

# 1-12 Partial Order and Lattice

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### SM Problem 14.44: Consistent Enumerations

Suppose the following are three consistent enumerations of an ordered set  $A = \{a, b, c, d\}$ :

$$A_1 : \quad a \quad \textcolor{red}{b} \quad \textcolor{red}{c} \quad d$$

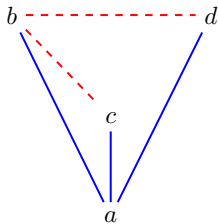
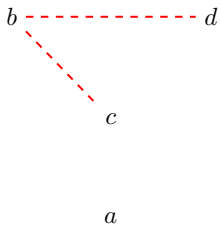
$$A_2 : \quad a \quad c \quad b \quad d$$

$$A_3 : \quad a \quad c \quad \textcolor{blue}{d} \quad \textcolor{blue}{b}$$

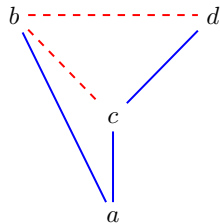
Assuming the Hasse diagram  $D$  of  $A$  is **connected**, draw  $D$ .

$$b \prec_{A_1} c \wedge c \prec_{A_2} b \implies b \parallel_A c$$

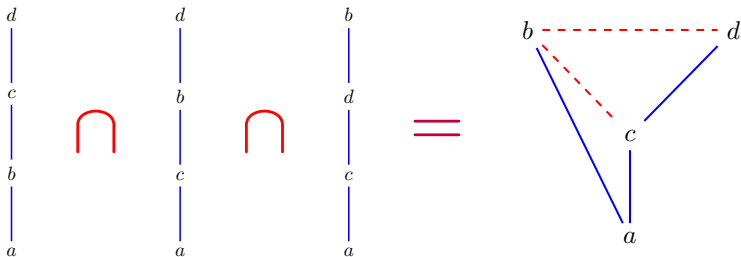
$$d \prec_{A_3} b \wedge b \prec_{A_2} d \implies b \parallel_A d$$



$$\# = 6$$



$$\# = 3$$



## Theorem

Every partial ordering on a set  $X$  is the *intersection* of the total orders on  $X$  *containing* it.

## Definition

Lattice A *lattice* is an algebra  $\mathcal{L} = (L, \wedge, \vee)$  satisfying,

$$\forall a, b, c \in L,$$

Idempotency:

$$a \wedge a = a \quad a \vee a = a$$

Commutativity:

$$a \wedge b = b \wedge a \quad a \vee b = b \vee a$$

Associativity:

$$(a \wedge b) \wedge c = a \wedge (b \wedge c) \quad (a \vee b) \vee c = a \vee (b \vee c)$$

Absorption:

$$a \wedge (a \vee b) = a \quad a \vee (a \wedge b) = a$$

$$a \wedge (a \vee b) = a \quad a \vee (a \wedge b) = a$$

(1) Very useful in lattice computations

$$a \wedge a = a \wedge (a \vee (a \wedge b)) = a$$

(2) The only laws connecting  $\wedge$  and  $\vee$

$\wedge$ -semilattice       $\vee$ -semilattice

(3) Ensure that  $\wedge$  and  $\vee$  induce the same order on  $L$

$$a \leq b \iff a \wedge b = a$$

$$a \leq b \iff a \vee b = b$$

$$a \wedge b = a \iff a \vee b = b$$

## SM Problem 14.72: “Weak” Distributive Laws

Prove that for any lattice  $L$ :

$$a \vee (b \wedge c) \leq (a \vee b) \wedge (a \vee c)$$

$$a \wedge (b \vee c) \geq (a \wedge b) \vee (a \wedge c)$$

$$a \vee (b \wedge c) \leq a \vee b$$

$$a \vee (b \wedge c) \leq a \vee c$$

$$a \leq b$$

$$c \leq d$$

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$$(a \vee c) \leq (b \vee d)$$

## 2017-1-final-exam (5): Lattice

假设  $(L, \leq)$  是格。

如果以下模律 (modular law) 成立, 则称  $L$  是模格 (modular lattice):

$$\forall x \in L : a \leq b \implies a \vee (x \wedge b) = (a \vee x) \wedge b.$$

以下均假设  $L$  是模格。

$$\text{vs. } a \vee (x \wedge b) = (a \vee x) \wedge (a \vee b)$$

Distributive Law  $\implies$  Modular Law



## 2017-1-final-exam (5): Lattice

$$\forall x \in L : a \leq b \implies a \vee (x \wedge b) = (a \vee x) \wedge b.$$

(1) 请证明模律与以下条件等价:

$$\forall x \in L : a \leq b \implies a \vee (x \wedge b) \geq (a \vee x) \wedge b.$$

$$\begin{aligned} & \forall x \in L : a \leq b \\ & \implies \\ & \left( (a \vee (x \wedge b) = (a \vee x) \wedge b) \iff (a \vee (x \wedge b) \geq (a \vee x) \wedge b) \right). \end{aligned}$$

$$a \leq b \implies a \vee (x \wedge b) \leq (a \vee x) \wedge b$$

$$a \leq b \implies a \vee (x \wedge b) \leq (a \vee x) \wedge b$$

$$(a \leq a \vee x) \wedge (a \leq b) \implies a \leq (a \vee x) \wedge b \quad (1)$$

$$(x \leq a \vee x) \wedge b \leq b \implies (x \wedge b) \leq (a \vee x) \wedge b \quad (2)$$

## 2017-1-final-exam (5): Lattice

$$\forall x \in L : a \leq b \implies a \vee (x \wedge b) = (a \vee x) \wedge b.$$

(2) 请证明:  $\forall a, b, c \in L$ ,

如果  $c \leq a$ ,  $a \wedge b = c \wedge b$ ,  $a \vee b = c \vee b$  成立, 则  $a = c$ .

