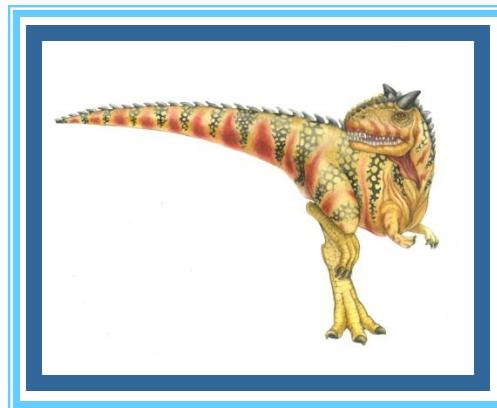


Chapter 13: File-System Interface



LRU : ? calculate effective bandwidth

? calculate average access time

Disk scheduling



Chapter 13: File System Interface

- File Concept *What is this?* *how os do*
- Access Methods *sequential/random access that?*
- Disk and Directory Structure
- File-System Mounting *(some methods) with trade-offs* *-low level*
- File Sharing *not protected*
- Protection *not sharing*





Objectives

- To explain the functions of file systems
- To describe the **interfaces to file systems**
- To discuss file-system design tradeoffs, including access methods, *file sharing*, and *directory structures*
- To explore file-system protection

why we need this?
- provide convenience
- efficient to manage!

- 7 - 這會發現

pro and cons

I understand
what is protection,
how it works

我們會 suffer from linux!





File Concept

- *logically continuous!*
Contiguous logical address space
- Types:
 - Data *instructions and this!*
 - numeric
 - character
 - binary
- Program
- Contents defined by the file's creator *不同 files!*
 - Many types, consider **text file**, source file, executable file *high flexibility!*





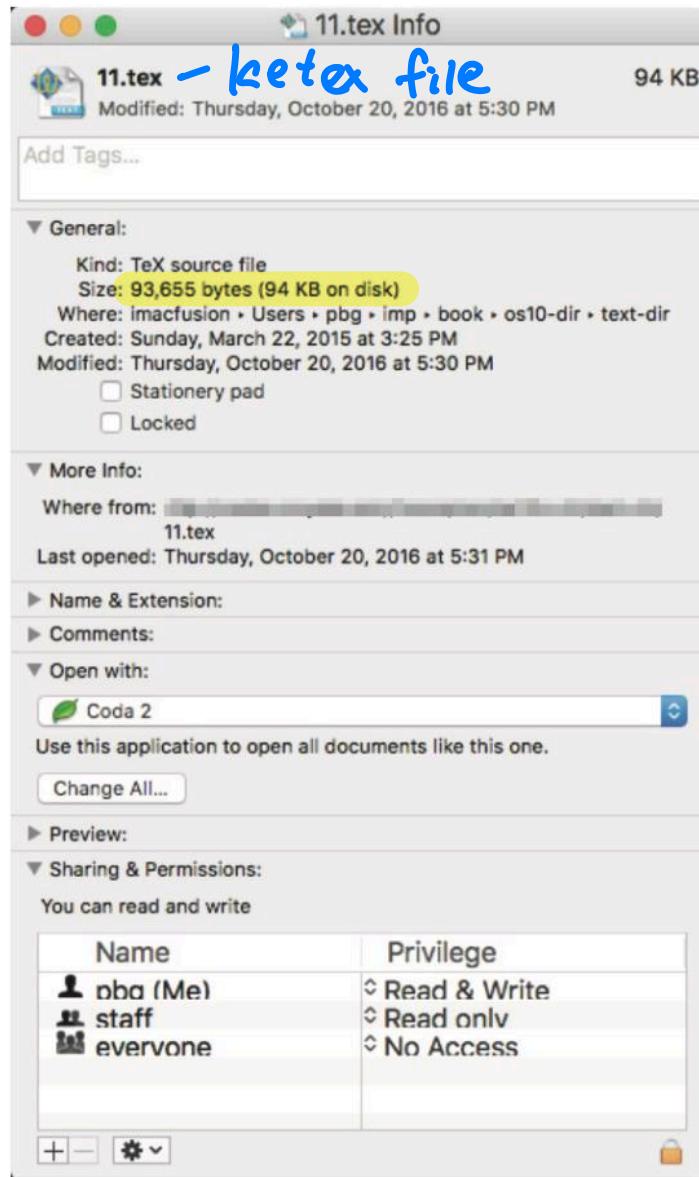
File Attributes

- Name – information kept in human-readable form
- Identifier – unique tag (number) identifies files within a file system
- Type – needed by systems that support different types *it b s know,
how to open?
Sometimes encoded!*
- Location – pointer to file location on device
- Size – current file size
- Protection – controls who can do reading, writing, executing, etc.
- Time, date, and user identification – data for protection, security, and usage monitoring *not Embedded in content itself!*
- Information about files are kept in a directory structure, maintained on the disk - part of which currently in use can be cached in main memory for fast access *Cache this in memory!*
- Many variations, including extended file attributes such as file checksum

*/
ensure that
not changed by someone
checksum in Mac OS*



File info Window on Mac OS X



file information





File Operations

- File is an **ADT** or **abstract data type**
 - **Create** – create a file
 - **Write** – at **write pointer** location
 - **Read** – at **read pointer** location
 - **Reposition within file** - **seek**
 - **Delete**
 - **Truncate** 缩短： :
 - **Open(F_i)** – search the directory structure on disk for entry F_i , and move the content of entry to memory, preparing file for subsequent access
 - **Close (F_i)** – move the content of entry F_i in memory to directory structure on disk
 - Such operations involve the changes of various OS kernel data structures
-] important!
- structure maintained by
the kernel!





Open Files

fork!

- Several data structures are needed to manage open files:
 - Open-file tables: tracks open files, system-wide open-file table, and per-process open-file table *← copy this table!*
 - File pointer: pointer to last read/write location, per process that has the file open *→ how open change the table?*
 - File-open count: counting the number of processes that the file has been opened – to allow removal of data from the open-file table when the last processes closes it (when file-open count is zero)
 - Disk location of a file: cache of data access information (*Fast access*)
 - Access rights: per-process access mode information

change the files : get the pointer!





