Optimizing a reed-solomon encoder (polynomial division)

I am trying to optimize a Reed-Solomon encoder, which is in fact simply a polynomial division operation over Galois Fields 2⁸ (which simply means that values wrap-around over 255). The code is in fact very very similar to what can be found here for Go: http://research.swtch.com/field

The algorithm for polynomial division used here is a synthetic division (also called Horner's method).

I tried everything: numpy, pypy, cython. The best performance I get is by using pypy with this simple nested loop:

```
def rsenc(msg in, nsym, gen):
          Reed-Solomon encoding using polynomial division, better explained at
http://research.swtch.com/field''
      msg_out = bytearray(msg_in) + bytearray(len(gen)-1)
      lgen = bytearray([gf_log[gen[j]] for j in xrange(len(gen))])
      for i in xrange(len(msg_in)):
            coef = msg_out[i]
           \#\ coef\ =\ gf\_mul(msg\_out[i],\ gf\_inverse(gen[0]))\ \ //\ for\ general\ polynomial
division (when polynomials are non-monic), we need to compute: coef = msg_out[i] / gen[0]

if coef != 0: # coef 0 is normally undefined so we manage it manually here (and it
also serves as an optimization btw)
                 lcoef = gf_log[coef] # precaching
\label{eq:coefficient} \begin{array}{ll} \text{for j in xrange(1, len(gen)): \# optimization: can skip g0 because the first coefficient of the generator is always 1! (that's why we start at position 1) \\ & \text{msg\_out[i + j] } ^= \text{gf\_exp[lcoef + lgen[j]] \# equivalent (in Galois Field)} \end{array}
2^8) to msg_out[i+j] += msg_out[i] * gen[j]
      # Recopy the original message bytes
     msg_out[:len(msg_in)] = msg_in
      return msg out
```

Can a Python optimization wizard guide me to some clues on how to get a speedup? My goal is to get at least a speedup of 3x, but more would be awesome. Any approach or tool is accepted, as long as it is cross-platform (works at least on Linux and Windows).

