1-12 Partial Order and Lattice

魏恒峰

hfwei@nju.edu.cn

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SM Problem 14.44: Consistent Enumerations

Suppose the following are three consistent enumerations of an ordered set $A = \{a, b, c, d\}$:

 $A_1: \ a \ b \ c \ d$ $A_2: \ a \ c \ b \ d$ $A_3: \ a \ c \ d \ b$

Assuming the Hasse diagram D of A is connected, draw D.

SM Problem 14.44: Consistent Enumerations

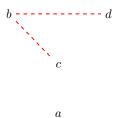
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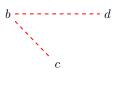
$$A_1: \ a \ b \ c \ d$$
 $A_2: \ a \ c \ b \ d$
 $A_3: \ a \ c \ d \ b$

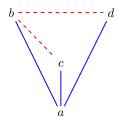
Assuming the Hasse diagram D of A is connected, draw D.

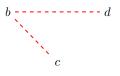
$$b \prec_{A_1} c \land c \prec_{A_2} b \implies b \parallel_A c$$
$$d \prec_{A_2} b \land b \prec_{A_2} d \implies b \parallel_A d$$

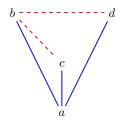




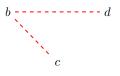


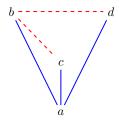




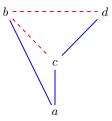


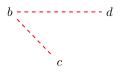
$$\# = 6$$

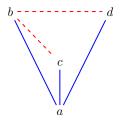




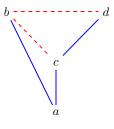
$$\# = 6$$





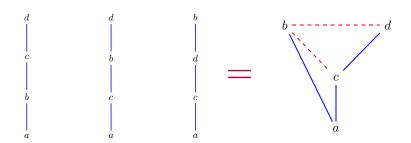


$$\# = 6$$

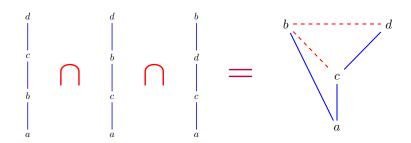


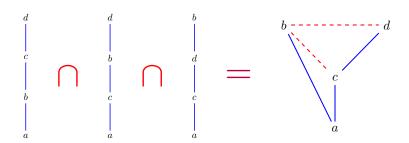
$$\# = 3$$

3 / 20



4/20





Theorem

Every partial ordering on a set X is the intersection of the total orders on X containing it.

Definition

Lattice A lattice is an algebra $\mathcal{L} = (L, \wedge, \vee)$ satisfying,

$$\forall a, b, c \in L$$
,

Idempotency:

$$a \wedge a = a \quad a \vee a = a$$

Commutativity:

$$a \wedge b = b \wedge a \quad a \vee b = b \vee a$$

Associativity:

$$(a \wedge b) \wedge c = a \wedge (b \wedge c) \quad (a \vee b) \vee c = a \vee (b \vee c)$$

Absorption:

$$a \wedge (a \vee b) = a \quad a \vee (a \wedge b) = a$$

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(2) The only laws connecting \wedge and \vee

$$a \land (a \lor b) = a \quad a \lor (a \land b) = a$$

- (1) Very useful in lattice computations
- $a \wedge a = a \wedge (a \vee (a \wedge b)) = a$
- (2) The only laws connecting \wedge and \vee

∧-semilattice ∨-semilattice

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(3) Ensure that \wedge and \vee induce the same order on L



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$$a \le b \iff a \land b = a$$

$$a \le b \iff a \lor b = b$$

$$a \land (a \lor b) = a \quad a \lor (a \land b) = a$$

$$a \wedge a = a \wedge (a \vee (a \wedge b)) = a$$

(2) The only laws connecting \wedge and \vee

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(3) Ensure that \wedge and \vee induce the same order on L

$$a \le b \iff a \land b = a$$
$$a \le b \iff a \lor b = b$$

$$a \wedge b = a \iff a \vee b = b$$



SM Problem 14.72: "Weak" Distributive Laws

Prove that for any lattice L:

$$a \vee (b \wedge c) \leq (a \vee b) \wedge (a \vee c)$$

$$a \wedge (b \vee c) \geq (a \wedge b) \vee (a \wedge c)$$

SM Problem 14.72: "Weak" Distributive Laws

Prove that for any lattice L:

$$a \lor (b \land c) \le (a \lor b) \land (a \lor c)$$
$$a \land (b \lor c) \ge (a \land b) \lor (a \land c)$$

