(十五) 离散数学: 复习 (Review)

魏恒峰

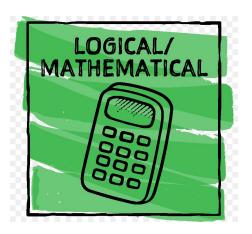
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Theorem

$$\Sigma \vdash \alpha \Longleftrightarrow \Sigma \models \alpha$$



" \rightarrow " and " \leftrightarrow " are used in a single formula.

" \Longrightarrow " and " \Longleftrightarrow " are used to connect two formulas.

$$x \in A \setminus B$$

$$\iff x \in A \land x \notin B$$

$$\iff x \in A \land (x \in U \land x \notin B)$$

$$\iff x \in A \land x \in \overline{B}$$

$$\iff x \in A \cap \overline{B}$$

$$p \oplus q \triangleq (p \lor q) \land \neg (p \land q)$$
$$= (p \land \neg q) \lor (\neg q \land q)$$

p	q	$p\oplus q$
0	0	0
0	1	1
1	0	1
1	1	0

$$p \oplus q = q \oplus r$$
$$(p \oplus q) \oplus r = p \oplus (q \oplus r)$$

Theorem

$\sqrt{2}$ is irrational.



The First Crisis in Mathematics



Theorem (Pigeonhole Principle)

If n objects are placed in r boxes, where r < n, then at least one of the boxes contains ≥ 2 ($\geq \lceil \frac{n}{r} \rceil$) object.

Numbers

Consider the numbers 1, 2, ..., 2n, and take any n + 1 of them. There are two among these n + 1 numbers which are relatively prime.

There must be two numbers which are only 1 apart.

Numbers

Consider the numbers $1, 2, \ldots, 2n$, and take any n+1 of them. There are two among these n+1 numbers such as one divides the other.

$$a = 2^k m$$
, $(1 \le m \le 2n - 1 \text{ is odd})$

There n+1 numbers have only n different odd parts.

There must be two numbers with the same odd part.

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Hand-shaking

If there are n > 1 people who can shake hands with one another, there are two people who shake hands with the same number of people.

$$0 \sim n-1$$

Either the '0' hole or the 'n -1' hole or both must be empty.

Sums

Suppose we are given n integers a_1, a_2, \ldots, a_n . Then there is a set of consecutive numbers $a_{k+1}, a_{k+2}, \ldots, a_l$ whose sum $\sum_{i=k+1}^{l} a_i$ is a multiple of n.

$$A_i = \sum_{k=1}^{k=i} a_i$$

$$A_0, A_1, A_2, \ldots, A_n$$

$$\exists 0 \le i < j \le n. \ A_i = A_j \mod n$$

$$A_j - A_i = a_{i+1} + \dots + a_j = 0 \mod n$$

Championship Match

"胡司令" (胡荣华) 要安排一次长达 77 天的象棋练习赛。

他想每天至少要有一场比赛,但是总共不超过132场比赛。

请证明, 无论如何安排, 他都要在连续的若干天内恰好完成 21 场比赛。

Let a_i denote the number of games he plays up through the *i*-th day.

$$a_1, a_2, \ldots, a_{76}, a_{77}, a_1 + 21, a_2 + 21, \ldots, a_{76} + 21, a_{77} + 21$$

There must be ≥ 2 elements having the same value.

It must be
$$a_i + 21 = a_i$$
.

Sequences

In any sequence $a_1, a_2, \ldots, a_{mn+1}$ of mn+1 distinct numbers, there exists an increasing subsequence

$$a_{i_1} < a_{i_2} < \dots < a_{i_{m+1}} \quad (i_1 < i_2 < \dots < i_{m+1})$$

of length m+1, or a decreasing subsequence

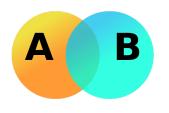
$$a_{j_1} > a_{j_2} > \dots > a_{j_{n+1}} \quad (j_1 > i_2 < \dots > j_{n+1})$$

of length n+1, or both.

