

Update/Outline

- Previously we have discussed
 - Probability of Error for Orthogonal M-ary modulation
 - Union Bound
 - Some signal types
- We will now consider
 - M-ary Modulation Types
 - MFSK, MPSK
 - Tradeoffs

Another



M-ary Modulation Types

- We have seen how to:
 - Design an optimum M-ary Receiver
 - Calculating the probability of bit and symbol errors
 - Considered some general transmission signal types
- Now we wish to discuss specific but popular modulation formats and determine their properties and when to use which one

M-ary Phase-Shift Keying (MPSK)

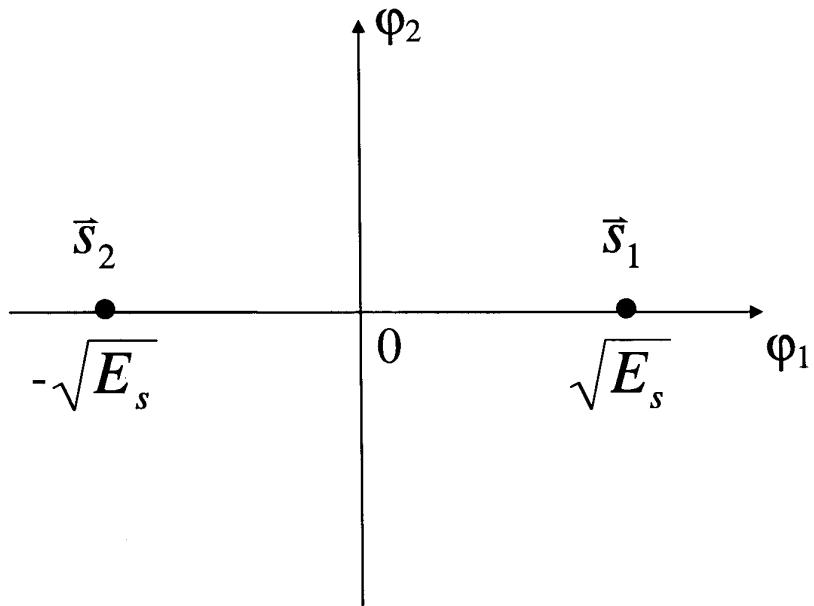
$$s_k(t) = \sqrt{\frac{2E_s}{T_s}} \cos \left[\omega_c t + \underbrace{\frac{2\pi(k-1)}{M}}_{\theta_k} \right] \quad 0 \leq t \leq T_s, k = 1, 2, \dots, M$$

phase angle'.

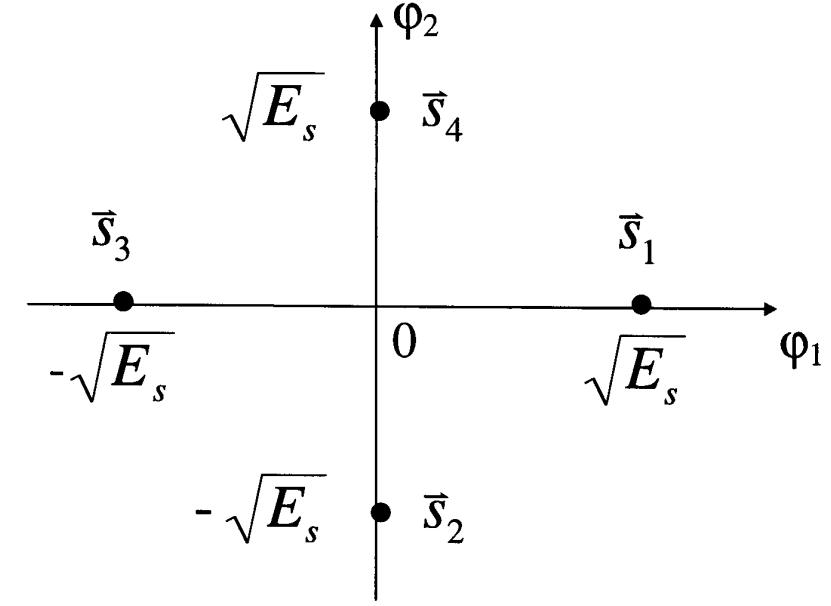
- Assuming that $\omega_c = \frac{\text{integer} \times 2\pi}{T_s}$

$$\Rightarrow \begin{cases} \phi_1(t) = \sqrt{\frac{2}{T_s}} \cos(\omega_c t) \\ \phi_2(t) = \sqrt{\frac{2}{T_s}} \sin(\omega_c t) \end{cases}$$

