# 3-10 Traversability

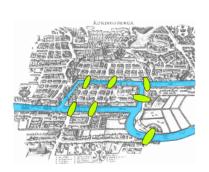
(Part I: Eulerian Graphs)

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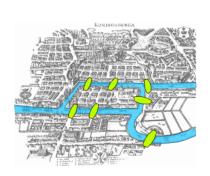
December 03, 2018







Leonhard Euler (1707 – 1783)





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Graph Theory Topology













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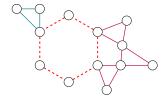
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$$H = G - E(T) = \bigcup H_i$$



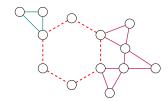
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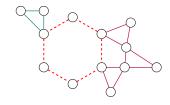
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(II) 
$$\forall i : E(H_i) < m$$



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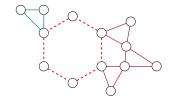
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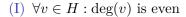
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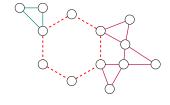
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Combine each  $T_i$  with T to get an Eulerian circuit of G.

1: **procedure** HIERHOLZER(G)

1: **procedure** FLEURY(G, w, s)





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