

# 1-9 关系及其基本性质 (II)

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# The Relational Data Model

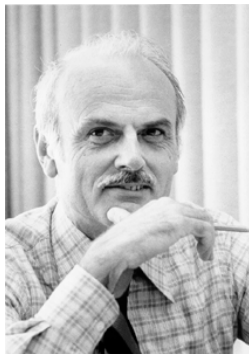
## A Relational Model of Data for Large Shared Data Banks

E. F. Codd

*IBM Research Laboratory, San Jose, California*

Future users of large data banks must be protected from having to know how the data is organized in the machine (the internal representation). A prompting service which supplies such information is not a satisfactory solution. Activities of users at terminals and most application programs should remain unaffected when the internal representation of data is changed and even when some aspects of the external representation are changed. Changes in data representation will often be needed as a result of changes in query, update, and report traffic and natural growth in the types of stored information.

Existing noninferential, formatted data systems provide users with tree-structured files or slightly more general network models of the data. In Section 1, inadequacies of these models are discussed. A model based on  $n$ -ary relations, a normal form for data base relations, and the concept of a universal data sublanguage are introduced. In Section 2, certain operations on relations (other than logical inference) are discussed and applied to the problems of redundancy and consistency in the user's model.



Codd@CACM'1970

Edgar F. Codd (1923 – 2003)

# The Relational Data Model

## — 如何靠“关系”赢得图灵奖？

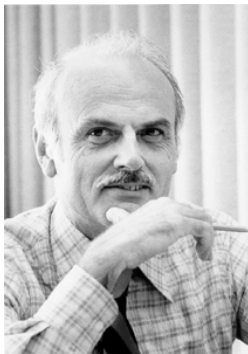
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1

2

3



# 1

## A Relational View of Data

$$(a, b) = \{\{a\}, \{a, b\}\}$$

$$R \subseteq A \times B$$

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$$(a, b, c) = ((a, b), c)$$



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$$R \subseteq A \times B$$

$$(a, b, c) = ((a, b), c)$$

$$(x_1, x_2, \dots, x_n) = ((x_1, x_2, \dots, x_{n-1}), x_n)$$

$$(a, b) = \{\{a\}, \{a, b\}\}$$

$$R \subseteq A \times B$$

$$(a, b, c) = ((a, b), c)$$

$$(x_1, x_2, \dots, x_n) = ((x_1, x_2, \dots, x_{n-1}), x_n)$$

$$R \subseteq X_1 \times X_2 \times \dots \times X_n$$

Course = {CS101, EE200, PH100}  
StudentId = {12345, 67890, 22222, 33333}  
Grade = {A, B, C, D, A<sup>-</sup>, B<sup>+</sup>, C<sup>+</sup>}

$\text{Course} = \{\text{CS101}, \text{EE200}, \text{PH100}\}$   
 $\text{StudentId} = \{12345, 67890, 22222, 33333\}$   
 $\text{Grade} = \{A, B, C, D, A^-, B^+, C^+\}$

$\text{CSG} \subseteq \text{Course} \times \text{StudentId} \times \text{Grade}$

$\text{Course} = \{\text{CS101}, \text{EE200}, \text{PH100}\}$   
 $\text{StudentId} = \{12345, 67890, 22222, 33333\}$   
 $\text{Grade} = \{A, B, C, D, A^-, B^+, C^+\}$

$\text{CSG} \subseteq \text{Course} \times \text{StudentId} \times \text{Grade}$

$\text{CSG} = \{(\text{CS101}, 12345, A),$   
 $(\text{CS101}, 67890, B),$   
 $(\text{CS101}, 33333, C),$   
 $(\text{EE200}, 12345, B^+),$   
 $(\text{EE200}, 22222, A^-),$   
 $(\text{PH100}, 67890, C^+)\}$

Course	StudentId	Grade
CS101	12345	A
CS101	67890	B
EE200	12345	C
EE200	22222	B+
CS101	33333	A-
PH100	67890	C+

Course	StudentId	Grade
CS101	12345	A
CS101	67890	B
EE200	12345	C
EE200	22222	B+
CS101	33333	A-
PH100	67890	C+

Row : tuple

Column : attribute/component

Course	StudentId	Grade
CS101	12345	A
CS101	67890	B
EE200	12345	C
EE200	22222	B+
CS101	33333	A-
PH100	67890	C+

StudentId	Name	Address	Phone
12345	C. Brown	12 Apple St.	555-1234
67890	L. Van Pelt	34 Pear Ave.	555-5678
22222	P. Patty	56 Grape Blvd.	555-9999

Course	Day	Hour
CS101	M	9AM
CS101	W	9AM
CS101	F	9AM
EE200	Tu	10AM
EE200	W	1PM
EE200	Th	10AM

Course	Room
CS101	Turing Aud.
EE200	25 Ohm Hall
PH100	Newton Lab.

