

# Data Structure

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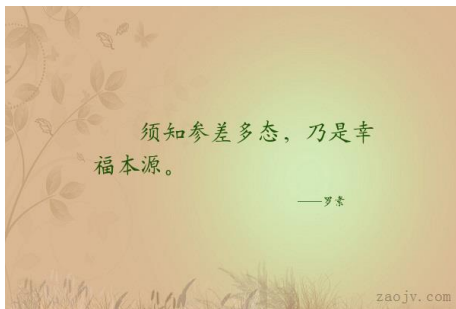
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# Data Structure

- (i) What is (and why) data structure?
- (ii) Common (simple) data structures:
  - (1) Variable, Pointer
  - (2) Linear data structures:
    - ▶ Array, List (Singly-linked List, Doubly-linked List, Circular List)
    - ▶ Stack, Queue (Deque, Priority Queue)
  - (3) Trees
    - ▶ Binary Search Tree (BST)
    - ▶ ...
  - (4) Hashes
  - (5) Graphs
  - (6) ...

Why are there so many data structures?



# Data Structure vs. Data Type

Data type: data + operations

Data structure: data type + structure

A data structure is an **implementation** of an abstract data type (ADT).

Example: Sequence of Data

OP: SEARCH, INSERT, DELETE

Array vs. List

# Variable and Pointer

# Memory

## Definition (Memory (K&R))

The memory is organized as a collection of consecutively **addressed** cells that may be manipulated **individually or in contiguous groups**.

**address of  
memory cell**    ***RAM (memory)***

000...000	00001101
000...001	00000011
000...010	00000000
000...011	00101101



# Variable

```
int x;
```

# Pointer

## Definition (Pointer (K&R))

A pointer is a **variable** that contains the **address** of a variable.

```
int a = 0;  
int *p = &a;
```

## Definition (Pointer in Memory (K&R))

A pointer is a group of cells (often two or four) that can hold an address.



## swap

```
swap(a, b);  
  
void swap(int a, int b) {  
    int temp = a;  
    a = b;  
    b = tmp;  
}
```

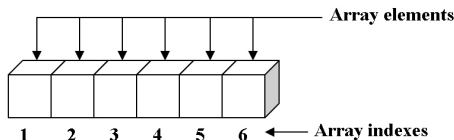
*Pointer arguments enable a function to access and change objects in the function that called it.* — K&R

```
swap(&a, &b);  
  
void swap(int *a, int *b) {  
    int temp = *a;  
    *a = *b;
```

# Array

# Array

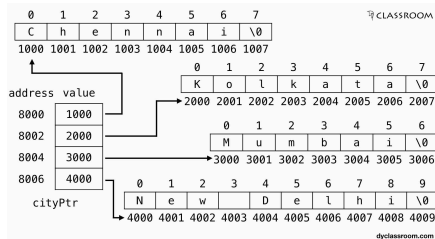
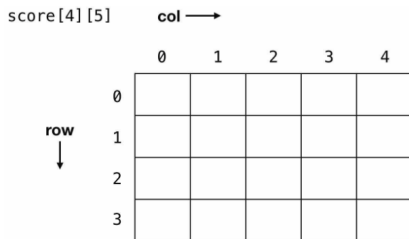
Array: A sequence of **contiguously** stored elements.



One-dimensional array with six elements

```
vector<int> array {1,5,7,9,10};  
  
array[1] = 3;    // offset  
array.insert(pos, val); // moving elements  
array.erase(pos) // moving elements
```

# 2D Array

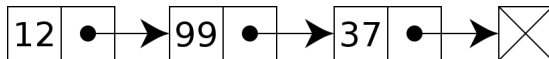


```
vector<int> array {1,5,7,9,10};  
matrix<int>
```

```
array[1] = 3;    // offset  
array.insert(pos, val); // moving elements  
array.erase(pos) // moving elements
```

