

Core concepts (recap):

- Eval everything recursively
 - eval(AST node, environment) -> string
 - text -> text
 - definitions
 - eval the name and parameters names
 - record that (do not eval the body)
 - invocations
 - look up the invocation name (eval args and name first)
 - create a new environment with param:arg bindings
 - exec the body in that environment
 - parameters {{{...}}}
 - look that up in the environment we just create
 - if we don't find it -> bound to itself (in the sense of the syntactic version of the parameter: "{{{...}}}")

With this, we can already do interesting things. We can do nested calls.

Ordering is implicit [see 14c302](#)

We can pass functions around (in a limited way) [see 14c302-1](#)

Recursive functions [see 14c302-2](#)

Arithmetic and Numbers

- Could define numbers and operators inside our grammar.
- We know we're "living on top of JS"
- We will assume that there is an already defined expr grammar {{ #expr | ... }}
 - The above is a special, always-recognized template
 - Its body is an arithmetic expression.
 - #expr evaluates its argument and return the result [see 14c302-3](#)
 - N.B. this is a glaring security hole since the input can be anything (should sanitize it)
 - Can also include logical operators

Conditionals

- [See 14c302-4](#)
- Notice "true" for us is a non-empty string, and "false" as an empty string.
- We need to treat #if and #ifeq differently:
 - eval the name
 - eval the condition
 - then we choose to evaluate the then or else, we only do one of them.

[See 14c302-5](#)

Scoping and environment

[See 14c302-6](#)

- Notice that we will need to keep track of the declared environment for templates
- For static scoping, this should be the declared environment ->

bar: param -> ... body -> environment ->
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Executing in an environment E0

Bar: param -> ...

Body ->

Environment -> E0

Property of Patrick Ghazal