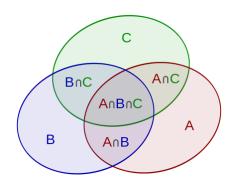


Paul Erdős (1913 \sim 1996)

Chapter 28 of "Proofs from THE Book"







$$|A \cup B \cup C| = |A| + |B| + |C|$$

$$-|A \cap B| - |A \cap C| - |B \cap C|$$

$$+|A \cap B \cap C|$$

Theorem (Inclusion-Exclusion Principle)

$$\left| \bigcup_{i=1}^{n} A_i \right| = \sum_{i=1}^{n} |A_i| - \sum_{1 \le i < j \le n} |A_i \cap A_j|$$

$$+ \sum_{1 \le i < j < k \le n} |A_i \cap A_j \cap A_k|$$

$$- \cdots$$

$$+ (-1)^{n-1} |A_1 \cap \cdots \cap A_n|.$$

$$\left| \bigcap_{i=1}^{n} \bar{A}_{i} \right| = \left| S - \bigcup_{i=1}^{n} A_{i} \right| = \left| S \right| - \sum_{i=1}^{n} \left| A_{i} \right| + \sum_{1 \leq i < j \leq n} \left| A_{i} \cap A_{j} \right| - \dots + (-1)^{n} \left| A_{1} \cap \dots \cap A_{n} \right|.$$

Counting Integers

How many integers in $1, \ldots, 100$ are not divisible by 2, 3 or 5?

$$100 - (50 + 33 + 20) + (16 + 10 + 6) - 3 = 26.$$

Counting Derangements (错排)

Suppose there is a deck of n cards numbered from 1 to n. Suppose a card numbered i is in the correct position if it is the i-th card in the deck. How many ways can the cards be shuffled without any cards being in the correct position?

 A_m : all of the orderings of cards with the m-th card correct

$$\left|\bigcap_{i=1}^{n} \overline{A_i}\right| = \left|S - \bigcup_{i=1}^{n} A_i\right| = n! - \sum_{i=1}^{n} |A_i| + \sum_{1 \le i < j \le n} |A_i \cap A_j|$$
$$- \dots + (-1)^n |A_1 \cap \dots \cap A_n|.$$

$$S_k \triangleq \sum_{1 \leq i_1 < i_2 < \dots < i_k \leq n} |A_{i_1} \cap A_{i_2} \cap \dots \cap A_{i_k}| = \binom{n}{k} (n-k)! = \frac{n!}{k!}$$

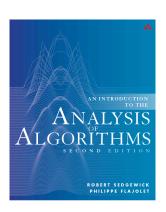
$$S_k = \frac{n!}{k!}$$

$$\left| \bigcap_{i=1}^{n} \overline{A_i} \right| = n! - \frac{n!}{1!} + \frac{n!}{2!} - \dots + (-1)^n \frac{n!}{n!}$$
$$= n! \sum_{k=0}^{n} \frac{(-1)^k}{k!}$$

$$n \to \infty \implies \sum_{k=0}^{n} \frac{(-1)^k}{k!} \to e^{-1} \approx 0.368$$



$$a_n = f(a_{n-1}, a_{n-2}, \dots, a_{n-t}) + g(n)$$



recurrence type	typical example	
first-order		
linear	$a_n = na_{n-1} - 1$	
nonlinear	$a_n = 1/(1 + a_{n-1})$	
second-order		
linear	$a_n = a_{n-1} + 2a_{n-2}$	
nonlinear	$a_n = a_{n-1}a_{n-2} + \sqrt{a_{n-2}}$	
variable coefficients	$a_n = na_{n-1} + (n-1)a_{n-2} + 1$	
tth order	$a_n = f(a_{n-1}, a_{n-2}, \dots, a_{n-t})$	
full-history	$a_n = n + a_{n-1} + a_{n-2} \dots + a_1$	
divide-and-conquer	$a_n = a_{\lfloor n/2 \rfloor} + a_{\lceil n/2 \rceil} + n$	

 Table 2.1
 Classification of recurrences

Homogeneous Linear Recurrence Relations with Constant Coefficients

$$a_n = c_1 a_{n-1} + c_2 a_{n-2} + \dots + c_t a_{n-t}$$



https://www.bilibili.com/video/BV1Cf4y187Cu?share_source=copy_web

$$R \subseteq A \times A$$

$$\begin{cases} R^0 = I_A \\ R^{n+1} = R \circ R^n \end{cases}$$

Representing Relations as Matrices/Digraphs

$$A = \{1, 2, 3, 4\}$$

$$R = \{(1, 1), (1, 2), (2, 1), (2, 2), (2, 3), (2, 4), (3, 4), (4, 1)\}$$

$$R^{2} \qquad R^{3}$$

$$R^{+} = \bigcup_{i=1}^{\infty} R \qquad R^{*} = \bigcup_{i=0}^{\infty} R$$

Definition (Reflexive Closure (自反闭包))