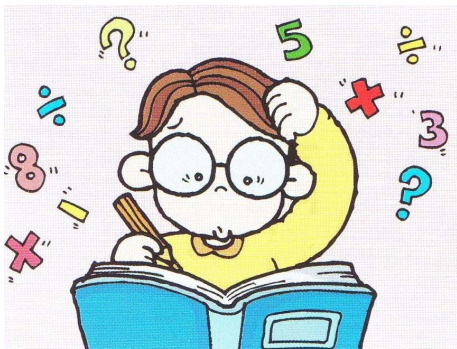


Too simple?

# 2

## Relational Algebra

# Relational Algebra



# Relational Algebra

# Relational Algebra

A set of relations

Operations on relations

Laws of operations

$\cup \quad \cap \quad \setminus$

## Selection ( $\sigma_C(R)$ )

$$\sigma_{\text{Course} = \text{"CS101"}}(\text{CSG})$$

Course	StudentId	Grade
CS101	12345	A
CS101	67890	B
EE200	12345	C
EE200	22222	B+
CS101	33333	A-
PH100	67890	C+

Course	StudentId	Grade
CS101	12345	A
CS101	67890	B
CS101	33333	A-

## Projection ( $\pi_L(R)$ )

$$\pi_{\text{StudentId}}(\sigma_{\text{Course} = \text{"CS101"}}(\text{CSG}))$$

Course	StudentId	Grade
CS101	12345	A
CS101	67890	B
EE200	12345	C
EE200	22222	B+
CS101	33333	A-
PH100	67890	C+

Course	StudentId	Grade
CS101	12345	A
CS101	67890	B
CS101	33333	A-

StudentId
12345
67890
33333



$$\text{Join } (R \bowtie_{R.A_i=S.B_j} S)$$

$$R \subseteq A_1 \times \cdots \times A_n$$

$$S \subseteq B_1 \times \cdots \times B_m$$

$$R \bowtie_{R.A_i=S.B_j} S \subseteq A_1 \times \cdots \times A_n \times B_1 \times \cdots \times B_{j-1} \times B_{j+1} \times \cdots \times B_m$$

$$\text{Join } (R \bowtie_{R.A_i=S.B_j} S)$$

$$R \subseteq A_1 \times \cdots \times A_n$$

$$S \subseteq B_1 \times \cdots \times B_m$$

$$R \bowtie_{R.A_i=S.B_j} S \subseteq A_1 \times \cdots \times A_n \times B_1 \times \cdots \times B_{j-1} \times B_{j+1} \times \cdots \times B_m$$

$$\text{Natural Join } (R \bowtie S)$$

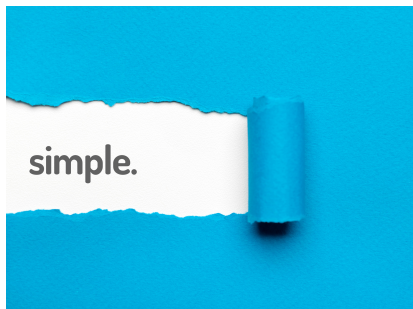
$$\text{CR} \quad \bowtie \quad \text{CDH} \quad (\text{CR} \bowtie \text{CDH})$$

$\text{CR.Course} = \text{CDH.Course}$

Course	Room
CS101	Turing Aud.
EE200	25 Ohm Hall
PH100	Newton Lab.

Course	Day	Hour
CS101	M	9AM
CS101	W	9AM
CS101	F	9AM
EE200	Tu	10AM
EE200	W	1PM
EE200	Th	10AM

Course	Room	Day	Hour
CS101	Turing Aud.	M	9AM
CS101	Turing Aud.	W	9AM
CS101	Turing Aud.	F	9AM
EE200	25 Ohm Hall	Tu	10AM
EE200	25 Ohm Hall	W	1PM
EE200	25 Ohm Hall	Th	10AM



Course	StudentId	Grade
CS101	12345	A
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EE200	12345	C
EE200	22222	B+
CS101	33333	A-
PH100	67890	C+