Examples of MPSK constellation



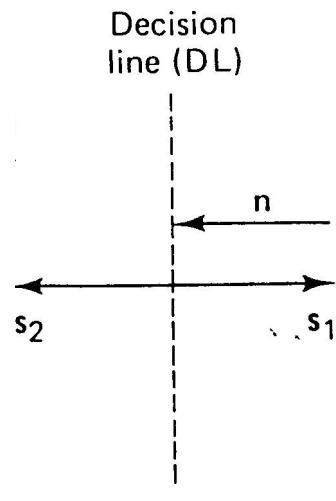
$M = 2$ (2-PSK \equiv BPSK)



$M = 4$ (4-PSK)

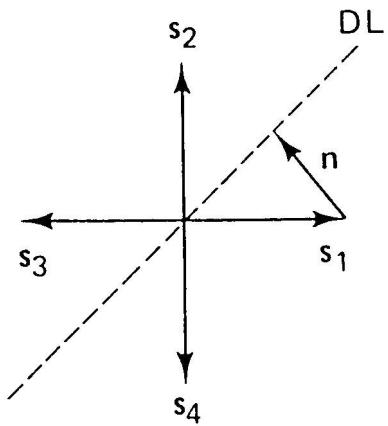
Also called **QPSK**

MPSK signal sets for $M=2,4,8,16$



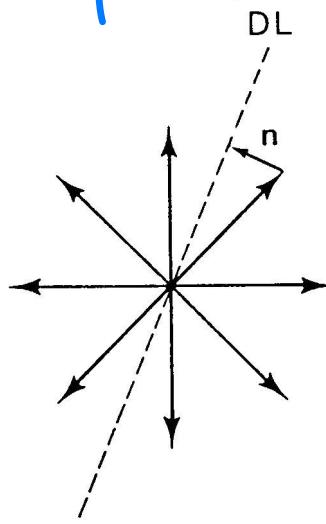
$M = 2$

(a)



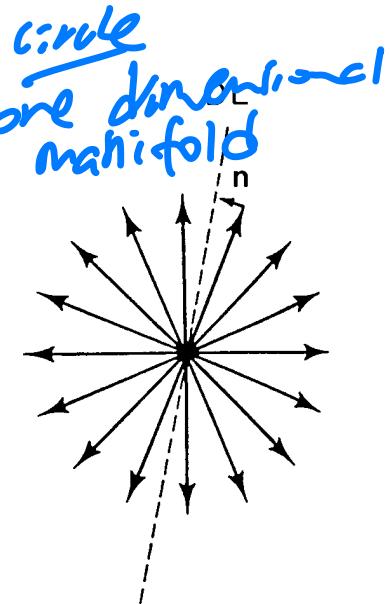
$M = 4$

(b)



$M = 8$

(c)



$M = 16$

(d)

partition of circle
one dimensional
manifold

Error Performance of MPSK

- It can be shown that

$$P_{eM} = \frac{1}{\pi} \int_0^{\pi(1-1/M)} \exp \left[-\frac{(E_s / N_0) \sin^2(\pi / M)}{\sin^2 \phi} \right] d\phi$$

i.e. Eqn 4-98 in Ziemer and Peterson

Except for $M = 2$ and 4, numerical integration is needed.