The reflexive closure $\operatorname{cl}_{\operatorname{ref}}(R)$ of a relation $R \subseteq X \times X$ is the smallest reflexive relation on X that contains R.

$$\mathsf{cl}_{\mathsf{ref}}(R) = R \cup I_X$$

Definition (Symmetric Closure (对称闭包))

The symmetric closure $\operatorname{cl}_{\operatorname{sym}}(R)$ of a relation $R \subseteq X \times X$ is the smallest symmetric relation on X that contains R.

$$\mathsf{cl}_{\mathrm{sym}}(R) = R \cup R^{-1}$$

Definition (Transitive Closure (传递闭包))

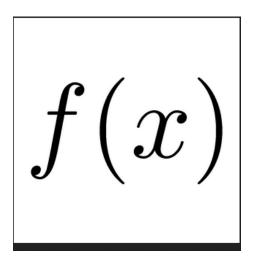
The transitive closure $\operatorname{cl}_{\operatorname{trn}}(R)$ of a relation $R \subseteq X \times X$ is the smallest transitive relation on X that contains R.

$$\mathsf{cl}_{\mathrm{trn}}(R) = R^+$$

- $ightharpoonup R^+$ contains R
- $ightharpoonup R^+$ is transitive
- $ightharpoonup R^+$ is minimal

If T is any transitive relation containing R, then $R^+ \subset T$.

By induction on i, we can show that $R^i \subseteq T$.



Injection (one-to-one; 1-1)

Surjection

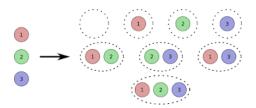
Bijection (one-to-one correspondence)

Definition (Characteristic Function (特征函数) of a Subset)

For a given subset $A \subseteq X$,

$$\chi_A: X \to \{0, 1\}$$

$$\chi_A(x) = 1 \iff x \in A.$$



$$\chi_A: X \to \{0,1\}$$
 vs. $\mathcal{P}(X)$

Definition (Natural Function)

Let $R \subseteq A \times A$ be an equivalence relation. The following function f

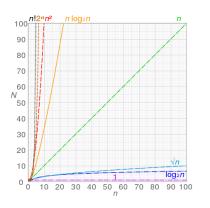
$$f: A \to A/R$$

$$f: a \mapsto R(a)$$

is called the <u>natural function</u> on A.



Asymptotic Growth Rates of Functions





https://www.bilibili.com/video/BV175411T7ph?share_source=copy_web



Definition (Order Isomorphism (同构))

Given two posets (S, \leq_S) and (T, \leq_T) , an order isomorphism from (S, \leq_S) to (T, \leq_T) is a bijection from S to T such that

$$\forall x, y \in S. \ x \leq_S y \leftrightarrow f(x) \leq_T f(y).$$

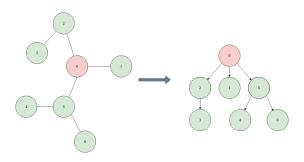
$$(\mathbb{R}, \leq) \xrightarrow{f: \mathbb{R} \to \mathbb{R}} (\mathbb{R}, \geq)$$

Definition (Order Automorphism (自同构))

An order isomorphism from a poset to itself is an order automorphism.

Definition (Rooted Tree (有根树))

A rooted tree is a tree where one vertex has been designated the root.



Definition (Directed Rooted Tree (有向有根树))

A directed rooted tree is a rooted tree where all edges directed away from or towards the root.

Definition

Parent, Child; Sibling; Ancestor, Descendant

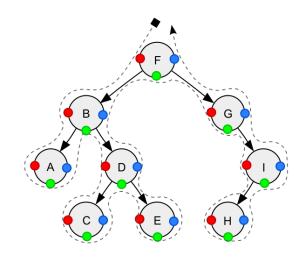
Definition (k-ary Trees (k-叉树))

A k-ary tree is a rooted tree in which each vertex has $\leq k$ children.

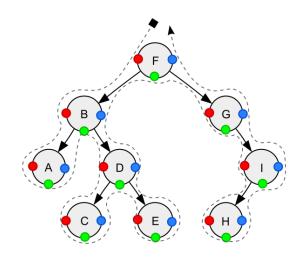
2-ary trees are often called binary trees.

