CS-GY 9223 D: Lecture something Linear Programming

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IMPROVING GRADIENT DESCENT

We now have a good understanding of gradient descent.

Number of iterations for ϵ error:

	G-Lipschitz	eta-smooth
R bounded start	$O\left(\frac{G^2R^2}{\epsilon^2}\right)$	$O\left(\frac{\beta R^2}{\epsilon}\right)$
lpha-strong convex	$O\left(\frac{G^2}{\alpha\epsilon}\right)$	$O\left(\frac{\beta}{\alpha}\log(1/\epsilon)\right)$

How do we use this understanding to design faster algorithms?



ACCELERATED GRADIENT DESCENT

Nesterov's accelerated gradient descent:

$$\begin{aligned} \cdot \ & \mathbf{x}^{(1)} = \mathbf{y}^{(1)} = \mathbf{z}^{(1)} \\ \cdot \ & \text{For } t = 1, \dots, T \\ \cdot \ & \mathbf{y}^{(t+1)} = \mathbf{x}^{(t)} - \frac{1}{\beta} \nabla f(\mathbf{x}^{(t)}) \\ \cdot \ & \mathbf{x}^{(t+1)} = \left(1 + \frac{\sqrt{\kappa} - 1}{\sqrt{\kappa} + 1}\right) \mathbf{y}^{(t+1)} + \frac{\sqrt{\kappa} - 1}{\sqrt{\kappa} + 1} \left(\mathbf{y}^{(t+1)} - \mathbf{y}^{(t)}\right) \end{aligned}$$

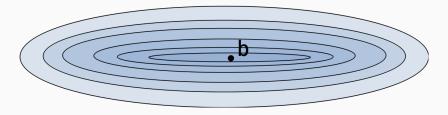
Theorem (AGD for β -smooth, α -strongly convex.)

Let f be a β -smooth and α -strongly convex function. If we run AGD for T steps we have:

$$f(\mathbf{x}^{(t)}) - f(\mathbf{x}^*) \le \kappa e^{-(t-1)\sqrt{\kappa}} \left[f(\mathbf{x}^{(1)}) - f(\mathbf{x}^*) \right]$$

Corollary: If $T = O(\sqrt{\kappa} \log(\kappa/\epsilon))$ achieve error ϵ .

INTUITION BEHIND ACCELERATION



Level sets of $\|\mathbf{A}\mathbf{x} - \mathbf{b}\|_2^2$.

Other terms for similar ideas:

- Momentum
- Heavy-ball methods

What if we look back beyond two iterates?



Main idea: Instead of minimizing f(x), find another function g(x) with the same minimum but which is better suited for first order optimization (e.g., has a smaller conditioner number).

Claim: Let
$$h(\mathbf{x}): \mathbb{R}^d \to \mathbb{R}^d$$
 be an invertible function. Let $g(\mathbf{x}) = f(h(\mathbf{x}))$. Then

$$\min_{\mathbf{x}} f(\mathbf{x}) = \min_{\mathbf{x}} g(\mathbf{x})$$
 and $\underset{\mathbf{x}}{\operatorname{arg min}} f(\mathbf{x}) = h \left(\underset{\mathbf{x}}{\operatorname{arg min}} g(\mathbf{x}) \right)$.

First Goal: We need g(x) to still be convex.

Claim: Let P be an (invertible) $d \times d$ matrix and let $g(\mathbf{x}) = f(P\mathbf{x})$.

 $g(\mathbf{x})$ is always convex.

If $y^* = \arg \min g(y)$, then $x^* = Py^* \min \operatorname{minimizes} f(x)$.

Second Goal:

 $g(\mathbf{x})$ should have better condition number κ than $f(\mathbf{x})$.

High dimensional chain rule:

If
$$g(\mathbf{x}) = f(\mathbf{P}\mathbf{x})$$
, $\nabla^2 g(\mathbf{x}) = \nabla^2 \mathbf{P}^T f(\mathbf{P}\mathbf{x}) \mathbf{P}$.

