1-5 Data Structures

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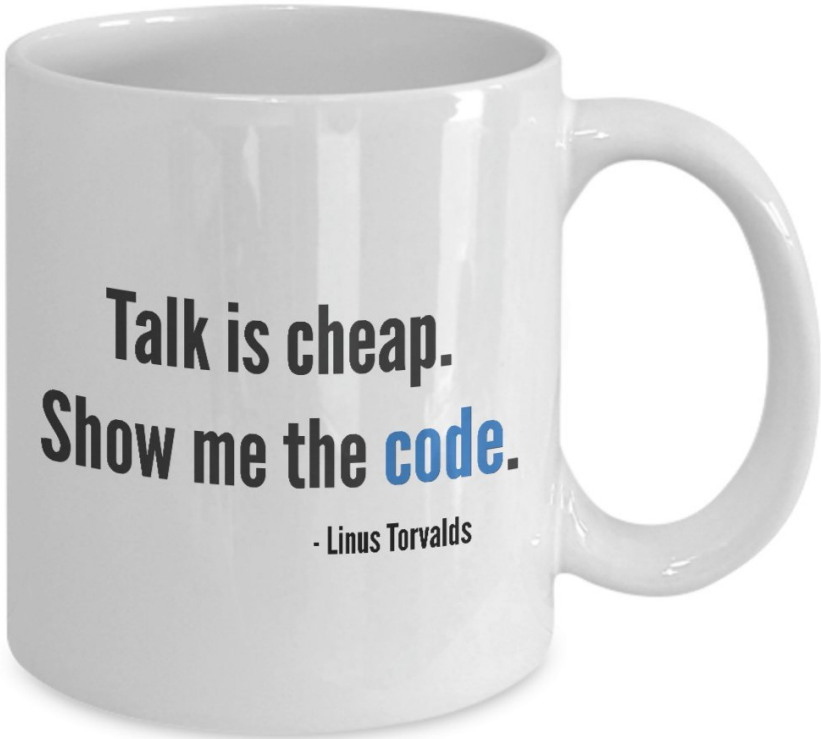
Permutations

Generating All Permutations
Stackable/Queueable Permutations

Generating All Permutations

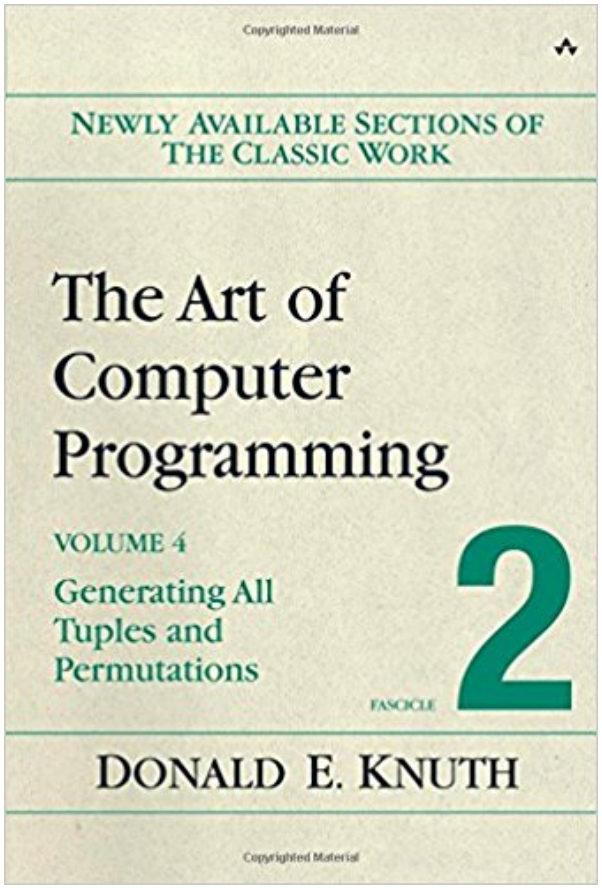
```
1: procedure PERMS(A[], l)
2:     if l = A.size − 1 then
3:         print A
4:     else
5:         for i ← l to A.size − 1 do
6:             SWAP(A[i], A[l])
7:             PERMS(A, l + 1)
8:             SWAP(A[i], A[l])
```