StudentId	Name	Address	Phone
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Course	Room
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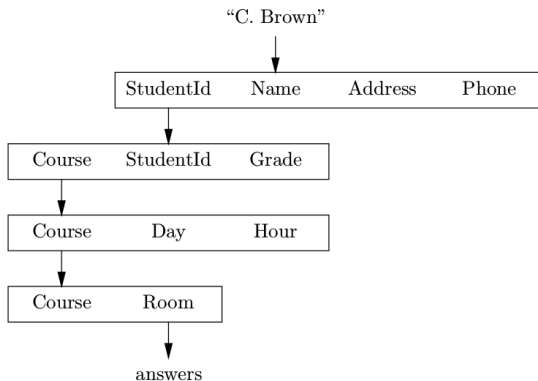
Course	StudentId	Grade
CS101	12345	A
CS101	67890	B
EE200	12345	C
EE200	22222	B+
CS101	33333	A-
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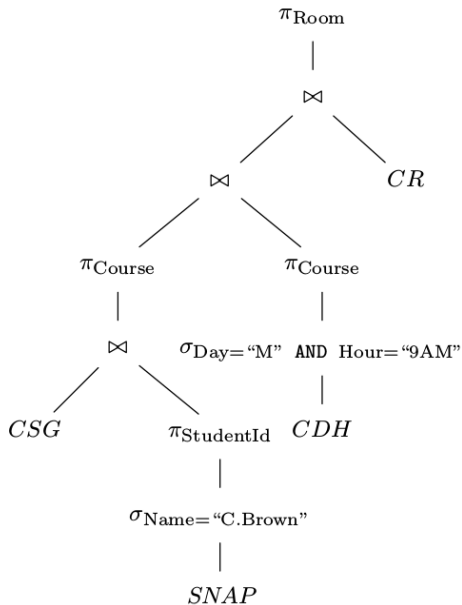
StudentId	Name	Address	Phone
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Course	Room
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*Q* : Where is “C. Brown” 9 AM on Mondays?







**TOO  
CLEVER IS  
DUMB.**

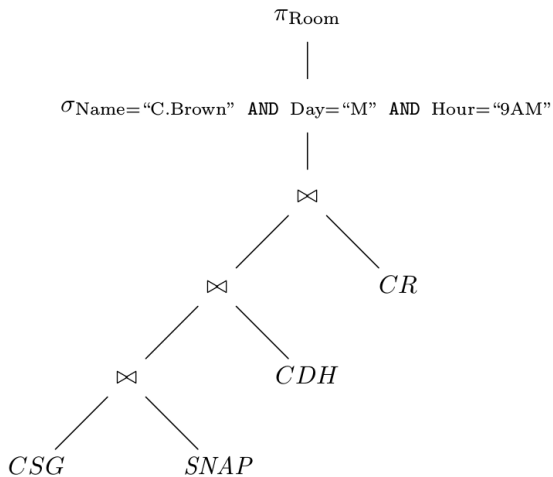


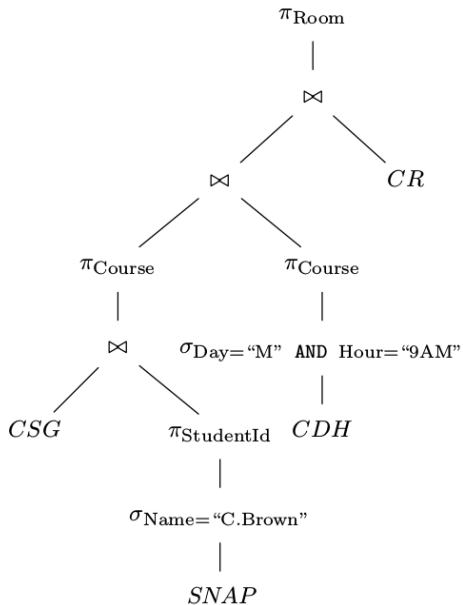
**Ogden Nash**  
American Poet

1902 - 1971



## Algebraic Laws for Relations





*“In fact, the **improvement in efficiency** made possible by **transforming expressions of relational algebra** is arguably the most striking example of **the power of algebra** that we find in computer science.”*

