

2-9 Sorting and Selection

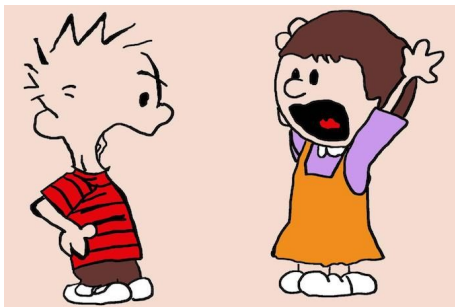
Hengfeng Wei

hfwei@nju.edu.cn

May 28, 2018

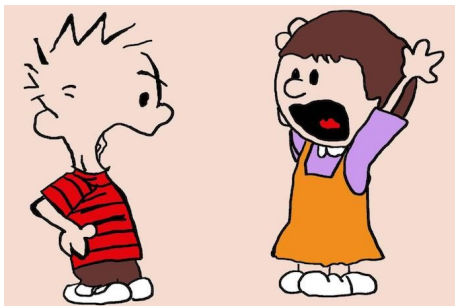


How to Argue?



Show that \dots , Argue that \dots , Explain why \dots

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Show that \dots , Argue that \dots , Explain why \dots
= Prove that \dots

不好，掉坑里了



不好，掉坑里了



“在千山万水人海相遇，喔，原来你也在这里”

入坑指南 (Coupon Collector Problem)

The Double Dixie Cup Problem

Donald J. Newman

The American Mathematical Monthly

Vol. 67, No. 1 (Jan., 1960), pp. 58-61

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$$1 \rightarrow 2 \rightarrow m$$

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The Coupon Collector's Problem

Marco Ferrante, Monica Saltalamacchia

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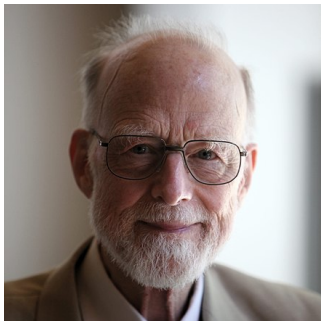


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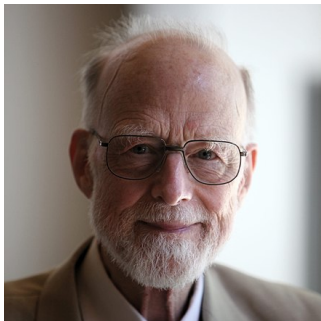
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$$n \log n + (m-1)n \log \log n + nC_m + o(n), \quad n \rightarrow \infty, m \text{ fixed}$$

QUICKSORT (Tony Hoare, 1959/1960)

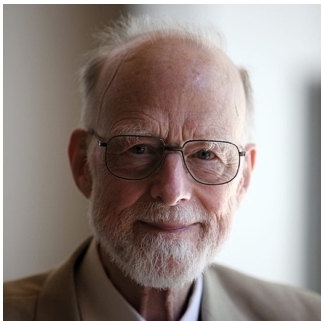


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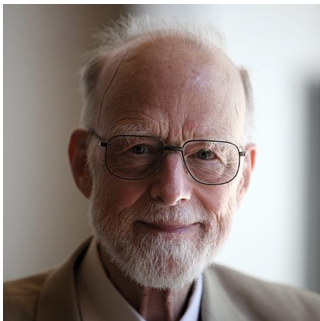
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By substitution. (Section 7.4.1, CLRS)

Sorting Almost-Sorted Inputs (Problem 7.2 – 4)

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The Analysis of Quicksort Programs*

Robert Sedgewick

Received January 19, 1976

Summary. The Quicksort sorting algorithm and its best variants are presented and analyzed. Results are derived which make it possible to obtain exact formulas describing the total expected running time of particular implementations on real computers of Quicksort and an improvement called the median-of-three modification. Detailed analysis of the effect of an implementation technique called loop unwrapping is presented. The paper is intended not only to present results of direct practical utility, but also to illustrate the intriguing mathematics which arises in the complete analysis of this important algorithm.