$$a \lor (b \land c) \le a \lor b$$
$$a \lor (b \land c) \le a \lor c$$

SM Problem 14.72: "Weak" Distributive Laws

Prove that for any lattice L:

$$a \lor (b \land c) \le (a \lor b) \land (a \lor c)$$
$$a \land (b \lor c) \ge (a \land b) \lor (a \land c)$$

$$a \lor (b \land c) \le a \lor b$$
$$a \lor (b \land c) \le a \lor c$$

$$a \le b$$

$$c \le d$$

$$(a \lor c) \le (b \lor d)$$

假设 (L, \leq) 是格。

如果以下模律 (modular law) 成立, 则称 L 是模格 (modular lattice):

$$\forall x \in L : a \le b \implies a \lor (x \land b) = (a \lor x) \land b.$$

以下均假设 L 是模格。

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vs.
$$a \lor (x \land b) = (a \lor x) \land (a \lor b)$$

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Distributive Law \implies Modular Law

$$\forall x \in L : a \le b \implies a \lor (x \land b) = (a \lor x) \land b.$$

(1) 请证明模律与以下条件等价:

$$\forall x \in L : a \le b \implies a \lor (x \land b) \ge (a \lor x) \land b.$$

$$\forall x \in L : a \le b \implies a \lor (x \land b) = (a \lor x) \land b.$$

(1) 请证明模律与以下条件等价:

$$\forall x \in L : a \le b \implies a \lor (x \land b) \ge (a \lor x) \land b.$$

$$\Longrightarrow \Big(\big(a \lor (x \land b) = (a \lor x) \land b \big) \iff \big(a \lor (x \land b) \ge (a \lor x) \land b \big) \Big).$$

 $\forall x \in L : \mathbf{a} \leq \mathbf{b}$

$$\forall x \in L : a \le b \implies a \lor (x \land b) = (a \lor x) \land b.$$

(1) 请证明模律与以下条件等价:

$$\forall x \in L : a \le b \implies a \lor (x \land b) \ge (a \lor x) \land b.$$

$$\Longrightarrow \Big(\big(a \vee (x \wedge b) = (a \vee x) \wedge b \big) \iff \big(a \vee (x \wedge b) \ge (a \vee x) \wedge b \big) \Big).$$

 $\forall x \in L : a \leq b$

$$a \le b \implies a \lor (x \land b) \le (a \lor x) \land b$$

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$$a \le b \implies a \lor (x \land b) \le (a \lor x) \land b$$

$$(a \le a \lor x) \land (a \le b) \implies a \le (a \lor x) \land b \tag{1}$$

$$(x \le a \lor x) \land b \le b \implies (x \land b) \le (a \lor x) \land b \tag{2}$$

$$\forall x \in L : a \le b \implies a \lor (x \land b) = (a \lor x) \land b.$$

$$\forall x \in L : a \le b \implies a \lor (x \land b) = (a \lor x) \land b.$$

$$\text{Modular Law}: [a \leftarrow c] \quad [b \leftarrow a]$$

$$\forall x \in L : c \le a \implies c \lor (x \land a) = (c \lor x) \land a.$$

$$\forall x \in L : a \le b \implies a \lor (x \land b) = (a \lor x) \land b.$$

$$\text{Modular Law}: [a \leftarrow c] \quad [b \leftarrow a]$$

$$\forall x \in L : c \le a \implies c \lor (x \land a) = (c \lor x) \land a.$$

$$[x := b]$$

$$c \le a \implies c \lor (\mathbf{b} \land a) = (c \lor \mathbf{b}) \land a$$

$$\forall x \in L : a \le b \implies a \lor (x \land b) = (a \lor x) \land b.$$

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$$c \lor (c \land b) = (a \lor b) \land a$$

$$c = a$$

$$\forall x \in L : a \leq b \implies a \vee (x \wedge b) = (a \vee x) \wedge b.$$

(3) 给定任意元素 $s,t \in L,$ 且 $s \le t,$ 构造集合 (称为区间 (interval)):

$$[s,t] \triangleq \{x \in L \mid s \leq x \leq t\}.$$

请证明 $([s,t],\leq)$ 是 L 的子格 (sublattice)。

2017-1-final-exam (5): Lattice

$$\forall x \in L : a \le b \implies a \lor (x \land b) = (a \lor x) \land b.$$

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 $a, b \in [s, t] \implies a \lor b, a \land b \in [s, t]$

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5. 格 (Lattice)

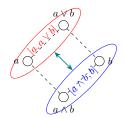
$$\forall x \in L : a \le b \implies a \lor (x \land b) = (a \lor x) \land b.$$