Modular Law :  $[a \leftarrow c] \quad [b \leftarrow a]$

$$\forall x \in L : c \leq a \implies c \vee (x \wedge a) = (c \vee x) \wedge a.$$

$$[x := b]$$

$$c \leq a \implies c \vee (b \wedge a) = (c \vee b) \wedge a$$

$$c \vee (c \wedge b) = (a \vee b) \wedge a$$

$$c = a$$

## 2017-1-final-exam (5): Lattice

$$\forall x \in L : a \leq b \implies a \vee (x \wedge b) = (a \vee x) \wedge b.$$

(3) 给定任意元素  $s, t \in L$ , 且  $s \leq t$ , 构造集合 (称为区间 (interval)):

$$[s, t] \triangleq \{x \in L \mid s \leq x \leq t\}.$$

请证明  $([s, t], \leq)$  是  $L$  的子格 (sublattice)。



$$a, b \in [s, t] \implies a \vee b, a \wedge b \in [s, t]$$

## 5. 格 (Lattice)

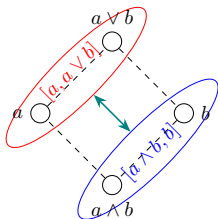
$$\forall x \in L : a \leq b \implies a \vee (x \wedge b) = (a \vee x) \wedge b.$$

(4) 给定任意元素  $a, b \in L$ , 定义函数

$$\varphi : [a \wedge b, b] \rightarrow [a, a \vee b] \quad \varphi(x) = x \vee a$$

$$\psi : [a, a \vee b] \rightarrow [a \wedge b, b] \quad \psi(y) = y \wedge b$$

请证明  $\varphi$  (类似地,  $\psi$ ) 是从  $[a \wedge b, b]$  到  $[a, a \vee b]$  的同构。



## Definition (Lattice Isomorphism)

$$(L, \vee_L, \wedge_L) \quad (M, \vee_M, \wedge_M)$$

A *lattice isomorphism* from  $L$  to  $M$  is a bijection

$$f : L \xleftrightarrow[\text{onto}]{1-1} M$$

such that  $\forall a, b \in L$ :

$$f(a \vee_L b) = f(a) \vee_M f(b)$$

$$f(a \wedge_L b) = f(a) \wedge_M f(b)$$

$f$  preserving  $\vee$  and  $\wedge$ .

$\varphi$  preserving  $\vee$  and  $\wedge$ .

$$\varphi : [a \wedge b, b] \rightarrow [a, a \vee b] \quad \varphi(x) = x \vee a$$

$$\varphi(x_1 \wedge x_2) = \varphi(x_1) \wedge \varphi(x_2)$$

$$\varphi(x_1 \wedge x_2) = (x_1 \wedge x_2) \vee a$$

$$\begin{aligned} \varphi(x_1) \wedge \varphi(x_2) &= (x_1 \vee a) \wedge (x_2 \vee a) \\ &= (a \vee x_1) \wedge (x_2 \vee a) \\ &=_{\text{modular law}} a \vee (x_1 \wedge (x_2 \vee a)) \\ &= \dots \end{aligned}$$

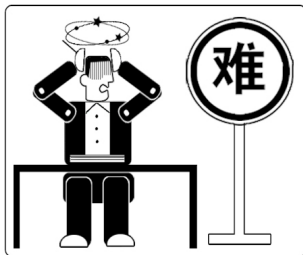




$$\varphi : [a \wedge b, b] \rightarrow [a, a \vee b] \quad \varphi(x) = x \vee a$$

$$\psi : [a, a \vee b] \rightarrow [a \wedge b, b] \quad \psi(y) = y \wedge b$$

$\varphi$  is bijective.



Theorem (UD Theorem 15.8 (iii))

$$f : A \rightarrow B$$

$$\exists g : B \rightarrow A \left( g \circ f = i_A \wedge f \circ g = i_B \right)$$

$$\implies$$

$$f : A \rightarrow B \text{ is bijective} \wedge g = f^{-1}$$

$$\psi \circ \varphi = id_{[a, a \vee b]} \quad \varphi \circ \psi = id_{[a \wedge b, b]}$$

$$(\psi \circ \varphi)(y) = \psi(\varphi(y)) = (y \wedge b) \vee a = a \vee (b \wedge y) = (a \vee b) \wedge y = y$$

$$(\varphi \circ \psi)(x) = \varphi(\psi(x)) = (x \vee a) \wedge b = x \vee (b \wedge a) = x$$

Back to  $\varphi$  preserving  $\vee$  and  $\wedge$ .

$\psi$  preserving  $\wedge$ :

$$\psi(y_1 \wedge y_2) = y_1 \wedge y_2 \wedge b = (y_1 \wedge b) \wedge (y_2 \wedge b) = \psi(y_1) \wedge \psi(y_2)$$

$$\psi(\varphi(x_1) \wedge \varphi(x_2)) = \psi(\varphi(x_1)) \wedge \psi(\varphi(x_2)) = x_1 \wedge x_2$$

$$\varphi(x_1 \wedge x_2) = \varphi(\psi(\varphi(x_1) \wedge \varphi(x_2))) = \varphi(x_1) \wedge \varphi(x_2)$$

Thank  
You!