Definition (Complete k-Tree (完全 k-叉树))

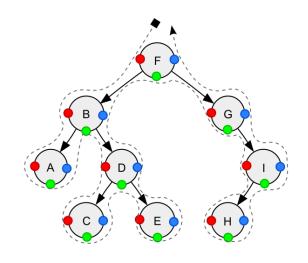
A complete k-tree is a k-ary tree in which each vertex, other than leaves, has = k children.



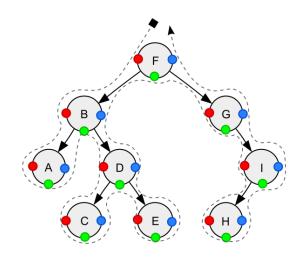
Depth-First Search (DFS)



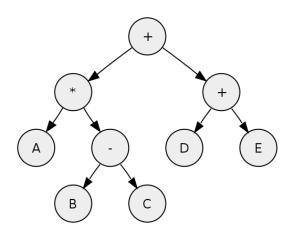
Pre-order (前序) Traversal: F, B, A, D, C, E, G, I, H



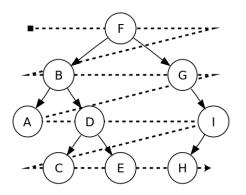
In-order (中序) Traversal: A, B, C, D, E, F, G, H, I



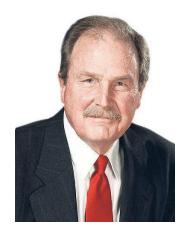
Post-order (后序) Traversal: A, C, E, D, B, H, I, G, F



Prefix Expression (前缀表达式): +*A - BC + DEInfix Expression (中缀表达式): A*(B-C) + (D+E)Postfix Expression (后缀表达式): ABC - *DE + +



Breadth-First Search (BFS): F, B, G, A, D, I, C, E, H



David A. Huffman (1925 $\sim 1999)$

$C[1 \dots n]$	a	b	c	d	e	f
$F[1 \dots n]$	45	13	12	16	9	5
Fixed Length Code	000	001	010	011	100	101
Variable Length Code	0	101	100	111	1101	1100

Prefix code (前缀码): No code is a prefix of some other code

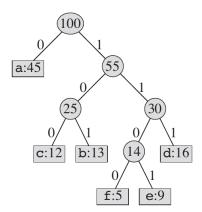
The Encoding Problem

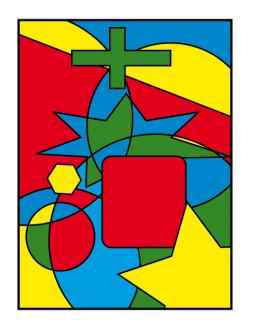
To find the optimal binary prefix code for C and F.

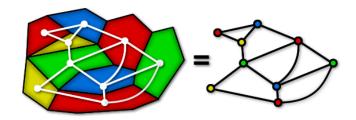
Let E be a binary prefix code for C and F. The length L(E) is

$$L(E) = \sum_{c \in C} f_c \cdot l_E(c)$$

$C[1 \dots n]$	a	b	c	d	e	f
$F[1 \dots n]$	45	13	12	16	9	5



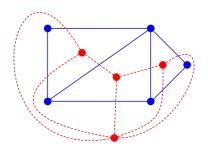


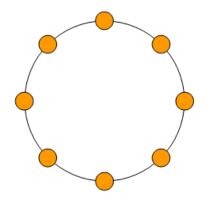


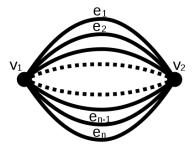
Definition (Dual Graph (对偶图))

The dual graph of a plane graph G is a graph G'

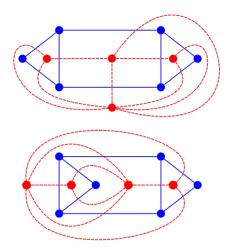
- ightharpoonup G' has a vertex for each face of G;
- ightharpoonup G' has an edge for each pair of faces in G that are separated from each other by an edge, and a self-loop when the same face appears on both sides of an edge.





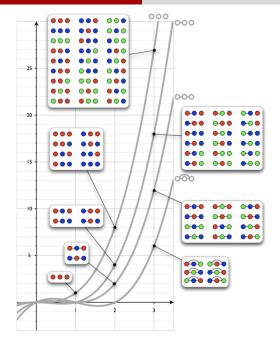


The dual graph G' depends on the choice of embedding of the graph G.



Theorem

G is a bipartite graph $\iff \chi(G) = 2 \iff G$ has no odd cycles.



