

# 1-5 数据与数据结构 (I)

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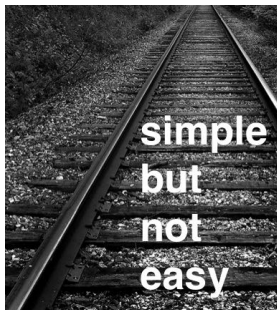


# Permutations

# Permutations

Generating All Permutations  
Stackable/Queueable Permutations

## Generating All Permutations



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For  $\dots$ :  $\dots$

Then, # of perms is

$$n \times (n - 1) \times \dots \times 1 = n!$$



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$$\underbrace{(n+1)}_{\text{1st choice}} \times \underbrace{n!}_{\text{I.H.}} = (n+1)!$$



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Design an algorithm which, given a positive integer  $n$ , generates/[prints](#) all the permutations of  $[0 \cdots n)$ .

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```
void perms (A[], n) {  
    if (n == 1)  
        print 'A[0] '  
    else  
        for (int i = 0; i < n; ++i)  
            print 'A[i] '  
            perms(A ← A \ A[i], n - 1)  
            print '\\n'  
}
```

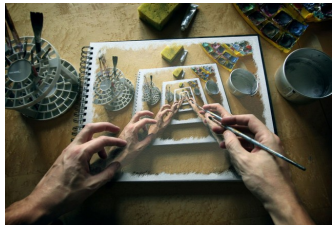
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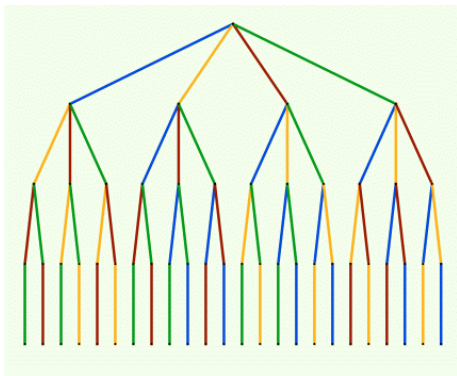
generate-perms.c







$$A = [0, 1, 2, 3] \quad n = 4$$



```

void perms (prifix, A[], n) {
    if (n == 1)
        print 'prifix ++ A[0]'
    else
        for (int i = 0; i < n; ++i)
            perms(prefix ← prefix ++ A[i],
                A ← A \ A[i], n - 1)
            print '\n'
}

```

```

void perms (prefix, A[], n) {
    if (n == 1)
        print ' 'prefix ++ A[0] ' '
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```
perms(' ', A, n);
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        for (int i = 0; i < n; ++i)
            perms(prefix ← prefix ++ A[i],
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        print ' '\n'
}

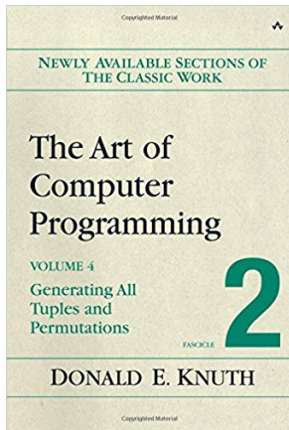
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# For more about “Generating All Permutations”:



## DH 2.10: Permutation Checking

- ▶ An integer  $n$
- ▶ An array of integers  $P$  of length  $n$

To check whether  $P$  is a permutation of  $1 \cdots n$ ?

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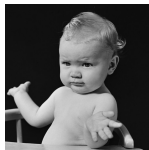
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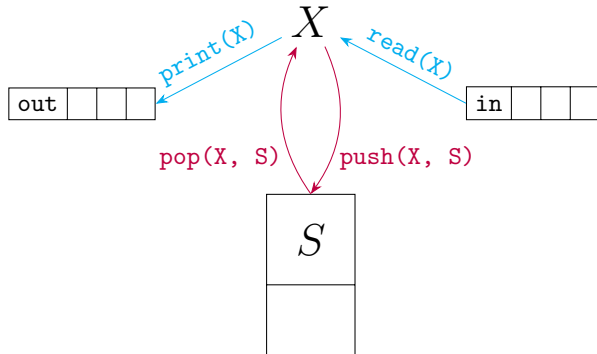


