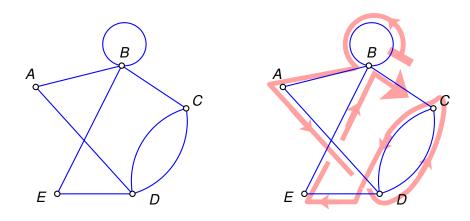
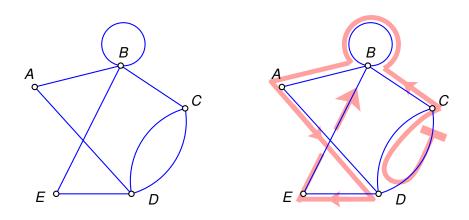
An **Euler path** is a path that uses every **edge** of a graph exactly once.

An **Euler circuit** is a circuit that uses every edge of a graph exactly once.

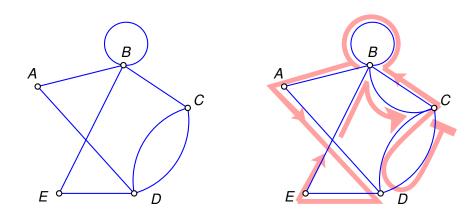
- An Euler path starts and ends at different vertices.
- An Euler circuit starts and ends at the same vertex.



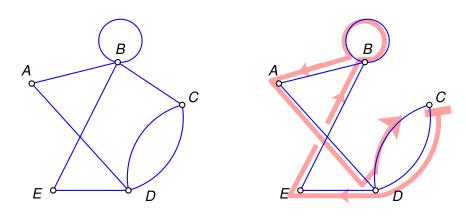
An Euler path: BBADCDEBC



Another Euler path: CDCBBADEB



An Euler circuit: CDCBBADEBC



Another Euler circuit: CDEBBADC

Is it possible to determine whether a graph has an Euler path or an Euler circuit, without necessarily having to find one explicitly?



Suppose that a graph has an Euler path P.

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For every vertex v other than the starting and ending vertices, the path P enters v the same number of times that it leaves v (say \boxed{s} times).

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Therefore, there are 2s edges having v as an endpoint.

Therefore, all vertices other than the two endpoints of P must be even vertices.

Suppose the Euler path P starts at vertex x and ends at y.

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Then *P* leaves *x* one more time than it enters, and leaves *y* one fewer time than it enters.

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Therefore, the two endpoints of P must be odd vertices.

The inescapable conclusion ("based on reason alone!"):

If a graph G has an Euler path, then it must have exactly two odd vertices.

Or, to put it another way,

If the number of odd vertices in G is anything other than 2, then G cannot have an Euler path.

▶ Suppose that a graph *G* has an Euler circuit *C*.

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- ▶ Suppose that a graph G has an Euler circuit C.
- ► For every vertex *v* in *G*, each edge having *v* as an endpoint shows up *exactly once* in *C*.
- ► The circuit *C* enters *v* the same number of times that it leaves *v* (say *s* times), so *v* has degree 2*s*.
- ► That is, v must be an even vertex.

The inescapable conclusion ("based on reason alone"):

If a graph *G* has an Euler circuit, then all of its vertices must be even vertices.

Or, to put it another way,

If the number of odd vertices in G is anything other than O, then G cannot have an Euler circuit.

Things You Should Be Wondering

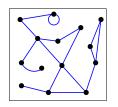
Does every graph with zero odd vertices have an Euler circuit?

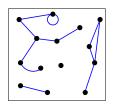
Does every graph with two odd vertices have an Euler path?

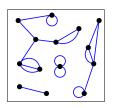
▶ Is it possible for a graph have just **one** odd vertex?

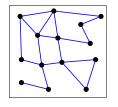


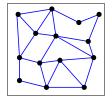
How Many Odd Vertices?

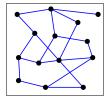




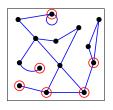


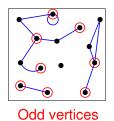


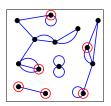


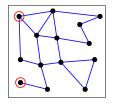


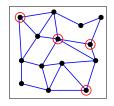
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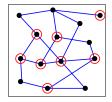




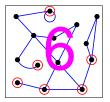


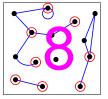


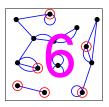




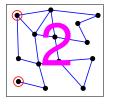
How Many Odd Vertices?



