

1-5 数据与数据结构 (II)

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温故而知新 — 孔子

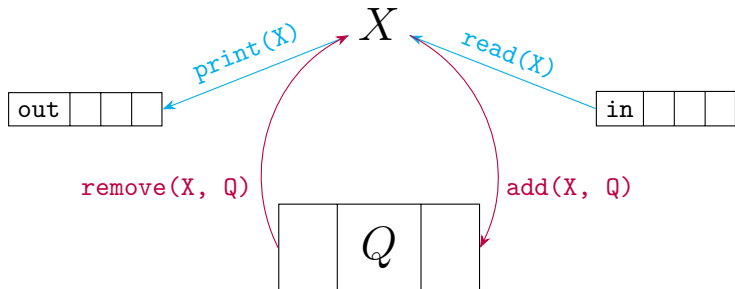
Stackable/Queueable Permutations

Treesort Algorithm

Queueable Permutations

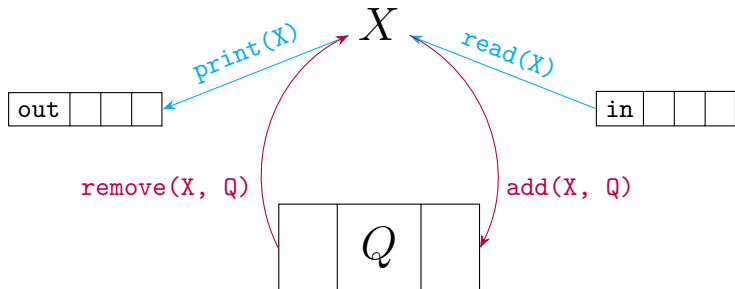


DH 2.14: Queueable Permutations



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$$\text{out} = (a_1, \dots, a_n) \xrightleftharpoons[X=0]{Q=\emptyset} \text{in} = (1, \dots, n)$$



DH 2.14: Queueable Permutations

(a) Show that the permutations given in Exercise 2.12(b) are queueable.

(i) $(3, 1, 2)$

(ii) $(4, 5, 3, 7, 2, 1, 6)$

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DH 2.14: Queueable Permutations

(b) Prove that every permutation are **queueable**.

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$X = 0$ $Q = \emptyset$ $in \neq EOF$

```
foreach 'a' ∈ out:
  if ('a' == in)
    read(X)
    print(X)
  else if ('a' > in)
    add-Q-till('a')
  else // ('a' < in)
    cycle-Q-till('a')
```

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```

```
add-Q-till('a'):
    while (('x' ∈ in) != 'a')
        add(X, Q)
    read(X)
    print(X)
```

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add-Q-till('a'):
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```

```
cycle-Q-till('a'):
    while (('x' ∈ Q) != 'a')
        remove(X, Q)
        add(X, Q)
    remove(X, Q)
    print(X)
```

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cycle-Q-till('a'):
    while (('x' ∈ Q) != 'a')
        remove(X, Q)
        add(X, Q)
    remove(X, Q)
    print(X)
```


DH 2.14: Queueable Permutations

(b) Prove that every permutation are **queueable**.

Proof.

```
foreach 'a' ∈ out:
  if ('a' >= in)
    add-Q-till('a')
  else // ('a' < in)
    cycle-Q-till('a')
```

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Proof.

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  else // ('a' ∈ Q)
    cycle-Q-till('a')
```



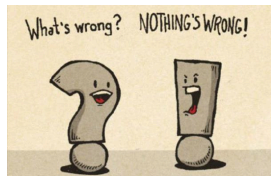
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```



Pseudocode

Pseudocode



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“Executable” at an **abstract** level.

DH 2.14: Queueable Permutations

(b) Prove that every permutation are **queueable**.

An “AHA!” Proof.

```
foreach 'a' ∈ in:
    add(X, Q)

foreach 'a' ∈ out:
    cycle-Q-till('a')
```

DH 2.14: Queueable Permutations

(b) Prove that every permutation are **queueable**.

An “AHA!” Proof.

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foreach 'a' ∈ in:
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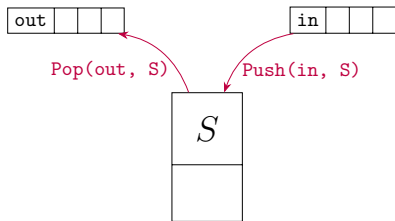
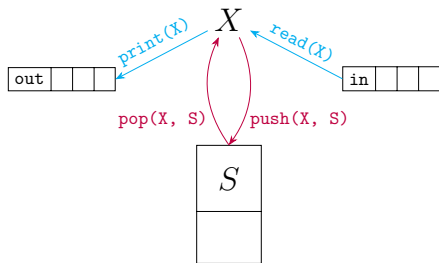


DH 2.14: Queueable Permutations

(c) Prove that every permutation can be obtained by **two stacks**.