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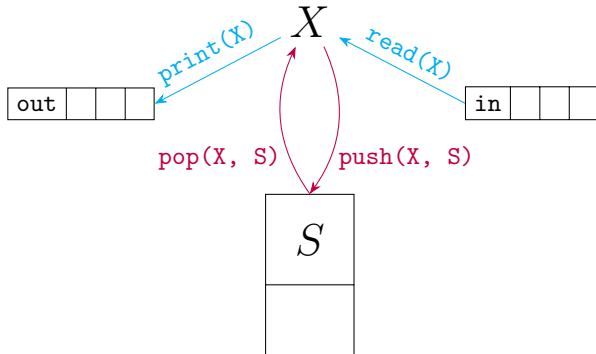


## Definition (Stackable Permutations)



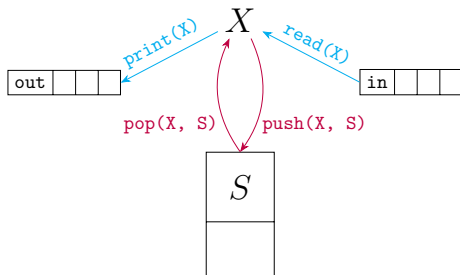
## Definition (Stackable Permutations)

$$\text{out} = (a_1, \dots, a_n) \xleftarrow[X=0]{S=\emptyset} \text{in} = (1, \dots, n)$$

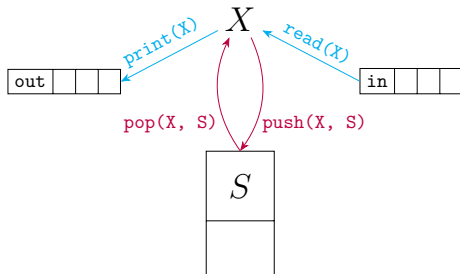




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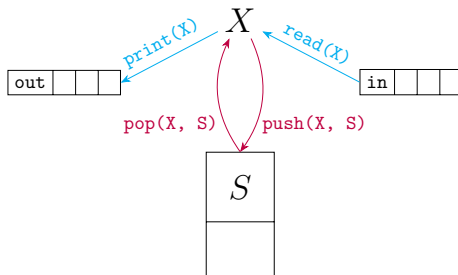


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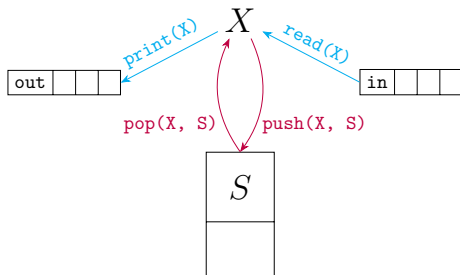


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$$a == X$$

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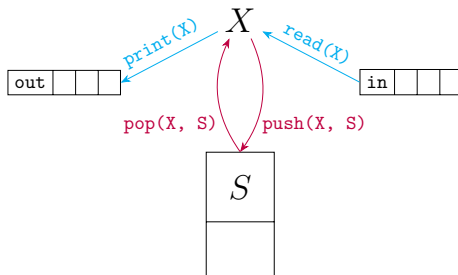


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$$a == X \quad a > X \ (a < X) \quad \text{top}(S)$$

## DH 2.12: Stackable Permutations

(a) **Show** that the following permutations *are* stackable:

(i)  $(3, 2, 1)$

(ii)  $(3, 4, 2, 1)$

(iii)  $(3, 5, 7, 6, 8, 4, 9, 2, 10, 1)$

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To check whether a given permutation can be obtained by a stack.

read   print   push   pop   is-empty

$X = 0$     $S = \emptyset$    in  $\neq$  EOF

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foreach 'a' in out:
    if (! is-empty(S)
        && 'a' == top(S))
        pop(S, X)
        print(X)
        continue
    else ... // T.B.C
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else // T.B.C
    while (in  $\neq$  EOF)
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        if (X == 'a')
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ERR
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(b) **Prove** that the following permutations are *not* stackable:

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$$\text{out} = \cdots a_i \cdots a_j \cdots a_k : i < j < k \wedge a_j < a_k < a_i$$



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$(\mathbf{3}, \mathbf{1}, \mathbf{2})$

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312-Pattern

## Theorem (Stackable Permutations)

A permutation  $(a_1, \dots, a_n)$  is stackable  $\iff$  it is not the case that

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Proof.



NO PROOF WARRANTY



## DH 2.12: Stackable Permutations

(c) How many permutations of  $A_4$  *cannot* be obtained by a stack?

