Functions

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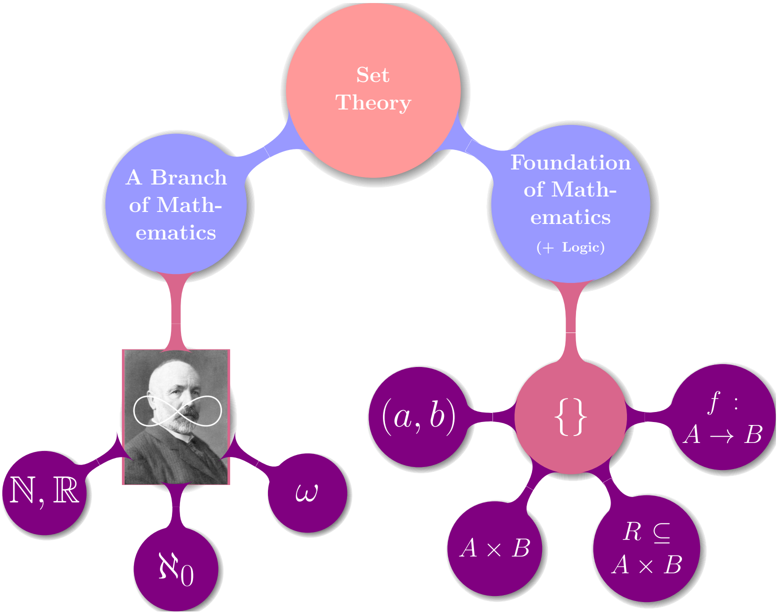
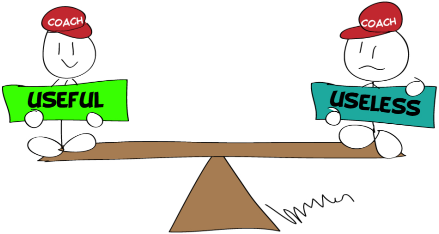
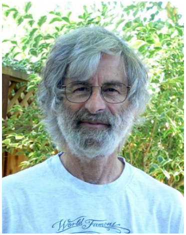

Set Theory
A Branch of Math-ematics
Foundation of Math-ematics
(+ Logic)
$\mathbb{N}, \mathbb{R}$
$\aleph_0$
ω
(a, b)
$\{\}$
$f : A \to B$
$A \times B$
$R \subseteq A \times B$


COACH
USEFUL
COACH
USELESS



Time, Clocks, and the Ordering of Events in a Distributed System

Leslie Lamport
Massachusetts Computer Associates, Inc.

The concept of one event happening before another in a distributed system is examined, and is shown to define a partial ordering of the events. A distributed algorithm is given for synchronizing a system of logical clocks which can be used to totally order the events.

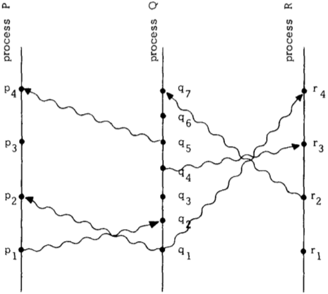


Figure 13. A selection of consistency axioms over an execution $(E, \mathsf{repl}, \mathsf{obj}, \mathsf{oper}, \mathsf{rval}, \mathsf{ro}, \mathsf{vis}, \mathsf{ar})$

Auxiliary relations

$\mathsf{sameobj}(e, f) \iff \mathsf{obj}(e) = \mathsf{obj}(f)$

Per-object causality (aka happens-before) order:

$\quad \mathsf{hbo} = ((\mathsf{ro} \cap \mathsf{sameobj}) \cup \mathsf{vis})^+$

Causality (aka happens-before) order: $\mathsf{hb} = (\mathsf{ro} \cup \mathsf{vis})^+$

Axioms

EVENTUAL:

$\quad \forall e \in E.\ \neg(\exists \text{ infinitely many } f \in E.\ \mathsf{sameobj}(e, f) \wedge \neg(e \xrightarrow{\mathsf{vis}} f))$

