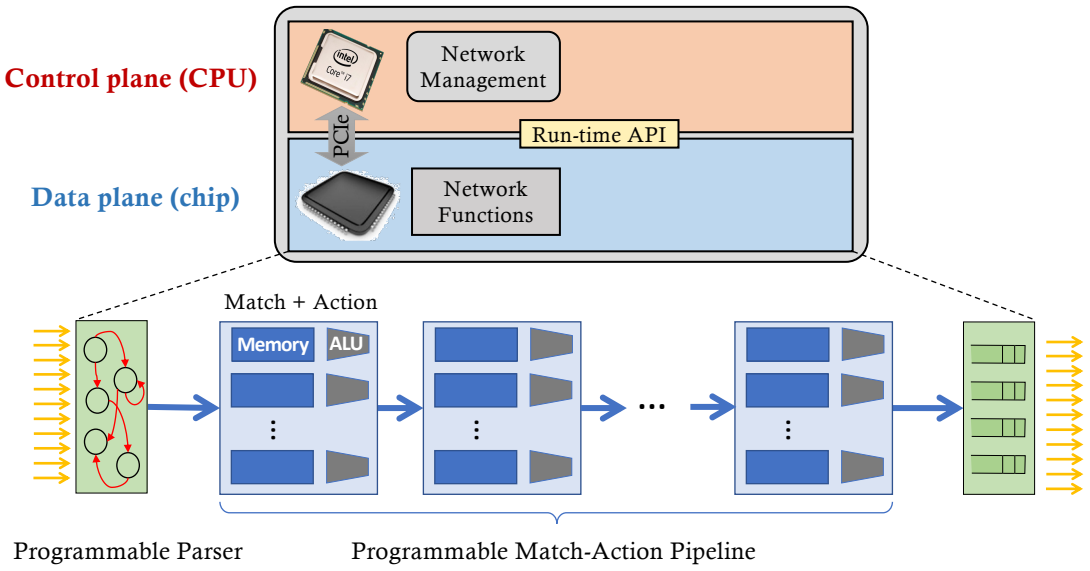
Control Plane of Programmable Switches



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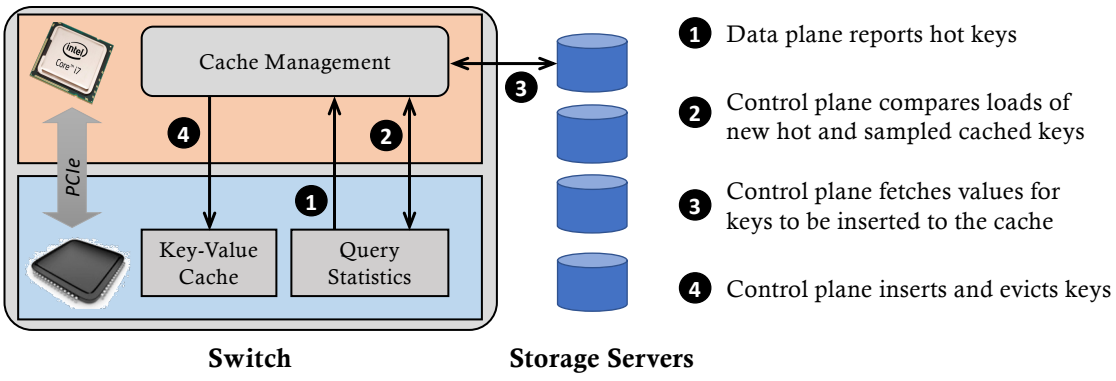
Cache Eviction and Replacement



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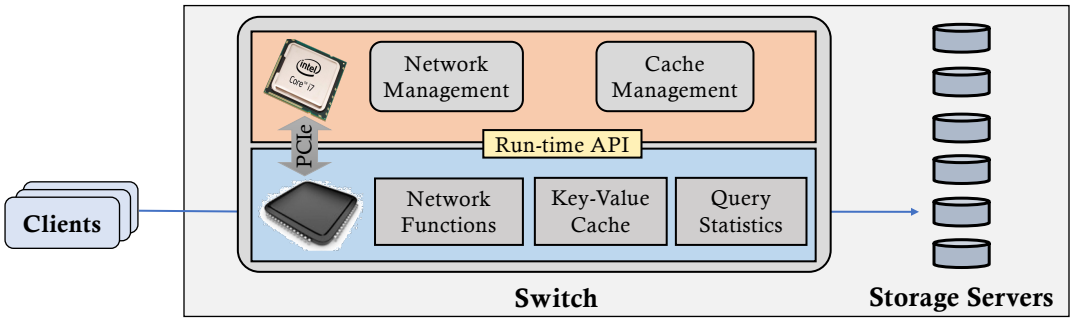
NetCache Architecture



- Switch data plane
  - Key-value store to serve queries for cached keys
  - Query statistics to enable efficient cache updates
- Switch control plane
  - Insert hot items into the cache and evict less popular items
  - Manage memory allocation for on-chip key-value store