大型“车祸”现场

Iteration Version of Perms

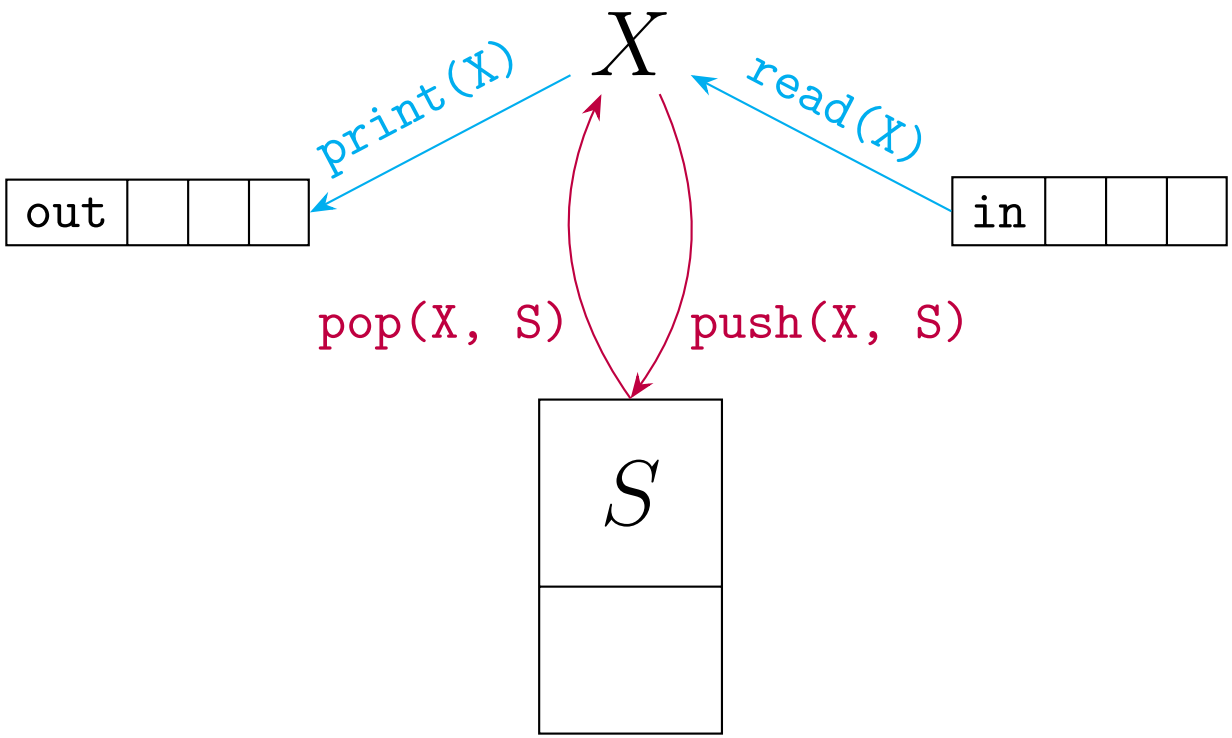
For more about "Generating All Permutations":