Here is a small test script with some of the other alternatives I tried (the cython attempt is not included since it was slower than native python!):

```
import random
from operator import xor
numpy_enabled = False
     import numpy as np
     numpy enabled = True
except ImportError:
102, 170, 229, 52, 92, 228, 55, 89, 235, 38, 106, 190, 217, 112
    144, 171, 236, 49, 83, 245, 4, 12, 20, 60, 68, 204, 79, 209, 104, 184, 211, 110, 178, 205, 76, 212, 103, 169, 224, 59, 77, 215, 98, 166, 241, 8, 24, 40, 120, 136, 131, 158, 185, 208, 107, 189, 220, 127, 129, 152, 179, 206, 73, 219, 118, 154, 181, 196, 87, 249, 16,
     48, 80, 240, 11, 29, 39, 105, 187, 214, 97, 163, 254, 25, 43, 125,
     135, 146, 173, 236, 47, 113, 147, 174, 233, 32, 96, 160, 251, 22, 58, 78, 210, 109, 183, 194, 93, 231, 50, 86, 250, 21, 63, 65, 195, 94, 226, 61, 71, 201, 64, 192, 91, 237, 44, 116, 156, 191, 218,
     117, 159, 186, 213, 100, 172, 239, 42, 126, 130, 157, 188, 223,
     122, 142, 137, 128, 155, 182, 193, 88, 232, 35, 101, 175, 234, 37,
     111, 177, 200, 67, 197, 84, 252, 31, 33, 99, 165, 244, 7, 9, 27,
     45, 119, 153, 176, 203, 70, 202, 69, 207, 74, 222, 121, 139, 134,
     145, 168, 227, 62, 66, 198, 81, 243, 14, 18, 54, 90, 238, 41, 123, 141, 140, 143, 138, 133, 148, 167, 242, 13, 23, 57, 75, 221, 124,
     132, 151, 162, 253, 28, 36, 108, 180, 199, 82, 246] * 2 + [1])
# Logarithm table, base 3
gf_log = bytearray([0, 0, 25, 1, 50, 2, 26, 198, 75, 199, 27, 104, 51, 238, 223, # BEWARE:
the first entry should be None instead of 0 because it's undefined, but for a bytearray we
can't set such a value
          3, 100, 4, 224, 14, 52, 141, 129, 239, 76, 113, 8, 200, 248, 105, 28, 193, 125, 194, 29, 181, 249, 185, 39, 106, 77, 228, 166, 114, 154, 201, 9, 120, 101, 47, 138, 5, 33, 15, 225, 36, 18, 240, 130,
           69, 53, 147, 218, 142, 150, 143, 219, 189, 54, 208, 206, 148, 19,
           92, 210, 241, 64, 70, 131, 56, 102, 221, 253, 48, 191, 6, 139
           179, 37, 226, 152, 34, 136, 145, 16, 126, 110, 72, 195, 163, 182,
           30, 66, 58, 107, 40, 84, 250, 133, 61, 186, 43, 121, 10, 21, 155, 159, 94, 202, 78, 212, 172, 229, 243, 115, 167, 87, 175, 88, 168,
           80, 244, 234, 214, 116, 79, 174, 233, 213, 231, 230, 173, 232, 44,
           215, 117, 122, 235, 22, 11, 245, 89, 203, 95, 176, 156, 169, 81,
          160, 127, 12, 246, 111, 23, 196, 73, 236, 216, 67, 31, 45, 164, 118, 123, 183, 204, 187, 62, 90, 251, 96, 177, 134, 59, 82, 161,
           108, 170, 85, 41, 157, 151, 178, 135, 144, 97, 190, 220, 252, 188,
           149, 207, 205, 55, 63, 91, 209, 83, 57, 132, 60, 65, 162, 109, 71,
           20, 42, 158, 93, 86, 242, 211, 171, 68, 17, 146, 217, 35, 32, 46,
          137, 180, 124, 184, 38, 119, 153, 227, 165, 103, 74, 237, 222, 197, 49, 254, 24, 13, 99, 140, 128, 192, 247, 112, 7])
```

```
if numpy_enabled:
    np_gf_exp = np.array(gf_exp)
    np\_gf\_log = np.array(gf\_log)
def gf_pow(x, power):
    return gf_exp[(gf_log[x] * power) % 255]
def gf poly mul(p, q):
    r = [0] * (len(p) + len(q) - 1)
    lp = [gf_log[p[i]] for i in xrange(len(p))]
    for j in range(len(q)):
        lq = gf_log[q[j]]
for i in range(len(p)):
            r[i + j] ^= gf_exp[lp[i] + lq]
    return r
def rs_generator_poly_base3(nsize, fcr=0):
    g all = \{\}
    g_all[0] = g_all[1] = g
    for i in range(fcr+1, fcr+nsize+1):
        g = gf_poly_mul(g, [1, gf_pow(3, i)])
        g_all[nsize-i] = g
    return g_all
# Fastest way with pypy
def rsenc(msg_in, nsym, gen):
    '''Reed-Solomon encoding using polynomial division, better explained at
http://research.swtch.com/field'''
    msg\_out = bytearray(msg\_in) + bytearray(len(gen)-1)
    lgen = bytearray([gf_log[gen[j]] for j in xrange(len(gen))])
    for i in xrange(len(msg_in)):
         coef = msg_out[i]
        \#\ coef\ =\ gf\_mul(msg\_out[i],\ gf\_inverse(gen[0]))\ \ \#\ for\ general\ polynomial\ division
(when polynomials are non-monic), the usual way of using synthetic division is to divide the divisor g(x) with its leading coefficient (call it a). In this implementation, this
means:we need to compute: coef = msg_out[i] / gen[0]
        if coef != 0: # coef 0 is normally undefined so we manage it manually here (and it
also serves as an optimization btw)
            lcoef = gf_log[coef] # precaching
             for j in xrange(1, len(gen)): # optimization: can skip g0 because the first
# Recopy the original message bytes
    msg\_out[:len(msg\_in)] = msg\_in
    return msg_out
# Alternative 1: the loops were completely changed, instead of fixing msg_out[i] and
updating all subsequent i+j items, we now fixate msg_out[i+j] and compute it at once using all couples msg_out[i] * gen[j] - msg_out[i+1] * gen[j-1] - \dots since when we fixate msg_out[i+j], all previous msg_out[k] with k < i+j are already fully computed.
def rsenc alt1(msg in, nsvm, gen):
    msg_in = bytearray(msg_in)
    msg_out = bytearray(msg_in) + bytearray(len(gen)-1)
    lgen = bytearray([gf_log[gen[j]] for j in xrange(len(gen))])