File Types – Name, Extension

file type	usual extension	function
executable	exe, com, bin or none	ready-to-run machine-language program
object	obj, o <i>like C program</i>	compiled, machine language, not linked
source code	c, cc, java, pas, asm, a	source code in various languages
batch	bat, sh <i>command line</i>	commands to the command interpreter
text	txt, doc	textual data, documents
word processor	wp, tex, rtf, doc <i>Linux Windows</i>	various word-processor formats
library	lib, a, so, dll <i>static dynamic</i>	libraries of routines for programmers
print or view	ps, pdf, jpg	ASCII or binary file in a format for printing or viewing
archive	arc, zip, tar <i>compression!</i>	related files grouped into one file, sometimes compressed, for archiving or storage
multimedia	mpeg, mov, rm, mp3, avi	binary file containing audio or A/V information

videos! modify by you

*ready to run → longer a program
extension of this file
↓
use this to distinguish*





Access Methods

Simple!

- Sequential Access – simplest access method

read next

write next

reset → ptr move to beginning
no read after last write

(rewrite)

beginning



→ read file more!

- Direct Access – implemented by sequential access method – file is fixed length logical records blocks.

read n

write n

position to n

read next

write next

rewrite n

Why need the blocks!

start index of
depends on OS!

Block so small ⇒ lots of overhead!

n = relative block number

- Relative block numbers allow OS to decide where file should be placed
 - See disk block allocation problem in Chapter 14



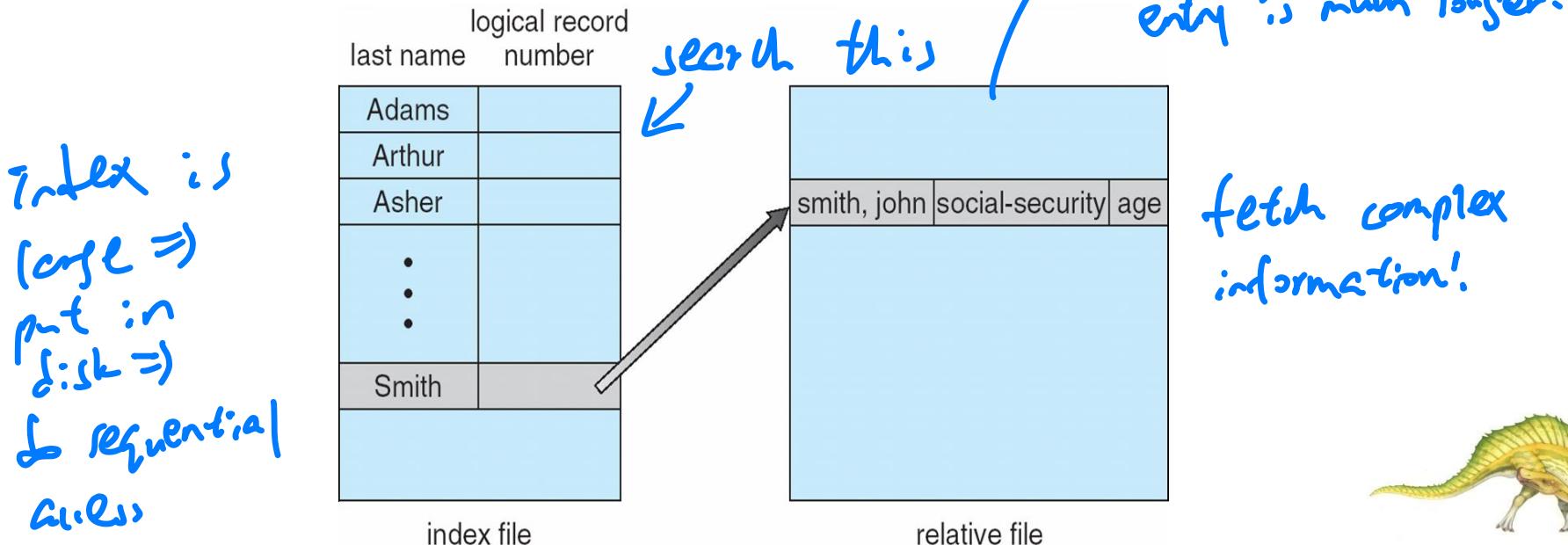


Other Access Methods

widely used nowadays

- Other file access methods can be built on top of direct-access method
- Generally, involve creation of an **index** for a file
 - Keep index in memory for fast location of the data to be operated on
 - If too large, index (in memory) of the index (on disk)
- IBM indexed sequential-access method (ISAM) is an example
 - Small master index, points to disk blocks of secondary index
 - File kept sorted on a defined key
 - All done by the OS
- VMS operating system provides index and relative files as another example