THINAIR: $\mathsf{ro} \cup \mathsf{vis}$ is acyclic

POCV (Per-Object Causal Visibility): $\mathsf{hbo} \subseteq \mathsf{vis}$

POCA (Per-Object Causal Arbitration): $\mathsf{hbo} \subseteq \mathsf{ar}$

COCV (Cross-Object Causal Visibility): $(\mathsf{hb} \cap \mathsf{sameobj}) \subseteq \mathsf{vis}$

COCA (Cross-Object Causal Arbitration): $\mathsf{hb} \cup \mathsf{ar}$ is acyclic

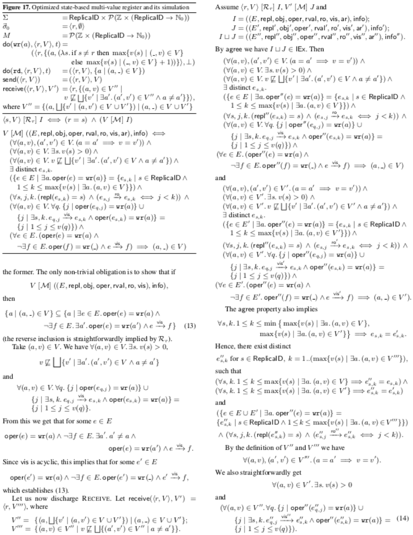


Figure 17. Optimized state-based multi-value register and its simulation

$\Sigma = \mathsf{ReplicaID} \times \mathcal{P}(\mathbb{Z} \times (\mathsf{ReplicaID} \to \mathbb{N}_0))$

$\vec{\sigma}_0 = \langle r, \emptyset \rangle$

$M = \mathcal{P}(\mathbb{Z} \times (\mathsf{ReplicaID} \to \mathbb{N}_0))$

$\mathsf{do}(\mathtt{wr}(a), \langle r, V \rangle, t) = (\langle r, \{(a, (\lambda s.\ \text{if } s \neq r \text{ then } \max\{v(s) \mid (\_, v) \in V\} \text{ else } \max\{v(s) \mid (\_, v) \in V\} + 1))\}\rangle, \bot)$

$\mathsf{do}(\mathtt{rd}, \langle r, V \rangle, t) = (\langle r, V \rangle, \{a \mid (a, \_) \in V\})$

$\mathsf{send}(\langle r, V \rangle) = (\langle r, V \rangle, V)$

$\mathsf{receive}(\langle r, V \rangle, V') = \langle r, \{(a, v) \in V'' \mid v \not\sqsubseteq \bigsqcup\{v' \mid \exists a'. (a', v') \in V'' \wedge a \neq a'\}\}\rangle$,

where $V'' = \{(a, \bigsqcup\{v' \mid (a, v') \in V \cup V'\}) \mid (a, \_) \in V \cup V'\}$

$\langle s, V \rangle\ [\mathcal{R}_r]\ I \iff (r = s) \wedge (V\ [\mathcal{M}]\ I)$

$V\ [\mathcal{M}]\ ((E, \mathsf{repl}, \mathsf{obj}, \mathsf{oper}, \mathsf{rval}, \mathsf{ro}, \mathsf{vis}, \mathsf{ar}), \mathsf{info}) \iff$
$(\forall (a, v), (a', v') \in V.\ (a = a' \implies v = v')) \wedge$
$(\forall (a, v) \in V.\ \exists s.\ v(s) > 0) \wedge$
$(\forall (a, v) \in V.\ v \not\sqsubseteq \bigsqcup\{v' \mid \exists a'. (a', v') \in V \wedge a \neq a'\}) \wedge$
$\exists$ distinct $e_{s,k}$.
$(\{e \in E \mid \exists a.\ \mathsf{oper}(e) = \mathtt{wr}(a)\} = \{e_{s,k} \mid s \in \mathsf{ReplicaID} \wedge 1 \le k \le \max\{v(s) \mid \exists a. (a, v) \in V\}\}) \wedge$
$(\forall s, j, k.\ (\mathsf{repl}(e_{s,k}) = s) \wedge (e_{s,j} \xrightarrow{\mathsf{ro}} e_{s,k} \iff j < k)) \wedge$
$(\forall (a, v) \in V.\ \forall q.\ \{j \mid \mathsf{oper}(e_{q,j}) = \mathtt{wr}(a)\} \cup \{j \mid \exists s, k.\ e_{q,j} \xrightarrow{\mathsf{vis}} e_{s,k} \wedge \mathsf{oper}(e_{s,k}) = \mathtt{wr}(a)\} = \{j \mid 1 \le j \le v(q)\}) \wedge$
$(\forall e \in E.\ (\mathsf{oper}(e) = \mathtt{wr}(a) \wedge \neg\exists f \in E.\ \mathsf{oper}(f) = \mathtt{wr}(\_) \wedge e \xrightarrow{\mathsf{vis}} f) \implies (a, \_) \in V)$