$(1, 4, 2, 3), (2, 4, 1, 3), (3, 1, 2, 4), (3, 1, 4, 2), (3, 4, 1, 2)$   
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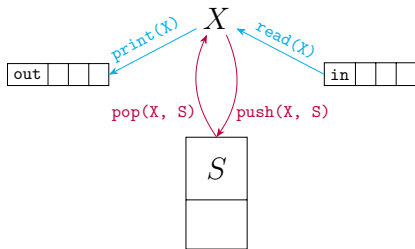
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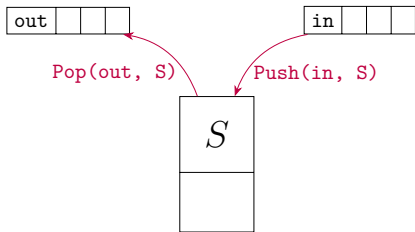
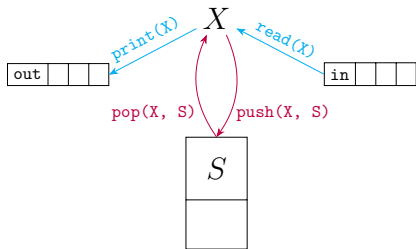
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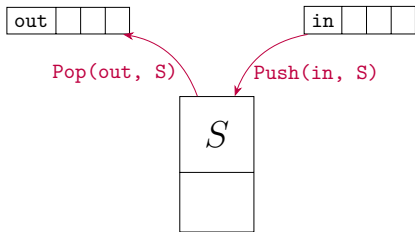
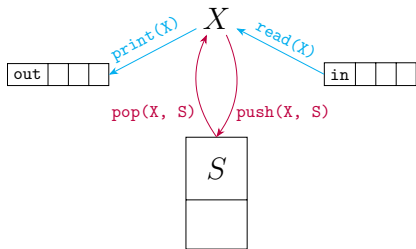
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*Q* : What about  $A_n$ ?

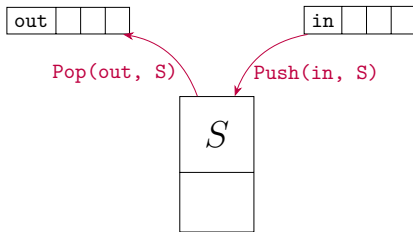
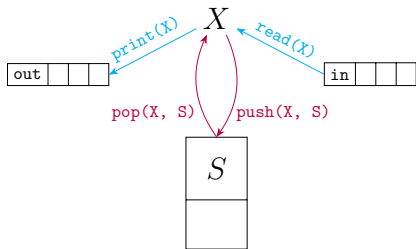




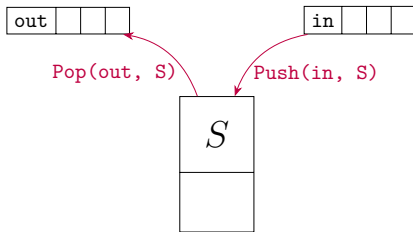
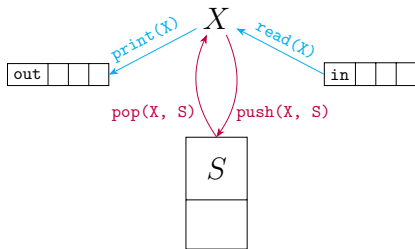




$Q$  : Are  $S + X$  and  $S$  are **equivalent**?

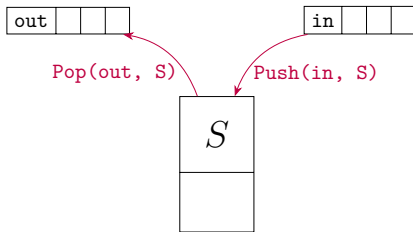
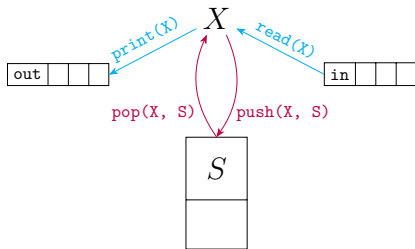


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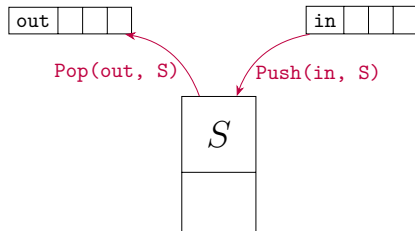
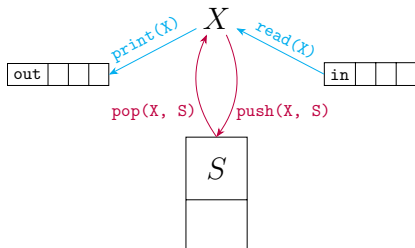
Producing the same set of permutations.

