

# 122com Data structures and types

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# Overview

- 1 Arrays
- 2 Linked lists
  - Array example
  - LL example
- 3 Data structures
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- 5 Queues
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## Arrays

## Linked lists

## Array example

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## End

A series of objects all of the same size and type.

```
char array[] = {'A', 'B', 'C', 'D', 'E'};
```

- Stored in contiguous blocks of memory.
- Python lists are functionally closest.
  - But are not arrays.
- Can't be resized.

The challenger for array's crown.

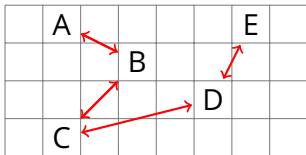
- Series of nodes, each of which points to the next element.
  - And to the previous element if it's a doubly linked list.

Singly linked       $\boxed{A} \rightarrow \boxed{B} \rightarrow \boxed{C} \rightarrow \boxed{D} \rightarrow$

Doubly linked       $\leftarrow \boxed{A} \rightarrow \boxed{B} \leftarrow \boxed{C} \rightarrow \boxed{D} \leftarrow$

Not in contiguous memory.

- Each node is separate.
- Scattered.
- Dynamic memory (pointers!).



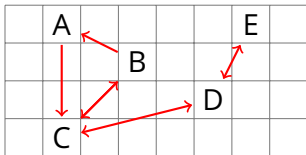
- Why would we use linked lists instead of arrays?
  - Can change size.
  - Can quickly insert and delete elements.

```
class Node:  
    __prev = None  
    __next = None  
    value = None
```

```
class Node  
{  
private:  
    Node *prev;  
    Node *next;  
  
public:  
    int value;  
};
```

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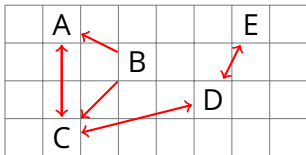
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# Removing array elements

C

```
char array[] = {'A', 'B', 'C', 'D', 'E'};
```

		A	B	C	D	E	

- Array in memory, multiple elements in a contiguous block.
- How do we remove elements from the middle?



## Removing array elements

C

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- Array in memory, multiple elements in a contiguous block.
- How do we remove elements from the middle?
  - 1 Remove element from the array.

## Removing array elements

C

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- Array in memory, multiple elements in a contiguous block.
- How do we remove elements from the middle?
  - 1 Remove element from the array.
  - 2 Move next element to occupy the empty space.

## Removing array elements

C

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  - 3 Repeat.

## Removing array elements

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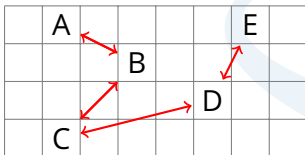
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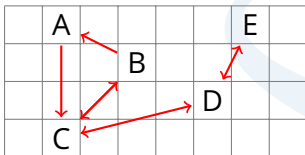
- Array in memory, multiple elements in a contiguous block.
- How do we remove elements from the middle?
  - 1 Remove element from the array.
  - 2 Move next element to occupy the empty space.
  - 3 Repeat.
- Is very slow with large arrays.

# Removing linked list elements



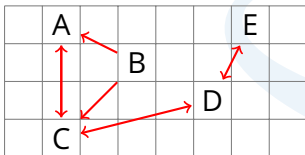
- Linked list, separate elements scattered in memory.
- Each pointing to the next/prev element.
- How do we remove elements?

# Removing linked list elements



- Linked list, separate elements scattered in memory.
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  - 1 Change pointers.

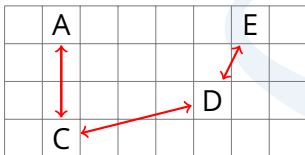
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# Removing linked list elements



- Linked list, separate elements scattered in memory.
- Each pointing to the next/prev element.
- How do we remove elements?
  - 1 Change pointers.
  - 2 Delete old element.

## Advantages

- Inserting and deleting elements is very fast.
  - $O(1)$ .
- No size limits, can keep adding new elements.
- Doesn't waste memory.

## Disadvantages

- Not indexed.
  - Can't ask for the 20<sup>th</sup> element etc.
  - Have to step through the list (slow).
- Needs more memory than an array to store the same number of elements.
  - Have to keep track of where the next/prev nodes are.

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Arrays and linked lists are data structures.

- A specific way of storing data.
- Can see how the various elements of the structure are laid out in memory.
- Direct access to the underlying memory.