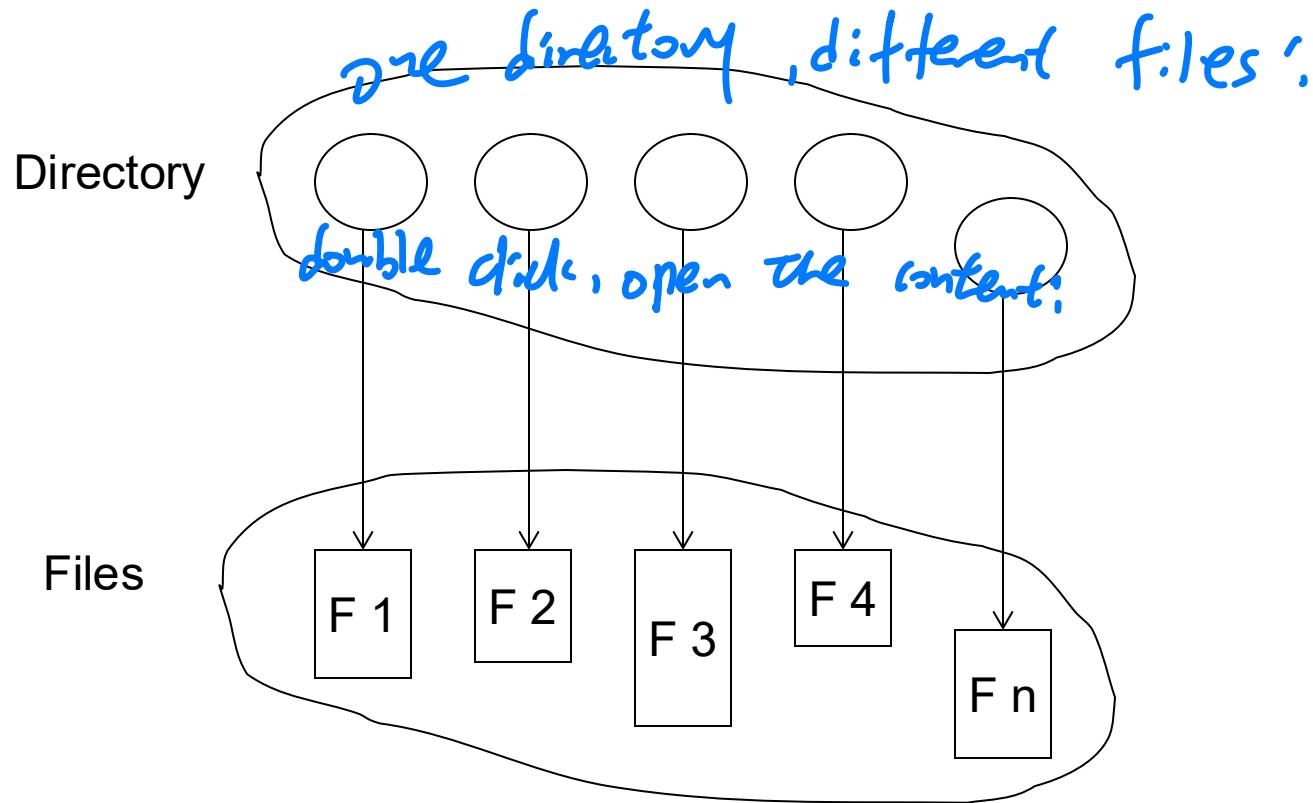
Memory I/O can
→ improve performance





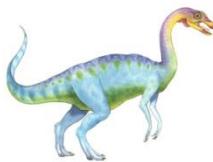
Directory Structure

- A collection of nodes containing information about all files



Both the directory structure and files reside on disk





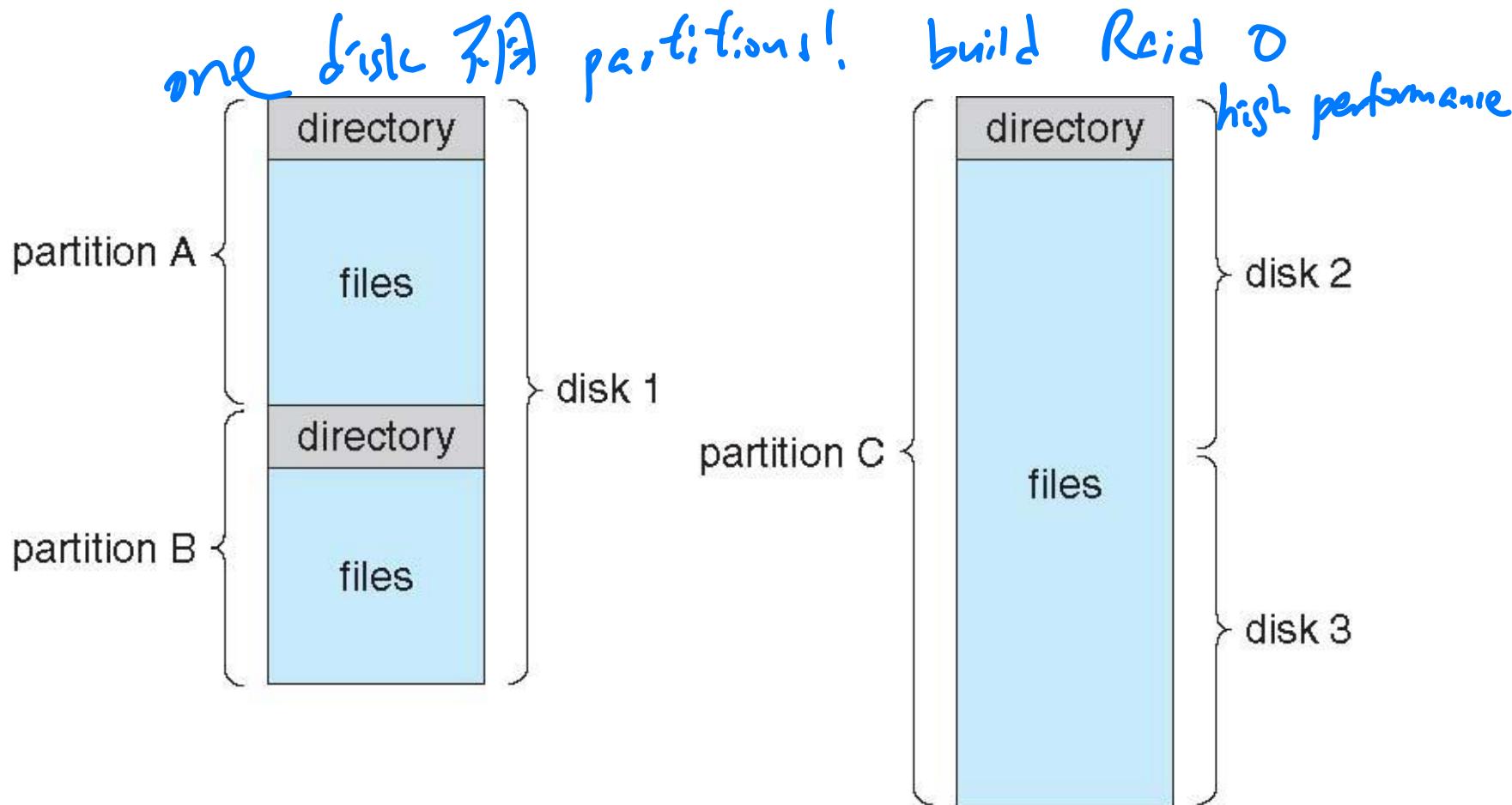
Physical Disk Structure

- Disk can be subdivided into partitions / CDE partitions
- Disks or partitions can be RAID protected against failure
- Disk or partition can be raw – without a file system, or formatted with a file system
- Partitions are also known as minidisks, slices *nick name Mac OS*
- An entity on a disk containing a file system known as a volume
- Each volume containing a file system also keeps track of the file system info in device directory or volume table of contents
- Other than general-purpose file systems, there are many special-purpose file systems, frequently within the same operating system or computing systems





A Typical File-system Organization





Operations Performed on Directory

- Search for a file ✓
- Create a file
- Delete a file 不想寫！
- List a directory ls ...
- Rename a file rn
- Traverse the file system dfn ...