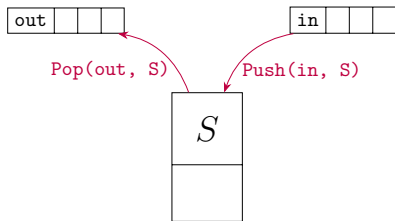
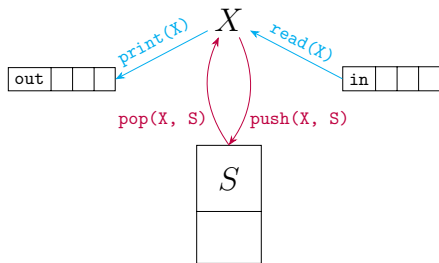
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DH 2.14: Queueable Permutations

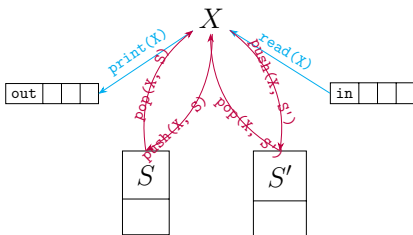
(c) Prove that every permutation can be obtained by **two stacks**.



We can similarly speak of a permutation obtained by **two stacks**, if we permit the **push** and **pop** operations on two stacks S and S' .

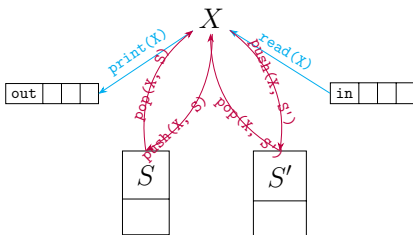
DH 2.14: Queueable Permutations

(c) Prove that every permutation can be obtained by **two stacks**.



DH 2.14: Queueable Permutations

(c) Prove that every permutation can be obtained by **two stacks**.



```
foreach 'a' ∈ in:
```

```
    read(X)
```

```
    push(X, S')
```

```
foreach 'a' ∈ out:
```

```
    if ('a' ≤ top(S')) // ∈ S'
```

```
        transfer-till(S', S, 'a')
```

```
    else // ∈ S
```

```
        transfer-till(S, S', 'a')
```

DH 2.15: Algorithm for Queueable Permutations

Extend the algorithm you were asked to design in Exercise 2.13, so that **if** the given permutation **cannot** be obtained by **a stack**, the algorithm will print the series of operations on **two stacks** that will generate it.

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`two-stackable-perm(in, X, S, S')`

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```
two-stackable-perm(in, X, S, S')
```

```
if (! stackable-perm(in, X, S))  
    two-stackable-perm(in, X, S, S')
```

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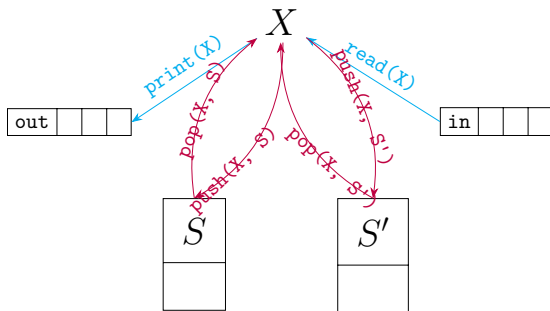
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two-stackable-perm(in, X, S, S')
```

```
if (! stackable-perm(in, X, S))  
    two-stackable-perm(in, X, S, S')
```

Embedding “transfer” into “stackable-perm”.

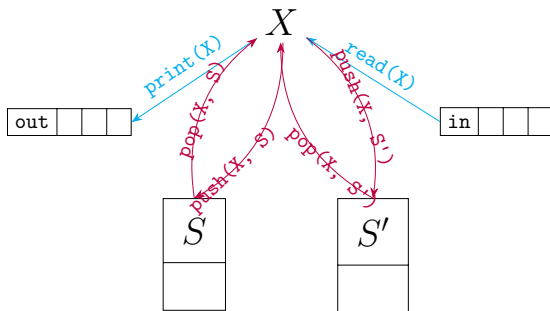
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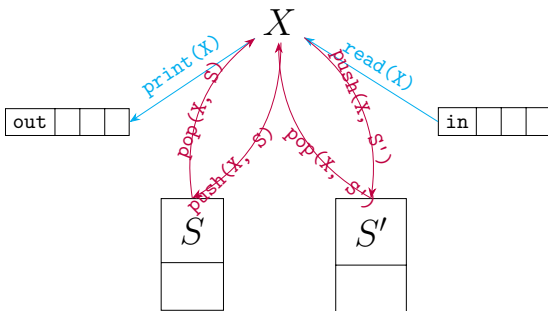
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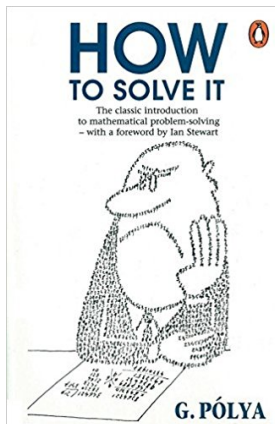
`transfer-till(S, S', top(S) == 'a')`