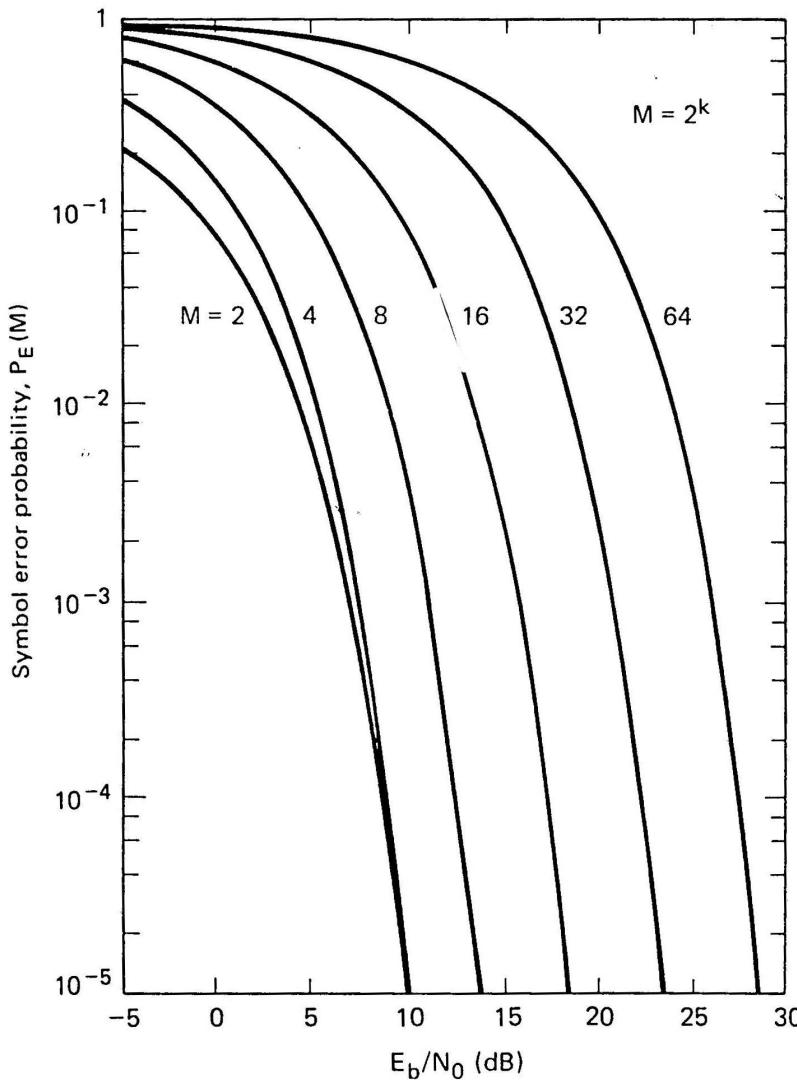
Alternatively, we can use the tight upper and lower bounds below.

confidence!

$$Q \left[\sqrt{\frac{2E_s}{N_0}} \sin(\pi / M) \right] \leq P_{e,M} \leq 2Q \left[\sqrt{\frac{2E_s}{N_0}} \sin(\pi / M) \right]$$

only have 2 bottleneck!

SER vs Eb/No for MPSK



Spreading
more and more
points in
Constellation
diagram

Figure 3.32 Symbol error probability for coherently detected multiple phase signaling. (Reprinted from W. C. Lindsey and M. K. Simon, *Telecommunication Systems Engineering*, Prentice-Hall, Inc., Englewood Cliffs, N.J., 1973, courtesy of W. C. Lindsey and Marvin K. Simon.)

For large Es/No:

$$P_{e,M} \approx 2Q\left[\sqrt{\frac{2E_s}{N_0}} \sin(\pi / M)\right]$$



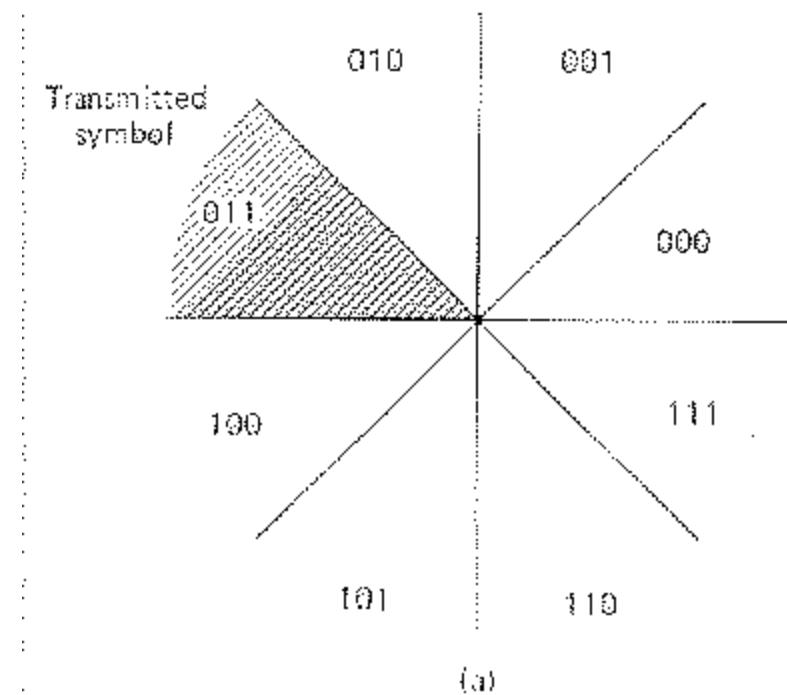
Very tight for fixed M as E_s/N₀ increases.