the former. The only non-trivial obligation is to show that if

$V\ [\mathcal{M}]\ ((E, \mathsf{repl}, \mathsf{obj}, \mathsf{oper}, \mathsf{rval}, \mathsf{ro}, \mathsf{vis}), \mathsf{info})$,

then

$$\{a \mid (a, \_) \in V\} \subseteq \{a \mid \exists e \in E.\ \mathsf{oper}(e) = \mathtt{wr}(a) \wedge \neg\exists f \in E.\ \exists a'.\ \mathsf{oper}(e) = \mathtt{wr}(a') \wedge e \xrightarrow{\mathsf{vis}} f\} \quad (13)$$

(the reverse inclusion is straightforwardly implied by $\mathcal{R}_r$).
Take $(a, v) \in V$. We have $\forall (a, v) \in V.\ \exists s.\ v(s) > 0$,

$$v \not\sqsubseteq \bigsqcup\{v' \mid \exists a'. (a', v') \in V \wedge a \neq a'\}$$

and

$$\forall (a, v) \in V.\ \forall q.\ \{j \mid \mathsf{oper}(e_{q,j}) = \mathtt{wr}(a)\} \cup \{j \mid \exists s, k.\ e_{q,j} \xrightarrow{\mathsf{vis}} e_{s,k} \wedge \mathsf{oper}(e_{s,k}) = \mathtt{wr}(a)\} = \{j \mid 1 \le j \le v(q)\}.$$

From this we get that for some $e \in E$

$$\mathsf{oper}(e) = \mathtt{wr}(a) \wedge \neg\exists f \in E.\ \exists a'.\ a' \neq a \wedge \mathsf{oper}(e) = \mathtt{wr}(a') \wedge e \xrightarrow{\mathsf{vis}} f.$$

Since vis is acyclic, this implies that for some $e' \in E$

$$\mathsf{oper}(e') = \mathtt{wr}(a) \wedge \neg\exists f \in E.\ \mathsf{oper}(e') = \mathtt{wr}(\_) \wedge e' \xrightarrow{\mathsf{vis}} f,$$

which establishes (13).
Let us now discharge RECEIVE. Let $\mathsf{receive}(\langle r, V \rangle, V') = \langle r, V''' \rangle$, where

$$V'' = \{(a, \bigsqcup\{v' \mid (a, v') \in V \cup V'\}) \mid (a, \_) \in V \cup V'\};$$
$$V''' = \{(a, v) \in V'' \mid v \not\sqsubseteq \bigsqcup\{(a', v') \in V'' \mid a \neq a'\}\}.$$

Assume $\langle r, V \rangle\ [\mathcal{R}_r]\ I$, $V'\ [\mathcal{M}]\ J$ and

$I = ((E, \mathsf{repl}, \mathsf{obj}, \mathsf{oper}, \mathsf{rval}, \mathsf{ro}, \mathsf{vis}, \mathsf{ar}), \mathsf{info})$;
$J = ((E', \mathsf{repl}', \mathsf{obj}', \mathsf{oper}', \mathsf{rval}', \mathsf{ro}', \mathsf{vis}', \mathsf{ar}'), \mathsf{info}')$;
$I \sqcup J = ((E'', \mathsf{repl}'', \mathsf{obj}'', \mathsf{oper}'', \mathsf{rval}'', \mathsf{ro}'', \mathsf{vis}'', \mathsf{ar}''), \mathsf{info}'')$.