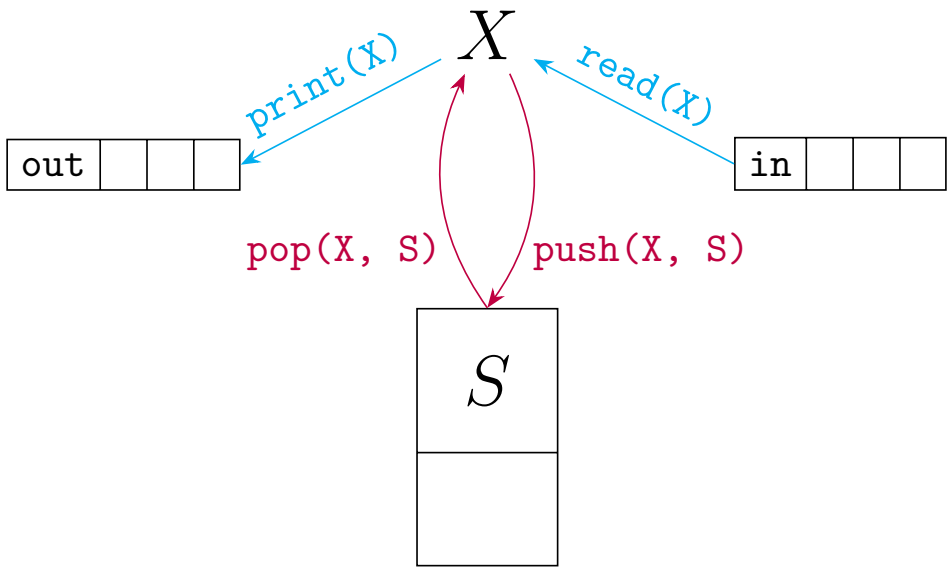
Stackable Permutations

Definition (Stackable Permutations)

$$\texttt{out} = (a_1, \cdots, a_n) \xleftarrow[X = 0]{S = \emptyset} \texttt{in} = (1, \cdots, n)$$

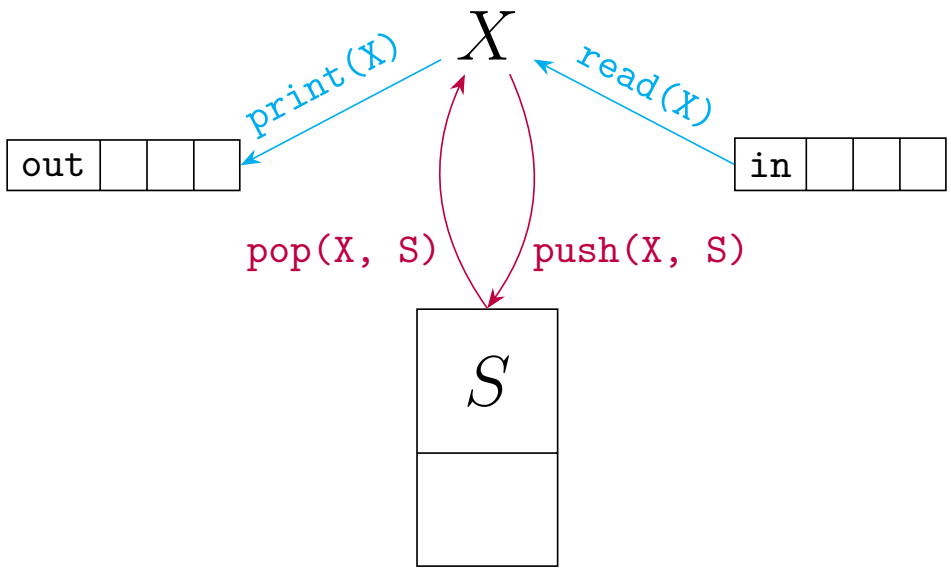


Definition (Stackable Permutations)



Q_2 : Using only "`read, print, push, pop`"?

$a == X$ $\quad$ $\texttt{top}(S)$ $\quad$ $a > X$ $(a < X)$



We can assume that X is always blank.

Proof.

What are the possible operations following `read(X)`/`pop(X, S)`?

DH 2.12: Stackable Permutations

(a) Show that the following permutations *are* stackable:

(i) $(3, 2, 1)$

(ii) $(3, 4, 2, 1)$

(iii) $(3, 5, 7, 6, 8, 4, 9, 2, 10, 1)$

DH 2.13: Stackable Permutations Checking Algorithm

To check whether a given permutation can be obtained by a stack.

`read` `print` `push` `pop` `is-empty`

```
X = 0    S = ∅    in != EOF
```

```
foreach 'a' in out:
  if (! is-empty(S)
       && 'a' == top(S))
    pop(S, X)
    print(X)
```

```
else // T.B.C
  while (in != EOF)
    read(X)
    if (X == 'a')
      print(X)
      break
    else
      push(X, S)
  if (in == EOF)
    ERR
```

DH 2.12: Stackable Permutations

(b) **Prove** that the following permutations are *not* stackable:

(i) $(3, 1, 2)$

(ii) $(4, 5, 3, 7, 2, 1, 6)$

$$(3, 1, 2)$$

$$(4, 5, 3, 7, 2, 1, 6)$$

$$\boxed{\texttt{out} = \cdots a_i \cdots a_j \cdots a_k : i < j < k \wedge a_j < a_k < a_i}$$

312-Pattern

Theorem (Stackable Permutations)

A permutation $(a_1, \cdots, a_n)$ *is stackable* $\iff$ *it is not the case that*

$$312\text{-}Pattern: \boxed{out = \cdots a_i \cdots a_j \cdots a_k : i < j < k \wedge a_j < a_k < a_i}$$

Proof.