# Abstract data types



As we move to storing more complex information in our software we will start to encounter Abstract Data Types (ADTs).

- Software engineering principal.

# Abstract data types



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...and how it does it separate.

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- Keep what a data type can do...  
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- Unlike data structure ADTs only concerned with the interface.

# Abstract data types



As we move to storing more complex information in our software we will start to encounter ADTs.

- Software engineering principal.
- Keep what a data type can do...  
...and how it does it separate.
- Unlike data structure ADTs only concerned with the interface.
- Internals of ADTs can vary widely between implementations.





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Imagine an ADT like a car.

- It has a set of supported operations, go faster, go slower, turn left, turn right.



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- Don't care how it achieves these.



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- Don't care how it achieves these.
- Don't care if, internally, it's using a combustion engine or an electric motor.



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Imagine an ADT like a car.

- It has a set of supported operations, go faster, go slower, turn left, turn right.
- Don't care how it achieves these.
- Don't care if, internally, it's using a combustion engine or an electric motor.
- Only care about the result.

## A First In First Out (FIFO) ADT.

- Ends of the queue called the front and back.
- New elements added to back of queue only.
  - Pushing - push(value)
- Old elements removed from front of queue only.
  - Popping - pop()
- No cutting in.

## A FIFO ADT.

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- Buffer to hold items for processing in the order in which they arrive.

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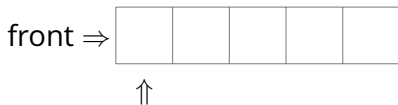
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- Which would be better for a queue? An array or a linked list?

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- 
- Which would be better for a queue? An array or a linked list?
    - Linked list.



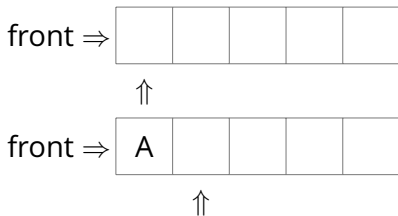
## Array as a queue.



- Very similar to stacks.
  - Keep track of next free space.
  - Limited size.

# Array as a queue.

1

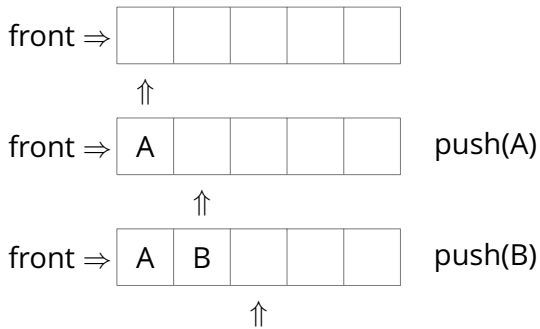


push(A)

- Very similar to stacks.
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# Array as a queue.

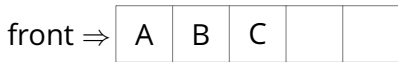
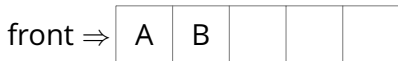
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push(A)

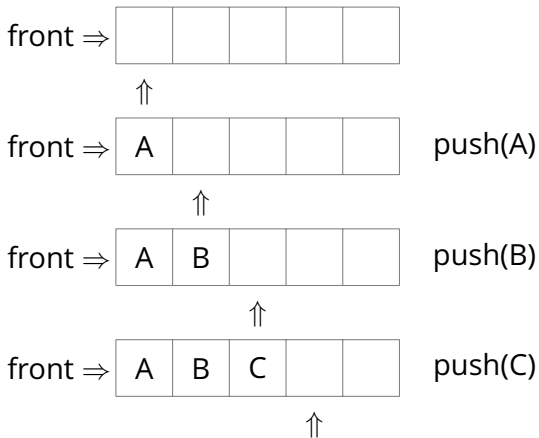
push(B)

push(C)

- Very similar to stacks.
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# Array as a queue.

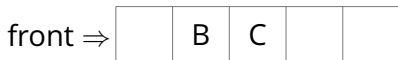
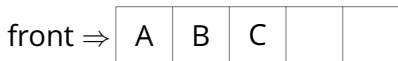
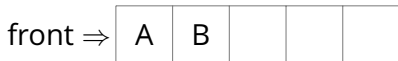
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- Very similar to stacks.
  - Keep track of next free space.
  - Limited size.
- What happens when we pop()?
  - Have to shuffle every element forward one space.
  - Inefficient.

