1-5 Data Structures

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1/26

Permutations

Generating All Permutations Stackable/Queueable Permutations

Generating All Permutations



```
1: procedure PERMS(A[], l)

2: if l = A.size - 1 then

3: print A

4: else

5: for i \leftarrow l to A.size - 1 do

6: SWAP(A[i], A[l])

7: PERMS(A, l + 1)
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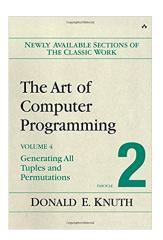


大型"车祸"现场



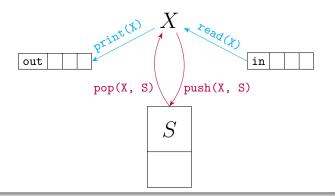
Iteration Version of PERMS

For more about "Generating All Permutations":

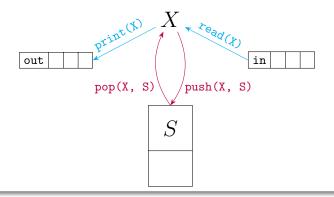


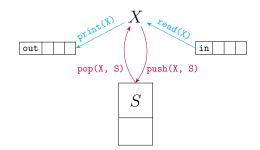


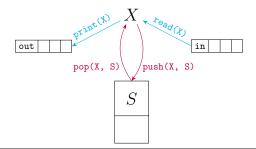
Stackable Permutations

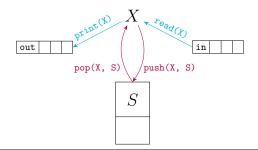


$$| \mathtt{out} = (a_1, \cdots, a_n) \underbrace{\frac{S = \emptyset}{X = 0}}_{\mathsf{in}} \mathtt{in} = (1, \cdots, n)$$

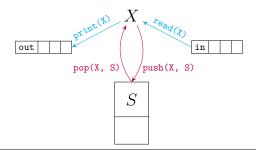




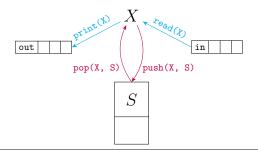




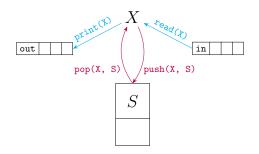
$$a == X$$



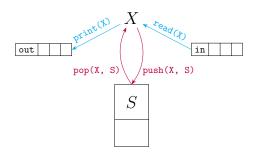
$$a == X \qquad \mathsf{top}(S)$$



$$a == X$$
 $top(S)$ $a > X (a < X)$



We can assume that X is always blank.



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Proof.

What are the possible operations following read(X)/pop(X, S)?

- (a) Show that the following permutations *are* stackable:
 - (i) (3,2,1)
 - (ii) (3,4,2,1)
 - ${\rm (iii)}\ (3,5,7,6,8,4,9,2,10,1)$

- (a) Show that the following permutations *are* stackable:
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DH 2.13: Stackable Permutations Checking Algorithm

To check whether a given permutation can be obtained by a stack.

read print push pop

X = 0 $S = \emptyset$ in != EOF

is-empty

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foreach 'a' in out:
   if (! is-empty(S)
        && 'a' == top(S))
      pop(S, X)
      print(X)
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foreach 'a' in out:
   if (! is-empty(S)
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      pop(S, X)
      print(X)
```

```
else // T.B.C
  while (in != EOF)
    read(X)
    if (X == 'a')
      print(X)
      break
    else
      push(X, S)
  if (in == EOF)
    ERR
```

- (b) **Prove** that the following permutations are not stackable:
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 - (ii) (4,5,3,7,2,1,6)

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(4, 5, 3, 7, 2, 1, 6)

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312-Pattern



Theorem (Stackable Permutations)

A permutation (a_1, \dots, a_n) is stackable \iff it is not the case that

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Proof.