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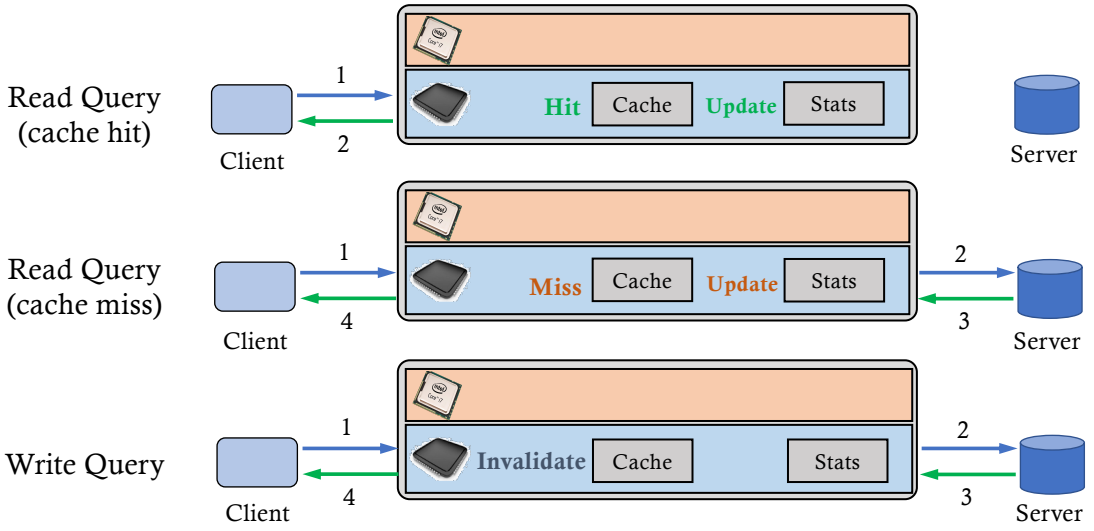
Data-Plane Query Handling



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NetCache Query Statistics

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Query Statistics in the Data Plane

- For cached keys
  - Count the number of accesses to each cached key
  - … so the control plane can decide which key-value pairs to *evict*
- For uncached keys
  - Identify new “hot” keys
  - … so the control plane can decide which key-value pairs to *insert*
  - Problem: huge number of uncached keys

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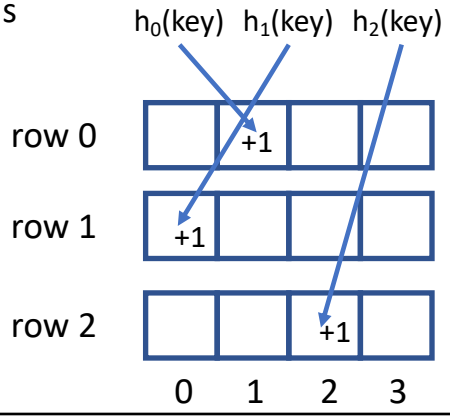
Compact Data Structure: CountMin Sketch

- Approximate answers with less memory
  - *r* rows, each with a hash function $h_i()$
  - *c* columns



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Compact Data Structure: CountMin Sketch

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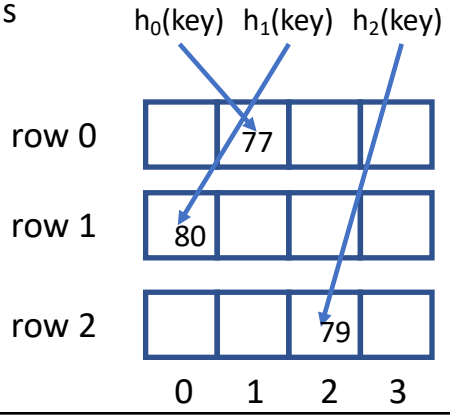

Estimated count:
min{77, 80, 79} = 77

Report this hot key!

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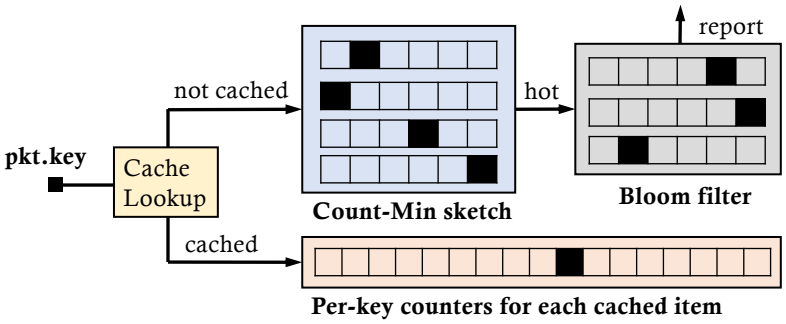
NetCache Query Statistics



- Cached key: per-key counter array
- Uncached key
  - Count-Min sketch: identify new hot keys
  - Bloom filter: remove duplicated hot key reports

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NetCache is a **key-value store** that leverages **in-network data-plane caching** to achieve

billions QPS throughput & **~10 µs latency**

even under

highly-skewed & **rapidly-changing**

workloads.

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Conclusion

- Key-value stores
  - Distributed, in-memory hash table
  - Simple, fast, and scalable
- Skewed workloads are still a challenge
  - Some keys are much more popular than others
  - So, some storage servers are overloaded
- In-network caching to the rescue
  - Store a small number of popular keys
  - … to reduce load (and variability of load) on the servers
- Programmable data-plane hardware
  - An opportunity for new network functionality
  - … but a challenge due to limited memory and processing

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