

# Assignment 1

## Question 1

- a) Given an array of integers of any size,  $n \geq 1$ , write an algorithm as a pseudo code (not a program!) that would reverse every two consecutive elements of the left half of the array (i.e. reverse elements at index 0 & 1, at index 2 & 3, etc.). Additionally, for the right half of the array, the algorithm must change the second element of every two consecutive elements to have to the sum of the two elements. For instance, if the right half starts at index 14, and the values at index 14 & 15 are 72 and 11, then the value at index 15 is changed to 83. If the given array is of an odd size, then the exact middle element should not be manipulated by the algorithm in any way. Finally, your algorithm must use the smallest auxiliary/additional storage to perform what is needed.

**Algorithm** rearrangeArray(arr,n)

**Input** --> Array arr with n elements

**Output** --> Resultant array

```
if ((n % 4 = 0 AND n % 2 = 0) OR n % 4 = 3) then
    mid ← ceiling(size/2)
else if (n % 2 = 0 OR n % 2 = 1) then
    mid ← ceiling((size + 2)/2)
endif
i ← 0
while (i is less than or equal to n) then
    if (i is greater than or equal to (mid - 1)) then
        break the loop
    end if
    a ← arr[i]
    b ← arr[i+1]
    a ← a + b
    b ← a - b
    a ← a - b
    increment i by 2
end while
i ← mid
while (i is less than n) then
    if (i is greater than or equal to (size - 1)) then
        break the loop
    end if
    a ← arr[i]
    b ← arr[i+1]
    b ← a + b
    increment i by 2
end while
```

- b) What is the time complexity of your algorithm, in terms of Big-O?  
-  $O(n)$ , n is number of elements in the array.
- c) What is the space complexity of your algorithm, in terms of Big-O?  
-  $O(1)$

## Question 2

Given a string of random length and random contents of characters, that do not include digits (0-9), write an algorithm, using pseudo code that will shorten the representation of that string by adding the number of consecutive characters, then finds out the first found character that occurred most. For instance, given str as "gggN@@@@@KKeeeejjjjdsmmu" the algorithm should show that the character that occurred most was @ and returns the following representation of the string: "g3N@5K2e4j5dsm2u".

**Algorithm** stringManipulation(str)

**Input** --> String with multiple characters

**Output** --> Resultant string

```
length ← LEN(str)
s ← ""
count ← 1
max ← 1
c ← 0
i ← 1
while (i is less than the length of the string) then
    c ← character at index (i-1)
    if (character at i = c) then
        increment count
    else
        if (max < count) then
            max ← count
            character ← c
        end if
        s ← s + c + count
        count ← 1
    end if
    increment i
end while
s ← s + c + count
Display (s)
Display (character)
Display (max)
```

- a) What is the time complexity of your algorithm, in terms of Big-O?
  - $O(n)$ , n is the length of the string
- b) What is the space complexity of your algorithm, in terms of Big-O?
  - $O(1)$

### Question 3

- i) Develop a well-documented pseudo code that finds the two consecutive elements in the array with the smallest difference in between, and the two consecutive elements with the biggest difference in between. The code must display the values and the indices of these elements. For instance, given the following array [20, 52, 400, 3, 30, 70, 72, 47, 28, 38, 41, 53, 20] your code should find and display something similar to the following (notice that this is just an example. Your solution must not refer to this particular example.):

The two consecutive indices with smallest difference between their values are: index 5 and index 6, storing values 70 and 73.

The two consecutive indices with largest difference between their values are: index 2 and index 3, storing values 400 and 3.

**Algorithm** minMaxDifference(arr,n)

**Input** --> Array with n elements

**Output** --> Indices with smallest difference between their values and the difference  
Indices with largest difference between their values and the difference

Initialize min\_i, min\_j, max\_i, max\_j to 0

min  $\leftarrow$  + infinite

max  $\leftarrow$  - infinite

diff  $\leftarrow$  0

for i  $\leftarrow$  0 to i  $\leftarrow$  n-1

    diff  $\leftarrow$  absolute(arr[i] – arr[i+1])

    if (diff is less than or equal to min) then

        min  $\leftarrow$  diff

        min\_i  $\leftarrow$  i

        min\_j  $\leftarrow$  i + 1

    end if

    if (diff is greater than or equal to max) then

        max  $\leftarrow$  diff

        max\_i  $\leftarrow$  i

        max\_j  $\leftarrow$  i + 1

    end if

    increment i

end for

display ("Indices with maximum difference: ", max\_i, max\_j)

display ("The maximum difference is:", max)

display ("Indices with minimum difference: ", min\_i, min\_j)

display ("The minimum difference is:", min)

- ii) Briefly justify the motive(s) behind your design.
- Initially, I have created two variables (min & max) to store minimum and maximum value. These variables have been initialized to minus (-) infinity and plus (+) infinity, respectively. Then, I have loop through the array. During each iteration, I have found the absolute difference (diff) between the values of successive indices. Then I have compared this difference with "min" variable. If "diff"

is less than “min” then this value is assigned to “min” and corresponding index values are stored in variables. Then, I have compared this difference with “max” variable. If “diff” is greater than “max” then this value is assigned to “max” and corresponding index values are stored in variables. This process is followed for all the values in the array.

- iii) What is the time complexity of your solution?
  - $O(n)$ ,  $n$  is the number of elements in the array. We loop through the array only one time.
- iv) What is the maximum size of stack growth of your algorithm?
  - $O(1)$ . The function call occurs only one time as the algorithm is non-recursive.