Organize the Directory (Logically) to Obtain

- **Efficiency** – locating a file quickly

for ↓ human !
directory convenient
to me, 請把 every
file to one
single file !

- **Naming** – convenient to users

- Two users can have same name for different files
 - The same file can have several different names

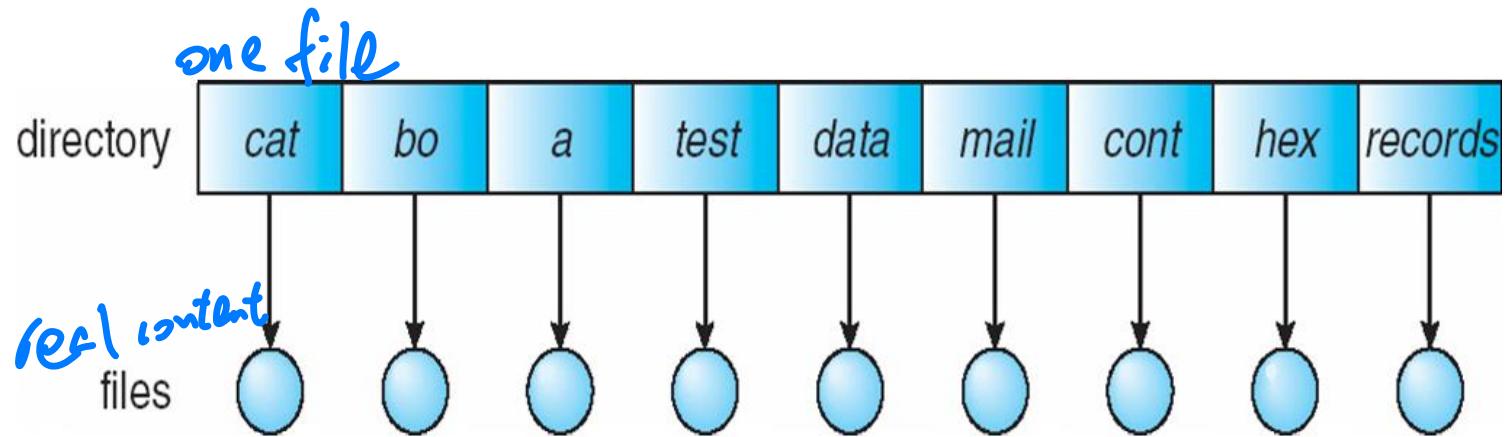
- **Grouping** – logical grouping of files by properties, (e.g., all Java programs, all games, my comp3511, ...)





Single-Level Directory

- A single directory for all users



Naming problem 不同內有 cat!

Grouping problem 不同 group, 33.000 group animals!

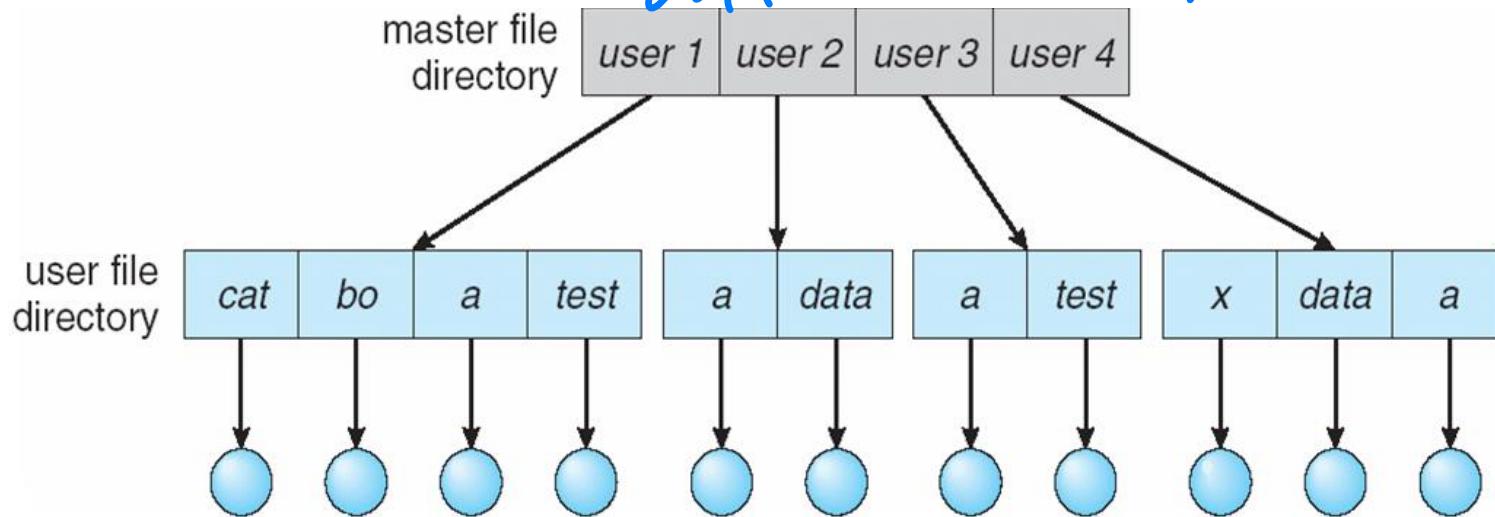




Two-Level Directory

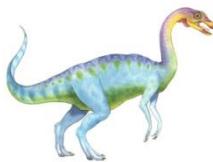
- Separate directory for each user

different files!