# Laws Involving Selection ( $\sigma_C(R)$ )

Selection Splitting

$$\sigma_{C \wedge C'}(R) \equiv \sigma_C(\sigma_{C'}(R))$$

Commutativity

$$\sigma_C(\sigma_{C'}(R)) \equiv \sigma_{C'}(\sigma_C(R))$$

# Laws Involving Selection ( $\sigma_C(R)$ )

## Selection Splitting

$$\sigma_{C \wedge C'}(R) \equiv \sigma_C(\sigma_{C'}(R))$$

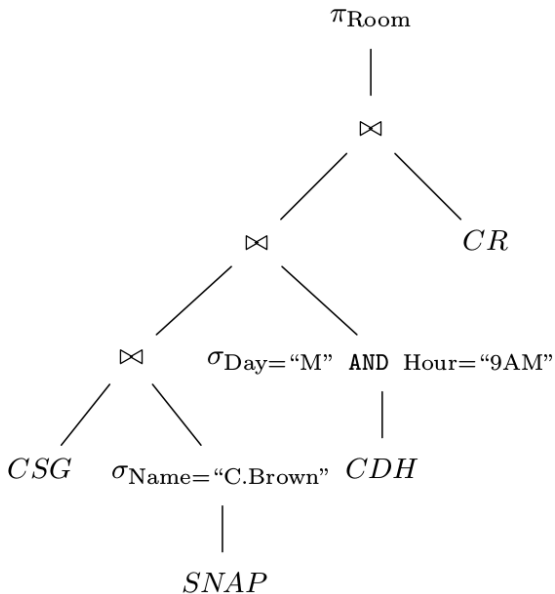
## Commutativity

$$\sigma_C(\sigma_{C'}(R)) \equiv \sigma_{C'}(\sigma_C(R))$$

## Selection Pushing

$$(\sigma_C(R \bowtie S)) \equiv (\sigma_C(R) \bowtie S)$$

$$(\sigma_C(R \bowtie S)) \equiv (R \bowtie \sigma_C(S))$$



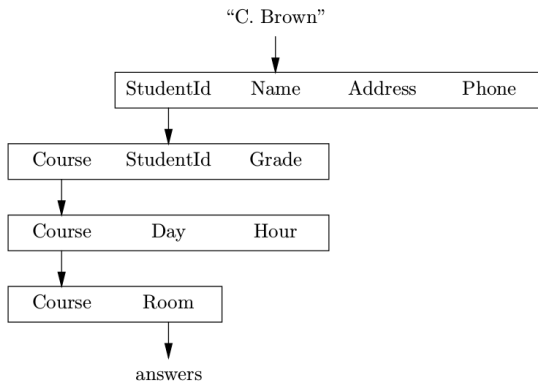


## Laws Involving Projection ( $\pi_L(R)$ )

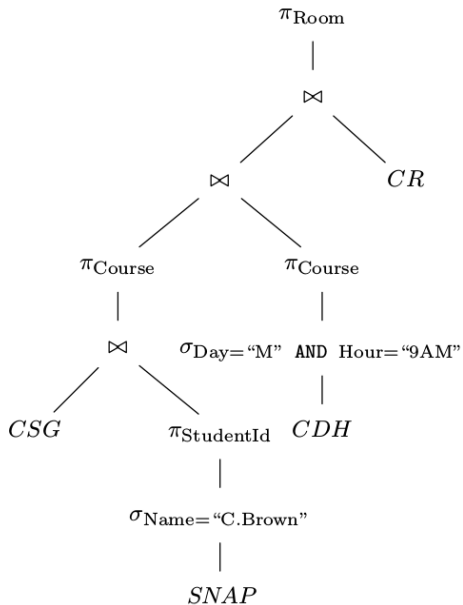
$$\left( \pi_L(R \bowtie_{A_i=B_j} S) \right) \equiv \left( \pi_L(\pi_M(R) \bowtie_{A_i=B_j} \pi_N(S)) \right)$$

$$M = (L \cup \{A_i\}) \cap \{A_i\}_i$$

$$N = (L \cup \{B_j\}) \cap \{B_j\}_j$$

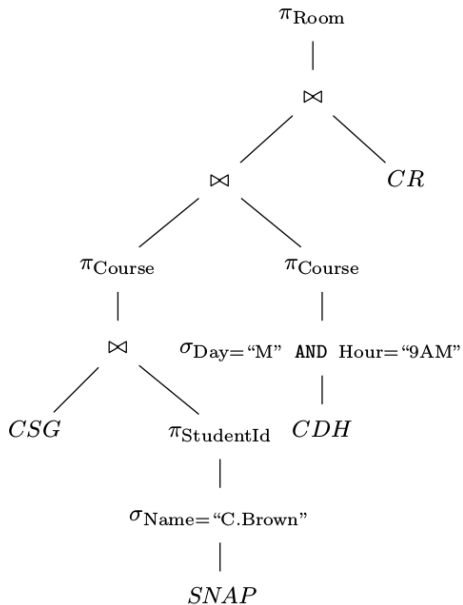


$$\begin{aligned}
 & \pi_{\text{Course}} \left( \text{CSG} \bowtie_{\text{StudentId}} \text{SNAP} \right) \\
 \equiv & \pi_{\text{Course}} \left( \pi_{\text{Course, StudentId}}(\text{CSG}) \bowtie_{\text{StudentId}} \pi_{\text{StudentId}}(\text{SNAP}) \right)
 \end{aligned}$$

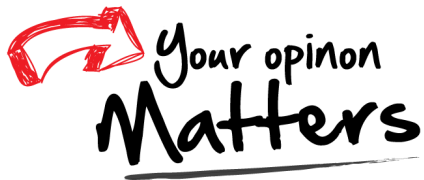


# 3

## Algebra as Language



Thank  
You!



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