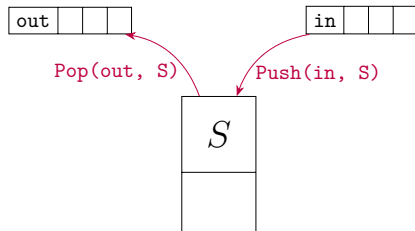
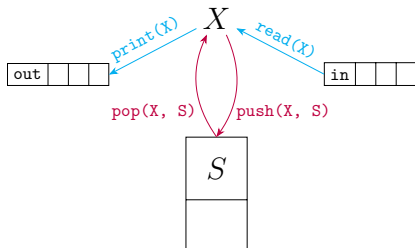


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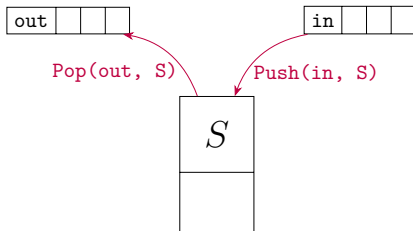
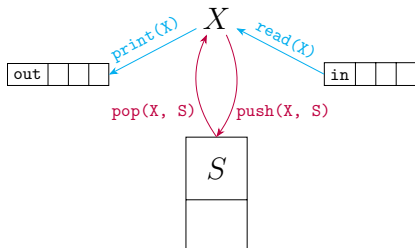
Producing the same set of permutations.

Accepting the same set of *admissible* operation sequences.





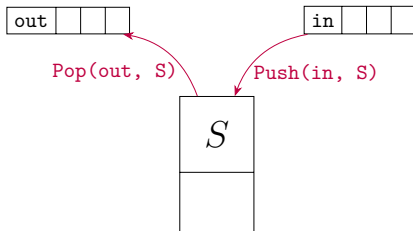
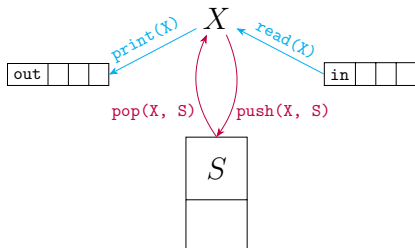
By simulations.



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Simulate  $S$  by  $S + X$ :

- Push
- Pop



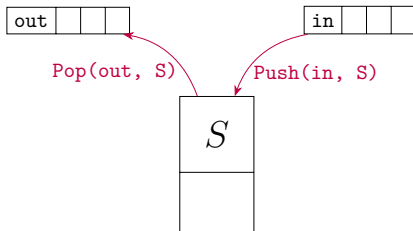
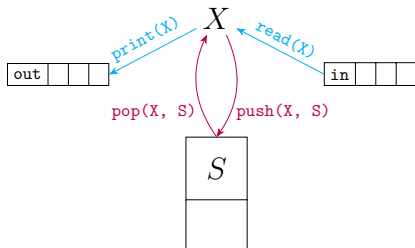
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Simulate  $S + X$  by  $S$ :

By iterative transformations.



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$(3, 2, 5, 6, 1, 4) : + + + - - + + - + - - -$

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## Theorem (Reflection Method)

*The number of stackable permutations is  $\binom{2n}{n} - \binom{2n}{n-1}$ .*

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# Catalan Number



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## Parenthesis

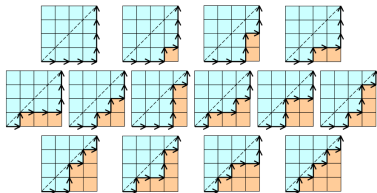
$(3, 2, 1) : ((( )))$        $(1, 2, 3) : ()()()$

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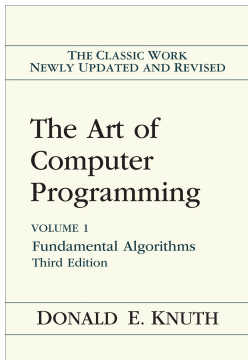
## Parenthesis

$(3, 2, 1) : ((( )))$        $(1, 2, 3) : ()()()$

## Grid Paths Not above the diagonal:



# For more about “Stackable Permutations”:



Thank  
You!