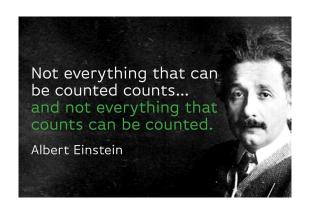
2-3 Counting

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所以, 学好 "2-3 组合与计数" 是多么重要!

Paring up (CS : 1.2 - 15)

A tennis club has 2n members. We want to pair up the members by twos for singles matches.

- (a) In how many ways can we pair up all the members of the club?
- (b) Suppose that we also determine who serves first for each pairing. In how many ways can we specify our pairs?

$$\frac{1}{n!} \binom{2n}{2, 2, \cdots, 2} = \frac{(2n)!}{\underbrace{2^n \cdot n!}_{\text{intra-pair inter-pair}}}$$

$$\frac{(2n)!}{2^n \cdot n!} \cdot 2^n = \frac{(2n)!}{n!}$$

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Passing out Apples to Children



k-Permutation (CS : 1.2 - 5)

We need to pass out k distinct apples (pieces of fruit) to n children such that each child may get at most one apple.

- (a) $k \leq n$?
- (b) What if k > n?

$$n^{\underline{k}} \triangleq n(n-1)\cdots(n-k+1)$$

0

Multisets (CS: 1.5-4)

Use multisets to determine the number of ways to pass out k identical apples to n children. Assume that a child may get more than one apple.

$$x_i$$
: the # of apples the *i*-th child gets

$$x_1 + x_2 + \dots + x_n = k, \quad x_i \ge 0$$

Integer composition (The order matters!)

$$y_i \triangleq x_i + 1$$

$$y_1 + y_2 + \dots + y_n = n + k, \qquad y_i \ge 1$$

$$\binom{n+k-1}{n-1} = \binom{n+k-1}{k}$$

Multisets (CS: 1.5-4)

Use **multisets** to determine the number of ways to pass out k identical apples to n children. Assume that a child may get more than one apple.

Q: k-multiset of $[1 \cdots n]$ vs. n-multiset of $[1 \cdots k]$

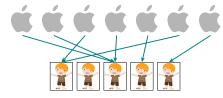
$$k = 7$$
 $n = 5$



Integer Partition (CS : 1.5 - 4 Extended)

What is the number of ways to pass out k identical apples to n- \hbar \hbar . Assume that a child may get more than one apple.





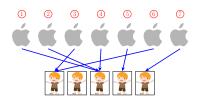
Integer partition of k into $\leq n$ parts (The order does not matter!)

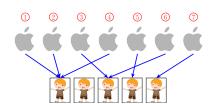
Theorem (G. H. Hardy, Ramanujan (1918))

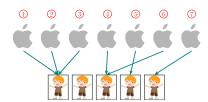
$$p(k) \triangleq \sum_{x=1}^{x=k} p_x(k) \sim \frac{1}{4\sqrt{3}k} \exp\left(\pi\sqrt{\frac{2k}{3}}\right)$$

Set Partition (CS : 1.5 - 4 Extended)

What is the number of ways to pass out k distinct apples to n-胞胎. Assume that a child may get more than one apple.







Set partition of $[1 \cdots k]$ into $\leq n$ parts

Set Partition (CS: 1.5 - 12)

$$S(n,k)$$
 $(\begin{Bmatrix} n \\ k \end{Bmatrix})$: # of set partitions of $[1\cdots n]$ into k classes

Stirling number of the second kind

Theorem (Recurrence for S(n,k))

$$S(0,0) = 1, \quad S(n,0) = S(0,n) = 0 \ (n > 0)$$

$$S(n,k) = S(n-1,k-1) + kS(n-1,k), \quad n > 0, k > 0$$

Proof.

$$S(n,k) = \underbrace{S(n-1,k-1)}_{n \text{ is alone}} + \underbrace{kS(n-1,k)}_{n \text{ is not alone}}$$



Bell number:
$$B_n = \sum_{k=0}^{k=n} \begin{Bmatrix} n \\ k \end{Bmatrix}$$

Theorem (Berend & Tassa (2010))

$$B_n < \left(\frac{0.792n}{\ln(n+1)}\right)^n, n \in \mathbb{Z}^+$$

Theorem (de Bruijn (1981))

As $n \to \infty$,

$$\frac{\ln B_n}{n} = \ln n - \ln \ln n - 1 + \frac{\ln \ln n}{\ln n} + \frac{1}{\ln n} + \frac{1}{2} \left(\frac{\ln \ln n}{\ln n} \right)^2 + O\left(\frac{\ln \ln n}{(\ln n)^2} \right)$$

