

# 3-11 Matchings and Factors

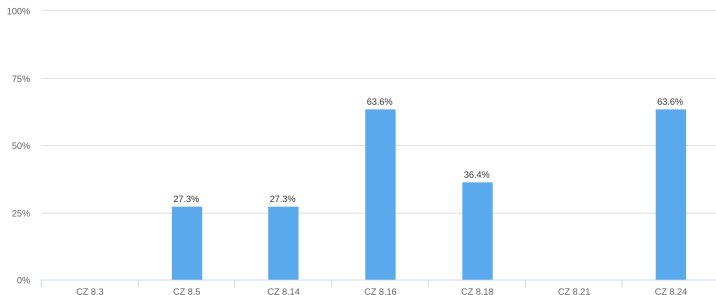
## (Part I: Matchings and Covers)

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8.5      8.14      8.16      Chinese Postman Problem (The Last Class?)  
8.18      8.24      (The Last Class)

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比较大的定理（证明比较长的）都不是很理解，想知道期末考什么

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点覆盖边覆盖那里只知道有这些性质，了解不是很深

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都理解

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图的分解的形象意义

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无

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定理8.3的证明

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$\alpha\beta$ 、 $\alpha'\beta'$ 的定义和几个定理推论

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为什么中英文书上的定义中 $\alpha$ 和 $\beta$ 反了。。

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定理8.10的证明看不懂；一些比较几何的构造法证明（比如把顶点排成正多边形，一个点放中间）是怎么保证这些分解不重不漏的？

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Kirkman三元系

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$$\alpha, \quad \beta, \quad \alpha', \quad \beta'$$

Theorem 8.10 (Tutte's Theorem)    (The Last Class)

## Theorem (Hall's Theorem, 1935; Theorem 8.3)

Let  $G$  be a *bipartite graph* with partite sets  $U$  and  $W$  such that  $r = |U| \leq |W|$ .

$G$  contains a matching of cardinality  $r \iff G$  satisfies *Hall's Condition*:

$$\forall X \subseteq U : |N(X)| \geq |X|$$

## TONCAS

(The Obvious Necessary Conditions are Also Sufficient)

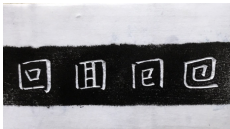


Other TONCAS?

## Perfect Matching on Trees (Problem 8.5)

Prove that every tree  $T$  has  $\leq 1$  perfect matching.

“这题有四样证法, 你知道吗?”



## Perfect Matching on Trees (Problem 8.5)

Prove that every tree  $T$  has  $\leq 1$  perfect matching.

By strong mathematical induction on the order  $n$  of trees.

Inductive step: Consider a tree of order  $n$ .

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1: if  $n$  is odd then
2:   # Perfect Matching = 0
3: else                                ▷  $n$  is even
4:   Consider  $T - r$                   ▷  $r$  : root of  $G$ 
5:   if  $k_o(T - r) > 1$  then
6:     # Perfect Matching = 0
7:   else                              ▷  $k_o(T - r) = 1$ 
8:     By Induction Hypothesis.
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## Perfect Matching on Trees (Problem 8.5)

Prove that every tree  $T$  has  $\leq 1$  perfect matching.

By strong mathematical induction on the order  $n$  of trees.

Inductive step: Consider a tree of order  $n$ .

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- 1: **if**  $T$  has no perfect matchings **then**
  - 2:      $\# \text{ Perfect Matching} = 0$
  - 3: **else**  $\triangleright T$  has perfect matchings
  - 4:     Consider a leaf  $v$
  - 5:      $v$  **must** be matched with its parent  $u$
  - 6:     **By Induction Hypothesis** on each component of  $G - \{u, v\}$
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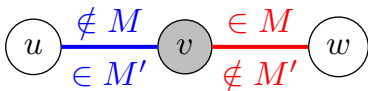
## Perfect Matching on Trees (Problem 8.5)

Prove that every tree  $T$  has  $\leq 1$  perfect matching.

By Contradiction.

Suppose that there are two different perfect matchings  $M$  and  $M'$  on  $T$ .

$\exists v : v$  is matched with different vertices in  $M$  and  $M'$ .



$Q$  : What about  $u$  and  $w$ ?

Contradiction: Cycle



## Perfect Matching on Trees (Problem 8.5)

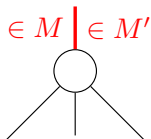
Prove that every tree  $T$  has  $\leq 1$  perfect matching.