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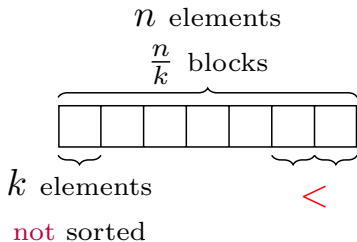
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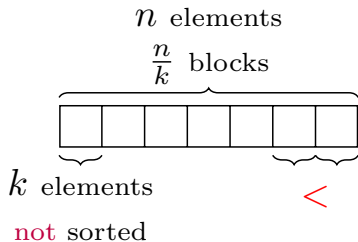
$$B_N = (N+1) \left(\frac{1}{3} H_{N+1} - \frac{1}{3} H_{M+2} + \frac{1}{6} - \frac{1}{M+2} \right) + \frac{1}{2} \quad \text{exchanges}$$

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Sorts a $\frac{n}{k}$ -sorted Array (Problem 8.1 – 4)

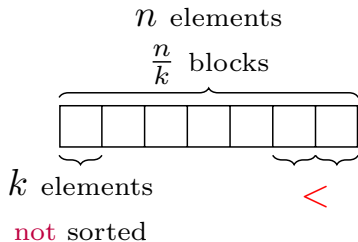


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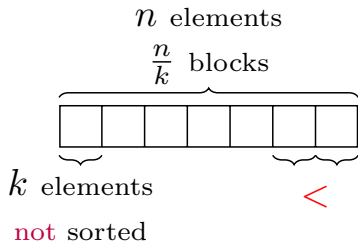
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$$\Omega(n \log k) \quad O(n \log k)$$

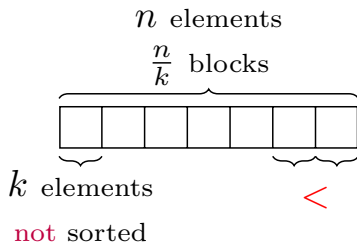
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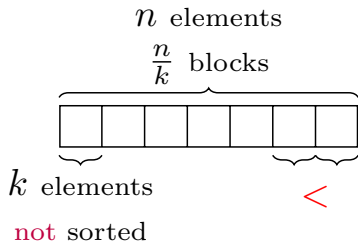


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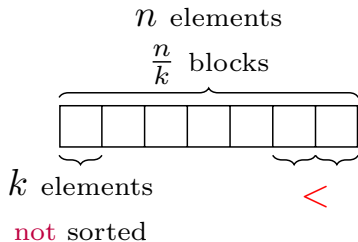
$$\Omega : \frac{n}{k} (k \log k)$$

$$(k!)^{\frac{n}{k}} \leq L \leq 2^H$$

$\frac{n}{k}$ -sorts an arbitrary array

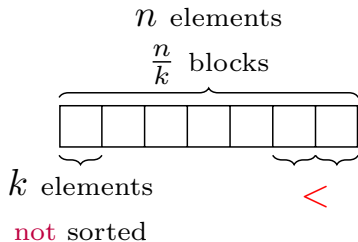


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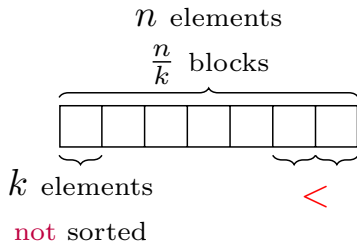
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$$O(?) \quad \Omega(?)$$

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$$L \geq \binom{n}{\underbrace{k, \dots, k}_{\frac{n}{k}}} = \frac{n!}{(k!)^{\frac{n}{k}}} \implies \Omega(n \log(n/k))$$

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Problem 3.1 – 8

When n and m go to ∞ **independently at different rates**:

$$O(g(n, m)) = \{f(n, m) \mid \exists c > 0, \exists n_0 > 0, \exists m_0 > 0 : \\ \forall n \geq n_0 \forall m \geq m_0, 0 \leq f(n, m) \leq cg(n, m)\}$$

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$$k = O(1) \quad k = \Theta(n); \quad n \rightarrow \infty$$

$$\Theta(1 + \alpha), \quad \alpha = \frac{n}{m}$$

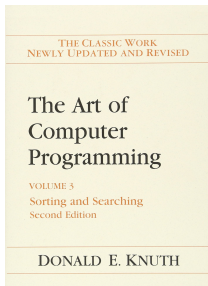
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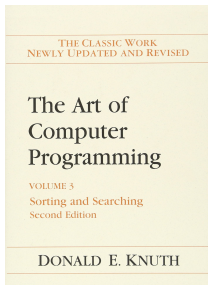
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Two ways of understanding

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$$\Theta(1 + \alpha)$$

$$n \rightarrow \infty, m \rightarrow \infty$$

$$n = f(m), m \rightarrow \infty$$

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$$n\text{-ary} \quad d = 3$$

$$n = 5 : [15, 39, 20, 123, 98] = \{030, 124, 040, 443, 343\}$$

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Any other costs?

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$$3n \cdot T(\square \div n)$$

Sorting in Place in Linear Time (Problem 8 – 2 (e))

Suppose that the n records have keys in the range $[0, k]$.