How about Bit error prob?

Different from M-ary orthogonal, it depends on the bit labeling scheme used.

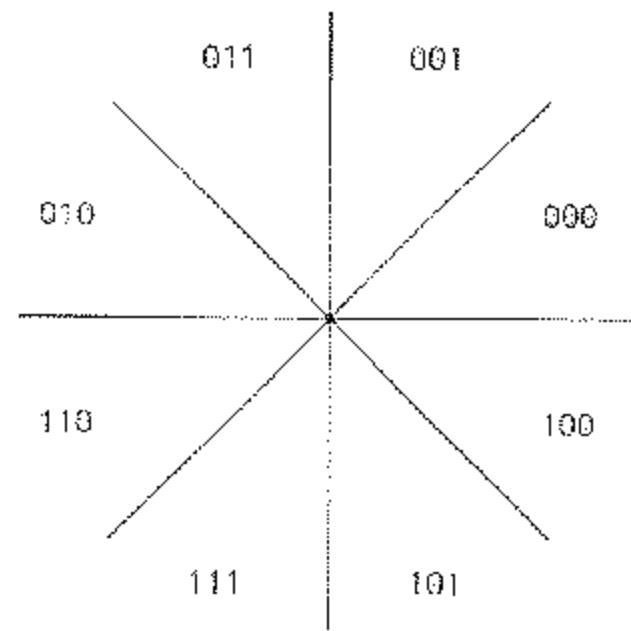
Typically Gray Coding is used for MPSK.
8

Binary Code



(a)

only have 1 bit
error!!!
Gray Code SNR ↑



(b)

Figure 3.36 Binary-coded versus Gray-coded decision regions in an MPSK signal space. (a) Binary coded. (b) Gray coded.

Error depends on
how you label
the symbols!

Gray Coding

- Allows representation of symbols or **bit-to-symbol mapping**
- In going from one symbol to an **adjacent** symbol, **only one bit** out of the k (or n bits in text) bits **changes**.
- An adjacent symbol error (i.e. the most likely symbol error) will therefore be accompanied by one and only one bit error.
- Thus, the **bit error** probability of Gray-coded MPSK can be well approximated by

$$P_b \cong \frac{P_{e,M}}{\log_2 M} \quad \text{至多1bit error!}$$

M-QAM Modulation

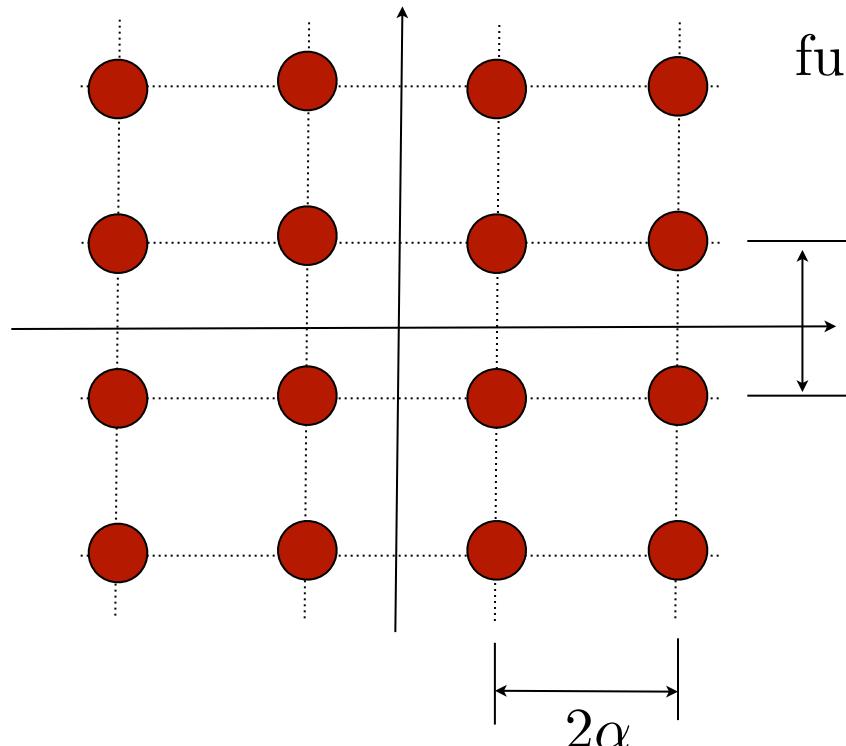
2-degrees of freedom:

