

1-5 Data Structures

魏恒峰

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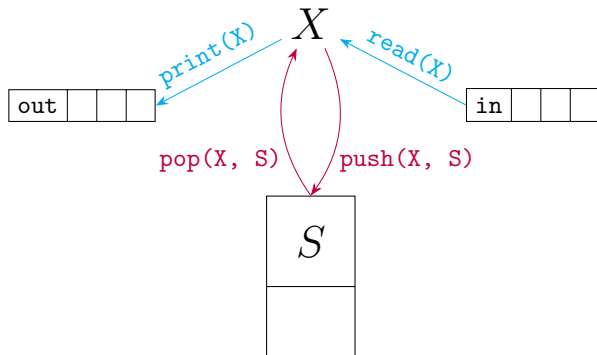


Permutations

Generating All Permutations
Stackable/Queueable Permutations

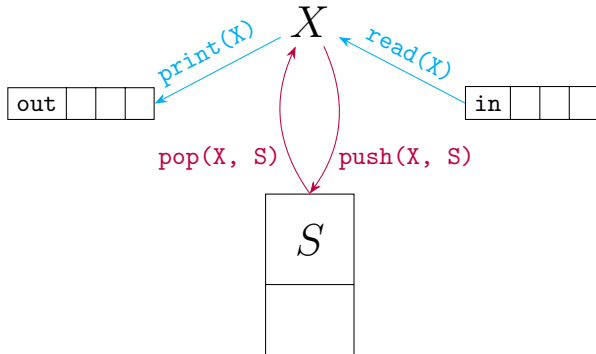
Stackable Permutations

Definition (Stackable Permutations)

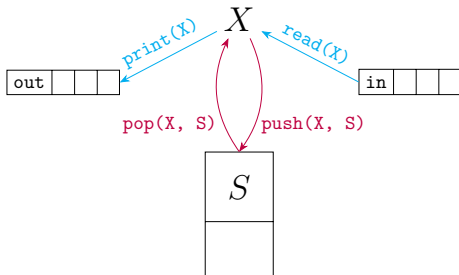


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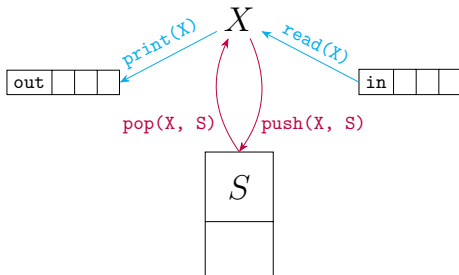
$$\text{out} = (a_1, \dots, a_n) \xleftarrow[X=0]{S=\emptyset} \text{in} = (1, \dots, n)$$



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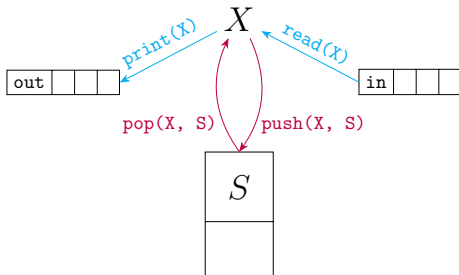


Definition (Stackable Permutations)



Q_2 : Using **only** “read, print, push, pop”?

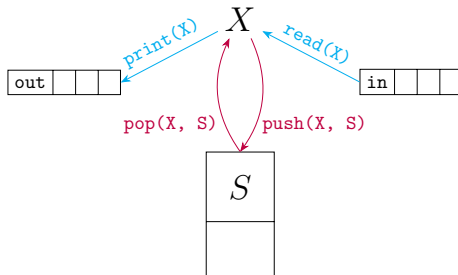
Definition (Stackable Permutations)



Q_2 : Using **only** “read, print, push, pop”?

$$a == X$$

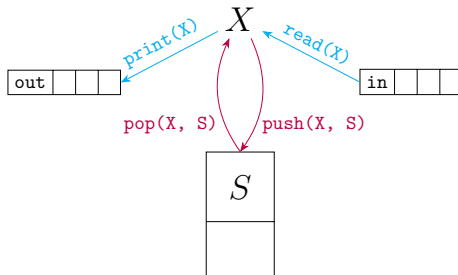
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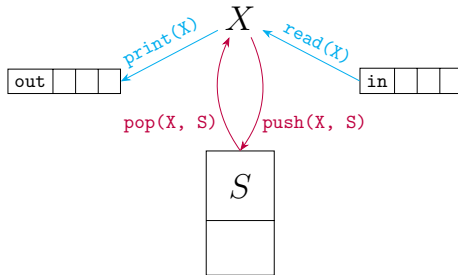
$$a == X \quad \text{top}(S)$$

