

## 2-11 Heapsort

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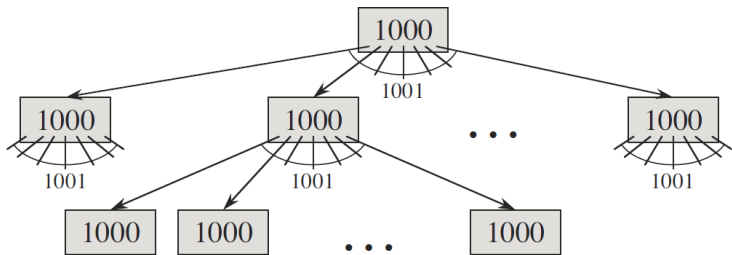
# Organization and Maintenance of Large Ordered Indexes

R. BAYER and E. McCREIGHT

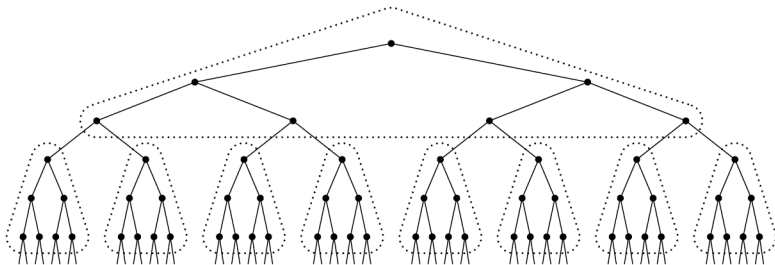
Received September 29, 1971

*Summary.* Organization and maintenance of an index for a dynamic random access file is considered. It is assumed that the index must be kept on some pseudo random access backup store like a disc or a drum. The index organization described allows retrieval, insertion, and deletion of keys in time proportional to  $\log_k I$  where  $I$  is the size of the index and  $k$  is a device dependent natural number such that the performance of the scheme becomes near optimal. Storage utilization is at least 50% but generally much higher. The pages of the index are organized in a special data-structure, so-called *B-trees*. The scheme is analyzed, performance bounds are obtained, and a near optimal  $k$  is computed. Experiments have been performed with indexes up to 100000 keys. An index of size 15000 (100000) can be maintained with an average of 9 (at least 4) transactions per second on an IBM 360/44 with a 2311 disc.

*“Bayer and McCreight introduced B-trees in 1972;  
they **did not** explain their choice of name.”*



2-way *vs.* multi-way

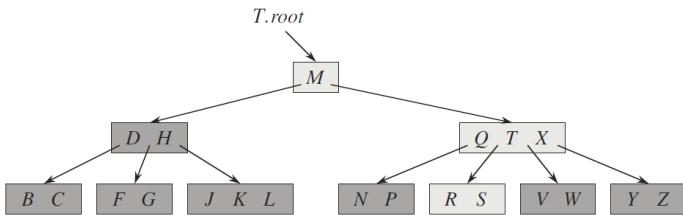


**Fig. 29.** A large binary search tree can be divided into “pages.”

node *vs.* pages

## Minimum (TC 18.2-3)

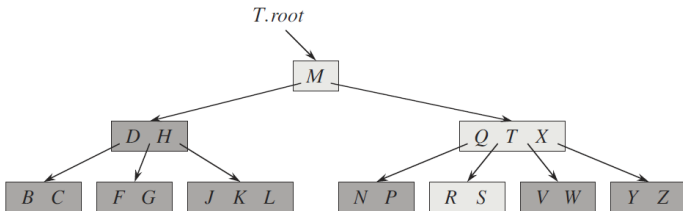
Explain how to find the **minimum** key stored in a B-tree.



the leftmost key in the leftmost node

## Predecessor (TC 18.2-3)

Explain how to find the predecessor of a given key stored in a B-tree.









### Insertion (TC 18.2-4 ★)

Suppose that we insert the keys  $\{1, 2, \dots, n\}$  in increasing order into an empty B-tree with minimum degree 2.

How many nodes, denoted  $X_n$ , does the final B-tree have?