Time Domain Description

$$s_k(t) = a_k \cos(\omega_c t) + b_k \sin(\omega_c t) \text{ for } t \in [0, T_s], k = \{1, 2, \dots, M\}$$

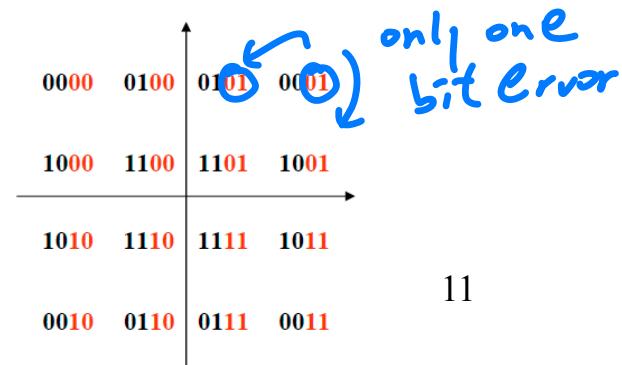
where $a_k, b_k \in \{\pm\alpha, \pm 3\alpha, \dots, \pm(\sqrt{M} - 1)\alpha\}$ and $M = \{4, 16, 64, 256, \dots\}$

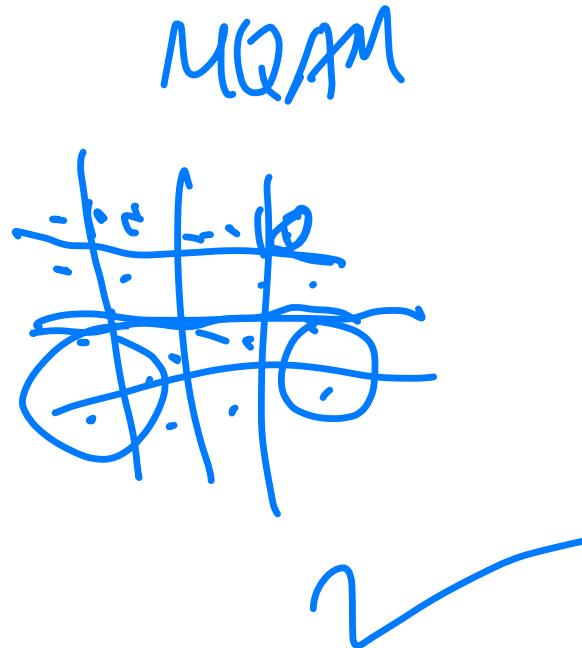
Geometric Domain Description



2-dim signal space with basis functions $\{\cos(\omega_c t), \sin(\omega_c t)\}$

Rectangular M-QAM signals can be generated easily using I-Q modulator





$$4 \cdot PQM = 4 \cdot QAM$$

$$\underline{16 \text{-PAM}} \rightarrow \underline{16 \text{ Cells}}$$

MQAM Performance

- Bit Rate: $R_b = \frac{1}{T_s} \log_2 M$
- Average Symbol Energy: $E_s = \frac{2}{\sqrt{M}} (2\alpha^2 + 2(3\alpha)^2 + \dots)$
 - e.g. for 16QAM, $E_s = \frac{2}{4} (2\alpha^2 + 2(3\alpha)^2) = 10\alpha^2$
- Average Transmit Power: $P_s = \frac{E_s}{T_s}$
- Remarks:
 - Not all M points have the same energy.
 - Information can be visualized as carried by the amplitude and phase information of the points