Definition (Stackable Permutations)

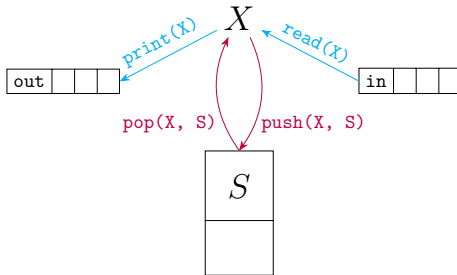


Q_2 : Using **only** “read, print, push, pop”?

$a == X$ $\text{top}(S)$ $a > X$ ($a < X$)



We can assume that X is always blank.



We can assume that X is always blank.

Proof.

What are the possible operations following $read(X)/pop(X, S)$?



DH 2.12: Stackable Permutations

(a) Show that the following permutations *are* stackable:

(i) $(3, 2, 1)$

(ii) $(3, 4, 2, 1)$

(iii) $(3, 5, 7, 6, 8, 4, 9, 2, 10, 1)$

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DH 2.13: Stackable Permutations Checking Algorithm

To check whether a given permutation can be obtained by a stack.

read print push pop is-empty

X = 0 S = \emptyset in \neq EOF

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foreach 'a' in out:
    if (! is-empty(S)
        && 'a' == top(S))
        pop(S, X)
        print(X)
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```

```
else // T.B.C
    while (in != EOF)
        read(X)
        if (X == 'a')
            print(X)
            break
        else
            push(X, S)
    if (in == EOF)
        ERR
```

DH 2.12: Stackable Permutations

(b) **Prove** that the following permutations are *not* stackable:

(i) $(3, 1, 2)$

(ii) $(4, 5, 3, 7, 2, 1, 6)$

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$$\text{out} = \cdots a_i \cdots a_j \cdots a_k : i < j < k \wedge a_j < a_k < a_i$$

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312-Pattern

Theorem (Stackable Permutations)

A permutation (a_1, \dots, a_n) is stackable \iff it is not the case that

312-Pattern : $out = \dots a_i \dots a_j \dots a_k : i < j < k \wedge a_j < a_k < a_i$

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Proof.



NO PROOF WARRANTY



DH 2.12: Stackable Permutations

(c) How many permutations of A_4 *cannot* be obtained by a stack?

$(1, 4, 2, 3), (2, 4, 1, 3), (3, 1, 2, 4), (3, 1, 4, 2), (3, 4, 1, 2)$
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DH 2.12: Stackable Permutations

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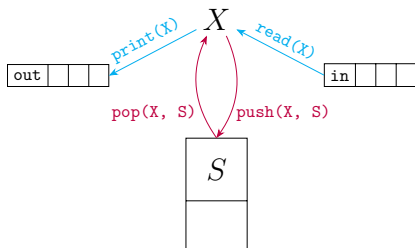
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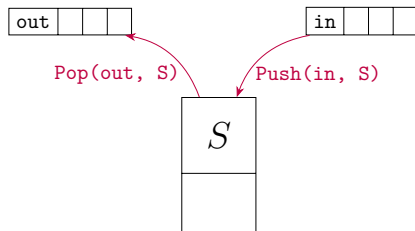
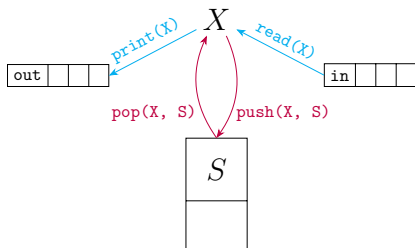
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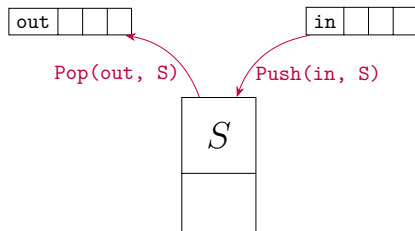
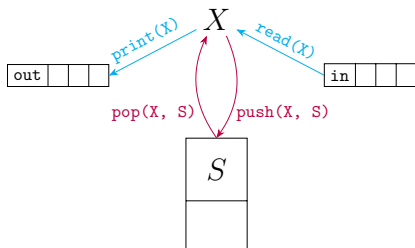
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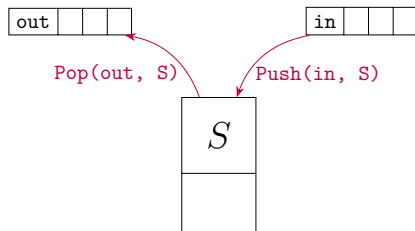
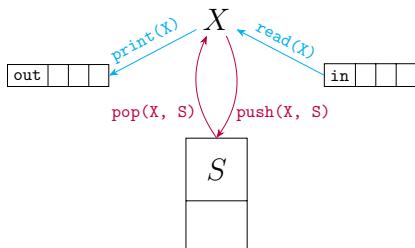
Q : What about A_n ?