By agree we have $I \sqcup J \in \mathsf{IEx}$. Then

$(\forall (a, v), (a', v') \in V.\ (a = a' \implies v = v')) \wedge$
$(\forall (a, v) \in V.\ \exists s.\ v(s) > 0) \wedge$
$(\forall (a, v) \in V.\ v \not\sqsubseteq \bigsqcup\{v' \mid \exists a'. (a', v') \in V \wedge a \neq a'\}) \wedge$
$\exists$ distinct $e_{s,k}$.
$(\{e \in E \mid \exists a.\ \mathsf{oper}''(e) = \mathtt{wr}(a)\} = \{e_{s,k} \mid s \in \mathsf{ReplicaID} \wedge 1 \le k \le \max\{v(s) \mid \exists a. (a, v) \in V\}\}) \wedge$
$(\forall s, j, k.\ (\mathsf{repl}''(e_{s,k}) = s) \wedge (e_{s,j} \xrightarrow{\mathsf{ro}} e_{s,k} \iff j < k)) \wedge$
$(\forall (a, v) \in V.\ \forall q.\ \{j \mid \mathsf{oper}''(e_{q,j}) = \mathtt{wr}(a)\} \cup \{j \mid \exists s, k.\ e_{q,j} \xrightarrow{\mathsf{vis}} e_{s,k} \wedge \mathsf{oper}''(e_{s,k}) = \mathtt{wr}(a)\} = \{j \mid 1 \le j \le v(q)\}) \wedge$
$(\forall e \in E.\ (\mathsf{oper}''(e) = \mathtt{wr}(a) \wedge \neg\exists f \in E.\ \mathsf{oper}''(f) = \mathtt{wr}(\_) \wedge e \xrightarrow{\mathsf{vis}} f) \implies (a, \_) \in V)$

and

$(\forall (a, v), (a', v') \in V'.\ (a = a' \implies v = v')) \wedge$
$(\forall (a, v) \in V'.\ \exists s.\ v(s) > 0) \wedge$
$(\forall (a, v) \in V'.\ v \not\sqsubseteq \bigsqcup\{v' \mid \exists a'. (a', v') \in V' \wedge a \neq a'\}) \wedge$
$\exists$ distinct $e_{s,k}$.
$(\{e \in E' \mid \exists a.\ \mathsf{oper}''(e) = \mathtt{wr}(a)\} = \{e_{s,k} \mid s \in \mathsf{ReplicaID} \wedge 1 \le k \le \max\{v(s) \mid \exists a. (a, v) \in V'\}\}) \wedge$
$(\forall s, j, k.\ (\mathsf{repl}''(e_{s,k}) = s) \wedge (e_{s,j} \xrightarrow{\mathsf{ro}'} e_{s,k} \iff j < k)) \wedge$
$(\forall (a, v) \in V'.\ \forall q.\ \{j \mid \mathsf{oper}''(e_{q,j}) = \mathtt{wr}(a)\} \cup \{j \mid \exists s, k.\ e_{q,j} \xrightarrow{\mathsf{vis}'} e_{s,k} \wedge \mathsf{oper}''(e_{s,k}) = \mathtt{wr}(a)\} = \{j \mid 1 \le j \le v(q)\}) \wedge$
$(\forall e \in E'.\ (\mathsf{oper}''(e) = \mathtt{wr}(a) \wedge \neg\exists f \in E'.\ \mathsf{oper}''(f) = \mathtt{wr}(\_) \wedge e \xrightarrow{\mathsf{vis}'} f) \implies (a, \_) \in V')$.

The agree property also implies

$$\forall s, k.\ 1 \le k \le \min\{\max\{v(s) \mid \exists a. (a, v) \in V\}, \max\{v(s) \mid \exists a. (a, v) \in V'\}\} \implies e_{s,k} = e'_{s,k}.$$

Hence, there exist distinct

$$e''_{s,k} \text{ for } s \in \mathsf{ReplicaID},\ k = 1..(\max\{v(s) \mid \exists a. (a, v) \in V'''\}),$$

such that

$(\forall s, k.\ 1 \le k \le \max\{v(s) \mid \exists a. (a, v) \in V\} \implies e''_{s,k} = e_{s,k}) \wedge$
$(\forall s, k.\ 1 \le k \le \max\{v(s) \mid \exists a. (a, v) \in V'\} \implies e''_{s,k} = e'_{s,k})$

and

$(\{e \in E \cup E' \mid \exists a.\ \mathsf{oper}''(e) = \mathtt{wr}(a)\} = \{e''_{s,k} \mid s \in \mathsf{ReplicaID} \wedge 1 \le k \le \max\{v(s) \mid \exists a. (a, v) \in V'''\}\})$
$\wedge (\forall s, j, k.\ (\mathsf{repl}(e''_{s,k}) = s) \wedge (e''_{s,j} \xrightarrow{\mathsf{ro}''} e''_{s,k} \iff j < k))$.

