# **CL Tutorial 5**

### Exercise 1

• Reduce the sequent  $(a \land \neg b) \lor (\neg a \land b) \models$  and use the result to derive the rule  $(\bigoplus L)$ .

$$\frac{a, \neg b \vDash}{a \land \neg b \vDash} \land L \qquad \frac{\neg a, b \vDash}{\neg a \land b \vDash} \land L (a \land \neg b) \lor (\neg a \land b) \vDash$$

Right-hand rule:

$$\frac{\Gamma, a, \neg b \vDash \Delta \qquad \Gamma, \neg a, b \vDash \Delta}{\Gamma, a \oplus b \vDash \Delta} \oplus L$$

### Exercise 2

• Reduce the sequents  $\vDash \neg a \lor b$  and  $\neg a \lor b \vDash$  and use the results to write rules  $(\rightarrow R)$  and  $(\rightarrow L)$ , respectively.

1.

$$\frac{a \vDash b}{\vDash \neg a, b} \neg R$$

$$\vDash \neg a \lor b$$

Right-hand rule:

$$\frac{\Gamma, a \vDash b, \Delta}{\Gamma \vDash a \to b, \Delta} \to R$$

2.

$$\frac{b \vDash a}{\neg a, b \vDash} \neg L$$

$$\neg a \lor b \vDash} \lor L$$

Left-hand rule:

$$\frac{\Gamma, b \vDash a, \Delta}{\Gamma, a \to b \vDash \Delta} \to L$$

## Exercise 3

• Reduce the sequents  $\vDash (a \land b) \lor (\neg a \land \neg b)$  and  $(a \land b) \lor (\neg a \land \neg b) \vDash$  and use the results to write rules  $(\leftrightarrow R)$  and  $(\leftrightarrow L)$ , respectively.

1.

$$\frac{\overline{a \vDash a}^{I}}{\vDash a, \neg a} \neg R \qquad \frac{b \vDash a}{\vDash a, \neg b} \neg R}{\vDash a, \neg b} \land R \qquad \frac{a \vDash b}{\vDash b, \neg a} \neg R \qquad \frac{\overline{b} \vDash \overline{b}^{I}}{\vDash b, \neg b} \neg R}{\vDash b, (\neg a \land \neg b)} \land R}$$

$$\frac{\vDash (a \land b), (\neg a \land \neg b)}{\vDash (a \land b) \lor (\neg a \land \neg b)} \lor R$$

Right-hand rule:

$$\frac{\Gamma, a \vDash b, \Delta \qquad \Gamma, b \vDash a, \Delta}{\Gamma \vDash a \leftrightarrow b, \Delta} \leftrightarrow R$$

2.

$$\frac{\overline{a \vDash a}^{I}}{a, \neg a \vDash} \neg L \qquad \frac{a \vDash b}{a, \neg b \vDash} \neg L \qquad \frac{b \vDash a}{b, \neg a \vDash} \neg L \qquad \frac{\overline{b \vDash b}^{I}}{b, \neg b \vDash} \neg L \\
\underline{a, (\neg a \land \neg b) \vDash} \qquad \land L \qquad \frac{b, (\neg a \land \neg b) \vDash}{b, (\neg a \land \neg b) \vDash} \land L \\
\underline{(a \land b), (\neg a \land \neg b) \vDash} \qquad \lor L$$

Left-hand rule:

$$\frac{\Gamma, a \vDash b, \Delta \qquad \Gamma, b \vDash a, \Delta}{\Gamma, a \leftrightarrow b \vDash \Delta} \leftrightarrow L$$

# Exercise 4

1.

$$\frac{a \vDash b \quad b \vDash a}{\vDash a \leftrightarrow b} \leftrightarrow R$$

$$\frac{a \vDash b}{\vDash (a \to b)} \to R \qquad \frac{b \vDash a}{\vDash (b \to a)} \to R$$

$$\vDash (a \to b) \land (b \to a) \land R$$

 $a \leftrightarrow b = (a \rightarrow b) \land (b \rightarrow a)$  is universally valid as both sides reduce to  $a \models b, b \models a$ .

