

# Starvation

## Starvation problem - disadvantage

- ∴ process in the ready state waiting indefinitely to get to CPU.
  - ↳ event occurrence keeps being postponed
- ∴ any priority based scheduling algorithm has a chance of starvation (ex. SJF, SRTF).
  - = low priority process can be starving

# Convoy effect

## convoy effect - disadvantage

- ∴ a smaller process (process with very small execution time) waiting for one big process to get off (release) CPU: (ex. FCFS, SJF)
- ⇒ higher waiting time, higher TAT

# Practical Implementation

- SJF, SRTF가 FCFS에 비해 throughput이 높다. throughput ↑ faster, more efficient computer.
- 그러나 SJF, SRTF를 실현하려면 정확하기는 어렵다. 모든 process의 burst time을 알아야 하기 때문  
burst time (execution time)이 복잡하기! BD도 많기 힘들기 ex. system power 등

# Throughput

: number of processes executed per unit time.

How will you get the maximum throughput? burst time 이 작은 것을 먼저 실행시키는 것이 좋다.

(조건: processes are infinite)

① SRTF ② SJF ③ FCFS

# LJF Longest Job First scheduling algorithm

process with the longest burst time will be scheduled first # non-preemptive # priority based

\* disadvantages

- 1) starvation
- 2) convey effect
- 3) low throughput
- 4) practically very difficult to implement

(ex)

process Id	1	2	3	4	5
arrival time	0	1	2	3	4
burst time	12	15	18	2	16

CT	12	61	30	63	46
TAT	12	60	28	60	42
WT	0	45	10	58	26
RT	0	45	10	58	26

P1	P3	P5	P2	P4
0	12	30	44	61
				63

Schedule length = 63

throughput =  $5/63$

# LRTF Longest Remaining Time First Scheduling algorithm

process with the longest burst time will be scheduled first # preemptive # priority-based

\* disadvantages

- 1) starvation
- 2) convey effect
- 3) low throughput
- 4) practically very difficult to implement

(ex)

process Id	1	2	3
arrival time	0	0	0
burst time	2	8	8

$\begin{matrix} 8 \\ 2 \end{matrix}$ 
 $\begin{matrix} 8 \\ 2 \end{matrix}$

(2T) burst time 0' 같다면 arrival time 4인, arrival time이 같다면 process Id 4인

P3	P2	P3	P2	P3	P1	P3	P1	P2	P3
0	4	5	6	7	8	9	10	11	12
									13
									14

(ex)

process Id	1	2	3
arrival time	1	2	3
burst time	2	4	4

(2T) burst time 0' 같다면 arrival time 4인, arrival time이 같다면 process Id 4인

///	P1	P2	P3	P2	P3	P2	P3	P1	P2	P3
0	1	2	3	4	5	6	7	8	9	10
										11

CT	9	10	11
TAT	8	8	8
WT	6	4	4
RT	0	0	0

Schedule length =  $11 - 1 = 10$

throughput =  $3/10$

# RR

## Round Robin Scheduling algorithm

windows, mac 등 여러 운영체제에 사용되는 algorithm

**Time Quantum** : Maximum allowable time a process can run without getting preempted.

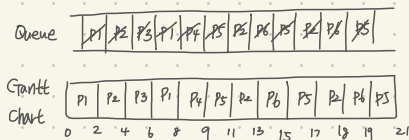
**RR** : TQ + FCFS - FCFS이지만 일정 time unit 동안 실행

- Works on basis of a particular time quantum
- uses queue data structure
- very popular & used in most of the OS today

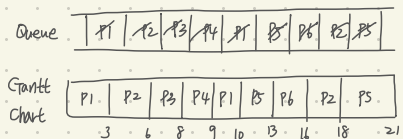
(EX)

process Id	1	2	3	4	5	6
arrival time	0	1	2	3	4	5
burst time	4	8	2	1	6	3
	✓	8	0	✓	✓	✓
		✓		2		

\* time quantum : 2



\* time quantum : 3



→ as time quantum increases, the number of context switching may decrease  
 → , the response time increases

### Advantages

- No starvation : not a priority based
- No Convey-effect
- Practically implementable
- Response time ↓

### Limitation

- Throughput is good but not as good as SJF, SRTF