By the definition of V'' and V''' we have

$$\forall (a, v), (a', v') \in V'''.\ (a = a' \implies v = v').$$

We also straightforwardly get

$$\forall (a, v) \in V'.\ \exists s.\ v(s) > 0$$

and

$$(\forall (a, v) \in V''.\ \forall q.\ \{j \mid \mathsf{oper}''(e''_{q,j}) = \mathtt{wr}(a)\} \cup \{j \mid \exists s, k.\ e''_{q,j} \xrightarrow{\mathsf{vis}''} e''_{s,k} \wedge \mathsf{oper}''(e''_{s,k}) = \mathtt{wr}(a)\} = \{j \mid 1 \le j \le v(q)\}). \quad (14)$$

Function

(1) from the perspective of set theory

(2) PROOF! PROOF! PROOF!

Definition of Function

Definition (Relation)

Let A and B be sets.
R is a (binary) relation if

$$R \subseteq A \times B = \{(a, b) \mid a \in A \land b \in B\}$$

Definition (Function)

Let A and B be sets.
A **function** f from A to B is a *relation* f from A to B such that

$$\forall a \in A \; \exists! b \in B \; (a, b) \in f.$$

For Proof:

$$\forall$$

$$\exists! : \forall b, b' \in B, (a, b) \in f \land (a, b') \in f \implies b = b'.$$

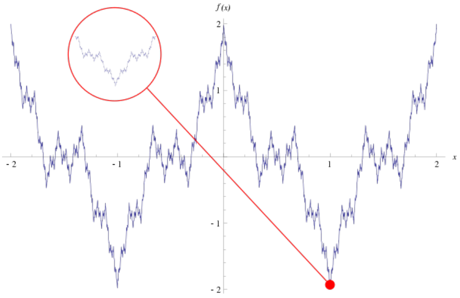
$$f : A \to B, \quad a \mapsto f(a) \qquad \Big(b = f(a)\Big)$$

$$A : \mathrm{dom}(f) \qquad B : \mathrm{cod}(f)$$

$$\mathrm{ran}(f) = f(A) = \{f(a) \mid a \in A\} \subseteq B$$

$$D(x) = \begin{cases} 1 & \text{if } x \in \mathbb{Q} \\ 0 & \text{if } x \in \mathbb{R} \setminus \mathbb{Q} \end{cases}$$

Dirichlet Function



Weierstrass Function (1872)

$$f(x) = \sum_{n=0}^{\infty} a^n \cos(b^n \pi x)$$

$$0 < a < 1,\ b \in 2\mathbb{N}+1,\ ab > 1 + \frac{3}{2}\pi$$

Problem $13.3\ (g)$

$$f : \mathbb{Q} \to \mathbb{R}$$

$$f(x) = \begin{cases} x+1 & \text{if } x \in 2\mathbb{Z} \\ x-1 & \text{if } x \in 3\mathbb{Z} \\ 2 & \text{otherwise} \end{cases}$$

Problem 13.4

$$f : \mathcal{P}(\mathbb{R}) \to \mathbb{Z}$$

$$f(A) = \begin{cases} \min(A \cap \mathbb{N}) & \text{if } A \cap \mathbb{N} \neq \emptyset \\ -1 & \text{if } A \cap \mathbb{N} = \emptyset \end{cases}$$

A function $f : A \to B$ is a set.

$$f \subseteq A \times B$$

$$A \times B = \{(a, b) \mid a \in A, b \in B\}$$

$$(a, b) = \{\{a\}, \{a, b\}\}$$

Definition (Axiom of Extensionality (集合的外延公理))

$$\forall A \forall B \forall x (x \in A \iff x \in B) \iff A = B.$$

Definition (函数的外延性原则)

$$f = g \iff \mathsf{dom}(f) = \mathsf{dom}(g) \wedge (\forall x \in \mathsf{dom}(f) : f(x) = g(x))$$