Definition (Chromatic Polynomial (色多项式; 非严格定义))

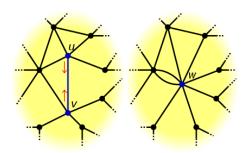
The chromatic polynomial P(G, k) counts the number of colorings of graph G as a function of the number k of colors.

Triangle K_3	x(x-1)(x-2)
Complete graph K_n	$x(x-1)(x-2)\cdots(x-(n-1))$
Edgeless graph \overline{K}_n	x^n
Path graph P_n	$x(x-1)^{n-1}$
Any tree on n vertices	$x(x-1)^{n-1}$
Cycle C_n	$(x-1)^n + (-1)^n(x-1)$
Petersen graph	$x(x-1)(x-2)\left(x^7-12x^6+67x^5-230x^4+529x^3-814x^2+775x-352\right)$

Theorem (Recurrence for Chromatic Polynomial)

Given a graph G and an edge $e \in E(G)$, then

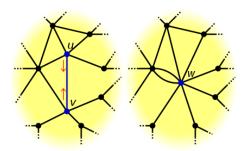
$$P(G,k) = P(G - e, k) - P(G/e, k)$$



G/e: 边的收缩

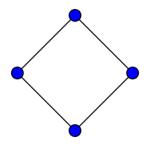
$$P(G, k) = P(G - e, k) - P(G/e, k)$$

$$P(G - e, k) = P(G/e, k) + P(G, k)$$



In $G - \{u, v\}$, Color(u) = Color(v) or $Color(u) \neq Color(v)$.

$$P(G,k) = P(G - e, k) - P(G/e, k)$$



$$P(C_4, k) = P(P_4, k) - P(K_3, k)$$

$$= k(k-1)^3 - k(k-1)(k-2)$$

$$= k(k-1)(k^2 - 3k + 3)$$

$$= (k-1)^4 + (-1)^4(k-1)$$

Cyclic Notation (轮换表示法) & Transposition (对换)

$$\sigma = \begin{pmatrix} 1 & 2 & 3 & 4 & 5 & 6 \\ 4 & 3 & 6 & 1 & 5 & 2 \end{pmatrix}$$

$$\sigma = (1 \ 4)(2 \ 3 \ 6)(5)$$

$$= (1 \ 4)(2 \ 3 \ 6)$$

$$= (2 \ 3 \ 6)(1 \ 4)$$

$$= (2 \ 3 \ 6)(4 \ 1)$$

$$= (3 \ 6 \ 2)(4 \ 1)$$

$$= (3 \ 6)(6 \ 2)(4 \ 1)$$

$$(i_1 \ i_2 \ \dots \ i_r) = (i_1 \ i_2)(i_2 \ i_3) \dots (i_{r-2} \ i_{r-1})(i_{r-1} \ i_r)$$

By induction on the length r.

$$\sigma = \begin{pmatrix} 1 & 2 & 3 & 4 & 5 & 6 & 7 \\ 7 & 3 & 6 & 2 & 5 & 4 & 1 \end{pmatrix} = (1\ 7)(2\ 3)(3\ 6)(6\ 4)$$
$$= (1\ 7)(3\ 6)(2\ 5)(6\ 4)(4\ 5)(2\ 5)$$

Theorem (Parity (奇偶性) of Permutations)

将一个置换表示成若干对换的乘积, 所用对换个数的奇偶性是唯一的。

Definition (Even/Odd Permutations (偶置换/奇置换))

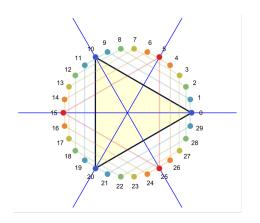
可表示为偶数个对换的乘积的置换称为偶置换; 否则, 称为奇置换。

Definition (Alternating Group (交错群; A_n))

由 S_n 的全体偶置换构成的子群称为 n 次交错群。

$$A_3 = \{(1), (1\ 2\ 3), (1\ 3\ 2)\}$$

$$\operatorname{sgn}: S_n \to \{1, -1\}$$
$$\operatorname{sgn}(\sigma_1 \sigma_2) = \operatorname{sgn}(\sigma_1) \operatorname{sgn}(\sigma_2)$$
$$S_n / A_n \cong \{1, -1\}$$



$$A_3 = \{(1), (1\ 2\ 3), (1\ 3\ 2)\}$$

 $(1\ 2)A_3 = \{(1\ 2), (2\ 3), (1\ 3)\}$

Thank You!



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