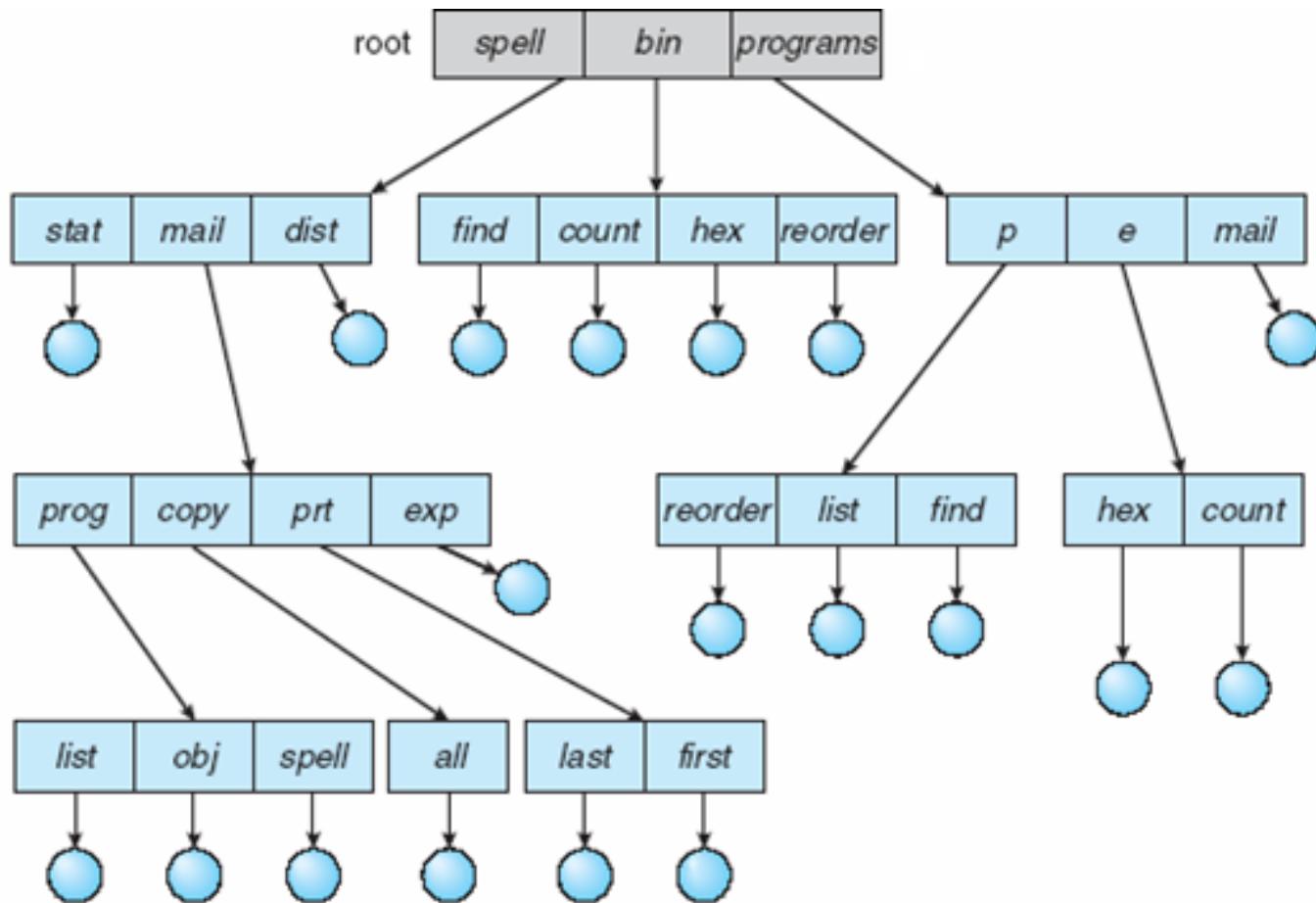


- Path name** – need a pathname to identify a file/dir, e.g., /user1/cat
- Can have the same file name under different users (paths)
- More efficient searching than single-level directory
- No grouping capability





more (level) Tree-Structured Directories





Tree-Structured Directories (Cont.)

- Efficient searching
- Grouping Capability
- Current directory (working directory)
 - `cd /spell/mail/prog` ✓ *change direction!*
 - `type list`





Tree-Structured Directories (Cont)

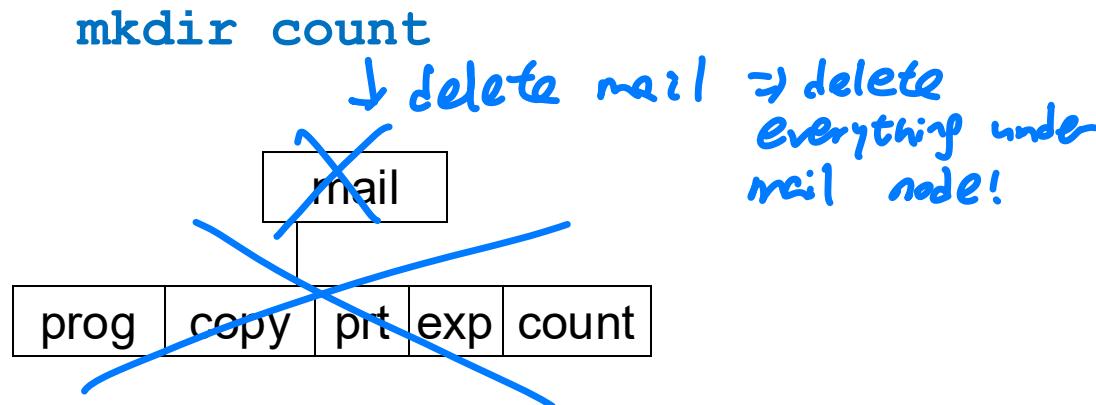
- Absolute or relative path name
- Creating a new file is done in the current directory
- Delete a file in the current directory