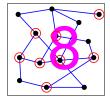




Number of odd vertices







The Handshaking Theorem

The Handshaking Theorem says that

In every graph, the sum of the degrees of all vertices equals twice the number of edges.

If there are n vertices V_1, \ldots, V_n , with degrees d_1, \ldots, d_n , and there are e edges, then

$$d_1 + d_2 + \cdots + d_{n-1} + d_n = 2e$$

Or, equivalently,

$$e = \frac{d_1 + d_2 + \dots + d_{n-1} + d_n}{2}$$

The Handshaking Theorem

Why "Handshaking"?

If n people shake hands, and the i^{th} person shakes hands d_i times, then the total number of handshakes that take place is

$$\frac{d_1+d_2+\cdots+d_{n-1}+d_n}{2}.$$

(How come? Each handshake involves two people, so the number $d_1 + d_2 + \cdots + d_{n-1} + d_n$ counts every handshake twice.)

▶ The number of edges in a graph is

$$\frac{d_1+d_2+\cdots+d_n}{2}$$

which must be an integer.

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- ▶ Therefore, $d_1 + d_2 + \cdots + d_n$ must be an **even number**.
- ▶ Therefore, the numbers d_1, d_2, \dots, d_n must include an **even number of odd numbers**.

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- ▶ Therefore, $d_1 + d_2 + \cdots + d_n$ must be an **even number**.
- ▶ Therefore, the numbers d_1, d_2, \dots, d_n must include an **even number of odd numbers**.
- Every graph has an even number of odd vertices!

Back to Euler Paths and Circuits

Here's what we know so far:

# odd vertices	Euler path?	Euler circuit?
0	No	Maybe
2	Maybe	No
4, 6, 8,	No	No
1, 3, 5,	No such graphs exist!	

Can we give a better answer than "maybe"?

Euler Paths and Circuits — The Last Word

Here is the answer Euler gave:

# odd vertices	Euler path?	Euler circuit?
0	No	Yes*
2	Yes*	No
4, 6, 8,	No	No
1, 3, 5,	No such graphs exist	

^{*} Provided the graph is connected.

Euler Paths and Circuits — The Last Word

Here is the answer Euler gave:

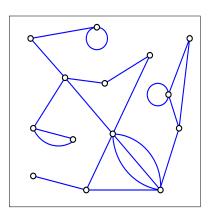
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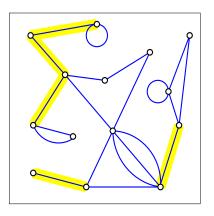
Next question: If an Euler path or circuit exists, how do you find it?

Removing a single edge from a connected graph can make it disconnected. Such an edge is called a **bridge**.

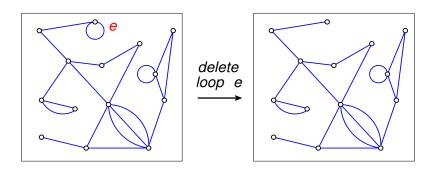
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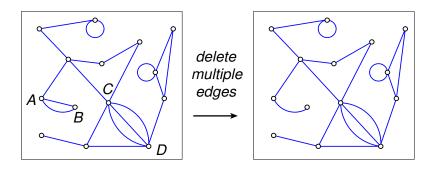


Loops cannot be bridges, because removing a loop from a graph cannot make it disconnected.



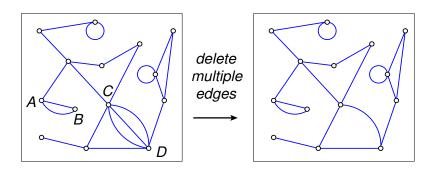
Bridges

If two or more edges share both endpoints, then removing any one of them cannot make the graph disconnected. Therefore, none of those edges is a bridge.



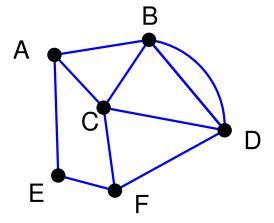
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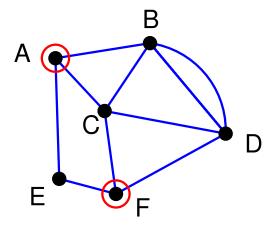


"Don't burn your bridges."