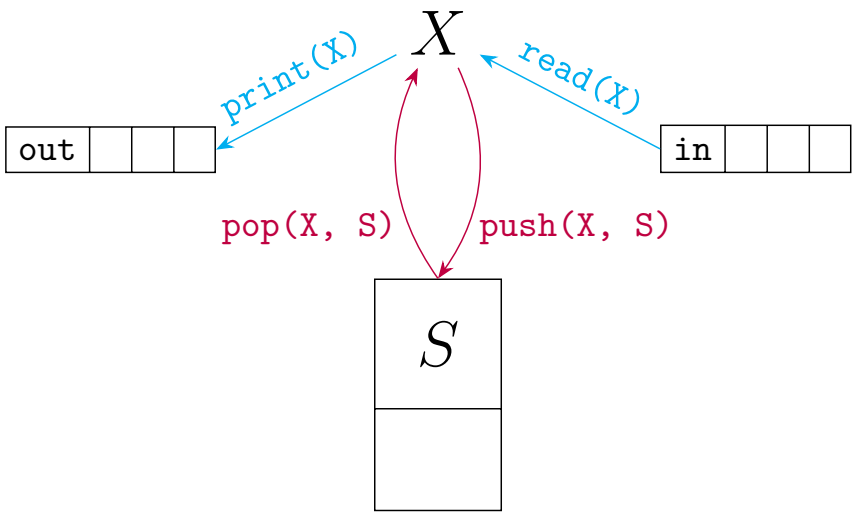
□

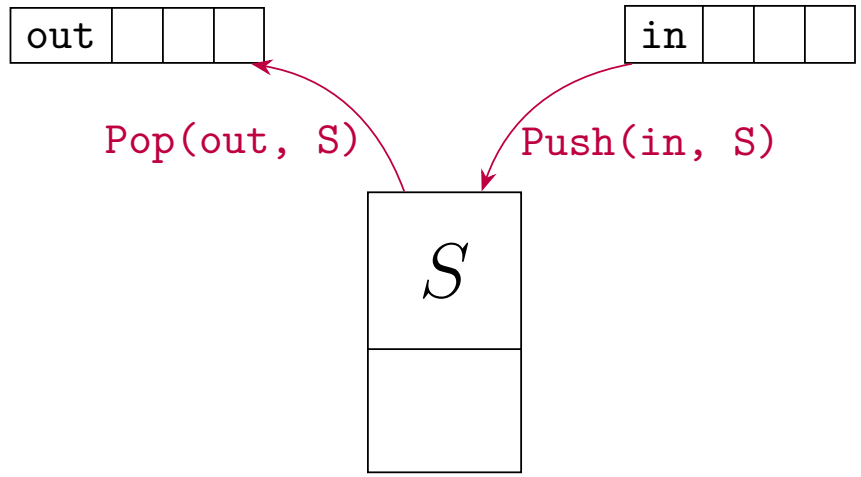
DH 2.12: Stackable Permutations

(c) How many permutations of A_4 *cannot* be obtained by a stack?

$$(1,4,2,3),(2,4,1,3),(3,1,2,4),(3,1,4,2),(3,4,1,2)$$
$$(4,1,2,3),(4,1,3,2),(4,2,1,3),(4,2,3,1),(4,3,1,2)$$

