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& 1.8 Linear Time Selection

1. One armed QUICKSORT : QUICKSELECT

QuickSelect (A[1..n], k): // Select Kun elem in A

if n=1
return A[1]
else
choose pivot element A[p]

r + Partition (A[1..n], p)

if K<r
return Quickselect (A[1..r-1], k)

else if k>r

return Quickselect (A[r+1..n], k-r)

else return A[r].

> Worst runtime:

$$T(n) \leq \max_{1 \leq r \leq n} \left\{ \frac{T \cdot (r-1) + T \cdot (n-r)}{l := length \text{ of recursive}} \right.$$

$$Subprob$$

$$T(n) \leq \max_{0 \leq l \leq n-1} T(l) + O(n) \text{, if always smallest,}$$

$$T(n) = O(n^2).$$

2. Choosing good pivots: l=an (a<1).

Idea: Choose recursively computing the median of a carefully-chosen subset of an array.

- How. Divide input array to [n/5] blocks, each containing exactly 5 (possible expept the last).
 - · Compute these blocks' median by brute force (possibly by a few ifs).
 - Collect those medians into a new Wst, $M[1... \lceil \frac{n}{5} \rceil]$,

· Use median of block medians as quick select pivot.

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MoMSELECT (A[1..n etc], k):

if n \le 25

use bruteforce

else

m \leftarrow \lceil n/5 \rceil

for i=1 to m

M[i] \leftarrow MedianOeFive(A[5i-4..5i])

y mom \leftarrow MomSelect(M[i..m], [m/2])

r \leftarrow Partition(A[i..n], mom)

if k < r

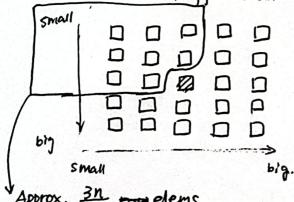
return MomSelect(A[i..r-i], k)

else if k > r

return MomSelect(A[r+i...n], k-r)

else return mom.
```

- · Analysis: MOM = Median Of Median
 - 1° Mom is a good pivot.
 - MOM is larger than $\lfloor \lceil h/5 \rceil / 2 \rfloor 1 \approx n/10$ block median,
 - · each block median is larger 2 other elems in its block
 - · Mom is bigger than at least 3n elems in input arr.



Assume sorted in both directions. (for demo)
THE ALGO PONT DO THIS!

2° The calls. $T(n) \leq T(\frac{n}{5}) + T(\frac{7n}{10}) = +0(n)$ $\sim \frac{9}{10}n$ of the prev layer, T(n) = 0(n).

About constant 5: First number result in exp. decay.