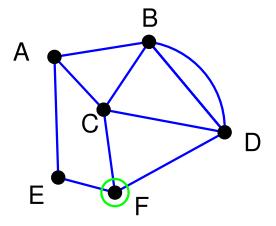
Problem: Find an Euler circuit in the graph below.



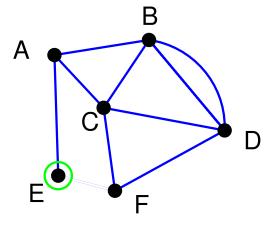
There are two odd vertices, A and F. Let's start at F.



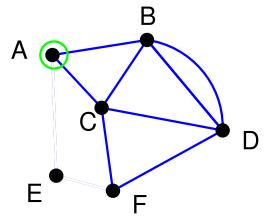
Start walking at F. When you use an edge, delete it.



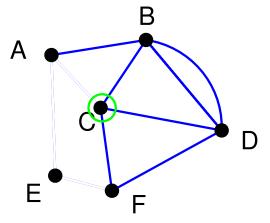
Path so far: FE



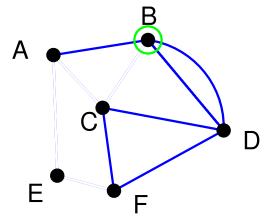
Path so far: FEA



Path so far: FEAC



Path so far: FEACB



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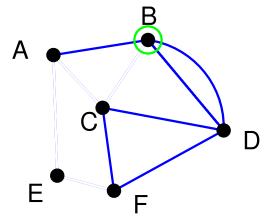
But now, crossing the edge BA would be a mistake, because we would be stuck there.

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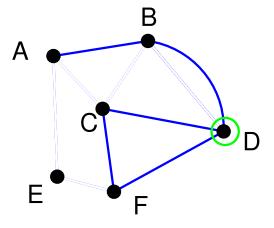
But now, crossing the edge BA would be a mistake, because we would be stuck there.

The reason is that BA is a **bridge**. We don't want to cross ("burn"?) a bridge unless it is the only edge available.

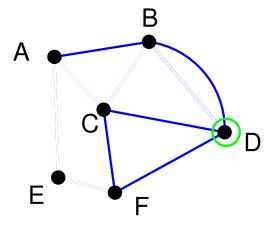
Path so far: FEACB



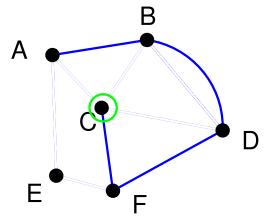
Path so far: FEACBD.



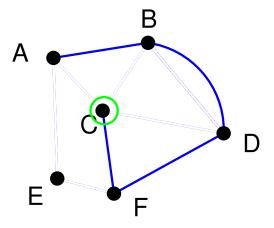
Path so far: FEACBD. Don't cross the bridge!



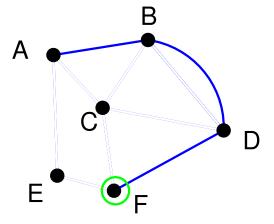
Path so far: FEACBDC



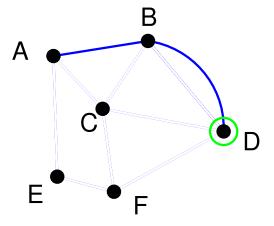
Path so far: FEACBDC Now we have to cross the bridge CF.



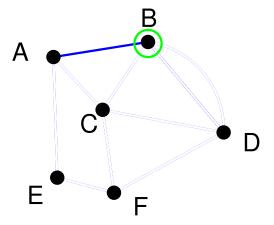
Path so far: FEACBDCF



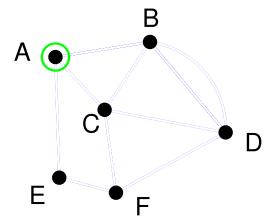
Path so far: FEACBDCFD



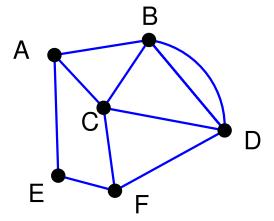
Path so far: FEACBDCFDB



Euler Path: FEACBDCFDBA



Euler Path: FEACBDCFDBA



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This is called **Fleury's algorithm**, and it always works!

Hierholzer?