

CS4341 Digital Logic & Computer Design

Lecture Notes 2

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Review: Number Representation

Remember the two examples:

$(\text{Hello})_{\text{English}} = (\text{Salut})_{\text{French}} = (\text{สวัสดี})_{\text{Thai}}$

They all mean the same thing but in different “language system”.

Different language systems may or may not share the same alphabet “symbols”

$(\text{Gift})_{\text{English}} =$  $(\text{Gift})_{\text{German}} =$  $(\text{Gift})_{\text{Norwegian}} =$ 

Sometimes same words have different meanings in the different “language systems”

Review: Positional Number System

Binary System:

$r = 2$ or base 2

Digit set: $\{0, 1\}$

Octal System:

$r = 8$ or base 8

Digit set: $\{0, 1, 2, 3, 4, 5, 6, 7\}$

Hexadecimal System:

$r = 16$ or base 16

Digit set: $\{0, 1, 2, 3, 4, 5, 6, 7, 8, 9, A, B, C, D, E, F\}$

Review: Numerical Values

An integer N of length n and base r is represented as:

$$(N)_r = (b_{n-1}, b_{n-2}, \dots, b_1, b_0)_r$$

And has a numerical value of:

$$b_{n-1}r^{n-1} + b_{n-2}r^{n-2} + \dots + b_1r^1 + b_0r^0$$

Can be representation using the summation formula:

$$\sum_{i=0}^{n-1} b_i r^i$$

Conversion Between Bases

- From any base r to Decimal: use the numerical value polynomial
- From Decimal to base t : use the repeated division method
- To convert $(N)_{10}$ to $(M)_t$, we repeatedly divide N by t . The remainder after each step is one digit of the required base t , starting from the right most digit.

Conversion Example

Convert $(37)_{10}$ to Binary

$$(37)_{10} = (100101)_2$$

Division	Quotient	Remainder
$37 / 2$	18	1
$18 / 2$	9	0
$9 / 2$	4	1
$4 / 2$	2	0
$2 / 2$	1	0
$1 / 2$	0	1

← Least significant bit

← Most significant bit

← Stop when 0

Conversion Example

Convert $(422)_{10}$ to Hexadecimal

$$(422)_{10} = (1A6)_{16}$$

Division	Quotient	Remainder
422 / 16	26	6
26 / 16	1	A
1 / 16	0	1

← Least significant bit

← Most significant bit

← Stop when 0

Binary, Octal & Hexadecimal Conversions

- Observation: Octal (8) is a power of 2 = 2^3 . Hence, each octal digit can be represented in exactly 3 binary bits.
 - To convert from binary to octal, divide the word into groups of 3 starting from the LSB and convert each into an octal digit.
 - To convert from octal to binary, simply reverse the process
- Observation: Hexadecimal (16) is a power of 2 = 2^4 . Hence, each hexadecimal digit can be represented in exactly 4 binary bits.
 - To convert from binary to hexadecimal, divide the word into groups of 4 starting from the LSB and convert each into a hexadecimal digit.
 - To convert from hexadecimal to binary, simply reverse the process

Conversion Example

Convert $(1101000101)_2$ to Octal

$$1 \mid 101 \mid 000 \mid 101 = 1 \ 5 \ 0 \ 5 = (1505)_8$$

Convert $(12A7F)_{16}$ to Binary

$$1 \ 2 \ A \ 7 \ F = 0001 \ 0010 \ 1010 \ 0111 \ 1111 = \\ (1001010100111111)_2$$

Famous Conversion Table

Decimal	Binary	Octal	Hex
0	0000	0	0
1	0001	1	1
2	0010	2	2
3	0011	3	3
4	0100	4	4
5	0101	5	5
6	0110	6	6
7	0111	7	7
8	1000	10	8
9	1001	11	9
10	1010	12	A
11	1011	13	B
12	1100	14	C
13	1101	15	D
14	1110	16	E
15	1111	17	F

Important Property

The largest possible value of a number of length n -digits in base system r is:

$$\mathbf{r^n - 1}$$

Representing Fractions

- A Number N_r can have a fraction part:

$$N_r = \underbrace{b_{n-1}b_{n-2} \dots b_1b_0}_{\text{Integer Part}} \cdot \underbrace{b_{-1}b_{-2} \dots b_{-m+1}b_{-m}}_{\text{Fraction Part}}$$

↑
Radix Point

- The numerical value is calculated as:

- $b_{n-1} \times r^{n-1} + \dots + b_1 \times r + b_0 +$

(Integer part)

- $b_{-1} \times r^{-1} + b_{-2} \times r^{-2} \dots + b_{-m} \times r^{-m}$

(Fraction part)

Representing Fractions

$$N_r = \sum_{i=0}^{i=n-1} b_i \times r^i + \sum_{j=-1}^{j=-m} b_j \times r^j$$

Examples

$$(2409.87)_{10} = 2 \times 10^3 + 4 \times 10^2 + 9 + 8 \times 10^{-1} + 7 \times 10^{-2}$$