By Contradiction.

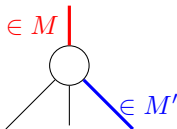
Suppose that there are two different perfect matchings  $M$  and  $M'$  on  $T$ .

$$M \Delta M' = (M - M') \cup (M' - M)$$

Consider the subgraph  $\mathcal{M}$  with  $V(T)$  and  $M \Delta M'$ .



CASE I



CASE II

$$\forall v \in V(\mathcal{M}) :$$

$$\deg(v) = 0 \vee \deg(v) = 2$$

$$T \text{ is a tree} \implies \deg(v) = 0$$

$$\deg(v) = 0 \implies \text{CASE I}$$

## Theorem (Gallai Identities, 1959; Theorem 8.7)

*If  $G$  is graph without isolated vertices, then*

$$\alpha'(G) + \beta'(G) = n(G).$$

$\alpha(G)$     Maximum size of independent set

$\beta(G)$     Minimum size of vertex cover

$\alpha'(G)$     Maximum size of matching

$\beta'(G)$     Minimum size of edge cover

## Theorem (Gallai Identities, 1959; Theorem 8.8)

*If  $G$  is graph without isolated vertices, then*

$$\alpha(G) + \beta(G) = n(G).$$

## Matching and Edge Cover (Problem 8.14)

A graph  $G$  without isolated vertices has a perfect matching if and only if  $\alpha'(G) = \beta'(G)$ .

“ $\implies$ ”

“ $\impliedby$ ”

$G$  has a perfect matching

$$\implies n \text{ is even} \wedge \alpha'(G) = n/2$$

$$\implies \beta'(G) = n/2$$

$$\alpha'(G) = \beta'(G)$$

$$\implies \alpha'(G) = n/2 \wedge n \text{ is even}$$

$$\implies G \text{ has a perfect matching}$$

Theorem (König, 1931; Egerváry, 1931)

If  $G$  is a *bipartite graph*, then

$$\alpha'(G) = \beta(G).$$

$\alpha(G)$  Maximum size of independent set

$\beta(G)$  Minimum size of vertex cover

$\alpha'(G)$  Maximum size of matching

$\beta'(G)$  Minimum size of edge cover

Theorem (König, 1931)

If  $G$  is a *bipartite graph*, then

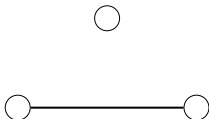
$$\alpha(G) = \beta'(G).$$

## Vertex Covering Number (Problem 8.16)

If  $G$  is a graph of order  $n$ , maximum degree  $\Delta$  and **having no isolated vertices**, then

$$\beta(G) \geq \frac{n}{\Delta + 1}.$$

By Contradiction:  $\beta < \frac{n}{\Delta + 1}$ .



$$n = 3, \Delta = 1, \frac{n}{\Delta + 1} = \frac{3}{2}, \beta = 1$$

$$\begin{aligned} \beta \cdot \Delta &< \frac{n\Delta}{\Delta + 1} \\ &= n - \frac{n}{\Delta + 1} \\ &\leq n - 1 \end{aligned}$$



## Vertex Covering Number (Problem 8.16)

If  $G$  is a graph of order  $n$ , maximum degree  $\Delta$  and having no isolated vertices, then

$$\beta(G) \geq \frac{n}{\Delta + 1}.$$

### Double Counting:

What is the number of **neighbors** of a vertex cover  $C$ ?

$$N(C) \leq |C| \cdot \Delta$$

$$N(C) = n - |C|$$

$$n - |C| \leq |C| \cdot \Delta \implies |C| \geq \frac{n}{\Delta + 1}$$

## Vertex Independence Number (Additional Problem)

If  $G$  is a graph of order  $n$ , maximum degree  $\Delta$ , then

$$\alpha(G) \geq \frac{n}{\Delta + 1}.$$

By Construction.

To construct an independent set  $S$  with  $|S| \geq \frac{n}{\Delta+1}$ .

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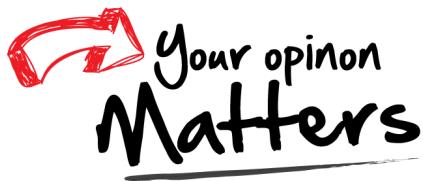
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1: while  $|V(G)| > 0$  do
2:   Choose  $v \in V(G)$ 
3:    $S \leftarrow S \cup \{v\}$ 
4:    $G \leftarrow G - \{v\} - N(v)$ 
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