

# Graphs

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# Graphs

- Set of vertices connected pairwise by edges.
- Collection of nodes/vertices connected to each other through set of edges
- Edges connected in any possible way
- $G = (V, E)$

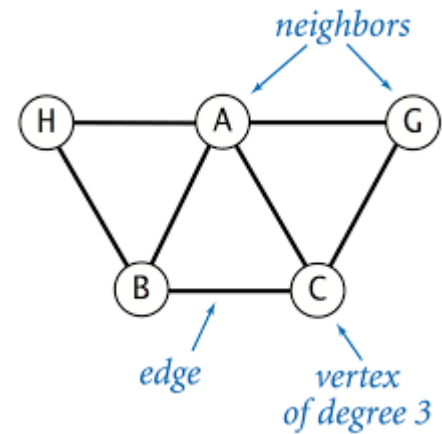


# Definitions

- **Path.** Sequence of vertices connected by edges.
- **Cycle.** Path whose first and last vertices are the same.
- Two vertices are **connected** if there is a path between them.

# A simple graph

Vertices	Neighbours
A	B C G H
B	A C H
C	A B G
G	A C
H	A B



# Graph types

- Acyclic graph has no cycles
  - A tree – an acyclic connected graph
- Forest – disjoint set of trees

# Tree

- $V-1$  edges and no cycles
- $V-1$  edges and is connected
- Connected and moving any edge disconnects it
- Acyclic but adding any edge creates a cycles
- Exactly one simple path connects each pair of vertices in  $G$

# Applications of Graphs

- Highway network, Flight network.
- Computer networks: Local area network, Internet, Web.
- Representing relationships between components in electronic circuits.
- Facebook Graph Search is a real-life example of application of graph algorithms for example friends suggestions



# Applications of Graphs

- Biology and conservation efforts where a vertex represents regions where certain species exist and the edges represent migration path or movement between the regions.
- Google Maps: Various locations are represented as vertices and the roads are represented as edges and graph theory is used to find shortest path between two nodes.
- E-commerce The “Recommendations for you” section on various e-commerce websites uses graph theory to recommend items of similar type to user’s choice.

# Graph continues

- Directed
- Undirected
- Weighted and unweighted

# Graph Representation

- Edge List representation
- Operations
  - Finding nodes adjacent to a given nodes  $O(|E|)$
  - Finding if two node are connected  $O(|E|)$
  - Number of edges can be very large
  - $O(|V|)$  better?

# Adjacency Matrix

0	1	2	3	4	5
1					
2					
3					
4					
5					

# Time Complexity

- Finding adjacent vertices
  - Go to the vertices list  $O(v)$
  - Then go to the row associated with the vertex in the two dimensional array  $O(v)$
  - Total  $O(v)$
- Finding if two nodes are connected
  - $O(1)$  if you know  $i$  and  $j$  indexes
  - $O(v + v + 1) = O(v)$

# Time Complexity

- Finding if two nodes are connected
  - $O(1) + O(V)$
  - How can the second  $O$  can be avoided?
- How would you represent weighted graph?

# Memory Big O?

- $V^2$
- Sparse vs dense graph
- Redundant information when it is not connected
- Most graph will not be anything close to  $V^2$
- Adjacency Matrix will not be a good fit

# What is the solution

- Only keep listed of nodes that are connected using indexes
- List
  - Array `v1 = new int[?]`
  - Linked list
  - Binary tree
- Space  $O(e)$
- $E < V^2$



# Memory usage

- $2 \times e$  (edges) (undirected graph)
- Exactly  $e$  edges (directed graph)
- Space is proportional to number of edges

# Time Complexity

- Two nodes connected?
  - $O(1)$   $A_{ij}$  in adjacency matrix
  - Linear search –  $O(v)$ 
    - Sort the binary search?
      - $O(\log v)$
- Adjacent nodes
  - $O(v)$  for both structure

# Simple Exercise

- Design a social network of friends in the class using a graph data structure
- Get inspirations from
  - <http://introcs.cs.princeton.edu/java/45graph/Graph.java>
  - [https://www.cs.duke.edu/courses/cps100e/fall10/class/11\\_Bacon/code/Graph.html](https://www.cs.duke.edu/courses/cps100e/fall10/class/11_Bacon/code/Graph.html)