    # Alternative 1
    jlist = range(1, len(gen))
    for k in xrange(1, len(msg_out)):
        for x in xrange(max(k-len(msg_in),0), len(gen)-1):
             if k-x-1 < 0: break
             msg_out[k] ^= gf_exp[msg_out[k-x-1] + lgen[jlist[x]]]
    # Recopy the original message bytes
    msg_out[:len(msg_in)] = msg_in
    return msg out
# Alternative 2: a rewrite of alternative 1 with generators and reduce
def rsenc_alt2(msg_in, nsym, gen):
    msg_in = bytearray(msg_in)
    msg_out = bytearray(msg_in) + bytearray(len(gen)-1)
    lgen = bytearray([gf_log[gen[j]] for j in xrange(len(gen))])
    # Alternative 1
    jlist = range(1, len(gen))
    for k in xrange(1, len(msg_out)):
        items_gen = ( gf_exp[msg_out[k-x-1] + lgen[jlist[x]]] if k-x-1 >= 0 else
next(iter(())) for x in xrange(max(k-len(msg_in),0), len(gen)-1) )
        msg_out[k] ^= reduce(xor, items_gen)
    # Recopy the original message bytes
    msg_out[:len(msg_in)] = msg_in
    return msg_out
# Alternative with Numpy
def rsenc numpy(msg in, nsym, gen):
    msg_in = np.array(bytearray(msg_in))
    msg_out = np.pad(msg_in, (0, nsym), 'constant')
    lgen = np_gf_log[gen]
    for i in xrange(msg in.size):
        msg_out[i+1:i+lgen.size] ^= np_gf_exp[np.add(lgen[1:], msg_out[i])]
    msg_out[:len(msg_in)] = msg_in
    return msg_out
gf mul arr = [bytearray(256) for in xrange(256)]
```

```
gf_add_arr = [bytearray(256) for _ in xrange(256)]
  # Precompute multiplication and addition tables
  def gf_precomp_tables(gf_exp=gf_exp, gf_log=gf_log):
         global gf_mul_arr, gf_add_arr
         for i in xrange(256):
                for j in xrange(256):
                      gf_mul_arr[i][j] = gf_exp[gf_log[i] + gf_log[j]]
gf_add_arr[i][j] = i ^ j
         return gf_mul_arr, gf_add_arr
 # Alternative with precomputation of multiplication and addition tables, inspired by zfec:
 https://hackage.haskell.org/package/fec-0.1.1/src/zfec/fec.c
 def rsenc_precomp(msg_in, nsym, gen=None):
        msg_in = bytearray(msg_in)
        {\tt msg\_out = bytearray(msg\_in) + bytearray(len(gen)-1)}
         for i in xrange(len(msg_in)): # [i for i in xrange(len(msg_in)) if msg_in[i] != 0]
                coef = msg_out[i]
                if coef != 0: # coef 0 is normally undefined so we manage it manually here (and
 it also serves as an optimization btw)
                       mula = gf mul arr[coef]
                       for j in xrange(1, len(gen)): # optimization: can skip g0 because the first
 coefficient of the generator is always 1! (that's why we start at position 1)
                              \# msg\_out[i + j] = gf\_add\_arr[msg\_out[i+j]][gf\_mul\_arr[coef][gen[j]]] \ \# sg\_out[i + j] = gf\_add\_arr[msg\_out[i+j]][gf\_mul\_arr[coef][gen[j]]] \ \# sg\_out[i+j][gf\_mul\_arr[coef][gen[j]]] \ \# sg\_out[i+j][gf\_mul\_
 slower...
                              \#msg\_out[i + j] \ ^= \ gf\_mul\_arr[coef][gen[j]] \ \# \ faster
                              msg_out[i + j] ^= mula[gen[j]] # fastest
        # Recopy the original message bytes
        msg\_out[:len(msg\_in)] = msg\_in \# equivalent to c = mprime - b, where mprime is msg\_in
 padded with [0]*nsvm
        return msg_out
 def randstr(n, size):
          ''Generate very fastly a random hexadecimal string. Kudos to jcdryer
 http://stackoverflow.com/users/131084/jcdyer'
         hexstr = '%0' + str(size) + 'x'
         for _ in xrange(n):
               yield hexstr % random.randrange(16**size)
 # Simple test case
  if __name__ == "__main_
         # Setup functions to test
         funcs = [rsenc, rsenc_precomp, rsenc_alt1, rsenc_alt2]
        if numpy_enabled: funcs.append(rsenc_numpy)
        gf_precomp_tables()
        # Setup RS vars
        n = 255
k = 213
        import time
        # Init the generator polynomial
        g = rs_generator_poly_base3(n)
        # Init the ground truth
        mes = 'hello world'
        mesecc_correct = rsenc(mes, n-11, g[k])
         # Test the functions
        for func in funcs:
                # Sanity check
               if func(mes, n-11, g[k]) != mesecc_correct: print func.__name__, ": output is
  incorrect!"
               # Time the function
                total time = 0
                for m in randstr(1000, n):
                       start = time.clock()
                       func(m, n-k, g[k])
               total_time += time.clock() - start
print func.__name__, ": total_time elapsed %f seconds." % total_time
And here is the result on my machine:
 With PyPy:
 rsenc : total time elapsed 0.108183 seconds.
 rsenc_alt1 : output is incorrect!
rsenc_alt1 : total time elapsed 0.164084 seconds.
 rsenc_alt2 : output is incorrect!
 rsenc_alt2 : total time elapsed 0.557697 seconds.
 Without PvPv:
 rsenc : total time elapsed 3.518857 seconds.
 rsenc_alt1 : output is incorrect!
 rsenc_alt1 : total time elapsed 5.630897 seconds.
 rsenc alt2 : output is incorrect!
 rsenc alt2: total time elapsed 6.100434 seconds.
 rsenc_numpy : output is incorrect!
 rsenc_numpy : total time elapsed 1.631373 seconds
```