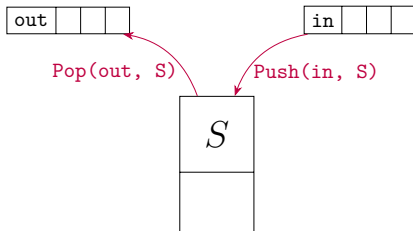
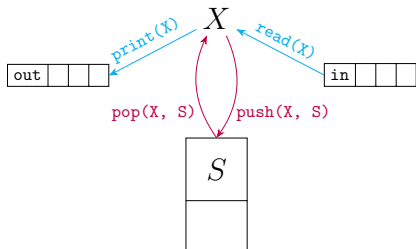


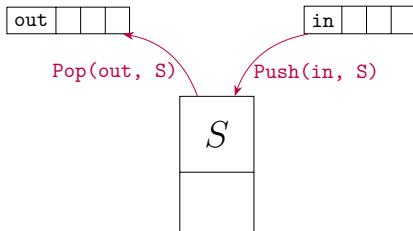
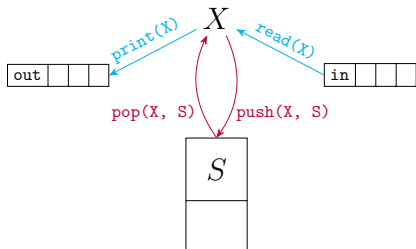
Q : Are $S + X$ and S are **equivalent**?



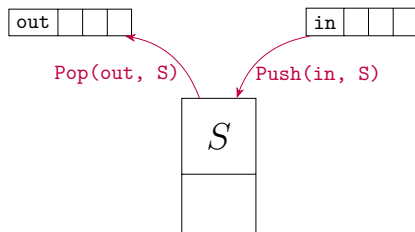
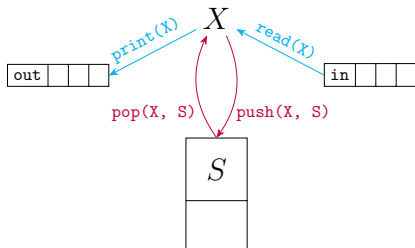
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Producing the same set of permutations.





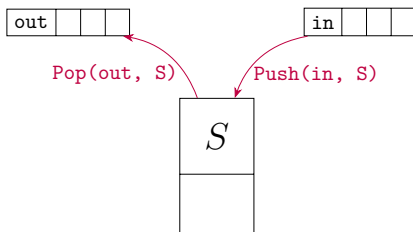
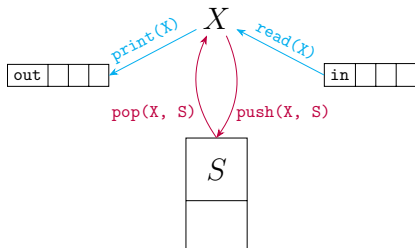
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Simulate S by $S + X$:

- Push
- Pop

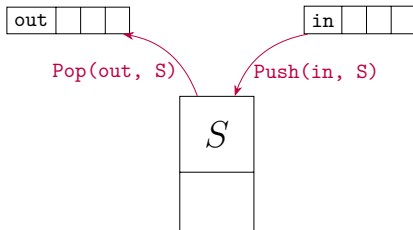
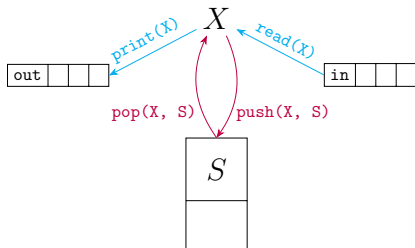