DH 2.15: Algorithm for Queueable Permutations

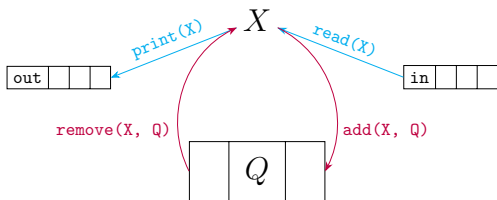
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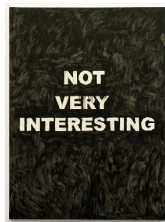
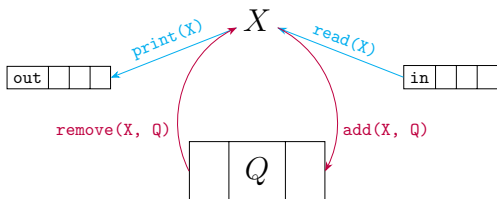


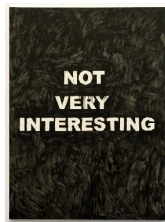
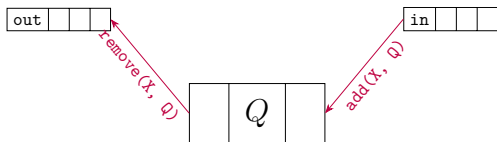
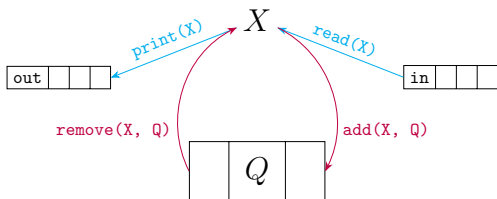
```
transfer-till(S, S', top(S) == 'a')  
transfer-till(S', S, S' ==  $\emptyset$ )
```

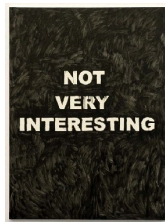
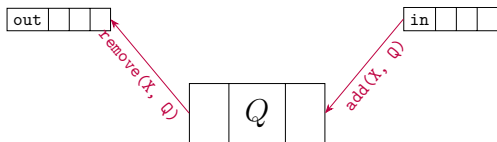
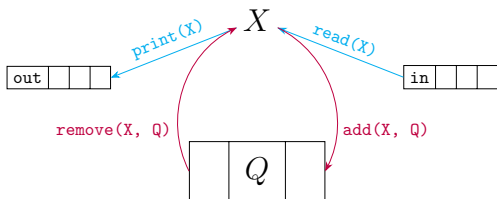


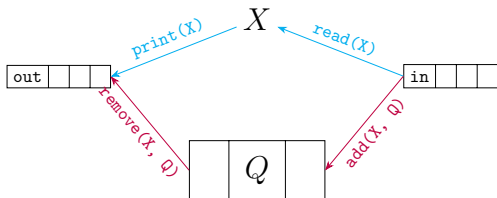
Step 4: Looking Back!

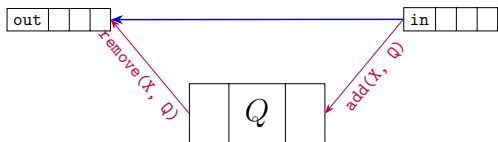
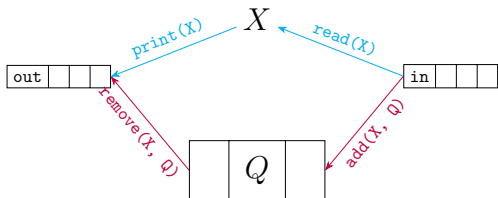


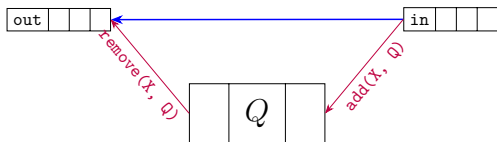
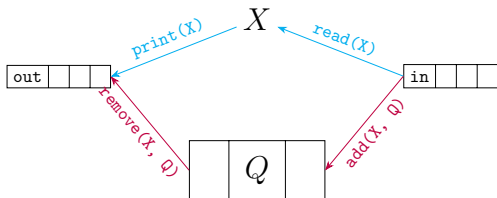