(Note: the alternatives should be correct, some index must be a bit off, but since they are slower anyway I did not try to fix them)

/UPDATE and goal of the bounty: I found a very interesting optimization trick that promises to speed up computations a lot: to precompute the multiplication table. I updated the code above with the new function rsenc_precomp(). However, there's no gain at all in my implementation, it's even a bit slower:

```
rsenc : total time elapsed 0.107170 seconds. rsenc precomp : total time elapsed 0.108788 seconds.
```

How can it be that arrays lookups cost more than operations like additions or xor? Why does it work in ZFEC and not in Python?

I will attribute the bounty to whoever can show me how to make this multiplication/addition lookup-tables optimization work (faster than the xor and addition operations) or who can explain to me with references or analysis why this optimization cannot work here (using Python/PyPy/Cython/Numpy etc.. I tried them all).

```
edited May 26 at 13:40

asked May 21 at 3:11

gaborous
1,933 3 17 31

Have you tried pypi.python.org/pypi/reedsolo? I'd be curious how it compares speed wise. – figs May 23 at 3:13

1 The code I posted is an enhancement of the reedsolo library you cite:) It's quite fast with pypy, but not as fast as the implementation I present above, but it's still quite far from the performance of zfec and other c/c++ libraries... – gaborous May 23 at 21:01
```

3 Answers

The following is 3x faster than pypy on my machine (0.04s vs 0.15s). Using Cython:

```
ctypedef unsigned char uint8_t # does not work with Microsoft's C Compiler: from
libc.stdint cimport uint8_t
cimport covthon.array as array
cdef uint8_t[::1] gf_exp = bytearray([1, 3, 5, 15, 17, 51, 85, 255, 26, 46, 114, 150, 161,
  lots of numbers omitted for space reasons
   ...])
cdef uint8_t[::1] gf_log = bytearray([0, 0, 25, 1, 50, 2, 26, 198, 75, 199, 27, 104,
    more numbers omitted for space reasons
    ...])
import cython
@cython.boundscheck(False)
@cython.wraparound(False)
@cython.initializedcheck(False)
def rsenc(msg_in_r, nsym, gen_t):

'''Reed-Solomon encoding using polynomial division, better explained at
http://research.swtch.com/field''
    make a memory view from a read only object
    cdef int[::1] gen = array.array('i',gen_t) # convert list to array
    cdef uint8_t[::1] msg_out = bytearray(msg_in) + bytearray(len(gen)-1)
    cdef int j
    cdef uint8_t[::1] lgen = bytearray(gen.shape[0])
    for j in xrange(gen.shape[0]):
        lgen[j] = gf_log[gen[j]]
    cdef uint8_t coef,lcoef
    cdef int i
    for i in xrange(msg_in.shape[0]):
    coef = msg_out[i]
        if coef != 0: # coef 0 is normally undefined so we manage it manually here (and it
also serves as an optimization btw)
            lcoef = gf_log[coef] # precaching
            for j in xrange(1, gen.shape[0]): # optimization: can skip g0 because the
first coefficient of the generator is always 1! (that's why we start at position 1)
msg\_out[i+j] \ ^= gf\_exp[lcoef + lgen[j]] \ \# \ equivalent \ (in \ Galois \ Field \ 2^8) \ to \ msg\_out[i+j] \ -= \ msg\_out[i] \ \# \ gen[j]
    # Recopy the original message bytes
    msg_out[:msg_in.shape[0]] = msg_in
```

It is just your fastest version with static types (and checking the html from cython -a until the loops aren't highlighted in yellow).