`rm <file-name>`

- Creating a new subdirectory is done in current directory

`mkdir <dir-name>`

Example: if in current directory `/mail`



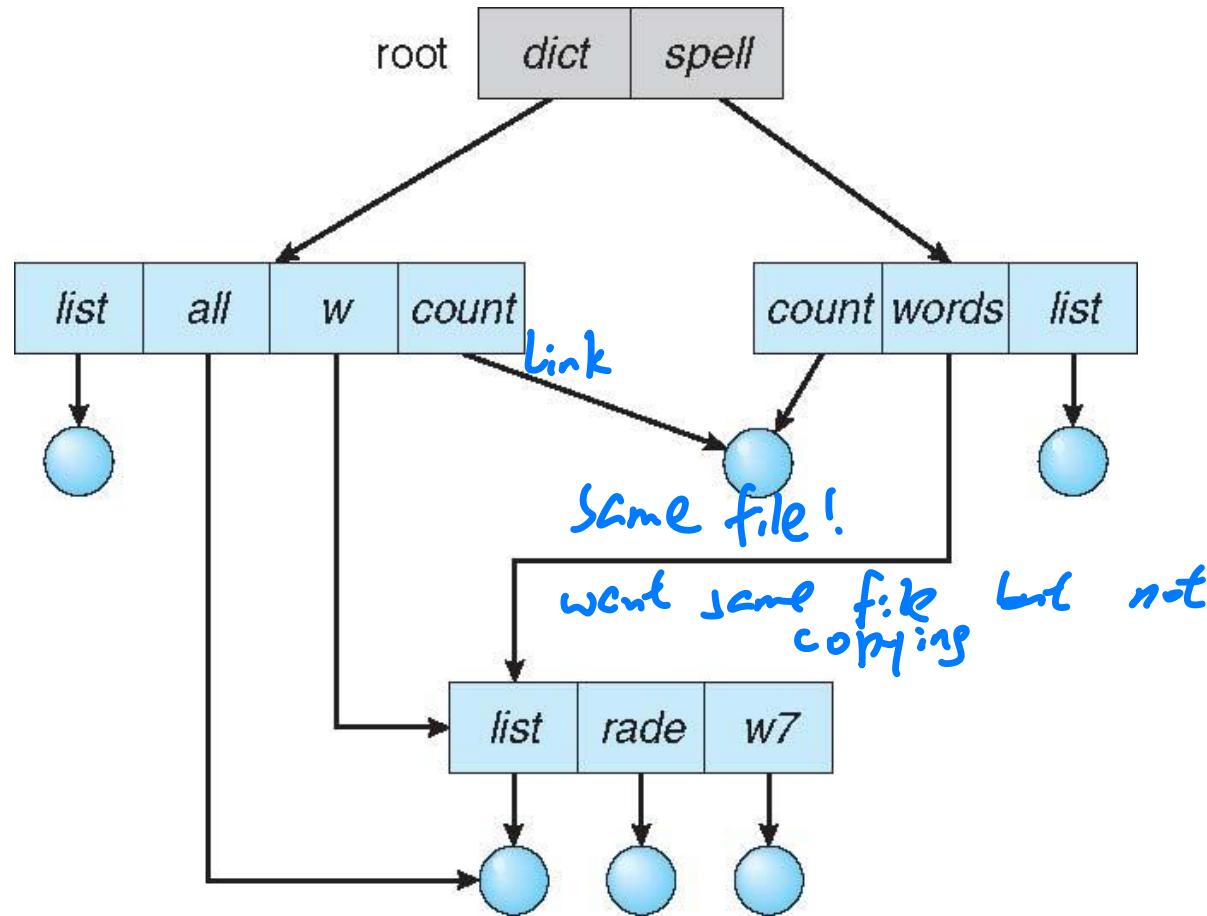
Deleting “mail” ⇒ deleting the entire subtree rooted by “mail”





Acyclic-Graph Directories

- Have shared subdirectories and files – more flexible and complex





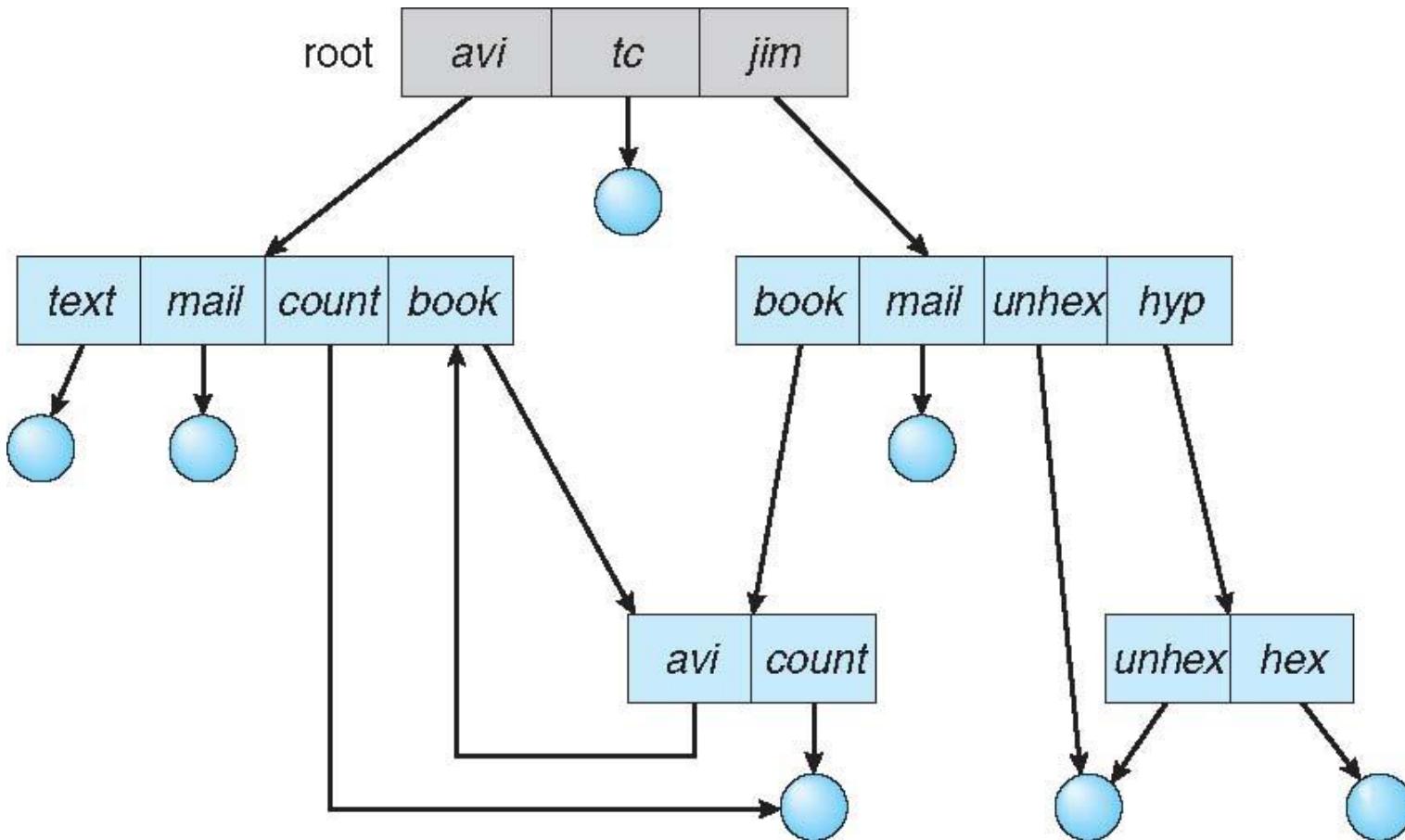
Acyclic-Graph Directories (Cont.)

