

2-10 Elementary Data Structures

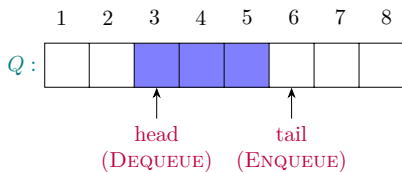
Hengfeng Wei

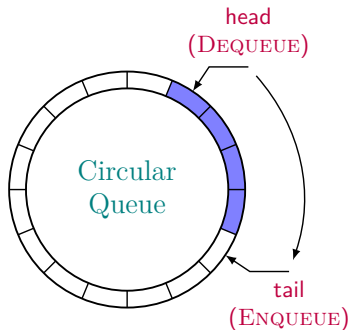
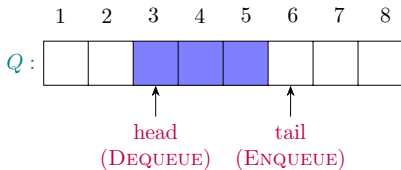
hfwei@nju.edu.cn

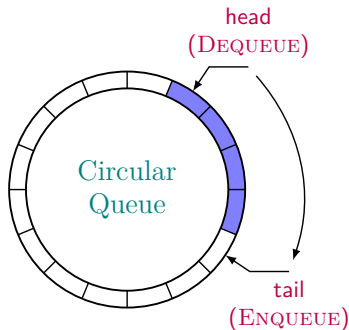
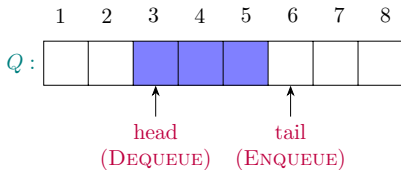
May 30, 2018



SO
EASY

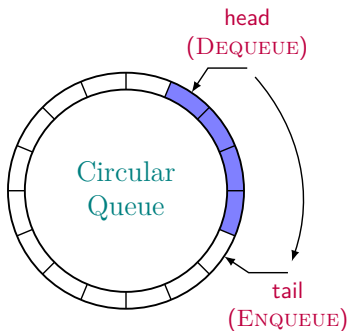




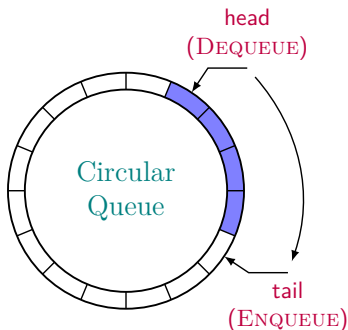


head & teal: following the same direction

Underflow and Overflow of a Circular Queue (Problem 10.1-4)



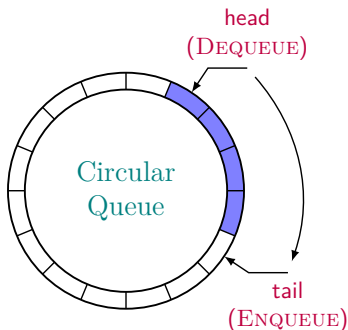
Underflow and Overflow of a Circular Queue (Problem 10.1-4)



```
procedure DEQUEUE( $Q$ )  
  if  $Q.head = Q.tail$  then  
    return "UNDERFLOW"
```

...

Underflow and Overflow of a Circular Queue (Problem 10.1-4)

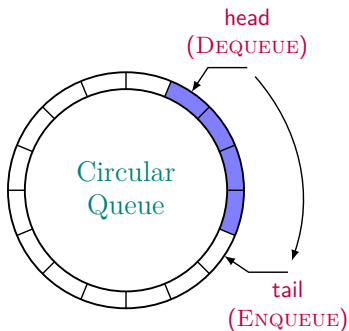


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procedure DEQUEUE(Q)
  if  $Q.head = Q.tail$  then
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```

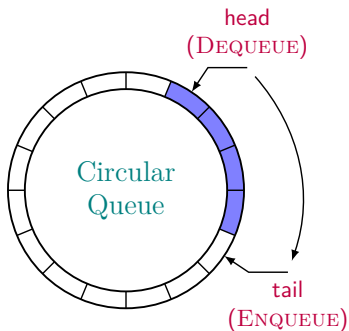
...

```
procedure ENQUEUE(Q, x)
  if  $Q.head = Q.tail + 1$  then
    return "OVERFLOW"
```

...



反馈: tail 为什么指向最后一个元素的后面? 这个太难受了。



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QUEUE-EMPTY

$[l, r)$ $(l, r]$ $[l, r]$ (l, r)

$[l, r)$ $(l, r]$ $[l, r]$ (l, r)

EWD831-0

Why numbering should start at zero EWD831.html

To denote the subsequence of natural numbers $2, 3, \dots, 12$ without the pernicious three dots, four conventions are open to us:

- a) $2 \leq i < 13$
- b) $1 < i \leq 12$
- c) $2 \leq i \leq 12$
- d) $1 < i < 13$



Why Numbering Should Start at Zero

A Queue, Two Stacks (Problem 10.1-6)

Show how to implement a queue using two stacks.