THE TWELVEFOLD WAY

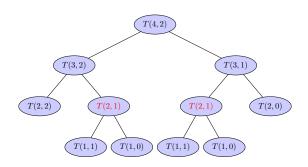
balls per urn	unrestricted	≤ 1	≥ 1
n labeled balls, m labeled urns	n-tuples of m things	n-permutations of m things	partitions of $\{1, \ldots, n\}$ into m ordered parts
n unlabeled balls, m labeled urns	n-multicombinations of m things	n-combinations of m things	compositions of n into m parts
n labeled balls, m unlabeled urns	partitions of $\{1, \dots, n\}$ into $\leq m$ parts	n pigeons into m holes	partitions of $\{1, \dots, n\}$ into m parts
n unlabeled balls, m unlabeled urns	partitions of n into $\leq m$ parts	n pigeons into m holes	partitions of n into m parts

Computing $\binom{n}{k}$ (CS 1.5:14)

1: **procedure** BINOM(n,k)

ightharpoonupRequired: $n \ge k \ge 0$

- 2: if $k = 0 \lor n = k$ then
- 3: **return** 1
- 4: **return** BINOM(n-1,k) + BINOM(n-1,k-1)



1: **procedure** BINOM(n,k)

 \triangleright Required: $n \ge k \ge 0$

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4: **return** BINOM(n-1,k) + BINOM(n-1,k-1)

(i) # of "+":

$$A(n,k) = 1 + A(n-1,k) + A(n-1,k-1)$$

(ii) # of recursive calls of BINOM:

$$R(n,k) = 2 + R(n-1,k) + R(n-1,k-1)$$

$$T(n,k) = \begin{cases} 0, & k = 0 \lor n = k \\ T(n-1,k) + T(n-1,k-1) + c, & o.w. \end{cases}$$

$$T(n,k) = \left\{ \begin{array}{ll} 0, & k = 0 \lor n = k \\ T(n-1,k) + T(n-1,k-1) + c, & \text{o.w.} \end{array} \right.$$

$$T(n,k) = T(n-1,k) + T(n-1,k-1) \implies T(n,k) = \frac{\alpha}{k} \binom{n}{k}$$

$$T(n,k) = \alpha \binom{n}{k} + \beta$$

$$\alpha \binom{n}{k} + \beta = \alpha \binom{n-1}{k} + \beta + \alpha \binom{n-1}{k-1} + \beta + c \implies \beta = -c$$

$$\alpha \binom{n}{0} - c = 0, \quad \alpha \binom{n}{n} - c = 0 \implies \alpha = c$$

$$T(n,k) = c \binom{n}{k} - c$$

```
    7. 斐波那契数列的定义如下: F₁=1, F₂=1, Fₙ=Fₙ₋₁+Fₙ₋₂(n≥3)。如果用下面的函数计算斐波那契数列的第 n 项,则其时间复杂度为 ( )。
        int F(int n)
        {
            if (n <= 2)
                 return 1;
            else
                 return F(n - 1) + F(n - 2);
        }
        A. O(1) B. O(n) C. O(n²) D. O(Fռ)</li>
```

时间复杂度 of what?

$$\begin{pmatrix}
0 \\
0
\end{pmatrix}$$

$$\begin{pmatrix}
1 \\
0
\end{pmatrix}
\begin{pmatrix}
1 \\
1
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\end{pmatrix}
\begin{pmatrix}
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1
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\begin{pmatrix}
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\end{pmatrix}
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\begin{pmatrix}
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4
\end{pmatrix}$$

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\end{pmatrix}
\begin{pmatrix}
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1
\end{pmatrix}
\begin{pmatrix}
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2
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\begin{pmatrix}
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4
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\begin{pmatrix}
5 \\
5
\end{pmatrix}$$

Q: How to calculate $\binom{5}{3}$?

```
1: procedure BINOM(n,k)
                                                                    \triangleright Required: n > k > 0
         for i \leftarrow 0 to n-k do
 2:
              B[i][0] \leftarrow 1
 3:
         for i \leftarrow 1 to k do
 4:
              B[i][i] \leftarrow 1
 5:
         for j \leftarrow 1 to k do
 6:
              for d \leftarrow 1 to n - k do
 7:
                  i \leftarrow j + d
 8:
                  B[i][j] \leftarrow B[i-1][j] + B[i-1][j-1]
 9:
         return B[n][k]
10:
```

$$(n-k+1) + (k) + k(n-k) = nk - k^2 + n + 1$$

Thank You!