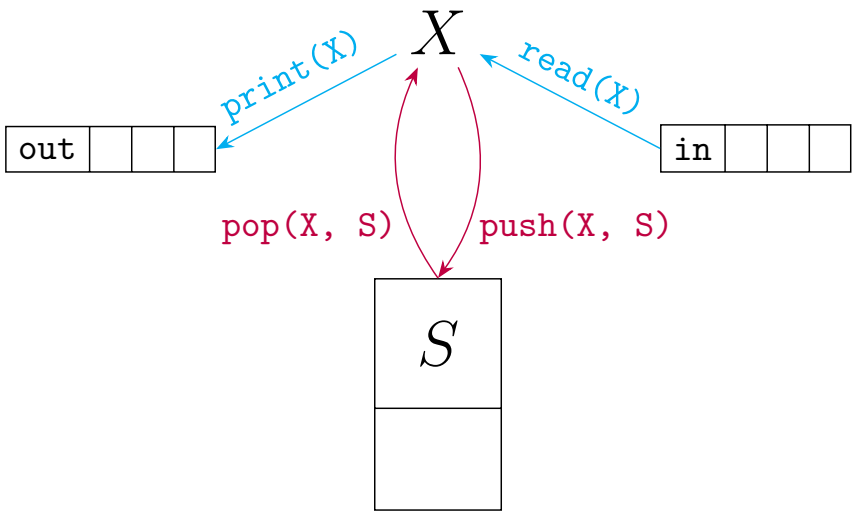
Q : What about A_n?

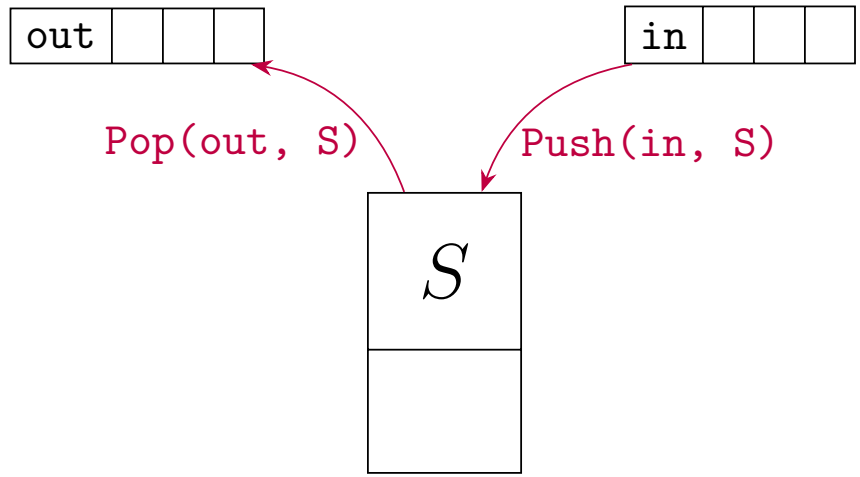




Q : Are $S + X$ and S are **equivalent**?

Producing the same set of permutations.





By simulations.

Simulate S by $S + X$:

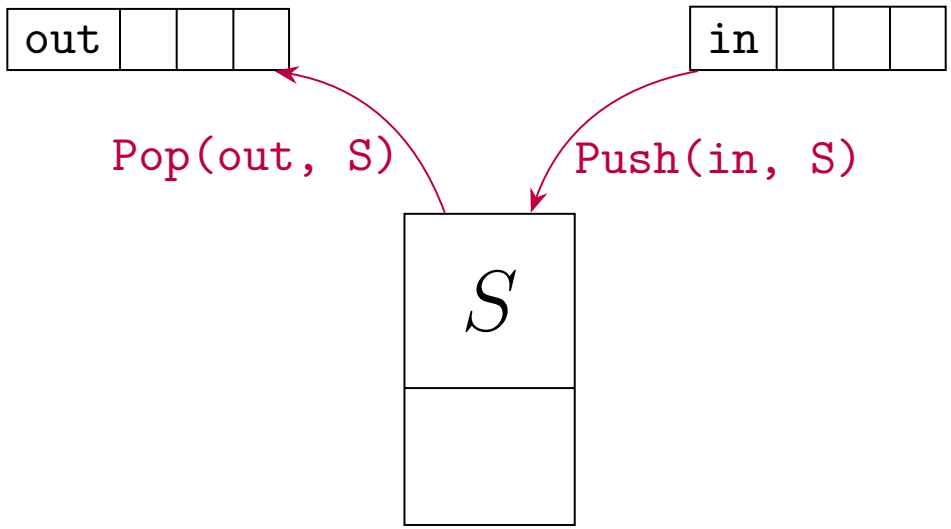
- `Push`
- `Pop`

Simulate $S + X$ by S:

By iterative transformations.