A few brief notes:

- Cython prefers x.shape[0] to len(shape)
- Defining the memoryviews as [::1] promises they are continuous in memory, which helps
- initializedcheck(False) is necesary for avoiding lots of existance checks on the globally defined gf_exp and gf_log. (You might find you can speed up your basic Python/PyPy

code by creating a local variable reference for these and using that istead)

 I had to copy a couple of the input arguments. Cython can't make a memoryview from a readonly object (in this case msg_in, a string. I could probably have just made it a char* though). Also gen (a list) needs to be in something with fast element access.

Other than that it's all fairly straight-forward. (I haven't tried any variations of it having got it faster). I'm really quite impressed at how well PyPy does.

edited May 21 at 21:20 gaborous 1,933 3 17 31 answered May 21 at 7:33 DavidW 1,645 3 13

Thank you very much for your answer, it seems very promising, but I can't get it to work because it crashes on the uint8_t type, which is part of C99 and unluckily on Windows, Microsoft's C compiler is C89. Could you suggest another type please? — gaborous May 21 at 18:39

1 Does unsigned char work? If you replace the cimport libc.stdint line with ctypedef unsigned char uint8 t I suspect it'll work fine (haven't tested it though). — DavidW May 21 at 18:43

Thank you very much, but it fails with invalid syntax, but there's something fishy. I tried to just ctypedef unsigned long ULong, which is the example given in the cython doc, and it fails with invalid syntax too. Googling... - gaborous May 21 at 18:49

Fixed: I was trying to directly run the .pyx. I just fixed it by creating a .py file that calls the .pyx. Silly me... This code indeed runs faster, and your explanations are really useful (I wondered how to create arrays of chars in Cython), thank's a lot! However, on my real setup, I only gain about a 2x speedup (if I account the rest of the Python code that is not JITed now, I go down to only a 0.25x speedup but I will cythonize later). Can you give me clues about the variations you would do to get even more speedup? — qaborous May 21 at 19:47

I'm honestly not sure what else you could do to make it faster (short of a clever redesign of the algorithm, and that may well not be possible!). The only obvious thing I can see is to avoid copying <code>msg_in</code> at the start (you could get a <code>const_uint8_t*</code> without a copy, but you need to change a few other things later), and I'd be surprised if it gained much. That's probably the best I can do. — <code>DavidW May 21</code> at 20:11

Yes indeed I already timed that part of the code, the copy isn't really a bottleneck, it's just that the nested loop is huge. I'll try to see if I can find another algorithm, but for now your solution is the one I'll use. Thank you very much for your detailed advices:)— gaborous May 21 at 20:51

I found a way to get a bit more juice: by initilizing directly the gen_t polynomial as a bytearray() instead of a sublist inside the generator function <code>rs_generator_poly_base3()</code>, there's no need to convert at encoding time, and in addition gen can then be declared as <code>uint8_t</code>. This makes timing go down from 0.042 to 0.037. — <code>qaborous May 22</code> at 17:05

Tried prange() to parallelize the inner loop, and it works great! This boosted quite a bit more. I added it to your answer. – gaborous May 22 at 17:56

Glad to here you got it faster. They're sensible ways to speed it up. It goes a bit beyond the scope of my original answer, so it's probably better to add an answer of your own saying that you got further improvements with these methods (or just leave it in the comments) rather than editing my answer. – DavidW May 22 at 20:46

Ok fair enough, I'll post a separate answer. I just did a fix to your answer, there was still one len(gen) instead of gen.shape[0]. After fixing that, there is another great speed boost! Also, I discovered that len() still calls Python code, while .shape does not (and in fact it is mandatory to use .shape when releasing the gil with with nogil:). — gaborous May 23 at 2:41

@DavidW Nice. Try replacing int i, j with Py_ssize_t i, j that should be faster (and loop could be potentially vectorized by the compiler), as illustrated in this guestion. — rth Jun 2 at 8:45

Alternatively, if you know C, I would recommend to rewrite this Python function in plain C and call it (say with CFFI). At least you know you reach the top performance in the inner loops of your functions without needing to be aware of either PyPy or Cython tricks.