Recall that the condition number is equal to:

$$\max_{\mathbf{x}} \frac{\lambda_{\max} \left(\nabla^2 g(\mathbf{x}) \right)}{\lambda_{\min} \left(\nabla^2 g(\mathbf{x}) \right)}$$

Example:

•
$$f(\mathbf{x}) = \|\mathbf{A}\mathbf{x} - \mathbf{b}\|_{2}^{2}$$
. $\nabla f(\mathbf{x}) = 2\mathbf{A}^{\mathsf{T}}\mathbf{A}$. $\kappa_{f} = \frac{\lambda_{1}(\mathbf{A}^{\mathsf{T}}\mathbf{A})}{\lambda_{d}(\mathbf{A}^{\mathsf{T}}\mathbf{A})}$.
• $g(\mathbf{x}) = \|\mathbf{A}\mathbf{P}\mathbf{x} - \mathbf{b}\|_{2}^{2}$. $\nabla g(\mathbf{x}) = 2\mathbf{P}^{\mathsf{T}}\mathbf{A}^{\mathsf{T}}\mathbf{A}\mathbf{P}$ $\kappa_{g} = \frac{\lambda_{1}(\mathbf{P}^{\mathsf{T}}\mathbf{A}^{\mathsf{T}}\mathbf{A}\mathbf{P})}{\lambda_{d}(\mathbf{P}^{\mathsf{T}}\mathbf{A}^{\mathsf{T}}\mathbf{A}\mathbf{P})}$.

Ideal preconditioner: Choose P so that $\mathbf{P}^T \mathbf{A}^T \mathbf{A} \mathbf{P} = \mathbf{I}$. For example, could set $P = \sqrt{(\mathbf{A}^T \mathbf{A})^{-1}}$. But obviously this is too expensive to compute.

DIAGONAL PRECONDITIONER

Third Goal: P should be easy to compute.

Many, many problem specific preconditioners are used in practice. There design is usually a heuristic process.

Example: Diagonal preconditioner for least squares problems.

- Let $D = diag(A^TA)$
- Want PA^TAP to be close to identity I.
- Let $P = \sqrt{D^{-1}}$

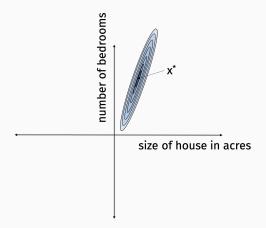
P is often called a Jacobi preconditioner. Often works very well in practice!

DIAGONAL PRECONDITIONER

A = -734 33 9111 0 -31 108 5946 -19 232 3502 101 10 426 0 -65 12503 -373 26 9298 -236 -2 -94 2398 2024 -132 -6904 -25 -2258 -1 92 -6516 2229 11921 -22 338 -5 -16118 -23

DIAGONAL PRECONDITIONER INTUITION

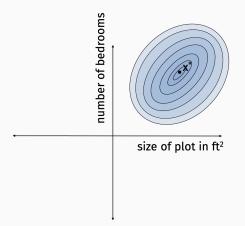
 $g(\mathbf{x}) = f(\|\mathbf{A}\mathbf{P}\mathbf{x} - \mathbf{b}\|_2^2)$ is the same least squares problem as $f(\mathbf{x}) = \|\mathbf{A}\mathbf{x} - \mathbf{b}\|_2^2$, but with each feature (column of **A**) scaled differently. The i^{th} column is scaled by P_{ii} .



Feature scaling can have a huge impact on conditioning.

DIAGONAL PRECONDITIONER INTUITION

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ADAPTIVE STEPSIZES

Another view: If g(x) = f(Px) then $\nabla g(x) = P^T \nabla f(Px)$.

 $\nabla g(\mathbf{x}) = \mathbf{P} \nabla f(\mathbf{P} \mathbf{x})$ when **P** is symmetric.