Modify COUNTING-SORT to sort them **in place** ($O(k)$) in $O(n + k)$ time.

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A:	2	5	3	0	2	3	0	3

	0	1	2	3	4	5
C:	2	0	2	3	0	1

C:	2	2	4	7	7	8
----	---	---	---	---	---	---

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B:	0	0	2	2	3	3	3	5

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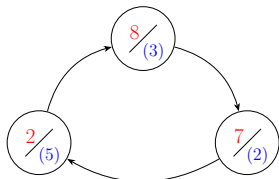
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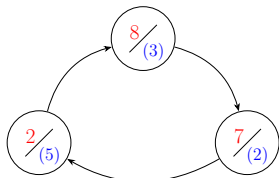
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for ($i \leftarrow n$ to 1):



Assigning elements according to C

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Q : What is the possible problem?

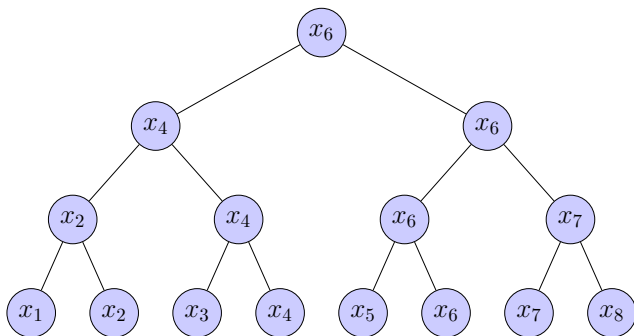
Finding the 2nd Smallest Element (Problem 9.1 – 1)

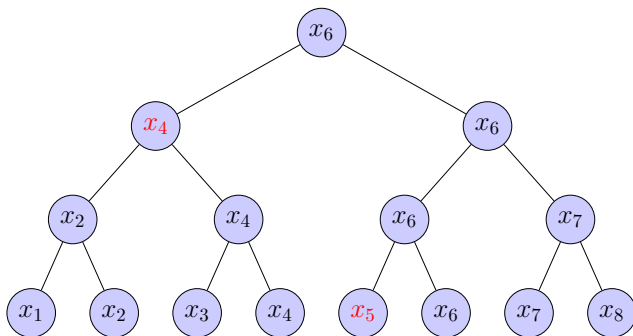
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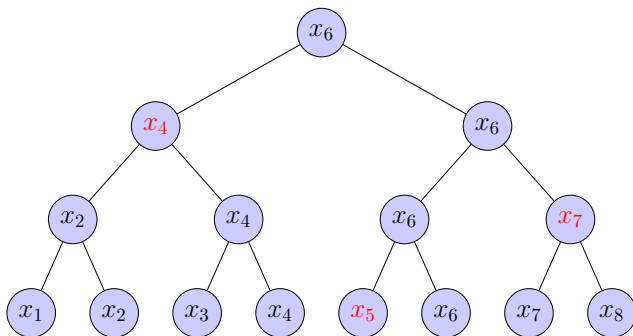
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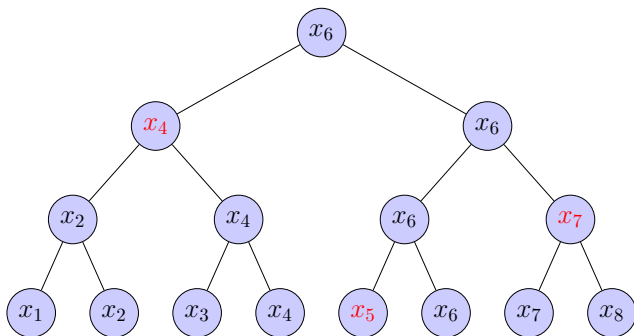
$$(n - 1) + (n - 1 - 1) = 2n - 3$$





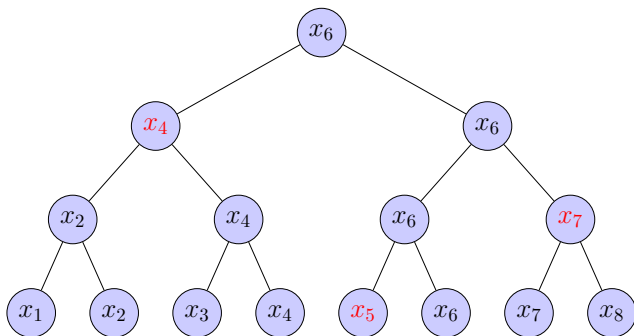


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Q : Can we do even better?