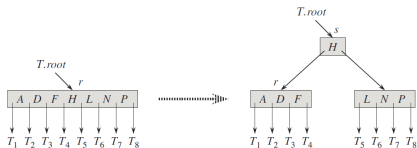


$$X_0 = 1$$

By [Yangjing Dong](#) (June 2018)

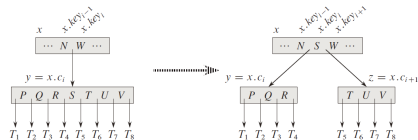
<https://maxmute.com/TC18.2-4.html>

Only **SPLIT** can create new nodes.



root SPLIT

+2



non-root SPLIT

+1

(I) Which nodes will SPLIT?  $S$

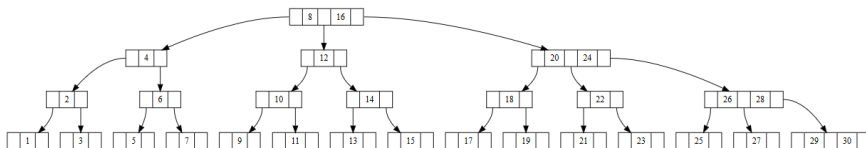
(II) When does each node  $s \in S$  SPLIT?  $T_s = \{s_1, s_2, \dots\}$

(III) How does it SPLIT, as a root or a non-root?  $T_s = s_R \uplus s_{NR}$

$$X_n = 1 + \sum_{s \in S} (2 |s_R| + |s_{NR}|)$$

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(I) Which nodes will SPLIT?

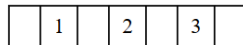


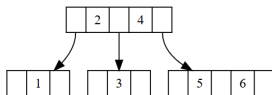
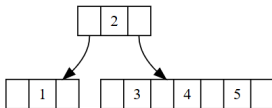
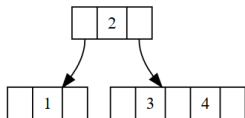
1, 2, ..., 30

$S = \left\{ \text{the nodes in the rightmost chain} \right\}$

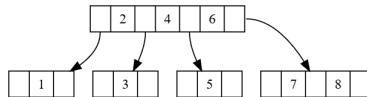
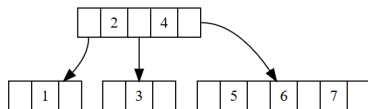
(II) When does each node SPLIT?

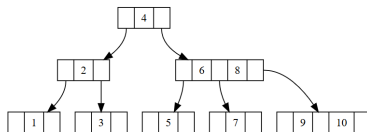
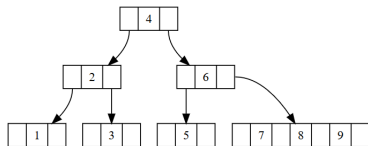
Let's focus the **rightmost** node first, denoted  $A$ .

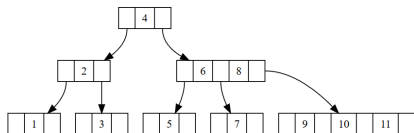










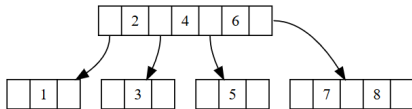
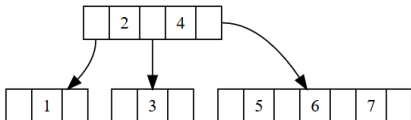
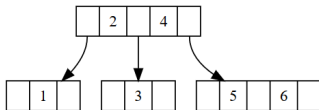
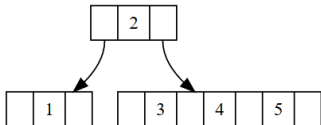
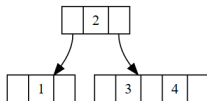
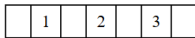
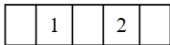


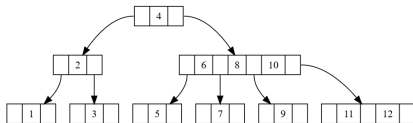
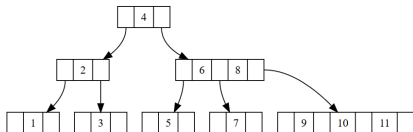
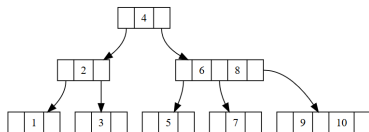
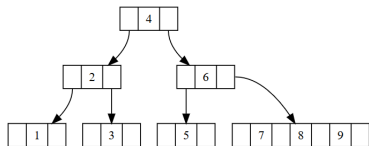
A SPLIT : 4, 6, 8, 10, 12, ...

