

2-7 The Algorithmic Methods

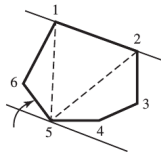
魏恒峰

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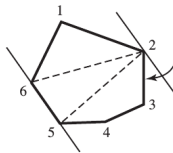
2018 年 05 月 07 日



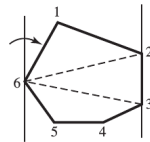
Convex Polygon Diameter



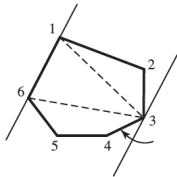
(a)



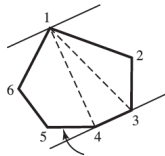
(b)



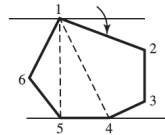
(c)



(d)



(e)



(f)

Convex Polygon Diameter (DH 6.8)

Show that the “Convex Polygon Diameter” algorithm is of **linear-time** complexity.

Q : Linear-time of **WHAT**?

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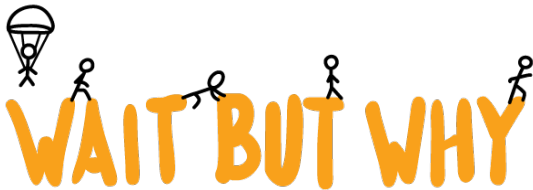
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$$\Theta(c \cdot n) = \Theta(n)$$

Correctness

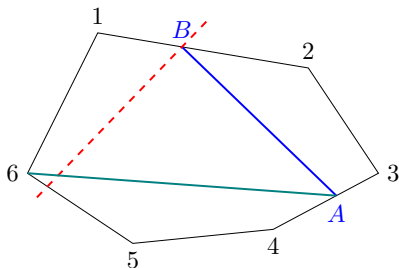


Theorem

For a convex polygon, a pair of vertices determine the diameter.

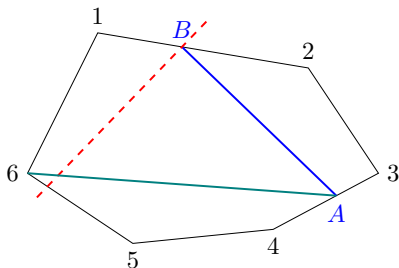
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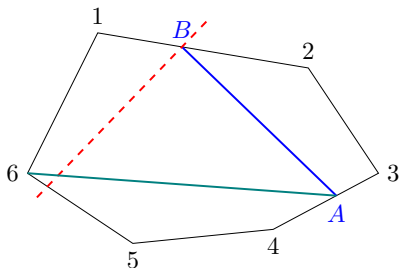
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BUT, we have *not* enumerated *all* pairs of vertices.

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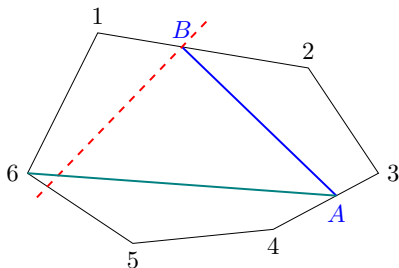


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Theorem

For a convex polygon, a pair of vertices determine the diameter.



BUT, we have *not* enumerated *all* pairs of vertices.

We have enumerated *all* pairs of vertices
that *admits parallel supporting lines*.

Definition (Line of Support)

A line L is a *line of support* of a convex polygon P if

$$L \cap P = \text{a vertex/an edge of } P.$$

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Definition (Antipodal)

An *antipodal* is a pair of points that admits parallel supporting lines.

We have enumerated *all* antipodals.

Theorem

If AB is a diameter of a convex polygon P , then AB is an antipodal.

Theorem

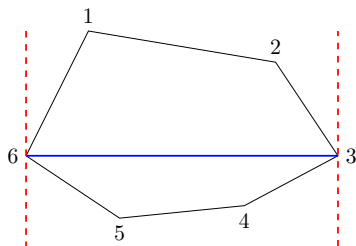
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Proof.

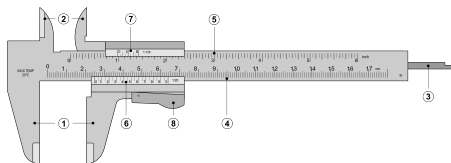
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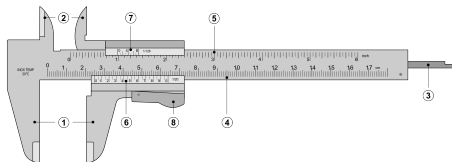
Proof.



Rotating Caliper



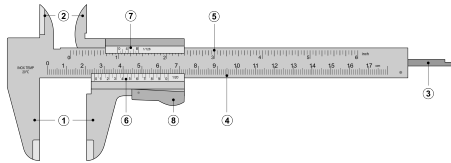
Rotating Caliper



“Computational Geometry”

Ph.D Thesis, Michael Shamos, 1978

Rotating Caliper

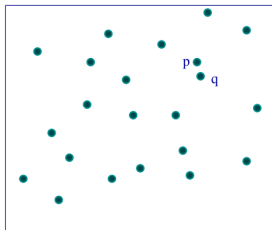


“Computational Geometry”
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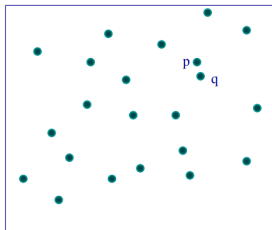


“Solving Geometric Problems with
the Rotating Calipers”, 1983

Finding the Closest Pair of Points

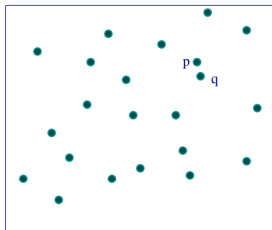


Finding the Closest Pair of Points



A Classic and Beautiful Divide-Conquer Algorithm:

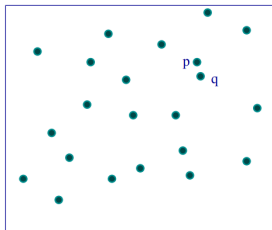
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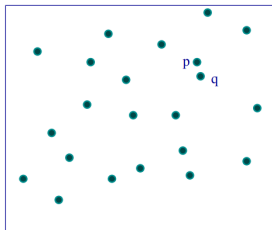
Finding the Closest Pair of Points



A Classic and Beautiful Divide-Conquer-Combine Algorithm:



Finding the Closest Pair of Points



A Classic and Beautiful Divide-Conquer-Combine Algorithm:



Section 33.4, CLRS

Thank
You!



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