# Array as a queue.

1



push(A)

push(B)

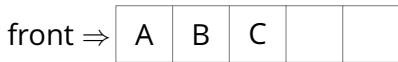
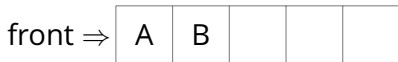
push(C)

pop()

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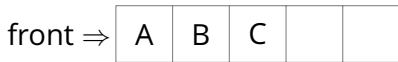
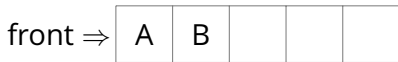
push(C)

pop()

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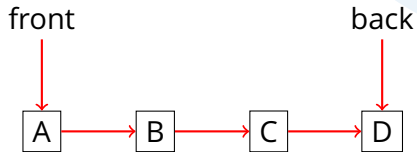
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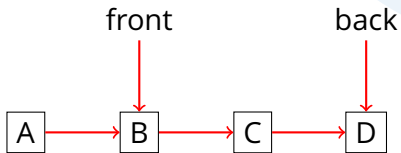
A



# Linked list as a queue.

A

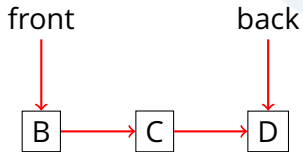
pop()



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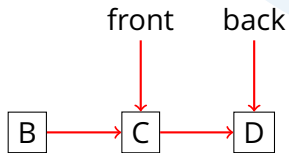
pop()



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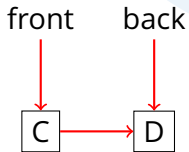
pop() , pop()



# Linked list as a queue.

A

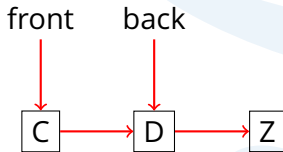
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## Linked list as a queue.

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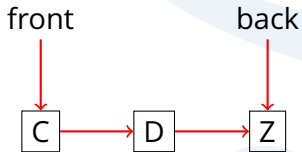
pop() , pop() , push(Z)



## Linked list as a queue.

A

pop() , pop() , push(Z)



## A First In Last Out (FILO) ADT.

- Ends of the stack are called the top and bottom.
- New elements add to top of stack only.
  - Pushing - push(value)
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## A FILO ADT.

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- Which would be better for a stack? An array or a linked list?

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- No cutting in.
  
- Which would be better for a stack? An array or a linked list?
  - Doesn't matter performance wise.
  - Linked list if  $n$  is unknown.

# Array as a stack.



- Keep track of position of the next free space in the array.
- Arrays have a fixed size.
  - Can't hold more values than we have space for.



push(A)

## Array as a stack.

1

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# Array as a stack.

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push(A)

push(B)

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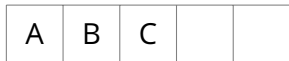
1



push(A)



push(B)



push(C)



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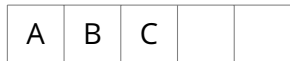
1



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push(B)



push(C)



pop()



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- An unordered ADT.
  - Items ordered by the set.
  - You have no control over it.
- Sets contain unique elements.
  - Can't contain duplicates.
- Can add items to a set.
- Can remove items from a set.
- Can see if an item is in a set.
- Can't get the  $n^{th}$  element.
  - It's unordered remember.





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add(A)

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add(A)



add(B)

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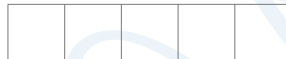
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add(A)



add(B)



add(A)

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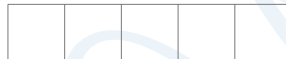
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add(A)



add(B)



add(A)

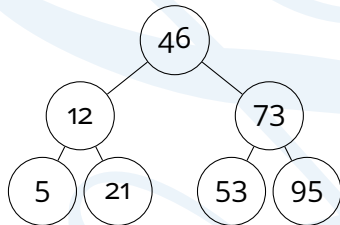


remove(A)

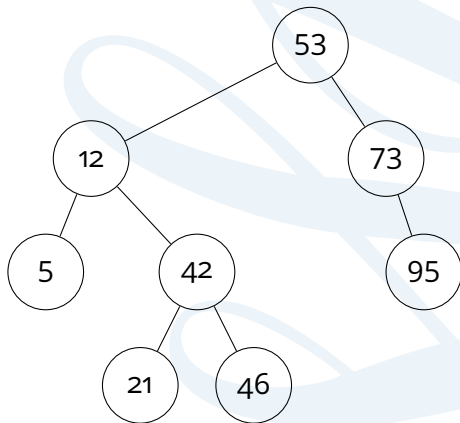
- Lots of other ADTs.
- Different names in different languages.
- Lists.
- Circular lists.
- Associative arrays.
  - Dictionaries/Maps.
- Double-ended queues.
- Trees.
- Graphs.