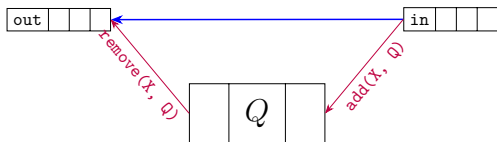
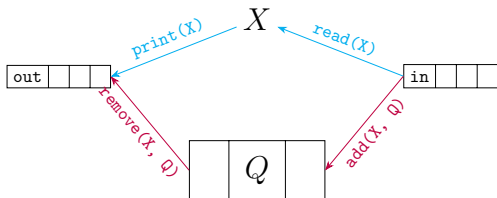








~~3 2 1~~



~~321~~



Theorem (Queueable Permutations)

A permutation (a_1, \dots, a_n) is **queueable** \iff it is not the case that

321-Pattern : $out = \dots a_i \dots a_j \dots a_k : i < j < k \wedge a_i > a_j > a_k$

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Proof.

Left as an exercise.

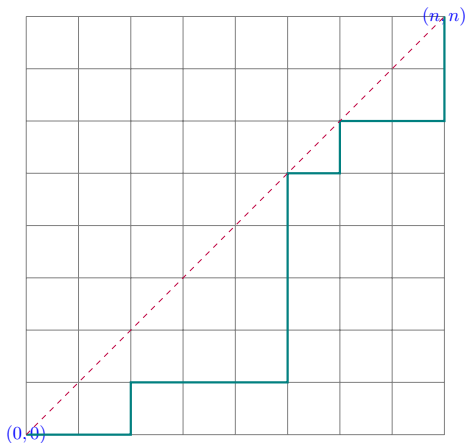


Theorem (# of Queueable Permutations)

The number of queueable permutations of $[1 \cdots n]$ is $\binom{2n}{n} - \binom{2n}{n-1}$.

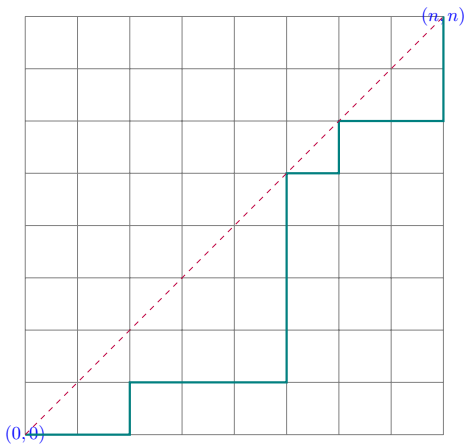
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Proof.

Left for your research.

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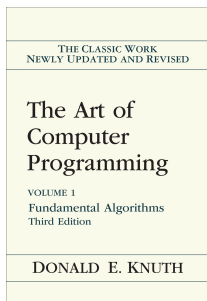
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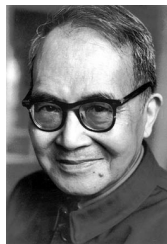
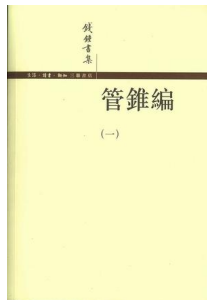
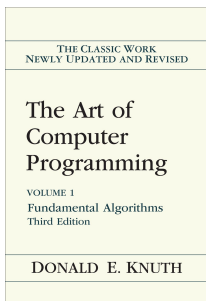
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Chapter 2—Information Structures	232
2.1. Introduction	232
2.2. Linear Lists	238
2.2.1. Stacks, Queues, and Deques	238
2.2.2. Sequential Allocation	244
2.2.3. Linked Allocation	254
2.2.4. Circular Lists	273
2.2.5. Doubly Linked Lists	280
2.2.6. Arrays and Orthogonal Lists	298
2.3. Trees	308
2.3.1. Traversing Binary Trees	318
2.3.2. Binary Tree Representation of Trees	334
2.3.3. Other Representations of Trees	348
2.3.4. Basic Mathematical Properties of Trees	362
2.3.4.1. Free trees	363
2.3.4.2. Oriented trees	372
*2.3.4.3. The “infinity lemma”	382
*2.3.4.4. Enumeration of trees	386
2.3.4.5. Path length	399
*2.3.4.6. History and bibliography	406
2.3.5. Lists and Garbage Collection	408
2.4. Multilinked Structures	424
2.5. Dynamic Storage Allocation	435
2.6. History and Bibliography	457





Thank
You!