- New directory entry type
 - Link – another name (pointer) to an existing file
 - Resolve the link – follow pointer to locate the file
- Two different (path) names (aliasing) *Z'ya dead link!*
 - Ensure not traversing shared structures more than once
- Deletion might lead to that dangling pointers that point to empty files or even wrong files
- There is also difficulty ensuring there is no cycles in a graph – complexity associated with it
Extra overhead.





General Graph Directory





General Graph Directory (Cont.)

- How do we guarantee no cycles?
 - Allow only links to file not subdirectories – sometime not convenient
 - Every time a new link is added use a cycle detection algorithm to determine whether there is a cycle or not – time consuming

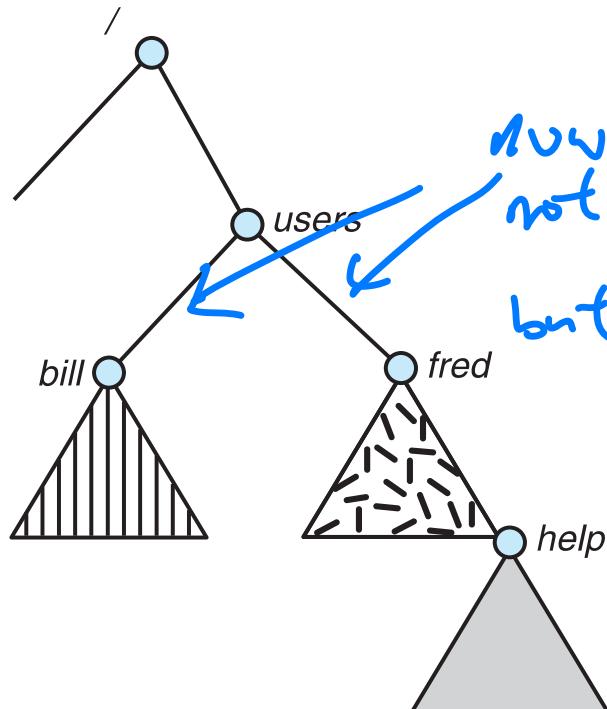
high overhead!..
indeed!



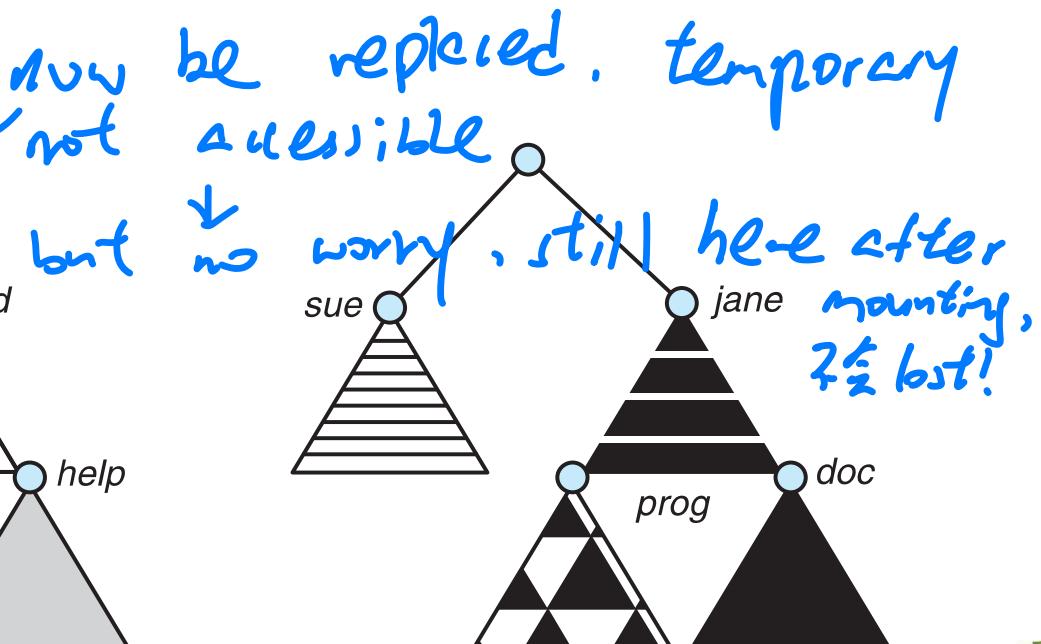


File System Mounting

- A file system must be **mounted** before it can be accessed – just like a file must be **opened** before it is used
- A unmounted file system (i.e., Fig. (b)), to be mounted at a **mount point**

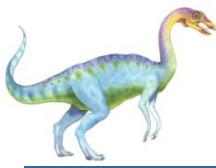


(a)