Gradient descent on g:

• For
$$t = 1, ..., T$$
,
• $\mathbf{x}^{(t+1)} = \mathbf{x}^{(t)} - \eta \mathbf{P} \left[\nabla f(\mathbf{P}\mathbf{x}^{(t)}) \right]$

Gradient descent on g:

• For
$$t = 1, ..., T$$
,
• $\mathbf{y}^{(t+1)} = \mathbf{y}^{(t)} - \eta \mathbf{P}^2 \left[\nabla f(\mathbf{y}^{(t)}) \right]$

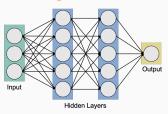
When **P** is diagonal, this is just gradient descent with a different step size for each parameter!

ADAPTIVE STEPSIZES

Less clear how to set P for general optimization problems where the Hessian is changing, but lots of heuristic algorithms based on this idea:

- · AdaGrad, AdaDelta
- · RMSprop
- · Adam optimizer

(Pretty much all of the most widely used optimization methods for training neural networks.)





STOCHASTIC METHODS

Main idea: Trade slower convergence (more iterations) for cheaper iterations.

Stochastic Gradient Descent: When $f(\mathbf{x}) = \sum_{i=1}^{n} f_i(\mathbf{x})$, approximate $\nabla f(\mathbf{x})$ with $\nabla f_i(\mathbf{x})$ for randomly chosen i.

STOCHASTIC METHODS

Main idea: Trade slower convergence (more iterations) for cheaper iterations.

Stochastic Coordinate Descent: Only compute a <u>single random</u> entry of $\nabla f(\mathbf{x})$ on each iteration:

$$\nabla f(\mathbf{x}) = \begin{bmatrix} \frac{\partial f}{\partial x_1}(\mathbf{x}) \\ \frac{\partial f}{\partial x_2}(\mathbf{x}) \\ \vdots \\ \frac{\partial f}{\partial x_d}(\mathbf{x}) \end{bmatrix} \qquad \nabla_i f(\mathbf{x}) = \begin{bmatrix} 0 \\ \frac{\partial f}{\partial x_i}(\mathbf{x}) \\ \vdots \\ 0 \end{bmatrix}$$

Update: $\mathbf{x}^{(t+1)} \leftarrow \mathbf{x}^{(t)} - \eta \nabla_i f(\mathbf{x}^{(t)})$.

COORDINATE DESCENT

When \mathbf{x} has d parameters, computing $\nabla_i f(\mathbf{x})$ sometimes costs just a 1/d fraction of what it costs to compute $\nabla f(\mathbf{x})$

Example: $f(\mathbf{x}) = \|\mathbf{A}\mathbf{x} - \mathbf{b}\|_2^2$ for $\mathbf{A} \in \mathbb{R}^{n \times d}, \mathbf{x} \in \mathbb{R}^d, \mathbf{b} \in \mathbb{R}^n$.

- $\cdot \nabla f(\mathbf{x}) = 2\mathbf{A}^{\mathsf{T}}\mathbf{A}\mathbf{x} 2\mathbf{A}^{\mathsf{T}}\mathbf{b}.$
- · $\nabla_i f(\mathbf{x}) = 2 \left[\mathbf{A}^\mathsf{T} \mathbf{A} \mathbf{x} \right]_i 2 \left[\mathbf{A}^\mathsf{T} \mathbf{b} \right]_i$

Computing full gradient takes O(nd) time. Can we do better here?

COORDINATE DESCENT

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$$\cdot \nabla f(\mathbf{x}) = 2\mathbf{A}^{\mathsf{T}}\mathbf{A}\mathbf{x} - 2\mathbf{A}^{\mathsf{T}}\mathbf{b}.$$

•
$$\nabla_i f(\mathbf{x}) = 2 [\mathbf{A}^T \mathbf{A} \mathbf{x}]_i - 2 [\mathbf{A}^T \mathbf{b}]_i$$
.