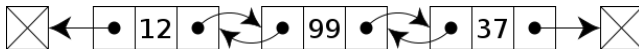
# List

# Singly-linked List



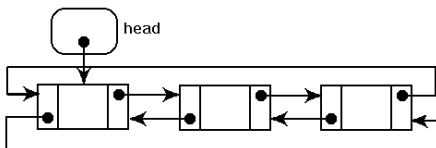
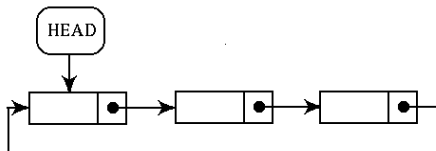
SEARCH, INSERT, DELETE

# Doubly-linked List



SEARCH, INSERT, DELETE

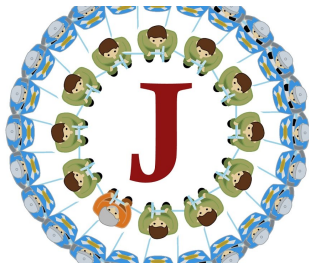
# Circular Linked List



Doubly Linked Circular list



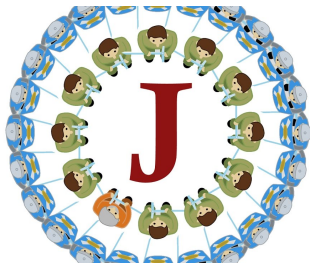
# The Josephus Puzzle



The Josephus Puzzle

$$J(n) = ?$$

# The Josephus Programming Task

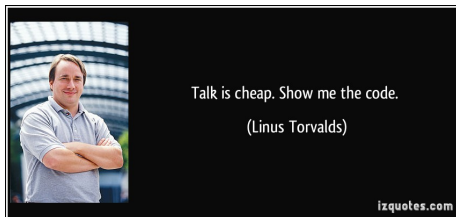


Input:  $n$   
Output:  $J(n)$

Input:  $n$   
Output:  $J(1), J(2), \dots, J(n)$

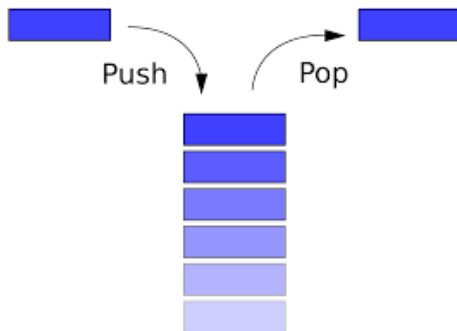
# The Josephus Programming Practice

Q: What data structure do you use?  
WHY?



# Stack

# Stack



PUSH, POP, EMPTY, PEEK

# Brackets Matching

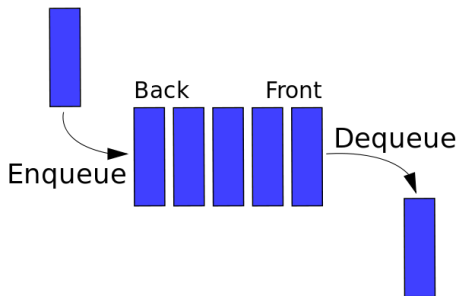
## Brackets Matching

判断给定字符串中的括号是否匹配。

- 输入格式:
- ▶ 首行是一个正整数 (记为  $n$ )。
  - ▶ 接着是  $n$  行字符串。
    - ▶ 每个字符串最多含有  $(, [, \{, ), ], \}$  六种不同字符。
    - ▶ 字符串可以为空。规定空字符串是“括号匹配的”。