MQAM Performance

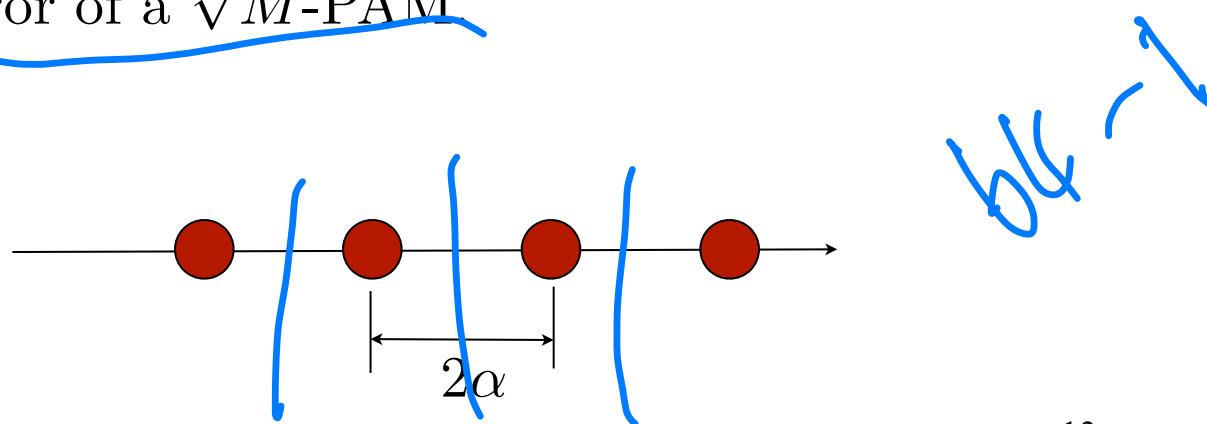
- Symbol Error Probability (SER)

Analysis:

$$4\text{-PAM} \times 4\text{-PAM} = 16\text{-QAM}$$

- M-QAM can be regarded as two independent M-PAM, each having $\sqrt{M} = 2^{k/2}$ points
- Probability of correct decision for M-QAM is given by:

$P_c = (1 - P_{\sqrt{M}})^2$ where $P_{\sqrt{M}}$ is the probability of error of a \sqrt{M} -PAM



M-QAM Performance

give directly in the question!

- It can be shown that

$$P_{\sqrt{M}} = 2 \left(1 - \frac{1}{\sqrt{M}} \right) Q \left(\sqrt{\frac{3}{M-1} \frac{E_s}{N_0}} \right)$$

symbol error rate! - edge!

- Hence, $1 - (1 - 2P_{\sqrt{M}} + P_{\sqrt{M}}^2)^2 =$

$$P_M = 1 - (1 - P_{\sqrt{M}})^2 \approx 2P_{\sqrt{M}} = 4 \left(1 - \frac{1}{\sqrt{M}} \right) Q \left(\frac{3}{M-1} \frac{E_s}{N_0} \right)$$

*if $\frac{E_s}{N_0} \gg 1$
exp. function:*

- Recall that the MPSK SER is given by:

$$P_M \approx 2Q \left(\sin \frac{\pi}{M} \sqrt{\frac{2E_s}{N_0}} \right)$$

$$\begin{aligned} E_s &= \log_2 M b \\ \bar{E}_s &= q \bar{b} b. \end{aligned}$$

- Compare M-QAM with M-PSK SER:

- The gain of M-QAM over M-PSK is given

by: $gain = \frac{3/(M-1)}{2 \sin^2(\pi/M)}$

M↑, gain ↑↑

M-QAM Union Bound

- Instead of working out the exact P_e , the union bound is simpler and is asymptotically accurate for high SNR.

$$P_e \leq 4Q\left(\sqrt{\frac{d_{min}^2}{2N_0}}\right) \approx 4Q\left(\sqrt{\frac{(2\alpha)^2}{2N_0}}\right) \approx 4Q\left(\sqrt{\frac{3E_s}{(M-1)N_0}}\right)$$

On average 4 nearest neighbors

- Compare with the exact P_e , they are very close especially for large Es/No.

Choices of M-ary Modulations

- There are a number of factors that one needs to consider in order to pick a modulation scheme.
For example:
 - Power efficiency (inversely proportional to **Eb/No**)
 - Bandwidth efficiency (bit rate/bandwidth)
- M-ary modulation allows us to trade in power,
bit rate and bandwidth

Error Probability Performance Curves

- Allow us to design and set an operating point for a system
- Consider MFSK
 - Increasing M can provide an improvement in P_b , or reduction in the E_b/N_0 required, at the cost of increased bandwidth
- Consider MPSK
 - Increasing M can provide a reduction in bandwidth requirement, at the cost of degraded P_b , or increase in the E_b/N_0 requirement

Theoretical Limits on Performance

- **Channel Capacity** is the theoretical upper bound for the maximum rate at which information could be transmitted without error (*Shannon 1948*)
 - For a bandlimited channel that is corrupted by AWGN ($N_0(f) = \frac{N_0}{2}$) the maximum rate achievable is given by

$$R = B \log_2 (1 + SNR) = B \log_2 \left(1 + \frac{P_s}{N_0 B}\right)$$

↑ Bandwidth: *↖ Power*

no need BER !