$$(1101.1001)_2 =$$

$$2^3 + 2^2 + 2^0 + 2^{-1} + 2^{-4} = 13.5625$$

$$(703.64)_8 =$$

$$7 \times 8^2 + 3 + 6 \times 8^{-1} + 4 \times 8^{-2} = 451.8125$$

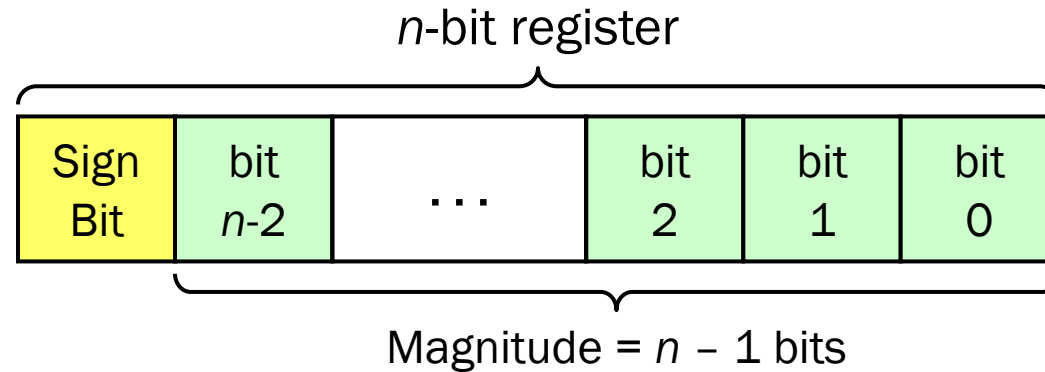
$$(A1F.8)_{16} =$$

$$10 \times 16^2 + 16 + 15 + 8 \times 16^{-1} = 2591.5$$

Signed Numbers

- *To carry out the various arithmetic operations, we need numbers as positive or negative.*
- *Three notations are commonly used:*
 - Sign-magnitude
 - $(r-1)$'s Complement
 - r 's complement
- *In Binary, 2's complement is widely used.*

Sign-Magnitude Representation



- *Independent representation of sign and magnitude.*
- *MSB is the sign bit: 0 means positive and 1 means negative*
- *Using n -bits, then largest represented magnitude is:*

$$2^{n-1} - 1$$

Sign-Magnitude Examples

Sign-magnitude
representation of +27
using 8-bit register

0	0	0	1	1	0	1	1
---	---	---	---	---	---	---	---

Sign-magnitude
representation of -27
using 8-bit register

1	0	0	1	1	0	1	1
---	---	---	---	---	---	---	---

Properties of Sign-Magnitude

- *Two representations of the zero: +0 and -0*
- *Symmetric range of represented positive and negative values*
 - For n-bit register, range is from $-(2^{n-1} - 1)$ to $+(2^{n-1} - 1)$
 - For example, using 8-bit register, range is -127 to +127
- *Hard to implement addition and subtraction*
 - Sign and magnitude parts must be processed separately
 - Sign bit should be examined to determine addition or subtraction
 - Increases the cost of the logic circuit

1's Complement Representation

- *Given a binary number N , the 1's complement of N is obtained by flipping each bit in N (0 becomes 1, and 1 becomes 0).*
- *Example: 1's complement of $(01101001)_2 = (10010110)_2$*
- *If N consists of n -bits then, 1's complement of $N = (2^n - 1) - N$*

2's Complement Representation

- *Used in most computers today*
- *Definition: given a binary number N , the 2's complement of N = 1's complement + 1*
- *Example: 1's complement of $(01101001)_2 = (10010110)_2$*
- *2's complement of $(01101001)_2 = (10010110 + 1)_2 = (10010111)_2$*
- *If N consists of n -bits then, 2's complement of $N = \mathbf{2^n - N}$*

2's Complement Example

starting value	$00100100_2 = +36$
step1: reverse the bits (1's complement)	11011011_2
step 2: add 1 to the value from step 1	$+ \quad \quad 1_2$
sum = 2's complement representation	$11011100_2 = -36$

2's Complement Representation

- *Another way to obtain the 2's complement: Starting at the least significant 1:*
- Leave all the 0s to its right unchanged
- Complement all the bits to its left

Binary Value

= 00100 **1** 00

2's Complement

= **11011** **1** 00

Ranges for Unsigned and Signed Numbers

- For n -bit unsigned integers: Range is 0 to $(2^n - 1)$
- For n -bit signed integers: Range is -2^{n-1} to $(2^{n-1} - 1)$:
 - Positive range: 0 to $(2^{n-1} - 1)$
 - Negative range: -2^{n-1} to -1

Note: in 2's Complement, there is only one zero: 2's complement of 0 = 0

Unsigned and Signed Values

➤ *Positive numbers:*

- Signed value = Unsigned value

➤ *Negative numbers:*

- To obtain the signed value, we assign a negative weight to MSB and add to the rest of the bit values.

$$N = b_{n-1} \times (-2^{n-1}) + b_{n-2} \times 2^{n-2} + \dots + b_0 \times 2^0$$

Example

Find the numerical value of signed number $(10110100)_2$

1	0	1	1	0	1	0	0
-128	64	32	16	8	4	2	1

- Value = $-128 + 32 + 16 + 4 = -76$
- If we store this number in 16-bit memory instead, would the numerical value change?

Example

Find the signed and unsigned ranges in 4-bit binary representation

- Unsigned: 0000_2 to $1111_2 = 0_{10} - 15_{10}$
- Signed (using 2's Complement):
 - Positive Range: $0111_2 (+7_{10})$ to $0000_2 (0_{10})$
 - Negative Range: $1000_2 (-8_{10})$ to $1111_2 (-1_{10})$

Ranges for Unsigned and Signed Numbers

8-bit Binary value	Unsigned value	Signed value
00000000	0	0
00000001	1	+1
00000010	2	+2
...
01111110	126	+126
01111111	127	+127
10000000	128	-128
10000001	129	-127
...
11111110	254	-2
11111111	255	-1

To Do List

- Review lecture notes
- Read chapter 1 until 1.5