Special Functions (*-jectivity*)

Definition (Injective (one-to-one; 1-1) 单射函数)

$$f : A \to B \qquad f : A \rightarrowtail B$$

$$\forall a_1, a_2 \in A : a_1 \neq a_2 \implies f(a_1) \neq f(a_2)$$

For Proof:

- To prove that f *is* 1-1:

$$\forall a_1, a_2 \in A : f(a_1) = f(a_2) \implies a_1 = a_2$$

- To show that f *is not* 1-1:

$$\exists a_1, a_2 \in A : a_1 \neq a_2 \wedge f(a_1) = f(a_2)$$

Definition (Surjective (onto) 满射函数)

$$f : A \to B \qquad f : A \twoheadrightarrow B$$

$$\text{ran}(f) = B$$

For Proof:

- To prove that f *is* onto:

$$\forall b \in B \left(\exists a \in A : f(a) = b \right)$$

- To show that f *is not* onto:

$$\exists b \in B \left(\forall a \in A : f(a) \neq b \right)$$

Theorem (Cantor Theorem (ES Theorem 24.4))

Let A be a set.
If $f : A \to 2^A$, then f is not onto.

Proof.

Proof. Let A be a set and let $f : A \to 2^A$. To show that f is not onto, we must find a $B \in 2^A$ (i.e., $B \subseteq A$) for which there is no $a \in A$ with $f(a) = B$. In other words, B is a set that f "misses." To this end, let

$$B = \{x \in A : x \notin f(x)\}.$$

We claim there is no $a \in A$ with $f(a) = B$.

Suppose, for the sake of contradiction, there is an $a \in A$ such that $f(a) = B$. We ponder: Is $a \in B$?

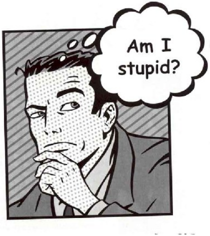
- If $a \in B$, then, since $B = f(a)$, we have $a \in f(a)$. So, by definition of B, $a \notin f(a)$; that is, $a \notin B$. $\Rightarrow\Leftarrow$
- If $a \notin B = f(a)$, then, by definition of B, $a \in B$. $\Rightarrow\Leftarrow$

Both $a \in B$ and $a \notin B$ lead to contradictions, and hence our supposition [there is an $a \in A$ with $f(a) = B$] is false, and therefore f is not onto. ■

□

Theorem (Cantor Theorem)

Let A be a set.
If $f : A \to 2^A$, then f is not onto.



Theorem (Cantor Theorem)

Let A be a set.
If $f : A \to 2^A$, then f is not onto.

Understanding this problem:

$$A = \{1, 2, 3\}$$

2^A

$$2^A = \Big\{\emptyset, \{1\}, \{2\}, \{3\}, \{1,2\}, \{1,3\}, \{2,3\}, \{1,2,3\}\Big\}$$

Onto

$$\forall B \in 2^A \, \Big(\exists a \in A \; f(a) = B\Big).$$

Not Onto

$$\exists B \in 2^A \, \Big(\forall a \in A \; f(a) \neq B\Big).$$

Theorem (Cantor Theorem)

Let A be a set.
If $f : A \to 2^A$, then f is not onto.

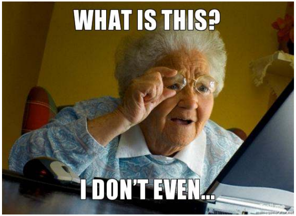
Proof.

- Constructive proof ($\exists$):

$$B = \{x \in A \mid x \notin f(x)\}.$$

- By contradiction ($\forall$):

$$\exists a \in A : f(a) = B.$$



$$Q : a \in B?$$

$\square$

Theorem (Cantor Theorem)

Let A be a set.
If $f : A \to 2^A$, then f is not onto.