DH 2.12: Stackable Permutations

How many permutations of $\{1 \cdots n\}$ are stackable on the model S?



Q : How many *admissible* operation sequences of "`Push`" and "`Pop`"?

Definition (Admissible Operation Sequences)

An operation sequence of "`Push`" and "`Pop`" is *admissible* if and only if

(i) # of "`Push`" $= n$ # of "`Pop`" $= n$

(ii) $\forall$ prefix : (# of "`Pop`") $\leq$ (# of "`Push`")

# of stackable perms = # of admissible operation sequences

Theorem

Different admissible operation sequences correspond to different permutations.

Proof.

```
Push Push Push Pop Pop Push···
Push Push Push Pop Pop Pop···
```

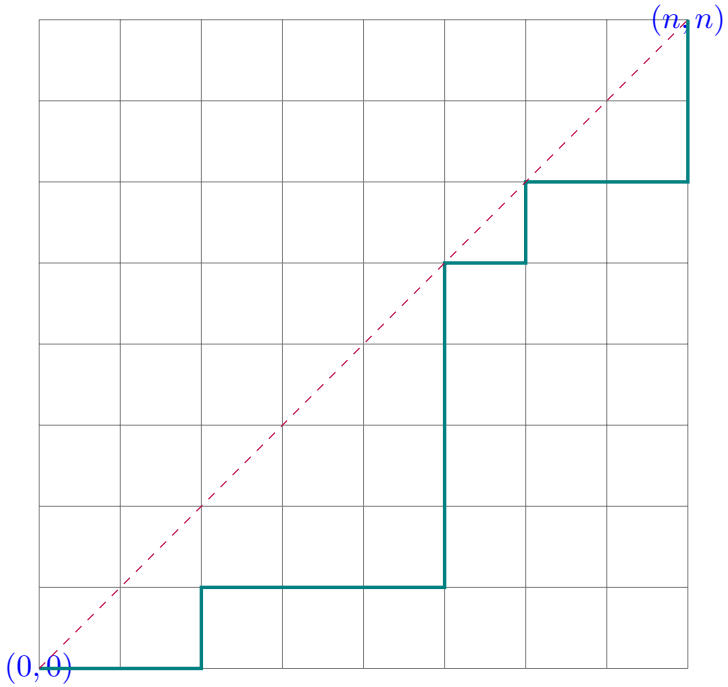
□

Theorem

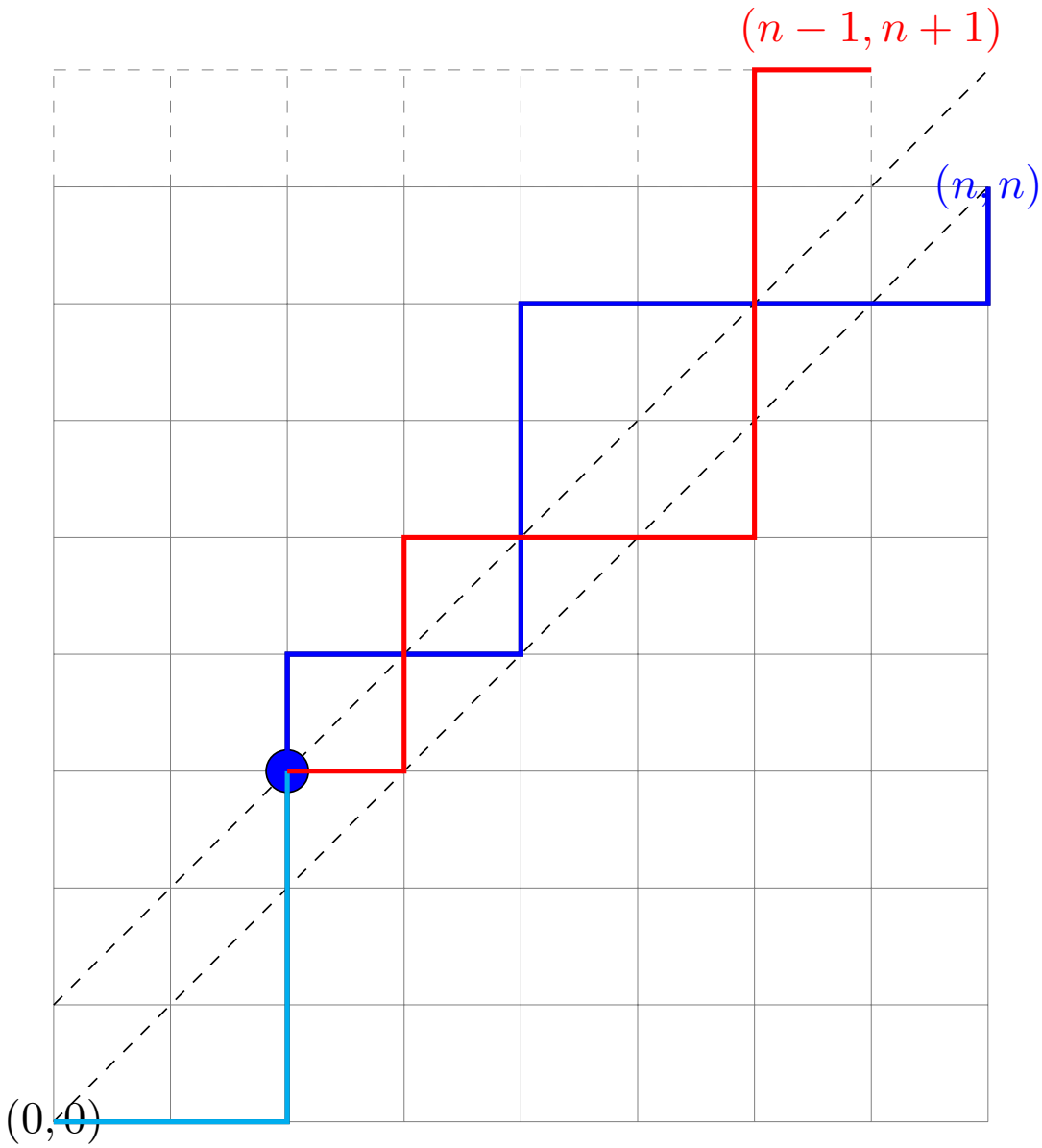
The number of admissible operation sequences of "Push" and "Pop" is $\binom{2n}{n} - \binom{2n}{n-1}$.