- $Ax^{(t+1)} = A(x^{(t)} + c \cdot e_i)$
- 2 $\left[\mathbf{A}^{\mathsf{T}}\left(\mathbf{A}\mathbf{x}^{(t+1)}-\mathbf{b}\right)\right]_{i}$

O(n) time

O(n) time

Stochastic Coordinate Descent:

- Choose number of steps T and step size η .
- For $t = 1, \ldots, T$:
 - Pick random $j \in 1, ..., d$ uniformly at random.
 - $\mathbf{x}^{(t+1)} = \mathbf{x}^{(t)} \eta \nabla_j f(\mathbf{x}^{(i)})$
- Return $\hat{\mathbf{x}} = \frac{1}{T} \sum_{t=1}^{T} \mathbf{x}^{(t)}$.

Theorem (Stochastic Coordinate Descent convergence)

Given a G-Lipschitz function f with minimizer \mathbf{x}^* and initial point $\mathbf{x}^{(1)}$ with $\|\mathbf{x}^{(1)} - \mathbf{x}^*\|_2 \le R$, SCD with step size $\eta = \frac{1}{Rd}$ satisfies the guarantee:

$$\mathbb{E}[f(\hat{\mathbf{x}}) - f(\mathbf{x}^*)] \le \frac{2GR}{\sqrt{T/d}}$$

IMPORTANCE SAMPLING

Often it doesn't make sense to sample *i* uniformly at random:

$$\mathbf{A} = \begin{bmatrix} 0 & 0 & 1 & 0 & 0 & 0 \\ 0 & 0 & 2 & 0 & 0 & 0 \\ 0 & 0 & -1 & 0 & 0 & 0 \\ 0 & 0 & -.5 & 0 & 0 & 0 \\ 0 & 0 & 3 & 0 & 0 & 0 \\ 0 & 0 & -2 & 0 & 0 & 0 \end{bmatrix} \qquad \mathbf{b} = \begin{bmatrix} 10 \\ 42 \\ -11 \\ -51 \\ 34 \\ -22 \end{bmatrix}$$

Select indices *i* proportional to $\|\mathbf{a}_i\|_2^2$:

$$\Pr[\text{select index } i \text{ to update}] = \frac{\|\mathbf{a}_i\|_2^2}{\sum_{i=1}^d \|\mathbf{a}_i\|_2^2} = \frac{\|\mathbf{a}_i\|_2^2}{\|\mathbf{A}\|_F^2}$$

Let's analyze this approach.

Specialization of SCD to $\|\mathbf{A}\mathbf{x} - \mathbf{b}\|_2^2$:

Randomized Coordinate Descent (Strohmer, Vershynin 2007 / Leventhal, Lewis 2018)

• For iterate $\mathbf{x}^{(t)}$, let $\mathbf{r}^{(t)}$ be the <u>residual</u>:

$$\mathbf{r}^{(t)} = \mathbf{A}\mathbf{x}^{(t)} - \mathbf{b}$$

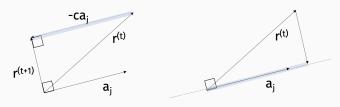
- $\mathbf{x}^{(t+1)} = \mathbf{x}^{(t)} c\mathbf{e}_{j}$.
- $\mathbf{r}^{(t+1)} = \mathbf{r}^{(t)} c\mathbf{a}_j$. Here \mathbf{a}_j is the i^{th} column of \mathbf{A} .

Typically *c* depends on fixed learning rate. Here we will choose it <u>optimally</u> – similar idea to gradient descent with <u>line search</u>.