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Simulate $S + X$ by S :



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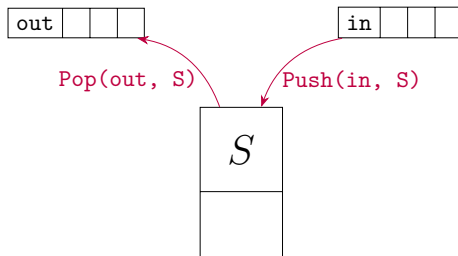
Simulate $S + X$ by S :

By iterative transformations.



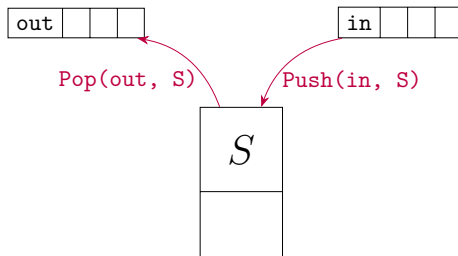
DH 2.12: Stackable Permutations

How many permutations of $\{1 \cdots n\}$ are stackable on the model S ?



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Q : How many *admissible* operation sequences of “Push” and “Pop”?

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of stackable perms = # of admissible operation sequences

Theorem

Different admissible operation sequences correspond to different permutations.

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Proof.

Push Push Push Pop Pop **Push**...

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Proof: The Reflection Method.

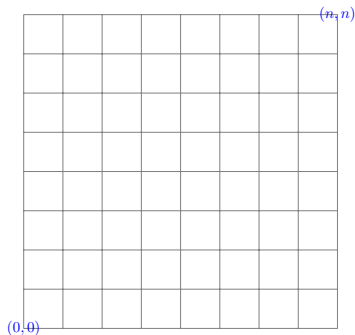
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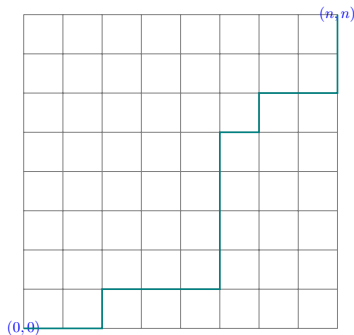


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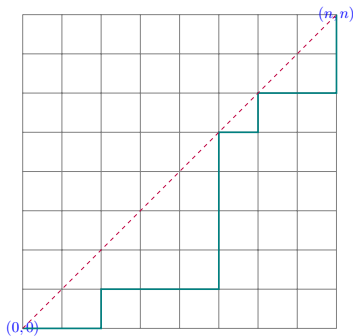


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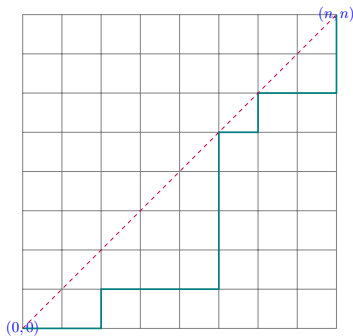