(c) How many permutations of A_4 cannot be obtained by a stack?

$$(1,4,2,3), (2,4,1,3), (3,1,2,4), (3,1,4,2), (3,4,1,2)$$

 $(4,1,2,3), (4,1,3,2), (4,2,1,3), (4,2,3,1), (4,3,1,2)$

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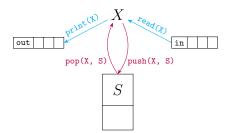
 $(4, 1, 2, 3), (4, 1, 3, 2), (4, 2, 1, 3), (4, 2, 3, 1), (4, 3, 1, 2)$

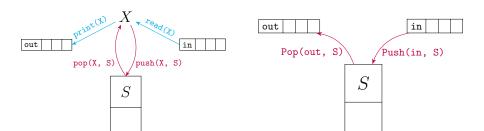
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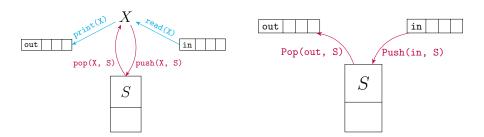
$$(1, 4, 2, 3), (2, 4, 1, 3), (3, 1, 2, 4), (3, 1, 4, 2), (3, 4, 1, 2)$$

 $(4, 1, 2, 3), (4, 1, 3, 2), (4, 2, 1, 3), (4, 2, 3, 1), (4, 3, 1, 2)$

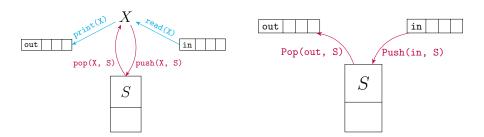
Q: What about A_n ?





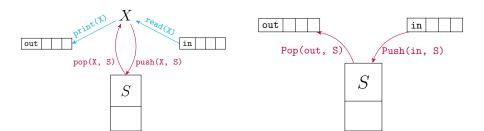


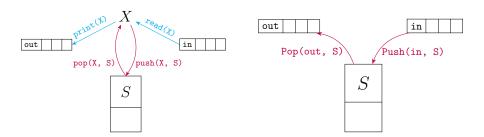
Q: Are S + X and S are equivalent?

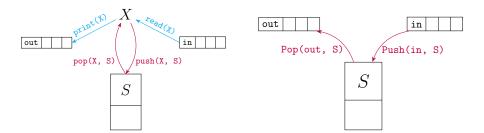


Q: Are S + X and S are equivalent?

Producing the same set of permutations.

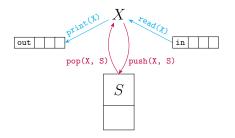


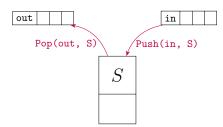




Simulate S by S + X:

- ▶ Push
- ► Pop

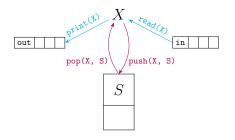


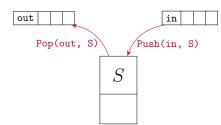


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Simulate S + X by S:





Simulate S by S + X:

- ▶ Push
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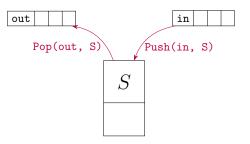
Simulate S + X by S:

By iterative transformations.



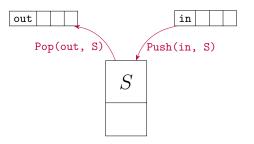
DH 2.12: Stackable Permutations

How many permutations of $\{1 \cdots n\}$ are stackable on the model S?



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Q: How many admissible operation sequences of "Push" and "Pop"?

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- (i) # of "Push" = n # of "Pop" = n
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of stackable perms = # of admissible operation sequences

Different admissible operation sequences correspond to different permutations.

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Proof.

```
Push Push Pop Pop Push · · · · Push Push Push Pop Pop Pop · · ·
```



The number of admissible operation sequences of "Push" and "Pop" is $\binom{2n}{n} - \binom{2n}{n-1}$.

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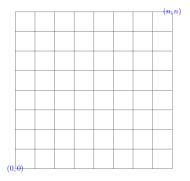
Proof: The Reflection Method.