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Show how to implement a queue using two stacks.

procedure ENQUEUE(x)

Push(S_1, x)

procedure DEQUEUE()

if $S_2 = \emptyset$ **then**

while $S_1 \neq \emptyset$ **do**

Push($S_2, \text{Pop}(S_1)$)

Pop(S_2)

A Queue, Two Stacks (Problem 10.1-6)

Show how to implement a queue using two stacks.

```
procedure ENQUEUE( $x$ )  
    Push( $S_1, x$ )
```

Correctness?

```
procedure DEQUEUE()  
    if  $S_2 = \emptyset$  then  
        while  $S_1 \neq \emptyset$  do  
            Push( $S_2, Pop(S_1)$ )  
    Pop( $S_2$ )
```

A Queue, Two Stacks (Problem 10.1-6)

Show how to implement a queue using two stacks.

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procedure ENQUEUE( $x$ )  
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     $Pop(S_2)$ 
```

$$\begin{aligned} &ENQ(x, t_1), ENQ(y, t_2) \wedge t_1 < t_2 \\ &\implies \\ &DEQ(x, t_3), DEQ(y, t_4) \wedge t_3 < t_4 \end{aligned}$$

A Queue, Two Stacks (Problem 10.1-6)

Show how to implement a queue using two stacks.

Analyze the running time of the queue operations.

procedure ENQUEUE(x)

Push(S_1, x)

Correctness?

procedure DEQUEUE()

if $S_2 = \emptyset$ **then**

while $S_1 \neq \emptyset$ **do**

Push($S_2, \text{Pop}(S_1)$)

Pop(S_2)

$\text{ENQ}(x, t_1), \text{ENQ}(y, t_2) \wedge t_1 < t_2$

\implies

$\text{DEQ}(x, t_3), \text{DEQ}(y, t_4) \wedge t_3 < t_4$

<i>item</i>	Push into S_1	Pop from S_1	Push into S_2	Pop from S_2
x	1	1	1	1

<i>item</i>	Push into S_1	Pop from S_1	Push into S_2	Pop from S_2
x	1	1	1	1

$$\hat{c}_{\text{ENQ}} = 4$$

$$\hat{c}_{\text{DEQ}} = 0$$

<i>item</i>	Push into S_1	Pop from S_1	Push into S_2	Pop from S_2
x	1	1	1	1

$$\hat{c}_{\text{ENQ}} = 4$$

$$\hat{c}_{\text{DEQ}} = 0$$

$$\hat{c}_{\text{ENQ}} = 3$$

$$\hat{c}_{\text{DEQ}} = 1$$

COMPACTIFY-LIST (Problem 10.3 – 5)

COMPACTIFY-LIST(L, F)

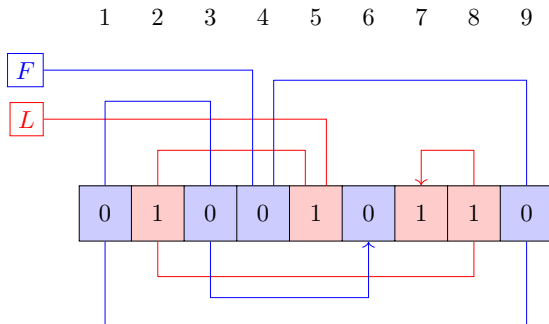
L : doubly linked list, $|L| = n$

F : doubly linked free list, $|F| = m - n$

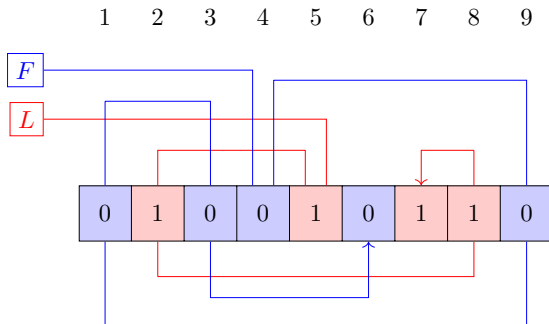
$\Theta(n)$

1	2	3	4	5	6	7	8	9
1	1	1	1	0	0	0	0	0

1	2	3	4	5	6	7	8	9
1	1	1	1	0	0	0	0	0

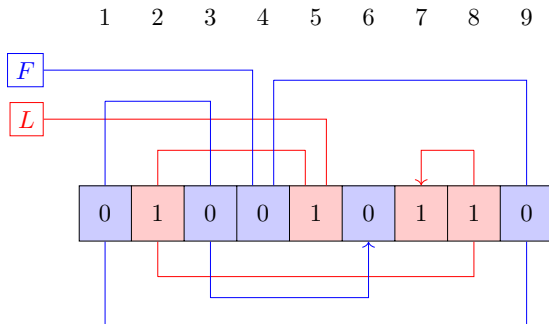


1	2	3	4	5	6	7	8	9
1	1	1	1	0	0	0	0	0



Swap (0, 1) pairs

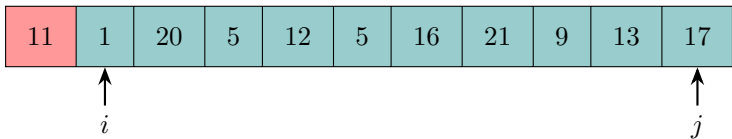
1	2	3	4	5	6	7	8	9
1	1	1	1	0	0	0	0	0



Swap (0, 1) pairs

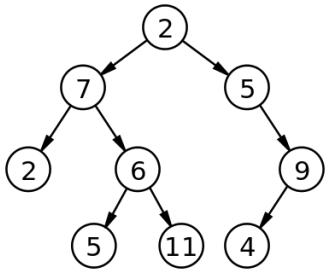
$$\Theta(n)$$

HOARE-PARTITION



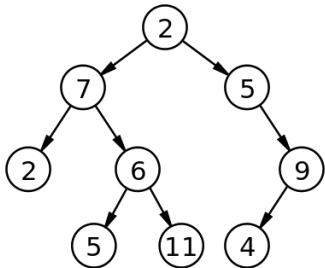
Recursive Binary Tree Traversal (Problem 10.4 – 2)

$O(n)$



Recursive Binary Tree Traversal (Problem 10.4 – 2)

$O(n)$



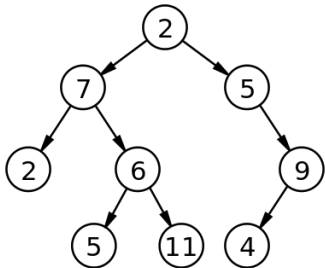
```
procedure RECURSIVE-DFS(t)  
  print t.key
```