See: http://cffi.readthedocs.org/en/latest/overview.html#performance

answered May 21 at 8:05

Armin Rigo 5,672 8 20

Thank you for the suggestion. CFFI is indeed very interesting and I know how to program in C, but compiling C code from the Python interpreter on Windows is always quite complicated. I could find a way to do that simply for Cython, but for CFFI it can be quite complicated from what I read, and I'd rather avoid the hassle for my users. But it could be used to make an optional high speed module for users willing to take the time to install everything that is needed, I'll give it a thought. — gaborous May 21 at 18-13.

Thank's to you, I discovered that CFFI relies on pycparser, a full C99 parser in pure python. Awesome! – gaborous May 21 at 18:46

Note that the very new CFFI 1.0 should fix most of the issues of installation (they are talking about CFFI 0.9 if the pages are more than a few days old). It is now like Cython: you need a C compiler, but you control when it is invoked—either when you run the build scripts, or in setuptools, or even manually on the generated C file—and you can then distribute the resulting DLL (.pyd file). — Armin Rigo May 21 at 23:19

Building on DavidW's answer, here's the implementation I am currently using, which is about 20% faster by u

```
from cython.parallel import parallel, prange
@cvthon.boundscheck(False)
@cython.wraparound(False)
@cython.initializedcheck(False)
cdef rsenc_cython(msg_in_r, nsym, gen_t) :
'''Reed-Solomon encoding using polynomial division, better explained at http://research.swtch.com/field'''
     cdef uint8_t[::1] msg_in = bytearray(msg_in_r) # have to copy, unfortunately - can't
make a memory view from a read only object
    #cdef int[::1] gen = array.array('i',gen_t) # convert List to array
cdef uint8_t[::1] gen = gen_t
     cdef uint8_t[::1] msg_out = bytearray(msg_in) + bytearray(len(gen)-1)
     cdef int i.
     cdef uint8_t[::1] lgen = bytearray(gen.shape[0])
     for j in xrange(gen.shape[0]):
         lgen[j] = gf_log_c[gen[j]]
     cdef uint8_t coef,lcoef
     with nogil:
         for i in xrange(msg in.shape[0]):
              coef = msg_out[i]
              if coef != 0: # coef 0 is normally undefined so we manage it manually here
(and it also serves as an optimization btw)
                  lcoef = gf_log_c[coef] # precaching
                   for j in prange(1, gen.shape[0]): # optimization: can skip g0 because the
first coefficient of the generator is always 1! (that's why we start at position 1)

msg_out[i + j] ^= gf_exp_c[lcoef + lgen[j]] # equivalent (in Galois
Field 2^8) to msg_out[i+j] -= msg_out[i] * gen[j]
     # Recopy the original message bytes
     msg_out[:msg_in.shape[0]] = msg_in
     return msg_out
```

I would still like it to be faster (on a real implementation, data is encoded at about 6.4 MB/s with n=255, n bein

The main lead to a faster implementation that I have found is to use a LUT (LookUp Table) approach, by prec Cython implementations, the LUT approach is slower than calculating XOR and addition operations.

There are other approaches to implement a faster RS encoder, but I don't have the abilities nor the time to try

- "Fast software implementation of finite field operations", Cheng Huang and Lihao Xu, Washington Univ http://citeseerx.ist.psu.edu/viewdoc/download?doi=10.1.1.99.6595&rep=rep1&type=pdf and a correct view=295
- Luo. Jiangiang, et al. "Efficient software implementations of large finite fields GF (2 n) for secure storac
- "A Performance Evaluation and Examination of Open-Source Erasure Coding Libraries for Storage.", F O'Hearn, Z, FAST. Vol. 9. 2009. https://www.usenix.org/legacy/events/fast09/tech/full_papers/plank/pl of Open-Source Erasure Coding Libraries for Storage Applications", Plank and Schuman.
- Sourcecode of the ZFEC library, with multiplication LUT optimization: https://hackage.haskell.org/pack
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However, I think the best lead is to use an efficient polynomial modular reduction instead of polynomial div

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/EDIT: in fact it seems "On computation of polynomial modular reduction" just uses the same approach as I d we precompute the couples of coefficients we will need, and reduce them all at once), and unluckily it's not fa for all since it depends on the message input).

/EDIT: I found a library with really interesting optimizations, lots that are not even found in any academic paper fastest software implementation of Reed-Solomon: the wirehair project and the related blog for more details. \ longhair with similar optimization tricks.

edited Jun 22 at 17:09



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