对角线论证 (Cantor's diagonal argument) (以下仅适用于可数集合 A).

a	$f(a)$					
	1	2	3	4	5	$\cdots$
1	1	1	0	0	1	$\cdots$
2	0	0	0	0	0	$\cdots$
3	1	0	0	1	0	$\cdots$
4	1	1	1	1	1	$\cdots$
5	0	1	0	1	0	$\cdots$
$\vdots$	$\vdots$	$\vdots$	$\vdots$	$\vdots$	$\vdots$	$\ldots$

$$B = \{0, 1, 1, 0, 1\}$$

□

Definition (Bijective (one-to-one correspondence) 一一对应)

$$f : A \to B \qquad f : A \xleftrightarrow[onto]{1-1} B$$

1-1 & onto

Problem 14.12

$$a, b, c, d \in \mathbb{R}, \; a < b, \; c < d$$

Define a bijective function:

$$f : [a, b] \xleftrightarrow[onto]{1-1} [c, d]$$

$$f : (a, b) \xleftrightarrow[onto]{1-1} (c, d)$$

Answer.

$$f(x) = c + \frac{d - c}{b - a}(x - a)$$

$\square$

Operations on Functions

Set

$\cup \quad \cap \quad \subseteq$

Relation

$\circ \quad f^{-1}(a) \quad f(A) \& f^{-1}(B)$

Definition (Intersection, Union)

$$f_1, f_2 : A \to B$$

(i) Q : Is $f_1 \cup f_2$ a function from A to B?

(ii) Q : Is $f_1 \cap f_2$ a function from A to B?

Definition (Restriction (Problem 15.20))

$$f : A \to B, \ A_0 \subseteq A$$

$$f \mid_{A_0} : A_0 \to B, \qquad f \mid_{A_0} (a) = f(a), \forall a \in A_0$$

Definition (Composition)

$$f : A \to B \qquad g : C \to D$$

$$\text{ran}(f) \subseteq C$$

The composition function

$$g \circ f : A \to D$$

$$(g \circ f)(x) = g(f(x))$$

Non-commutative:

$$f \circ g \neq g \circ f$$

Theorem (Associative Property for Composition)

$$f : A \to B \quad g : B \to C \quad h : C \to D$$

$$h \circ (g \circ f) = (h \circ g) \circ f$$

Proof.

(i)

$$\mathsf{dom}(h \circ (g \circ f)) = \mathsf{dom}((h \circ g) \circ f)$$

(ii)

$$(h \circ (g \circ f))(x) = ((h \circ g) \circ f)(x)$$

□

Theorem (Properties of Composition (UD Theorem 15.7))

$$f : A \to B \qquad g : B \to C$$

(i) *If* f, g *are injective, then* $g \circ f$ *is injective.*

(ii) *If* f, g *are surjective, then* $g \circ f$ *is surjective.*

(iii) *If* f, g *are bijective, then* $g \circ f$ *is bijective.*

Proof for (i).

$$\forall a_1, a_2 \in A \Big((g \circ f)(a_1) = (g \circ f)(a_2) \implies a_1 = a_2 \Big)$$

□

Theorem (Properties of Composition (UD Theorem 15.8))

$$f : A \to B \qquad g : B \to C$$

(i) *If* $g \circ f$ *is injective, then* f *is injective.*

(ii) *If* $g \circ f$ *is surjective, then* g *is surjective.*

(iii) *If* $g \circ f$ *is bijective, then* f *is injective and* g *is surjective.*

Proof.

Left as Exercise (15.9).

□

Cancellation Property for Composition (Problem 15.11)

$$f : A \to B \qquad g_1, g_2 : B \to A$$

$$f \circ g_1 = f \circ g_2 \land f \text{ is bijective} \implies g_1 = g_2$$

Remark:

f is one-to-one.

Proof.

$$\forall b \in B \Big(f \circ g_1(b) = f \circ g_2(b) \implies \cdots \Big)$$

□

Definition (Inverse)

Let $f : A \to B$ be a bijective function.
The inverse of f is the function $f^{-1} : B \to A$ defined by

$$f^{-1}(b) = a \iff f(a) = b.$$

Q: Why "Bijective"?