```
  if t.left  $\neq$  NIL then  
    RECURSIVE-DFS(t.left)  
  if t.right  $\neq$  NIL then  
    RECURSIVE-DFS(t.right)
```

```
RECURSIVE-DFS(T.root)
```

Recursive Binary Tree Traversal (Problem 10.4 – 2)

$O(n)$



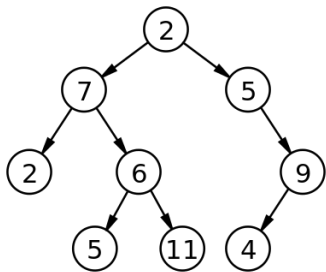
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procedure RECURSIVE-DFS(t)  
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```
  if t.left  $\neq$  NIL then  
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  if t.right  $\neq$  NIL then  
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```

```
RECURSIVE-DFS(T.root)
```

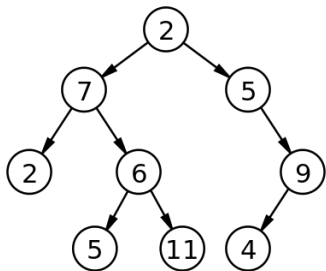
Non-recursive Binary Tree Traversal (Problem 10.4 – 2)

$O(n)$



Non-recursive Binary Tree Traversal (Problem 10.4 – 2)

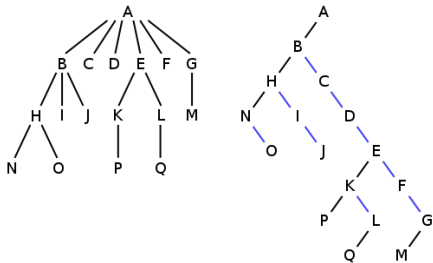
$O(n)$



procedure ITERATIVE-DFS(t) $S.PUSH(t)$ $\triangleright S : \text{stack}$ **while** $S \neq \emptyset$ **do** $v \leftarrow S.POP()$ **print** $v.key$ **if** $v.right \neq \text{NIL}$ **then** $S.PUSH(v.right)$ **if** $v.left \neq \text{NIL}$ **then** $S.PUSH(v.left)$ $\text{ITERATIVE-DFS}(T.root)$

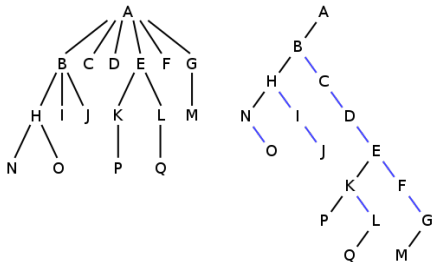
“LCRS” Tree Traversal (Problem 10.4 – 2)

$$O(n)$$



“LCRS” Tree Traversal (Problem 10.4 – 2)

$O(n)$



```
procedure RECURSIVE-DFS(t)  
  print t.key
```

```
  if t.lc  $\neq$  NIL then  
    RECURSIVE-DFS(t.lc)
```

```
  if t.rs  $\neq$  NIL then  
    RECURSIVE-DFS(t.rs)
```

```
RECURSIVE-DFS(T.root)
```

Thank
You!



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