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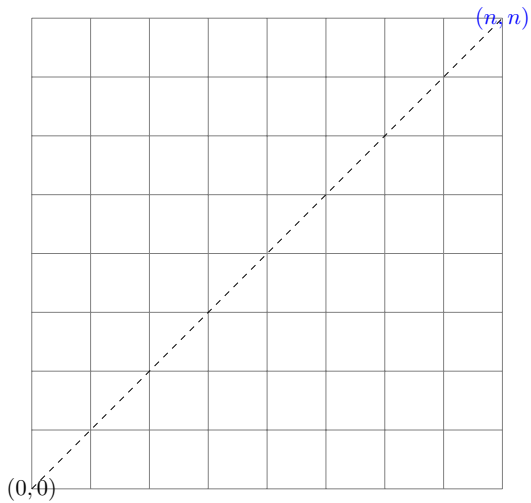
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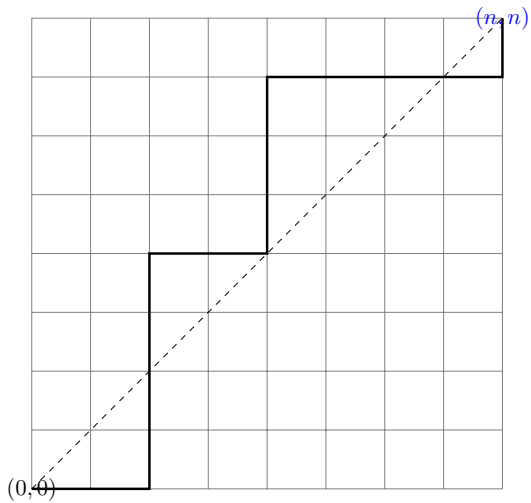
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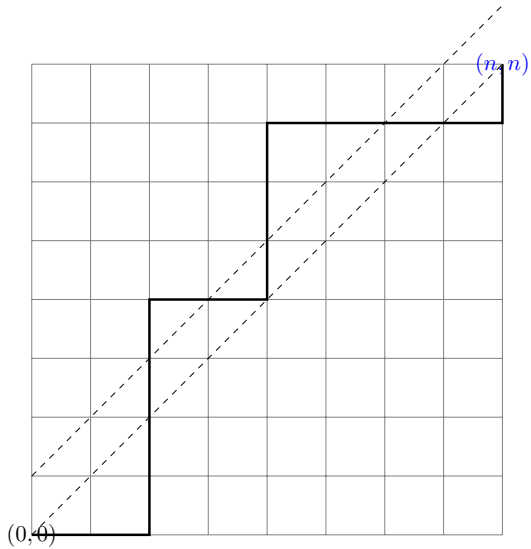


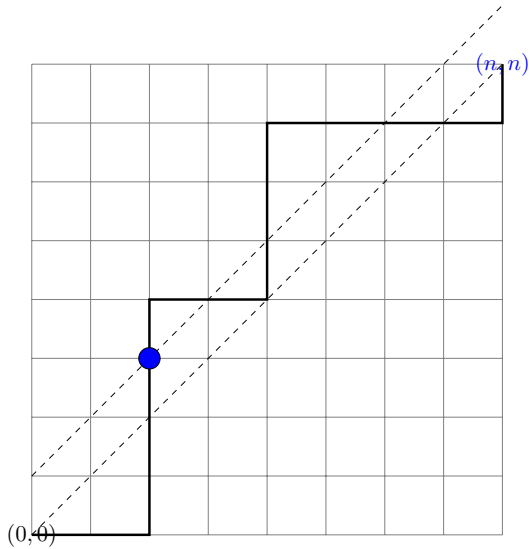
$$\underbrace{\binom{2n}{n}}_{\text{all}} - \underbrace{\binom{2n}{n-1}}_{\text{inadmissible}}$$