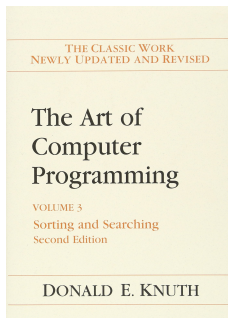
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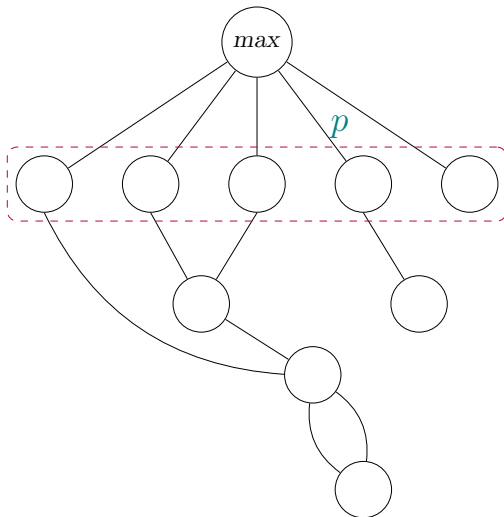


TAOCP, Vol. 3 (Page 209, Section 5.3.3)

Theorem

Any comparison-based algorithm to find secondLargest in n keys must do at least $n + \lceil \log n \rceil - 2$ comparisons in the worst case.

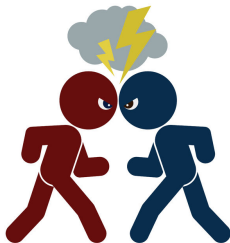
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$$p \geq \lceil \log n \rceil$$

Adversary \mathcal{A} :

$$x \succcurlyeq y$$



Algorithm \mathcal{A} :

COMPARE(x, y)

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Algorithm \mathbb{A} :

COMPARE(x, y)

$\exists \mathcal{A} \forall \mathbb{A} \left(\geq \lceil \log n \rceil \text{ distinct keys lost directly to } \max \text{ in } \mathbb{A} \right)$



$$\forall x : w(x) = 1$$

COMPARE(x, y)

CASE	REPLY	(LOCAL) ACTION
$w(x) > w(y)$	$x > y$	$w'(x) \leftarrow w(x) + w(y); w'(y) \leftarrow 0$
$w(x) = w(y) > 0$	$x > y$	$w'(x) \leftarrow w(x) + w(y); w'(y) \leftarrow 0$
$w(x) = w(y) = 0$	No Conflict	NOP
$w(y) > w(x)$	$y > x$	$w'(y) \leftarrow w(x) + w(y); w'(x) \leftarrow 0$

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贏者通吃; 敗者永不得翻身

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For any algorithm \mathbb{A} that finds *secondLargest*, when it stops,

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t has directly won against $\geq \lceil \log n \rceil$ distinct keys.

Proof.

K : # of comparisons x wins against previously undefeated keys

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units of WIN/LOSE info.

$$2n - 2 = 2 \times (n/2) + 1 \times (n - 2)$$

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Do we have to check every bit?

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Do we have to check every bit?

No, if n is odd;

Yes, if n is even.

1 2 3 4 5 6 7 8 9

--	--	--	--	--	--	--	--	--

First checking $A[2, 4, \dots, n - 1]$

1 2 3 4 5 6 7 8 9

--	--	--	--	--	--	--	--	--

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	0		0		0		0	
--	---	--	---	--	---	--	---	--

	1		1		1		1	
--	---	--	---	--	---	--	---	--

			0				1	
--	--	--	---	--	--	--	---	--

	1		1		0		0	
--	---	--	---	--	---	--	---	--

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CHECK(i)

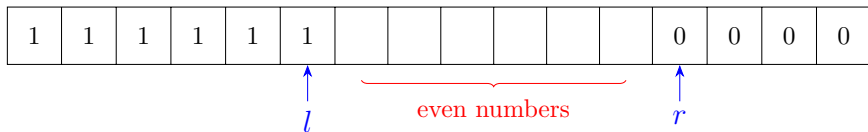
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median + subtraction + $(k + 1)$ -th smallest + partition + add back



The Coupon Collector's Problem

Marco Ferrante, Monica Saltalamacchia

In this note we will consider the following problem: how many coupons we have to purchase (on average) to complete a collection. This problem, which takes everybody back to his childhood when this was really “a problem”, has been considered by the probabilists since the eighteenth century and nowadays it is still possible to derive some new results, probably original or at least never published. We will present some classic results, some new formulas, some alternative approaches to obtain known results and a couple of amazing expressions.



Thank
You!



Office 302

Mailbox: H016

hfwei@nju.edu.cn