$b_1 b_2 \dots \rightarrow$ $\boxed{\text{Tx}}$ $\circlearrowleft (\text{Aug10})$

\circlearrowright $\boxed{\text{Rx}}$ $\xrightarrow{\text{11}} b_1 b_2 \dots$

$b_1 b_2 b_3 c_1 \dots$

$n(t) \sim \text{random}$ $\xrightarrow{T} p_e$

$\hat{b}_1 \hat{b}_2 \hat{b}_3$
 $m \downarrow \text{mining}$
o

Spectral Efficiency

Shannon Limit

R : bit rate
 B : bandwidth

- $N = N_0 B$, hence normalize for efficiency:

$$\frac{R_b}{B} = \log_2 \left(1 + \frac{P}{N_0 B} \right) = \log_2 \left(1 + \frac{P}{N} \right)$$

- Next note that

Energy Efficiency SNR

$$\frac{E_b}{N_0} = \frac{PT}{N_0} = \frac{P}{RN_0} = \frac{PB}{RN_0 B} = \frac{P}{N} \frac{B}{R}$$

- Then

$$\frac{E_b}{N_0} = \frac{B}{R} \left(2^{R/B} - 1 \right)$$

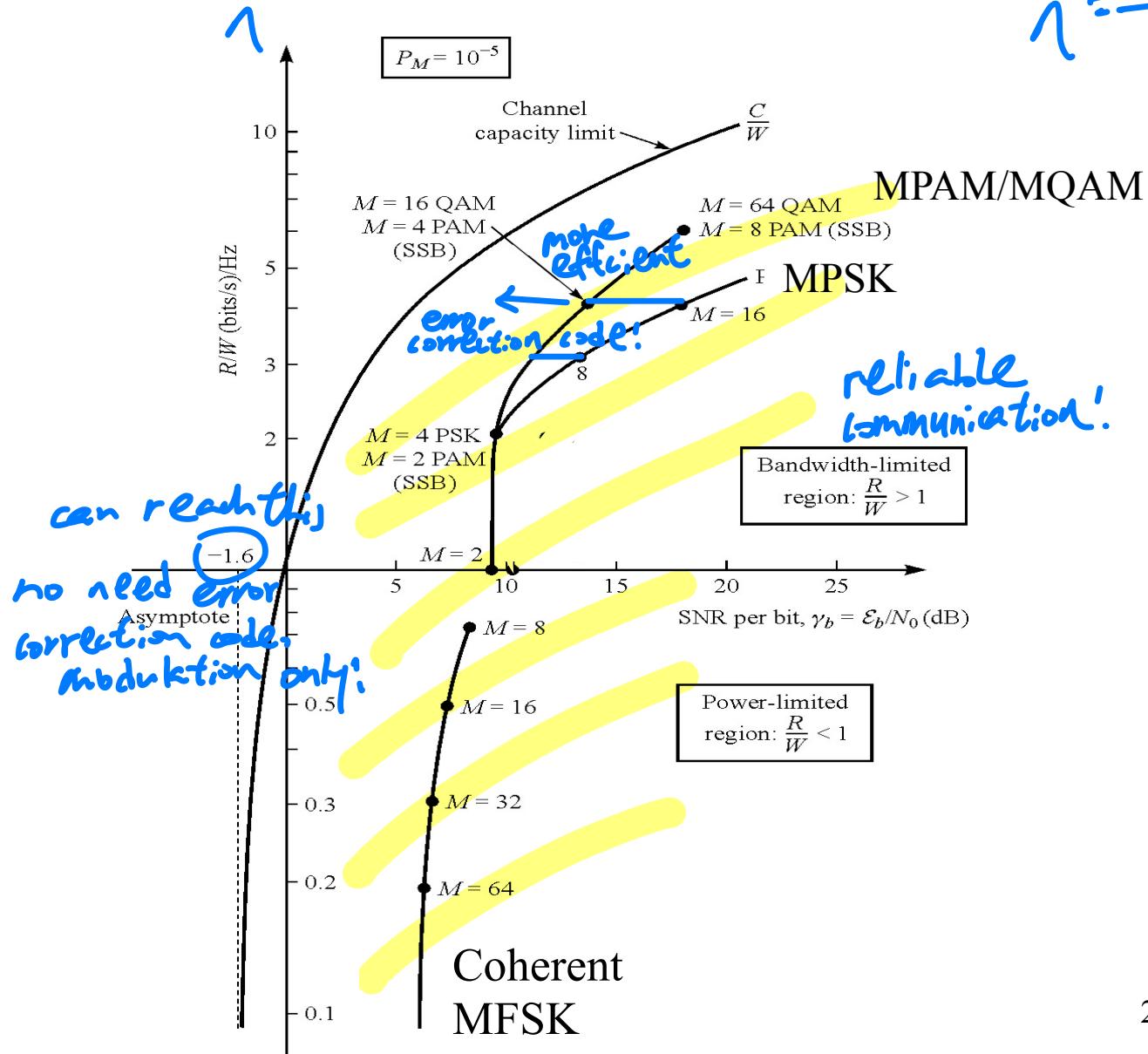
increasing function for η

$$\begin{aligned} \eta &= \frac{R_b}{B} \\ n &= \log_2 \left(1 + \eta \frac{E_b}{N_0} \right) \\ \frac{E_b}{N_0} &= \frac{n-1}{n^{19}} \end{aligned}$$

Another form!

Bandwidth- efficiency plane

$$N = \frac{\log_2 M}{M}$$



$$\left(\frac{\bar{E}_h}{N_0} \right)_{\min} = \lim_{n \rightarrow \infty} \frac{2^n - 1}{n}$$

MTJK, by very large M

Trade-Offs

- *Power-Limited Systems*: Power scarce but bandwidth available
 - Improved Pb by expanding bandwidth (for a given Eb/N0) or required Eb/N0 can be reduced by expanding bandwidth (for a given Pb)
- *Bandwidth-Limited Systems*: bandwidth scarce
 - Maximize R over the bandlimited channel at the expense of Eb/N0 (for a given Pb)

Shannon Limit

- In the limit as R/B goes to 0, we get