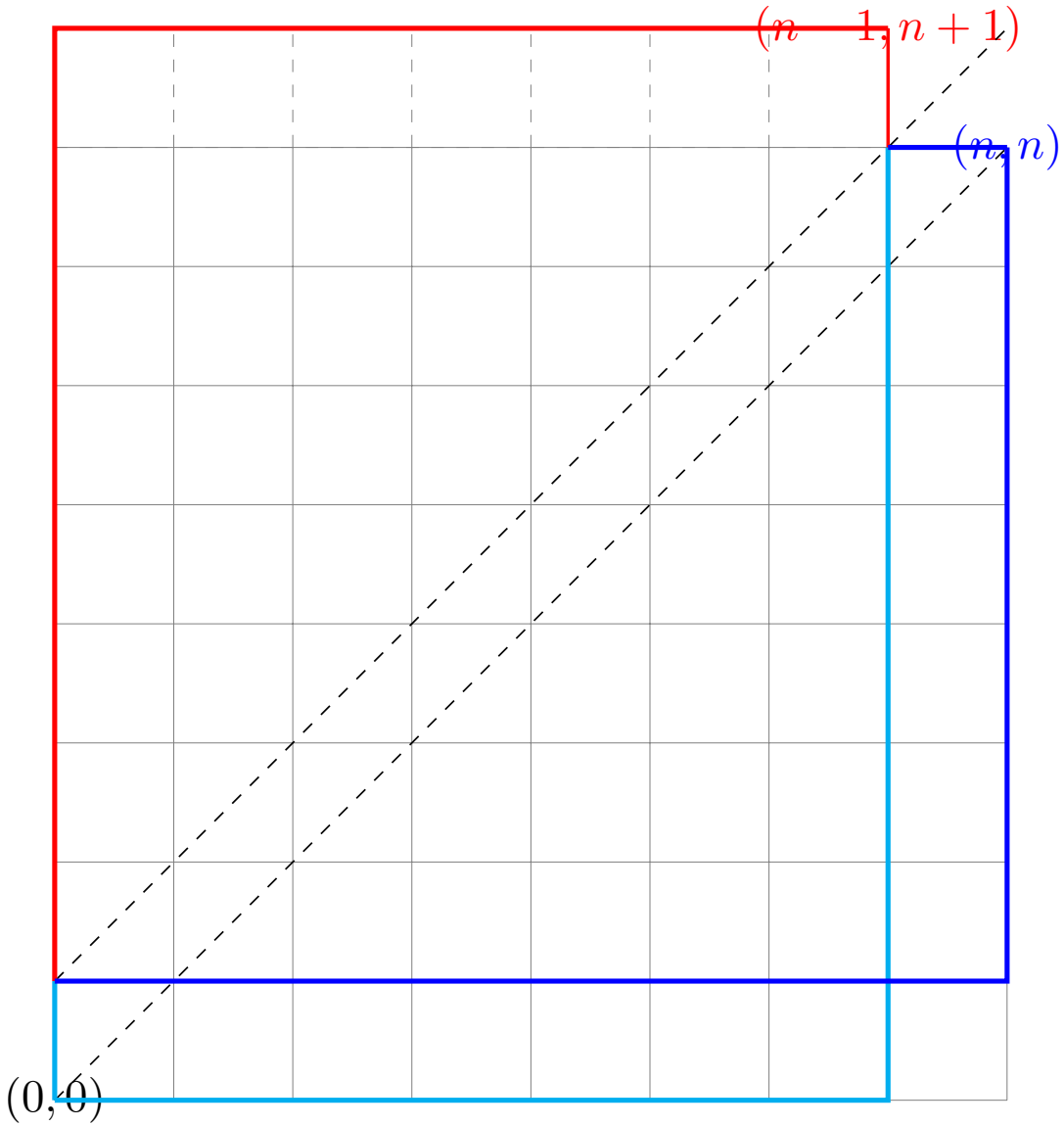
Proof: The Reflection Method.

`Push` : $\rightarrow$ `Pop` : $\uparrow$



$$\underbrace{\binom{2n}{n}}_{\text{all}} - \underbrace{\binom{2n}{n-1}}_{\text{inadmissible}}$$

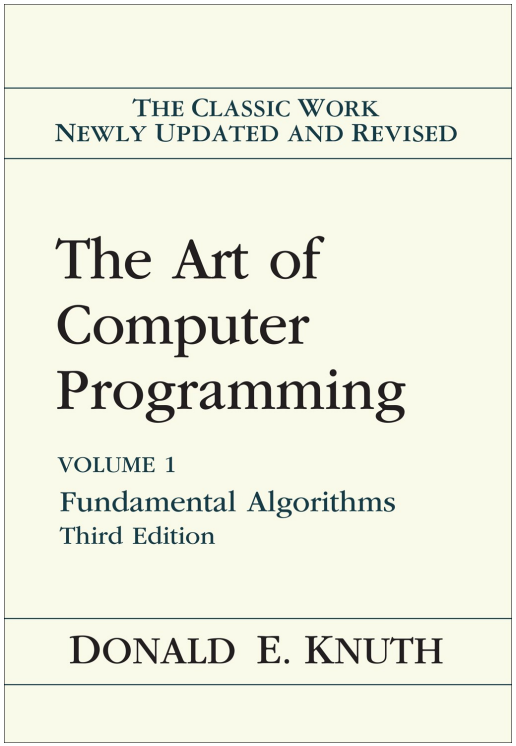

$(n-1, n+1)$
(n, n)
$(0, 0)$


$(n-1, n+1)$
(n, n)
$(0, 0)$

Catalan Number

$(3, 2, 1) : (((())))$ $\quad$ $(1, 2, 3) : ()()()$

For more about "Stackable Permutations" (Section 2.2.1)




Thank
You!