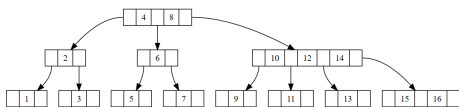
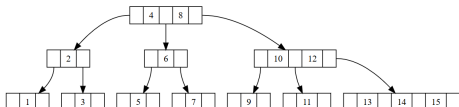
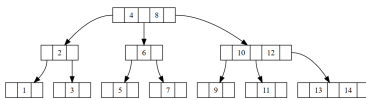
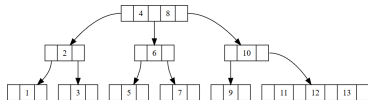
(II) When does each node SPLIT?

Let's consider the parent of  $A$ , denoted  $B \triangleq p(A)$ .

Every time  $A$  splits,  $B$  obtains a new key.







*B* SPLIT : 9, 13, 17, 21, 25, ...

*A* SPLIT : 4, 6, 8, 10, 12, ...

*B* SPLIT : 9, 13, 17, 21, 25, ...



(II) When does each node SPLIT?

Let's consider the parent of  $B$ , denoted  $C = p(B)$ .

$A$  SPLIT : 4, 6, 8, 10, 12, ...  
 $B$  SPLIT : 9, 13, 17, 21, 25, ...  
 $C$  SPLIT : 18, 26, 34, 42, 50, ...

---

$A : 1$      $B : 2$      $C : 3$

$T_i$  : the first time point the  $i$ -th node splits

$$T_1 = 4$$

$$T_i = \underbrace{T_{i-1}}_{\text{its right child first split}} + \underbrace{2 \times 2^{i-1}}_{\text{its right child split twice more}} + \underbrace{1}_{\text{insert one more}}$$

$$T_i = 2^{i+1} + i - 1$$

$$X_n = 1 + \sum_{s \in S} \left( 2 \left| s_R \right| + \left| s_{NR} \right| \right)$$

$$(T_s = s_R \uplus s_{NR})$$

(III) How does it SPLIT, as a root or a non-root?

$$s_R = \{s_1\} \quad s_{NR} = \{s_2, s_3, \dots\}$$

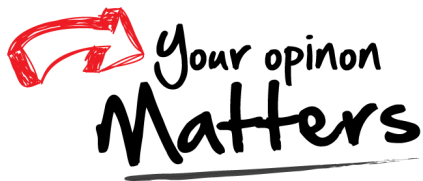
$$\left| s_R \right| = 1 \quad \left| s_{NR} \right| = \left| T_s \right| - 1$$

$$X_n = 1 + \sum_{s \in S} \left( 2 + \left| T_s \right| - 1 \right) = 1 + \sum_{s \in S} \left( \left| T_s \right| + 1 \right)$$

$$X_n = 1 + \sum_{s \in S} (|T_s| + 1) \quad T_i = 2^{i+1} + i - 1$$

$$\begin{aligned}
 X_n &= 1 + \sum_{i=1}^{\infty} [T_i \leq n] \left( \left( \left\lfloor \frac{n - T_i}{2^i} \right\rfloor + 1 \right) + 1 \right) \\
 &= 1 + \sum_{i=1}^{\infty} [T_i \leq n] \left( \left\lfloor \frac{n - T_i}{2^i} \right\rfloor + 2 \right) \\
 &= 1 + \sum_{i=1}^{\infty} [T_i \leq n] \left( \left\lfloor \frac{n - 2^{i+1} - i + 1}{2^i} \right\rfloor \right) \\
 &= 1 + \sum_{i=1}^{2^{i+1} + i - 1 \leq n} \left\lfloor \frac{n - i + 1}{2^i} \right\rfloor \\
 &= 1 + \sum_{i=0}^{2^{i+2} + i \leq n} \left\lfloor \frac{n - i}{2^{i+1}} \right\rfloor
 \end{aligned}$$

Thank  
You!



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