$$\frac{E_b}{N_0} = \frac{1}{\log_2 e} = 0.693 = -1.59 dB$$

This value is called the **Shannon Limit** (what is the relationship with k going to infinity)

Received Eb/N0 must be $> -1.6dB$ for reliable communications to be possible

Summary of M-ary Modulation Schemes

	M-FSK	M-PSK	M-QAM
Bit Rate	$R_b = \frac{1}{T_s} \log_2(M)$	$R_b = \frac{1}{T_s} \log_2(M)$	$R_b = \frac{1}{T_s} \log_2(M)$
BW (Bandpass)	$BW = \frac{M}{T_s}$	$BW = \frac{1}{T_s}$	$BW = \frac{1}{T_s}$ Size of letter
Average Transmit Power	$\frac{E_s}{T_s}$	$\frac{E_s}{T_s}$	$\frac{4\alpha^2}{\sqrt{M}T_s} \sum_{i=1}^{\frac{\log_2 M/2}{2}} (2^i - 1)^2$
Average Symbol Error Probability (SER)	$P_e \leq (M-1)Q\left(\sqrt{\frac{E_s}{N_0}}\right)$	$P_e \leq 2Q\left(\sqrt{\frac{2E_s}{N_0}} \sin(\pi/M)\right)$	$P_M \approx 4\left(1 - \frac{1}{\sqrt{M}}\right) Q\left(\sqrt{\frac{3}{M-1} \frac{E_s}{N_0}}\right)$ ↓ q nearest neighbours
Remarks	<ul style="list-style-type: none"> Orthogonal Signaling Schemes (Equi-energy points & Mutually orthogonal signals) Enhance Energy Efficiency at the expense of extra BW 	<ul style="list-style-type: none"> Equi-energy constellation (information carried by phase values only) Dimension of the signal set is always 2 (I-Q modulator) Enhance spectral efficiency at the expense of extra power 	<ul style="list-style-type: none"> Points are NOT equi-energy Information is carried by both amplitude and phase Enhance spectral efficiency at the expense of extra power Better than M-PSK for $M > 4$.