## Variation on linked lists.

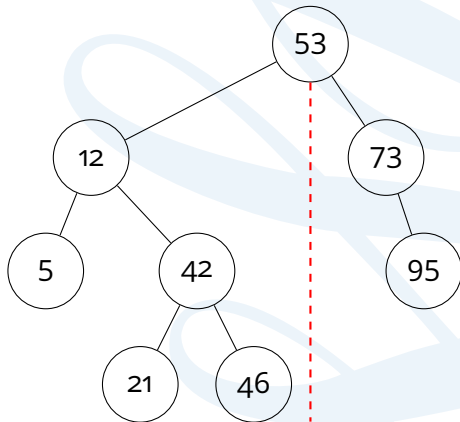
- Made of nodes and relationships.
- Root node at top.
- Each node can have  $\geq 0$  children.
- Binary search tree.
  - Very common type.
  - Ordered.
  - Max two children.
  - Binary searching.
  - Very good for sets.



- Trees can be balanced or unbalanced.
- Not required for all trees.
- Going to be talking about BSTs from here on.
- Unbalanced because more than a one node difference between the two halves.

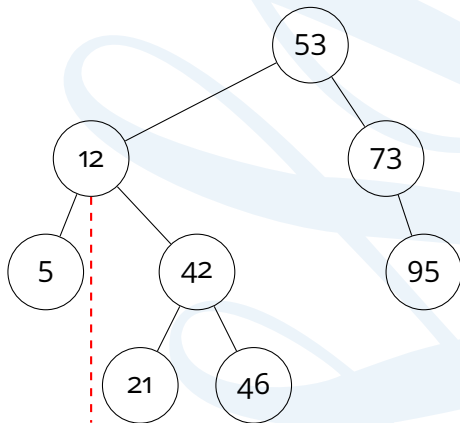


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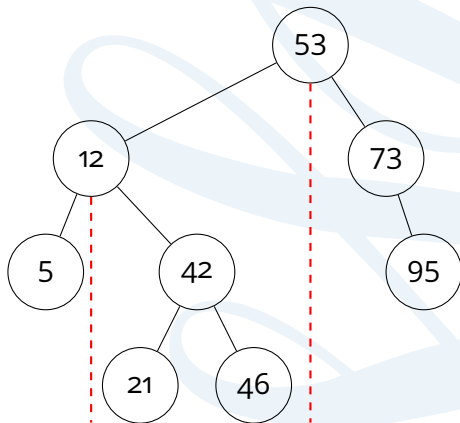




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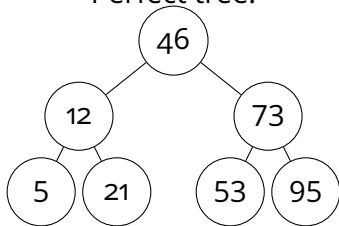


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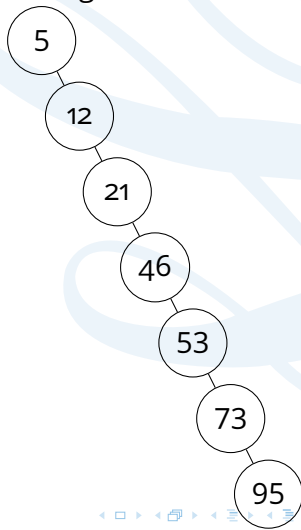


Important that you keep your BSTs balanced.

Perfect tree.



Degenerate tree.



Arrays

Linked lists

Array example

LL example

Data  
structures

Abstract data  
types

Queues

Stacks

Sets

Other

Trees

End

# Quiz

## Recap

- Arrays.

- Advantages/disadvantages.

- Linked lists .

- Advantages/disadvantages.
- How to insert/delete.

- Difference between data structure and ADTs.

- Stack.

- FILO.
- Using an array as one.
- Using a LL as one.

- Queue.

- FIFO.
- Using an array as one.
- Using a LL as one.

- Sets.

- No duplicates.
- Unordered.

- Trees.

- Balanced/unbalanced.

Arrays

Linked lists

Array example

LL example

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structures

Abstract data  
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Sets

Other

Trees

**End**

# The End