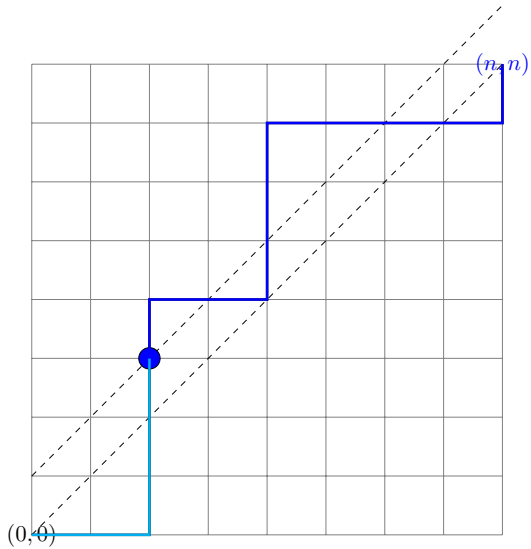


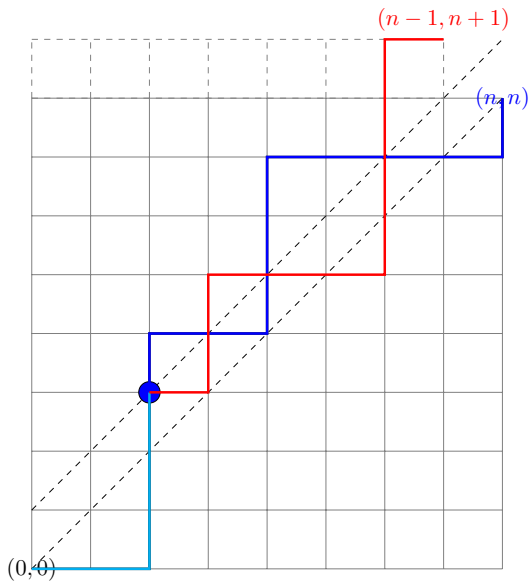


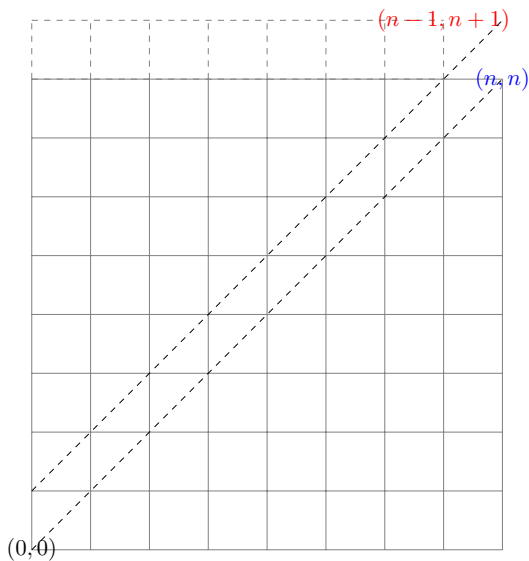


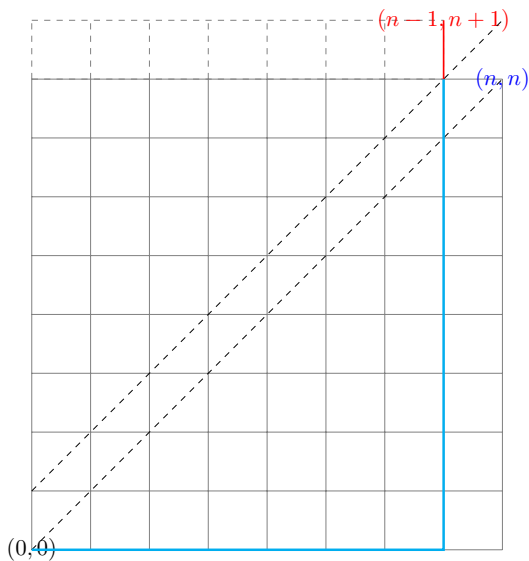


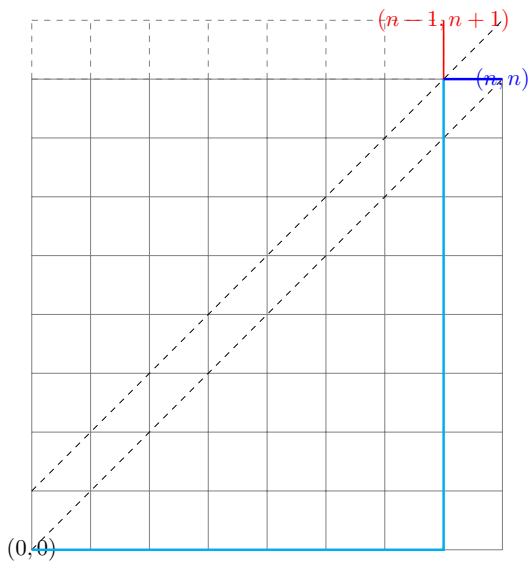


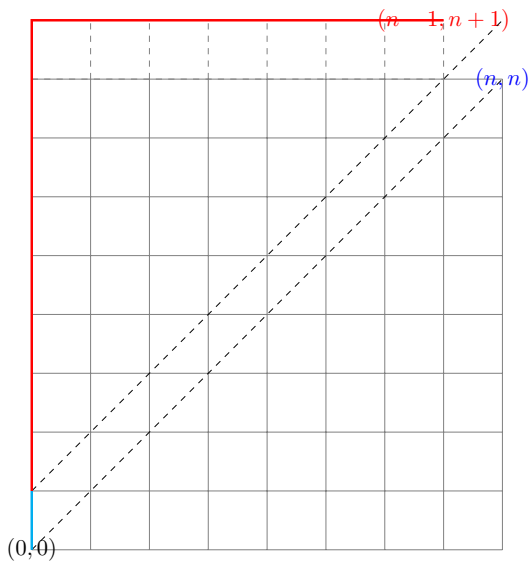


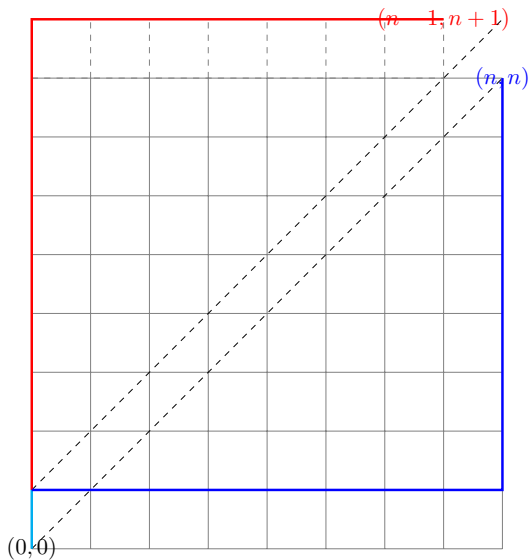








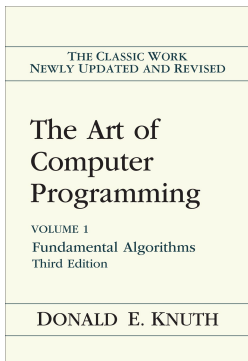




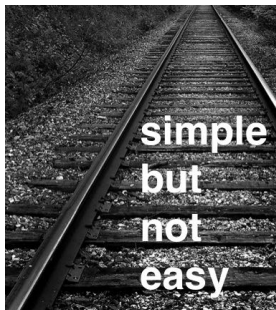
Catalan Number

$(3, 2, 1) : ((()))$ $(1, 2, 3) : ()()()$

For more about “Stackable Permutations” (Section 2.2.1)



Generating All Permutations



DH 2.11: Generate All Permutations

Design an algorithm which, given a positive integer n , generates/prints all the permutations of $[0 \cdots n)$.

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```
void perms (A[], n) {  
    if (n == 1)  
        print 'A[0] '  
    else  
        for (int i = 0; i < n; ++i)  
            print 'A[i] '  
            perms(A ← A \ A[i], n - 1)  