 $\mathtt{Push}: \rightarrow \qquad \mathtt{Pop}: \uparrow$

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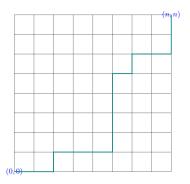




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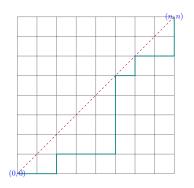




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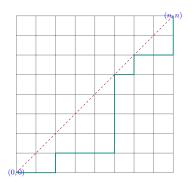
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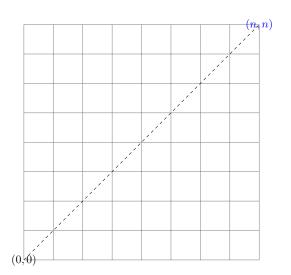
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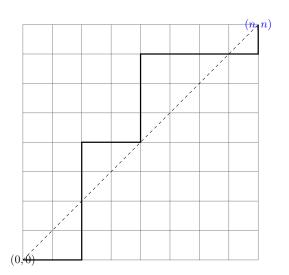
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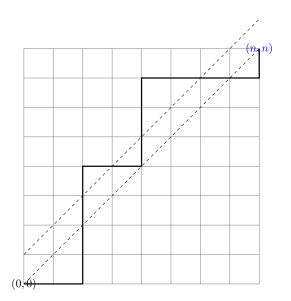
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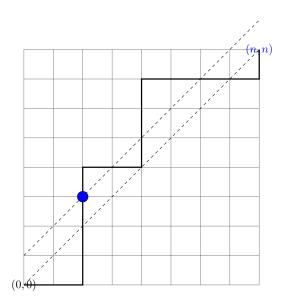


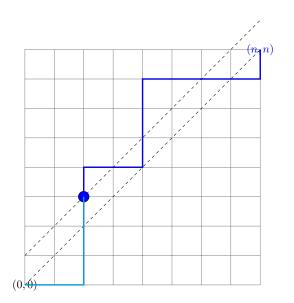
$$\underbrace{\binom{2n}{n}}_{\text{all}} - \underbrace{\binom{2n}{n-1}}_{\text{inadmissible}}$$

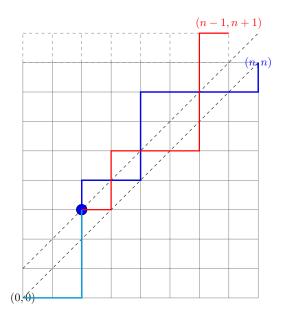


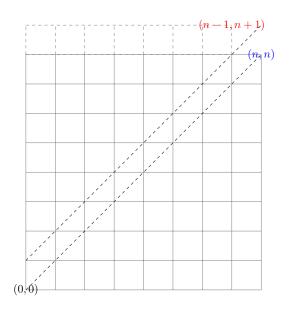


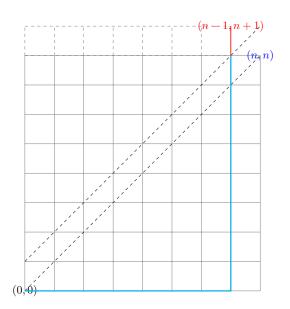


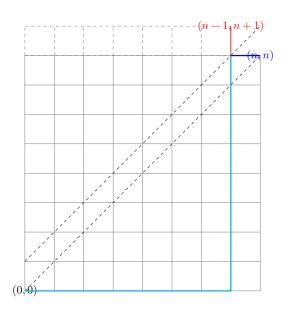


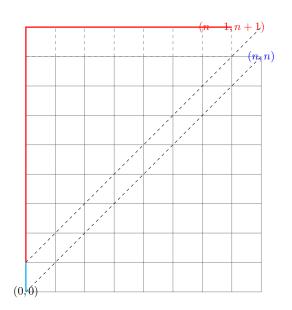


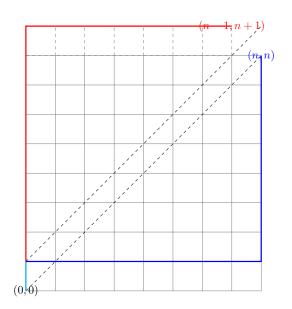












Catalan Number

$$(3,2,1):((()))$$
 $(1,2,3):()()()$

For more about "Stackable Permutations" (Section 2.2.1)

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Thank You!