1-5 数据与数据结构(Ⅱ)

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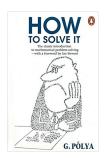


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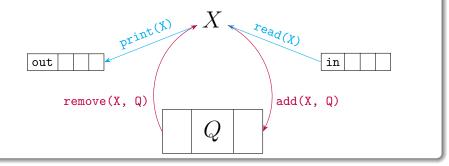
Looking Back!

Stackable/Queueable Permutations Treesort Algorithm

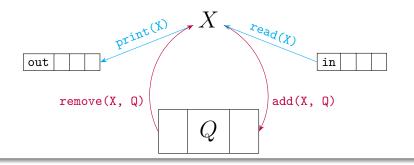
Queueable Permutations



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$$\mathtt{out} = (a_1, \cdots, a_n) \stackrel{Q=\emptyset}{\longleftarrow} \mathtt{in} = (1, \cdots, n)$$



- (a) Show that the permutations given in Excecise 2.12(b) are queueable.
 - (i) (3,1,2)
 - (ii) (4,5,3,7,2,1,6)

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```
X = 0 Q = \emptyset in != EOF
```

```
foreach 'a' ∈ out:
    if ('a' == in)
        read(X)
        print(X)
    else if ('a' > in)
        add-Q-till('a')
    else // ('a' < in)
        cycle-Q-till('a')</pre>
```

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```

```
add-Q-till('a'):
   while (('x' \in in) != 'a')
      add(X, Q)
   read(X)
   print(X)
```

```
X = 0
             Q = \emptyset
                         in != EOF
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```
add-Q-till('a'):
  while (('x' \in in) != 'a')
    add(X, Q)
  read(X)
  print(X)
```

```
cycle-Q-till('a'):
  while (('x' \in Q) != 'a')
    remove(X, Q)
    add(X, Q)
  remove(X, Q)
  print(X)
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add-Q-till('a'):
   while (('x' \in in) != 'a')
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   read(X)
   print(X)
```

```
cycle-Q-till('a'):
    while (('x' ∈ Q) != 'a')
        remove(X, Q)
        add(X, Q)
    remove(X, Q)
    print(X)
```

(b) Prove that every permutation are queueable.

Proof.

```
foreach 'a' \in out:
    if ('a' >= in)
        add-Q-till('a')
    else // ('a' < in)
        cycle-Q-till('a')</pre>
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Proof.

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foreach 'a' \in out:
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      cycle-Q-till('a')
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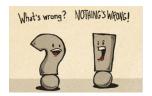


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foreach 'a' ∈ out:
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```



Pseudocode

Pseudocode



Pseudocode



"Executable" at an abstract level.

(b) Prove that every permutation are queueable.

An "AHA!" Proof.

```
foreach 'a' ∈ in:
   add(X, Q)

foreach 'a' ∈ out:
   cycle-Q-till('a')
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An "AHA!" Proof.

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foreach 'a' ∈ in:
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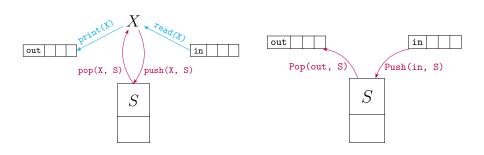
foreach 'a' ∈ out:
  cycle-Q-till('a')
```



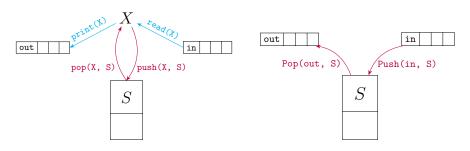


(c) Prove that every permutation can be obtained by two stacks.

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We can similarly speak of a permutation obtained by two stacks, if we permit the push and pop operations on two stacks S and S'.

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2 stacks model

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2 stacks model

```
foreach 'a' \in in:
    read(X)
    push(X, S')

foreach 'a' \in out:
    if ('a' <= top(S')) // \in S'
        transfer-till(S', S, 'a')
    else // \in S
        transfer-till(S, S', 'a')
```

(c) Prove that every permutation can be obtained by two stacks.

2 stacks model

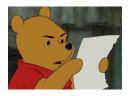
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    else // \in S
        transfer-till(S, S', 'a')
```

Can you PROVE it?

Extend the algorithm you were asked to design in Exercise 2.13, so that **if** the given permutation cannot be obtained by a stack, the algorithm will print the series of operations on two stacks that will generate it.

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two-stackable-perm(in, X, S, S')

if (! stackable-perm(in, X, S))
 two-stackable-perm(in, X, S, S')

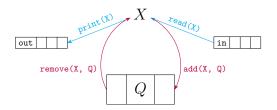
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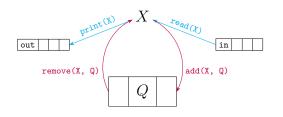


two-stackable-perm(in, X, S, S')

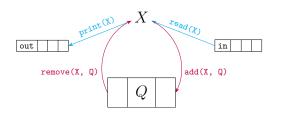
if (! stackable-perm(in, X, S))
 two-stackable-perm(in, X, S, S')

Embedding "transfer" into "stackable-perm".

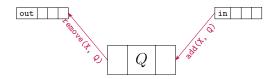


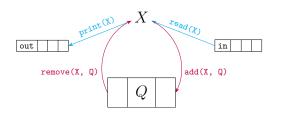




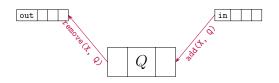




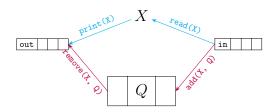


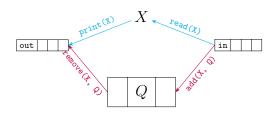


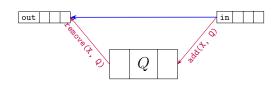


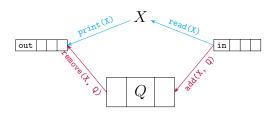


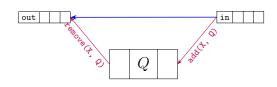




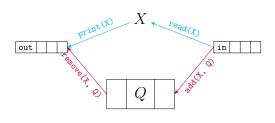


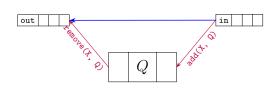






3 2 1







3 2 1

Theorem (Queueable Permutations)

A permutation (a_1, \cdots, a_n) is queueable \iff it is not the case that

321-Pattern :
$$out = \cdots a_i \cdots a_j \cdots a_k : i < j < k \land a_i > a_j > a_k$$

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Thank You!