(4) 给定任意元素 $a,b \in L$, 定义函数

$$\varphi:[a\wedge b,b]\to [a,a\vee b]\quad \varphi(x)=x\vee a$$

$$\psi: [a, a \lor b] \to [a \land b, b] \quad \psi(y) = y \land b$$

请证明 φ (类似地, ψ) 是从 $[a \land b, b]$ 到 $[a, a \lor b]$ 的同构。



Definition (Lattice Isomorphism)

$$(L, \vee_L, \wedge_L)$$
 (M, \vee_M, \wedge_M)

A *lattice isomorphism* from L to M is a bijection

$$f: L \stackrel{1-1}{\longleftrightarrow} M$$

such that $\forall a, b \in L$:

$$f(a \vee_L b) = f(a) \vee_M f(b)$$

$$f(a \wedge_L b) = f(a) \wedge_M f(b)$$

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f preserving \vee and \wedge .



 φ preserving \vee and \wedge .

$$\varphi: [a \wedge b, b] \to [a, a \vee b] \quad \varphi(x) = x \vee a$$

 φ preserving \vee and \wedge .

$$\varphi:[a\wedge b,b]\to [a,a\vee b]\quad \varphi(x)=x\vee a$$

$$\varphi(x_1 \wedge x_2) = \varphi(x_1) \wedge \varphi(x_2)$$

φ preserving \vee and \wedge .

$$\varphi : [a \land b, b] \to [a, a \lor b] \quad \varphi(x) = x \lor a$$

$$\varphi(x_1 \land x_2) = \varphi(x_1) \land \varphi(x_2)$$

$$\varphi(x_1 \land x_2) = (x_1 \land x_2) \lor a$$

$$\varphi(x_1) \land \varphi(x_2) = (x_1 \lor a) \land (x_2 \lor a)$$

$$= (a \lor x_1) \land (x_2 \lor a)$$

$$=_{\text{modular law}} a \lor (x_1 \land (x_2 \lor a))$$



$$\varphi: [a \wedge b, b] \to [a, a \vee b] \quad \varphi(x) = x \vee a$$

$$\psi: [a, a \vee b] \to [a \wedge b, b] \quad \psi(y) = y \wedge b$$

$$\varphi \text{ is bijective.}$$

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$$\psi: [a, a \vee b] \to [a \wedge b, b] \quad \psi(y) = y \wedge b$$

 φ is bijective.



Theorem (UD Theorem 15.8 (iii))

$$f: A \to B$$

$$\exists g: B \to A \ \Big(g \circ f = i_A \land f \circ g = i_B \Big)$$

$$f: A \to B$$
 is bijective $\land q = f^{-1}$

$$\psi \circ \varphi = id_{[a,a \vee b]} \qquad \varphi \circ \psi = id_{[a \wedge b,b]}$$

$$\psi \circ \varphi = id_{[a,a \vee b]} \qquad \varphi \circ \psi = id_{[a \wedge b,b]}$$

$$(\psi \circ \varphi)(y) = \psi(\varphi(y)) = (y \land b) \lor a = a \lor (b \land y) = (a \lor b) \land y = y$$

$$\psi \circ \varphi = id_{[a,a \lor b]} \qquad \varphi \circ \psi = id_{[a \land b,b]}$$

$$(\psi \circ \varphi)(y) = \psi(\varphi(y)) = (y \wedge b) \vee a = a \vee (b \wedge y) = (a \vee b) \wedge y = y$$

$$(\varphi \circ \psi)(x) = \varphi(\psi(x)) = (x \lor a) \land b = x \lor (b \land a) = x$$

 ψ preserving \wedge :

$$\psi(y_1 \wedge y_2) = y_1 \wedge y_2 \wedge b = (y_1 \wedge b) \wedge (y_2 \wedge b) = \psi(y_1) \wedge \psi(y_2)$$

 ψ preserving \wedge :

$$\psi(y_1 \wedge y_2) = y_1 \wedge y_2 \wedge b = (y_1 \wedge b) \wedge (y_2 \wedge b) = \psi(y_1) \wedge \psi(y_2)$$

$$\psi(\varphi(x_1) \wedge \varphi(x_2)) = \psi(\varphi(x_1)) \wedge \psi(\varphi(x_2)) = x_1 \wedge x_2$$

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$$\psi(\varphi(x_1) \wedge \varphi(x_2)) = \psi(\varphi(x_1)) \wedge \psi(\varphi(x_2)) = x_1 \wedge x_2$$

$$\varphi(x_1 \wedge x_2) = \varphi(\psi(\varphi(x_1) \wedge \varphi(x_2))) = \varphi(x_1) \wedge \varphi(x_2)$$

Thank You!