Theorem (UD Theorem 15.4 (ii))

$$f : A \to B \textit{ is bijective} \implies f^{-1} \textit{ is bijective.}$$

Theorem (Solving Equations (UD Theorem 15.4))

$$f : A \to B \textit{ is bijective}$$

(i) $f \circ f^{-1} = i_B$

(ii) $g : B \to A \wedge f \circ g = i_B \implies g = f^{-1}$

(iii) $f^{-1} \circ f = i_A$

(iv) $g : B \to A \wedge g \circ f = i_A \implies g = f^{-1}$

Solving the equations:

$$f \circ g = i_B \qquad g \circ f = i_A$$

Bijective $\implies$ Inverse:

$$f : A \to B \text{ is bijective}$$

$$\implies$$

$$\exists g : B \to A \left(f \circ g = i_B \land g \circ f = i_A \right) \land g = f^{-1}$$

Theorem (Inverse $\implies$ Bijective (UD Theorem 15.8 (iii)))

$$\exists g : B \to A \left(g \circ f = i_A \land f \circ g = i_B \right)$$

$$\implies$$

$$f : A \to B \textit{ is bijective} \land g = f^{-1}$$

Theorem (Inverse of Composition (UD Theorem 15.6))

$$f : A \to B, g : B \to C \text{ are bijective}$$

(i) $g \circ f$ *is bijective*

(ii) $(g \circ f)^{-1} = f^{-1} \circ g^{-1}$

Proof for (ii).

$$(f^{-1} \circ g^{-1}) \circ (g \circ f) = i_A$$

$$(g \circ f) \circ (f^{-1} \circ g^{-1}) = i_B$$

□

$$f : A \to B,\ A_0 \subseteq A,\ B_0 \subseteq B$$

Definition (Image)

The image of A_0 under f is the set

$$f(A_0) = \{f(a) \mid a \in A_0\}.$$

Definition (Inverse Image)

The inverse image of B_0 under f is the set

$$f^{-1}(B_0) = \{a \in A \mid f(a) \in B_0\}.$$

Theorem (Properties of f and f^{-1} (Theorem 16.7))

$$f : A \to B,\ A_0, A_1, A_2 \subseteq A,\ B_0, B_1, B_2 \subseteq B$$

(i) f, *when applied to subsets of* A, *preserves only "*$\subseteq$*" and* $\cup$*:*

(1) $A_1 \subseteq A_2 \implies f(A_1) \subseteq f(A_2)$

(2) $f(A_1 \cup A_2) = f(A_1) \cup f(A_2)$

(3) $f(A_1 \cap A_2) \subseteq f(A_1) \cap f(A_2)$

(4) $f(A \setminus A_0) \neq B \setminus f(A_0)$

(ii) f^{-1}, *when applied to subsets of* B, *preserves* $\subseteq, \cup, \cap$, *and* $\setminus$*:*

(5) $B_1 \subseteq B_2 \implies f^{-1}(B_1) \subseteq f^{-1}(B_2)$

(6) $f^{-1}(B_1 \cup B_2) = f^{-1}(B_1) \cup f^{-1}(B_2)$

(7) $f^{-1}(B_1 \cap B_2) = f^{-1}(B_1) \cap f^{-1}(B_2)$

(8) $f^{-1}(B \setminus B_0) = A \setminus f^{-1}(B_0)$

Theorem (Properties of f and f^{-1} (Theorem 16.7))

$$f : A \to B,\ A_0 \subseteq A,\ B_0 \subseteq B$$

(iii) *f and f^{-1}:*

(9) $A_0 \subseteq f^{-1}(f(A_0))$

Q: When is $A_0 = f^{-1}(f(A_0))$?

(10) $B_0 \subseteq f(f^{-1}(B_0))$

Q: When is $B_0 = f(f^{-1}(B_0))$?

Problem 16.20

$$f : A \to B, \quad A_1, A_2 \subseteq A$$

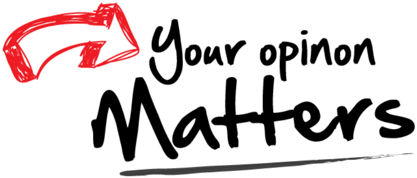
(i) When is $f(A_1) = f(A_2) \implies A_1 = A_2$?

Problem 16.21

$$f : A \to B, \quad B_1, B_2 \subseteq B$$

(i) When is $f^{-1}(B_1) = f^{-1}(B_2) \implies B_1 = B_2$?


Thank
You!



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