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Language theoretic comparison of LL and LR grammars

Asked 8 years, 10 months ago Active 3 years, 8 months ago Viewed 12k times



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People often say that $LR(k)$ parsers are more powerful than $LL(k)$ parsers. These statements are vague most of the time; in particular, should we compare the classes for a fixed k or the union over all k ? So how is the situation really? In particular, I am interested in how $LL(*)$ fits in.

As far as I know, the respective sets of grammars LL and LR parsers accept are orthogonal, so let us talk about the languages generated by the respective sets of grammars. Let $LR(k)$ denote the class of languages generated by grammars that can be parsed by an $LR(k)$ parser, and similar for other classes.

I am interested in the following relations:

- $LL(k) \stackrel{?}{\subseteq} LR(k)$
- $\bigcup_{i=1}^{\infty} LL(k) \stackrel{?}{\subseteq} \bigcup_{i=1}^{\infty} LR(k)$
- $\bigcup_{i=1}^{\infty} LL(k) \stackrel{?}{=} LL(*)$
- $LL(*) \stackrel{?}{\circ} \bigcup_{i=1}^{\infty} LR(k)$

Some of these are probably easy; my goal is to collect a "complete" comparison. References are appreciated.



edited Jun 6 '13 at 14:05



frafl

2,240

1

12

32

asked Mar 7 '12 at 0:32



Raphael ♦

68.9k

27

156

350

2 Maybe this can help you! [Grammar Hierarchy image](#) – Andrea Tucci Jul 11 '12 at 15:32

1 @AndreaTucci: Yes, but that only covers the grammars, not the generated languages. – Raphael ♦ Jul 11 '12 at 18:06

1 Answer

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There are numerous containments known. Let \subseteq denote containment and \subset proper containment. Let \times denote incomparability.

Let $LL = \bigcup_k LL(k)$, $LR = \bigcup_k LR(k)$.

Grammar level



For LL

- $LL(0) \subset LL(1) \subset LL(2) \subset LL(2) \subset \dots \subset LL(k) \subset \dots \subset LL \subset LL(*)$
- $SLL(1) = LL(1)$, $SLL(k) \subset LL(k)$, $SLL(k+1) \times LL(k)$



Most of these are proven in [Properties of deterministic top down grammars](#) by Rosenkrantz and Stearns. $SLL(k+1) \times LL(k)$ is a rather trivial exercise. [This presentation](#) by Terence Parr places $LL(*)$ on slide 13. The paper LL-regular grammars by Jarzabek and Krawczyk show $LL \subset LLR$, and their proof trivially extends to $LL \subset LL(*)$

For LR

- $LR(0) \subset SLR(1) \subset LALR(1) \subset LR(1)$
- $SLR(k) \subset LALR(k) \subset LR(k)$
- $SLR(1) \subset SLR(2) \subset \dots \subset SLR(k)$
- $LALR(1) \subset LALR(2) \subset \dots \subset LALR(k)$
- $LR(0) \subset LR(1) \subset LR(2) \subset \dots \subset LR(k) \subset \dots \subset LR$



LL versus LR

- $LL(k) \subset LR(k)$ ([Properties of deterministic top down grammars](#) plus any left recursive grammar)
- $LL(k) \times SLR(k), LALR(k), LR(k-1)$ (simple exercise)
- $LL \subset LR$ (any left recursive grammar)
- $LL(*) \times LR$ (left recursion versus arbitrary lookahead)

Language level

For LL

- $LL(0) \subset LL(1) \subset LL(2) \subset \dots \subset LL(k) \subset \dots \subset LL \subset LL(*)$
- $SLL(k) = LL(k)$

Most of these are proven in [Properties of deterministic top down grammars](#). [The equivalence problem for LL- and LR-regular grammars](#) by Nijholt makes references to papers showing $LL(k) \subset LL(*)$. The paper LL-regular grammars by Jarzabek and Krawczyk show $LL \subset LLR$, and their proof trivially extends to $LL \subset LL(*)$

For LR

- $LR(0) \subset SLR(1) = LALR(1) = LR(1) = SLR(k) = LALR(k) = LR(k) = LR$

Some of these were proven by Knuth in his paper [On the Translation of Languages from Left to Right](#) in which he introduced LR(k), the rest are proven in [Transforming LR\(k\) Grammars to LR\(1\), SLR\(1\), and \(1,1\) Bounded Right-Context Grammars](#) by Mickunas et al.

LL versus LR

- $LL \subset LR(1)$ (containment follows from the above, $\{a^i b^j \mid i \geq j\}$ is the canonical example for strict containment)
- $LL(*) \times LR$ (the language $\{a^i b^j \mid i \geq j\}$ shows half the claim, and the introduction of [The equivalence problem for LL- and LR-regular grammars](#) by Nijholt makes references to papers showing the other half)
- $LR(1) = DCFL$ (see e.g. reference [here](#)).

Excellent answer though, I had already upvoted. I would have thought Frank deRemer proved $LALR \leq LR$ in his original LALR paper? (1969?) – [user207421](#) Apr 11 '12 at 18:43

@EJP: $LALR(k) \subseteq LR(k)$ is immediate if you define $LALR$ by collapsing LR states. deRemer only proved that his construction created the same parser. The grammar classes are not equal, which is a nice exercise (or even a nice question for this site). The language class equality is a lot more tricky: read the paper if you really want the details ;) – [Alex ten Brink](#) Apr 11 '12 at 22:58

1 @AlextenBrink I did read the paper, and was taught by Frank de Remer, but it's 30+ years ago ;-) Thanks for all the detail. – [user207421](#) Apr 12 '12 at 0:13

It might be nice to collect example grammars for each inequality. – [o11c](#) May 31 '15 at 5:51

1 @o11c I think that would overburden a single answer. My impression is that Alex gave good references where necessary; he states "easy exercise" for some. I guess if a reader can not come up with a grammar, they can post a new question asking for that specific case. – [Raphael](#) ♦ Jun 18 '15 at 10:33