

# 1-9 关系及其基本性质 (II)

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# The Relational Data Model

## — 如何靠“关系”赢得图灵奖？

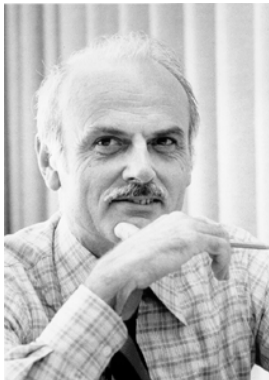
### A Relational Model of Data for Large Shared Data Banks

E. F. CODD

*IBM Research Laboratory, San Jose, California*

Future users of large data banks must be protected from having to know how the data is organized in the machine (the internal representation). A prompting service which supplies such information is not a satisfactory solution. Activities of users at terminals and most application programs should remain unaffected when the internal representation of data is changed and even when some aspects of the external representation are changed. Changes in data representation will often be needed as a result of changes in query, update, and report traffic and natural growth in the types of stored information.

Existing noninferential, formatted data systems provide users with tree-structured files or slightly more general network models of the data. In Section 1, inadequacies of these models are discussed. A model based on  $n$ -ary relations, a normal form for data base relations, and the concept of a universal data sublanguage are introduced. In Section 2, certain operations on relations (other than logical inference) are discussed and applied to the problems of redundancy and consistency in the user's model.



Edgar Frank Codd (1923 – 2003)

$$(a, b) = \{\{a\}, \{a, b\}\}$$

$$R \subseteq A \times B$$

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$$(a, b, c) = ((a, b), c)$$

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$$R \subseteq A \times B$$

$$(a, b, c) = ((a, b), c)$$

$$(x_1, x_2, \dots, x_n) = ((x_1, x_2, \dots, x_{n-1}), x_n)$$

$$(a, b) = \{\{a\}, \{a, b\}\}$$

$$R \subseteq A \times B$$

$$(a, b, c) = ((a, b), c)$$

$$(x_1, x_2, \dots, x_n) = ((x_1, x_2, \dots, x_{n-1}), x_n)$$

$$R \subseteq X_1 \times X_2 \times \dots \times X_n$$

Course = {CS101, EE200, PH100}  
StudentId = {12345, 67890, 22222, 33333}  
Grade = {A, B, C, D}



Course = {CS101, EE200, PH100}  
StudentId = {12345, 67890, 22222, 33333}  
Grade = {A, B, C, D}

$CSG \subseteq \text{Course} \times \text{StudentId} \times \text{Grade}$

$CSG = \{(CS101, 12345, A),$   
           $(CS101, 67890, B),$   
           $(CS101, 33333, C),$   
           $(EE200, 12345, B^+),$   
           $(EE200, 22222, A^-),$   
           $(PH100, 67890, C^+)\}$

Course	StudentId	Grade
CS101	12345	A
CS101	67890	B
EE200	12345	C
EE200	22222	B+
CS101	33333	A-
PH100	67890	C+

Course	StudentId	Grade
CS101	12345	A
CS101	67890	B
EE200	12345	C
EE200	22222	B+
CS101	33333	A-
PH100	67890	C+

Row : tuple

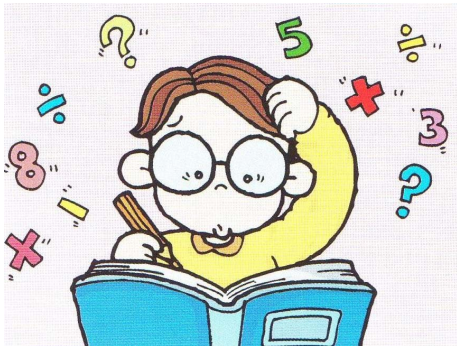
Column : attribute

StudentId	Name	Address	Phone
12345	C. Brown	12 Apple St.	555-1234
67890	L. Van Pelt	34 Pear Ave.	555-5678
22222	P. Patty	56 Grape Blvd.	555-9999

Course	Day	Hour
CS101	M	9AM
CS101	W	9AM
CS101	F	9AM
EE200	Tu	10AM
EE200	W	1PM
EE200	Th	10AM

Course	Room
CS101	Turing Aud.
EE200	25 Ohm Hall
PH100	Newton Lab.

# Relational Algebra



# Relational Algebra



## Operations on Relations

$\cup \quad \cap \quad \setminus$

## Selection ( $\sigma$ )

$$\sigma_{\text{Course}='CS101'}(\text{CSG})$$

Course	StudentId	Grade
CS101	12345	A
CS101	67890	B
EE200	12345	C
EE200	22222	B+
CS101	33333	A-
PH100	67890	C+

Course	StudentId	Grade
CS101	12345	A
CS101	67890	B
CS101	33333	A-











Thank  
You!



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