            print '\\n '  
}
```

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            perms(A ← A \ A[i], n - 1)
            print '\n'
}
```

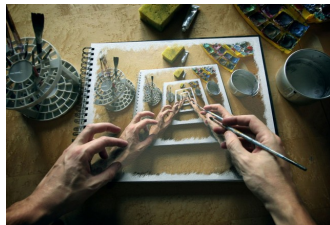
generate-perms.c



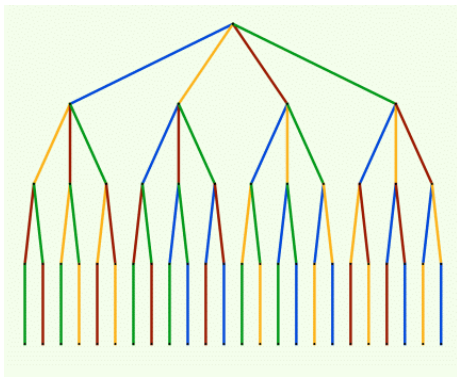
4perms.md



4perms.md



$$A = [0, 1, 2, 3] \quad n = 4$$



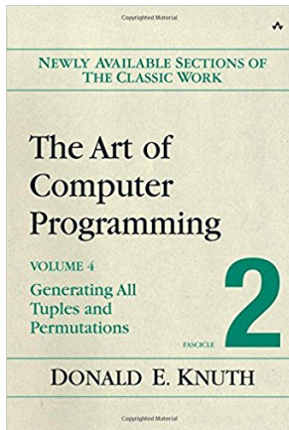
“手动单步调试”

```
void perms (prefix, A[], n) {  
    if (n == 1)  
        print ' 'prefix ++ A[0] '  
    else  
        for (int i = 0; i < n; ++i)  
            perms(prefix ← prefix ++ A[i],  
                A ← A \ A[i], n - 1)  
            print ' '\n'  
}
```

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void perms (prefix, A[], n) {  
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                A ← A \ A[i], n - 1)  
            print ' '\n'  
}
```

```
perms(' ', A, n);
```

For more about “Generating All Permutations”:



Thank
You!