输出格式: 如果某行字符串中的括号是匹配的, 则对应行输出 1, 否则输出 0。

# Queue



ENQUEUE, DEQUEUE



# Stutter

## Stutter

- ▶ Given a queue of integers
- ▶ To replace every element with two copies of itself

$$\{1, 2, 3\} \rightarrow \{1, 1, 2, 2, 3, 3\}$$

# Binary Search Tree

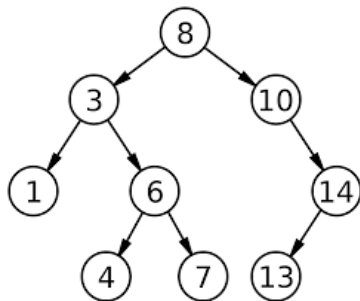
# Binary Search Tree

## Definition (BST)

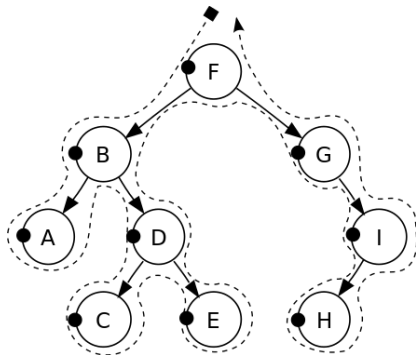
A binary search tree is a rooted binary tree,

1. each internal node stores a key/value
2. each internal node has two distinguished subtrees
  - left subtree** the key in each node must be  $\geq$  any key stored in the left subtree
  - right subtree** the key in each node must be  $\leq$  any key stored in the right subtree

# BST

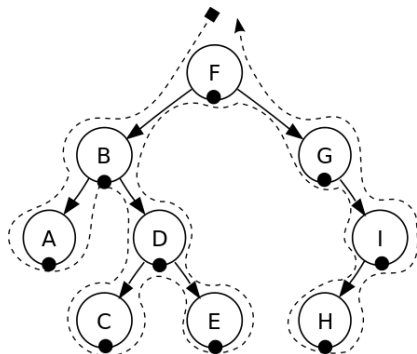


# Preorder Traversal



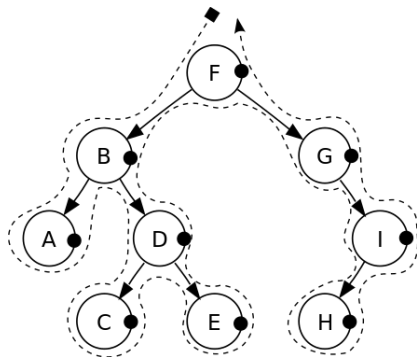
1. Check if the current node is a leaf.
2. Display the data part of the root (or current node).
3. Traverse the left subtree by **recursively** calling the preorder function.
4. Traverse the right subtree by **recursively** calling the preorder function.

# Inorder Traversal



1. Check if the current node is a leaf.
2. Traverse the left subtree by **recursively** calling the inorder function.
3. Display the data part of the root (or current node).
4. Traverse the right subtree by **recursively** calling the inorder function.

# Postorder Traversal



1. Check if the current node is a leaf.
2. Traverse the left subtree by **recursively** calling the postorder function.
3. Traverse the right subtree by **recursively** calling the postorder function.
4. Display the data part of the root (or current node).

## DH 2.16: Treesort

- (i) Construct an algorithm that transforms a given list of integers into a binary search tree.

```
procedure put-x-into-BST (t):  
    ... call put-x-into-BST (t's left subtree)  
    ... call put-x-into-BST (t's right subtree)  
end procedure
```



## DH 2.16: Treesort

- (i) Construct an algorithm that transforms a given list of integers into a binary search tree.

**Node:**

```
int val = NIL,  
Node left = NULL,  
Node right = NULL
```

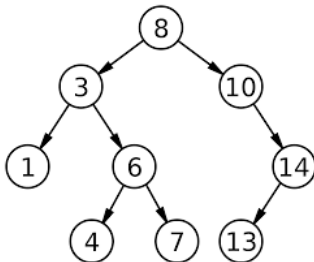
**buildBST**(int eles[]):  
Node root(eles[0])

```
foreach e ∈ eles[1..]:  
    insert(root, e)
```

```
insert(Node T, int e):  
    if (e < T.val)  
        if (T.left == NULL)  
            T.left = new Node(e)  
        else  
            insert(T.left, e)  
    else // e >= T.val  
        if (T.right == NULL)  
            T.right = new Node(e)  
        else  
            insert(T.right, e)
```

## DH 2.16: Treesort

(ii) right; val; left



Thank  
You!