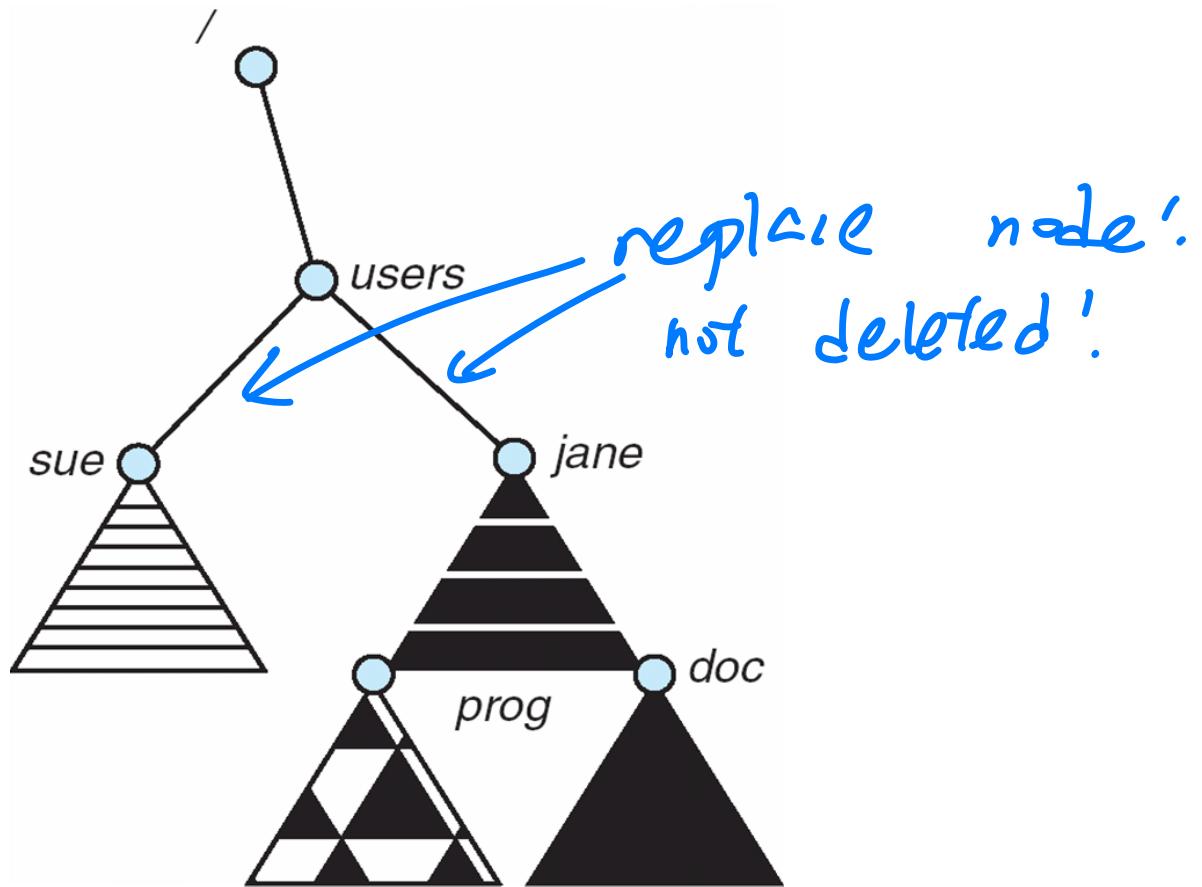
(b)





Mount Point

- Volume is mounted at `/users`





File Sharing

3 students SP

- Sharing of files in multi-user systems is desirable
- Sharing may be done through a protection scheme *not must be protected*
- In distributed systems, files may be shared across a network
 - Network File System (NFS) is a common distributed file-sharing method
- With a multi-user system *add permission*
 - User IDs identify users, allowing permissions and protections to be per-user *owner and groups*
 - Group IDs allow users to be in groups, permitting group access rights
 - Owner of a file / directory
 - Group of a file / directory





Protection

- File owner/creator of the file should be able to control:
 - what can be done
 - by whom
- Types of access
 - Read
 - Write
 - Execute
 - Append
 - Delete
 - List

control

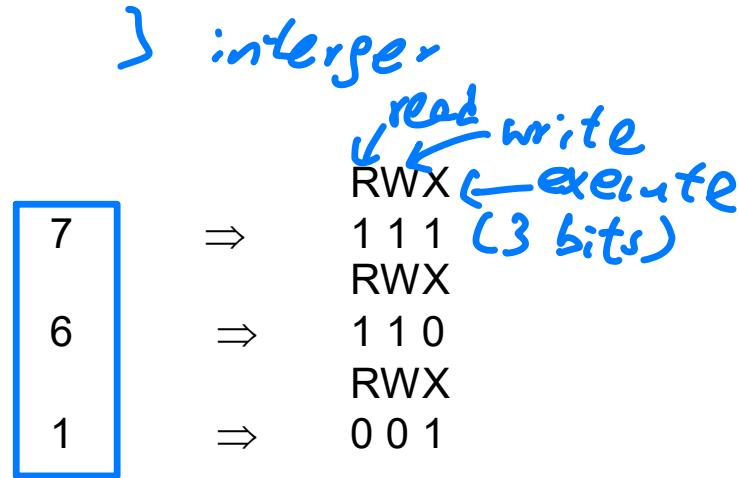




Access Lists and Groups

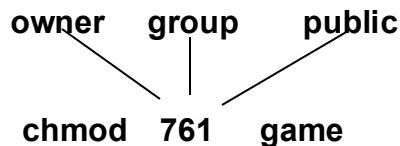
- Mode of access: read, write, execute
- Three classes of users on Unix / Linux

- a) owner access
- b) group access
- c) public access



which one else can do what

- Ask manager to create a group (unique name), say G, and add some users to the group.
- For a particular file or subdirectory, define an appropriate access.



Attach a group to a file

↳ can

chgrp G game
have different groups





Windows 10 Access-Control List Management

java Properties

General	Compatibility	Digital Signatures
Security	Details	Previous Versions

Object name: C:\Program Files\Java\jre1.8.0_92\bin\java.exe

Group or user names:

- ALL APPLICATION PACKAGES
- SYSTEM
- Administrators (LL11-SXIAO\Administrators) **(selected)**
- Users (LL11-SXIAO\Users)

To change permissions, click Edit.

Edit...

Permissions for Administrators	Allow	Deny
Full control	✓	
Modify	✓	
Read & execute	✓	
Read	✓	
Write	✓	
Special permissions		

For special permissions or advanced settings, click Advanced.

Advanced

OK Cancel Apply

Almost like
Linux, but
Linux only
have read
write execute!





A Sample UNIX Directory Listing

d = directory
- = folder

-rw-rw-r--	1	pbg	staff	31200	Sep 3 08:30	intro.ps
drwx-----	5	pbg	staff	512	Jul 8 09:33	private/
drwxrwxr-x	2	pbg	staff	512	Jul 8 09:35	doc/
drwxrwx---	2	pbg	student	512	Aug 3 14:13	student-proj/
-rw-r--r--	1	pbg	staff	9423	Feb 24 2003	program.c
-rwxr-xr-x	1	pbg	staff	20471	Feb 24 2003	program
drwx--x--x	4	pbg	faculty	512	Jul 31 10:31	lib/
drwx-----	3	pbg	staff	1024	Aug 29 06:52	mail/
drwxrwxrwx	3	pbg	staff	512	Jul 8 09:35	test/



End of Chapter 13

