Splay tree

Property: recently accessed elements are quick to access again

- is a binary search tree
- and all operations are combined with

splaying operation

- rearranges the tree so that the node of the (looked up) element is placed at the root of the tree
- use ~ rotations

Whenever an element is looked up in the tree,
move that element to the root of the tree

insert x: as with a normal binary search tree.
Splay the newly inserted node x to the top of the tree

delete a node x: as with a binary search tree
splay the parent of the removed node to the top of the tree

Splay operations

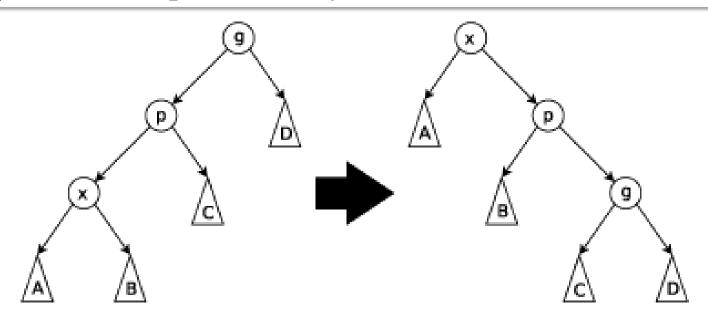
Assume p is the parent of x

Zig-zig step: when *x* and *p* are either both right children or they are both left children.

• The tree is rotated around g and then around p.

 $(\Rightarrow p \text{ is the new root of the subtree})$

e.g.: when **x** and **p** are both left children



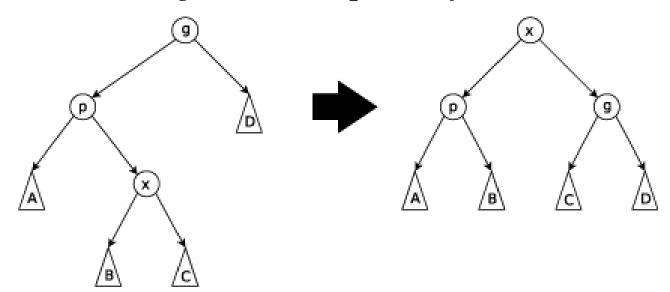
Splay operations

Assume p is the parent of x

Zig-zag step: when x is a right child and p is a left child or x is left and p is right child

The tree is rotated around p, and then rotated around q.

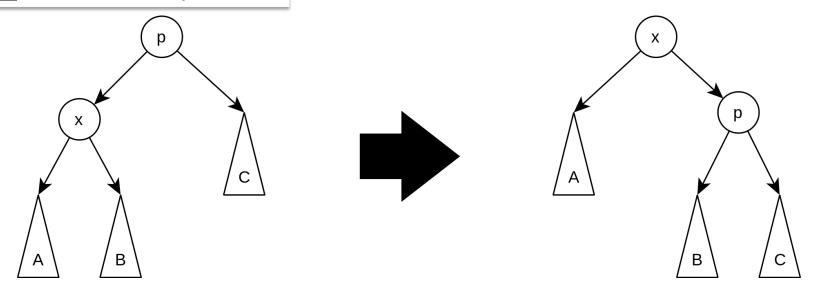
e.g.: when **x** is a right child and **p** is a left child



Splay operations

Zig step: this step is done when *p* is the root.

e.g.: when **x** is left child



This operation is performed:

- only as the last step in a splay operation.
- only if when x has odd depth at the beginning of the operation.

Splay tree. Insert / delete

To **insert** a value *x* into a splay tree:

- insert x as with a normal binary search tree;
- then a splay operation of the inserted node is performed.

 As a result, the newly inserted node x becomes the root of the tree.

To **delete** a node *x* from a splay tree:

- use the same method as with a binary search tree;
- then splay the parent of the removed node to the top of the tree.

Remark: we can find other "close" variants in the literature

Operation splay

End subalq.

```
Subalq. splay(x) {
                                                                  BinarySearchTree:
                                                                       root: ↑ BSTNode
0.
     if x=NIL then return:
    while ([x].parent <> NIL) execute
1.
        if ([[x].parent].parent = NIL ) then
2.
3.
               if ([[x].parent].left = x) then RightRotate([x].parent)
               else
                                               LeftRotate([x].parent)
4.
5.
               end if
6.
        else if ([[x].parent].left = x AND [[[x].parent].parent].left = [x].parent) then
                   RightRotate([[x].parent].parent);
                                                        RightRotate([x].parent)
7.
        else if ([[x].parent].right = x AND [[[x].parent].right = [x].parent) then
8.
9.
                   LeftRotate([[x].parent].parent);
                                                        LeftRotate([x].parent)
10.
       else if ([[x].parent].left = x AND [[[x].parent].right = [x].parent) then
11.
                   RightRotate([x].parent);
                                                        LeftRotate([x].parent)
12.
        else
13.
                   LeftRotate([x].parent);
                                                        RightRotate([x].parent)
14.
       end if
15.
16.
       end if
17.
     end while
```

BSTNode:

info: TComp

left: ↑ BSTNode

right: ↑ BSTNode

parent: ↑ BSTNode