2.

$$\frac{b \vDash a, c}{a \to b \vDash c} \to L$$

$$\vDash (a \to b) \to c$$

$$\frac{a, b \vDash c}{a \vDash b \to c} \to R$$

$$\vDash a \to (b \to c)$$

 $(a \rightarrow b) \rightarrow c = a \rightarrow (b \rightarrow c)$  is not universally valid.

3.

$$\frac{a \vDash b, c \quad b \vDash a, c}{a \leftrightarrow b \vDash c} \leftrightarrow L \quad \frac{b, c \vDash a \quad a, c \vDash b}{c \vDash a \leftrightarrow b} \leftrightarrow R$$

$$\vDash (a \leftrightarrow b) \leftrightarrow c$$

$$\frac{a.b \vDash c \quad b, c \vDash a}{a \vDash b \leftrightarrow c} \leftrightarrow R \quad \frac{b \vDash a, c \quad c \vDash a, b}{b \leftrightarrow c \vDash a} \leftrightarrow L$$

$$\vDash a \leftrightarrow (b \leftrightarrow c)$$

 $(a \leftrightarrow b) \leftrightarrow c = a \leftrightarrow (b \leftrightarrow c)$  is not universally valid.

$$\frac{a, \neg b \vDash c \quad \neg a, b \vDash c}{(a \oplus b), c \vDash} \oplus L \quad \frac{a, b \vDash c \quad \vDash a, b, c}{\vDash (a \oplus b), c} \oplus R$$
$$\vDash (a \oplus b) \oplus c$$

$$\frac{c, \neg b \vDash a \quad \neg c, b \vDash a}{a, (b \oplus c) \vDash} \oplus L \quad \frac{b, c \vDash a \quad \vDash a, b, c}{\vDash a, (b \oplus c)} \oplus R$$
$$\vDash a \oplus (b \oplus c)$$

 $(a \oplus b) \oplus c = a \oplus (b \oplus c)$  is not universally valid.

#### Exercise 5

 $1.r \leftrightarrow (a \land b)$ 

$$\frac{r \vDash a \quad r \vDash b}{r \vDash (a \land b)} \land R \qquad \frac{a, b \vDash r}{(a \land b) \vDash r} \land L$$
$$\vDash r \leftrightarrow (a \land b)$$

Shifting everything to the right:  $\models a, \neg r \models b, \neg r \models r, \neg a, \neg b$ 

CNF:  $(a \lor \neg r) \land (b \lor \neg r) \land (r \lor \neg a \lor \neg b)$ 

 $2.r \leftrightarrow (a \lor b)$ 

$$\frac{r \vDash a, b}{r \vDash (a \lor b)} \lor R \qquad \frac{a \vDash r \quad b \vDash r}{(a \lor b) \vDash r} \lor L$$
$$\vDash r \leftrightarrow (a \lor b)$$

Shifting everything to the right:  $\models a, b, \neg r \models r, \neg a \models r, \neg b$ 

CNF:  $(a \lor b \lor \neg r) \land (r \lor \neg a) \land (r \lor \neg b)$ 

 $3.r \leftrightarrow (a \rightarrow b)$ 

$$\frac{r, a \vDash b}{r \vDash (a \to b)} \to R \qquad \frac{b \vDash r, a}{(a \to b) \vDash r} \to L$$
$$\vDash r \leftrightarrow (a \to b)$$

Shifting everything to the right:  $\models b, \neg a, \neg r \qquad \models r, a, \neg b$ 

CNF:  $(b \lor \neg a \lor \neg r) \land (r \lor a \lor \neg b)$ 

 $4.r \leftrightarrow (\neg a)$ 

$$\frac{r \vDash \neg a \qquad \neg a \vDash r}{\vDash r \leftrightarrow (\neg a)} \leftrightarrow R$$

Shifting everything to the right:  $\vdash \neg a, \neg r \vdash a, r$ 

CNF:  $(\neg a \lor \neg r) \land (a \lor r)$