What choice for *c* minimizes $\|\mathbf{r}^{(t+1)}\|_2^2$?

$$\|\mathbf{r}^{(t+1)}\|_{2}^{2} = \|\mathbf{r}^{(t)} - c\mathbf{a}_{j}\|_{2}^{2}$$

· Requires projecting $\mathbf{r}^{(t)}$ onto perpendicular of \mathbf{a}_{j} .



$$\cdot c = \frac{\mathbf{a}_j^\mathsf{T} \mathbf{r}^{(t)}}{\|\mathbf{a}_j\|_2^2}$$

Note that
$$\|\mathbf{r}^{(t+1)}\|_2^2 = \|\mathbf{r}^{(t)}\|_2^2 - \|c\mathbf{a}_j\|_2^2 = \|\mathbf{r}^{(t)}\|_2^2 - \frac{(\mathbf{a}_j^{\mathsf{T}}\mathbf{r}^{(t)})^2}{\|\mathbf{a}_j\|_2^2}$$

Specialization of SCD to $\|\mathbf{A}\mathbf{x} - \mathbf{b}\|_2^2$:

Randomized Coordinate Descent

- · Choose number of steps T.
- Let $x^{(1)} = 0$ and $r^{(1)} = b$.
- For t = 1, ..., T:
 - Pick random $j \in 1, ..., d$. Index j is selected with probability proportional to $\|\mathbf{a}_i\|_2^2/\|\mathbf{A}\|_F^2$.
 - Set $c = \mathbf{a}_{i}^{T} \mathbf{r}^{(t)} / \|\mathbf{a}_{j}\|_{2}^{2}$
 - $\mathbf{x}^{(t+1)} = \dot{\mathbf{x}^{(t)}} c\mathbf{e}_i$
 - $\cdot \mathbf{r}^{(t+1)} = \mathbf{r}^{(t)} c\mathbf{a}_j$
- Return $\mathbf{x}^{(T)}$.

CONVERGENCE

Claim

$$\mathbb{E} \|\mathbf{r}^{(t+1)}\|_{2}^{2} = \|\mathbf{r}^{(t)}\|_{2}^{2} - \frac{1}{\|\mathbf{A}\|_{F}^{2}} \|\mathbf{A}^{T}\mathbf{r}^{(t)}\|_{2}^{2}$$

CONVERGENCE

Any residual \mathbf{r} can be written as $\mathbf{r} = \mathbf{r}^* + \overline{\mathbf{r}}$ where $\mathbf{r}^* = \mathbf{A}\mathbf{x}^* - \mathbf{b}$ and $\overline{\mathbf{r}} = \mathbf{A}(\mathbf{x}^t - \mathbf{x}^*)$. Note that $\mathbf{A}^T \mathbf{r}^* = \mathbf{0}$ and $\overline{\mathbf{r}} \perp \mathbf{r}^*$.

Claim

$$\mathbb{E}\|\bar{\mathbf{r}}^{(t+1)}\|_{2}^{2} \leq \|\bar{\mathbf{r}}^{(t)}\|_{2}^{2} - \frac{\lambda_{\min}(\mathbf{A}^{T}\mathbf{A})}{\|\mathbf{A}\|_{F}^{2}}$$

$$\mathbb{E}\|\overline{\mathbf{r}}^{(t+1)}\|_{2}^{2} + \|\mathbf{r}^{*}\|_{2}^{2} \leq \|\overline{\mathbf{r}}^{(t)}\|_{2}^{2} + \|\mathbf{r}^{*}\|_{2}^{2} - \frac{1}{\|\mathbf{A}\|_{F}^{2}}\|\mathbf{A}^{T}\overline{\mathbf{r}}^{(t)}\|_{2}^{2}\|\overline{\mathbf{r}}^{(t)}\|_{2}^{2}$$

Exercise: Because \bar{r} is in the column span of A,

$$\|\mathbf{A}^{\mathsf{T}}\mathbf{\bar{r}}^{(t)}\|_{2}^{2} \geq \lambda_{\min}(\mathbf{A}^{\mathsf{T}}\mathbf{A})\|\mathbf{\bar{r}}^{(t)}\|_{2}^{2}$$

CONVERGENCE

Theorem (Randomized Coordinate Descent convergence)

After T steps of RCD with importance sampling run on $f(\mathbf{x}) = \|\mathbf{A}\mathbf{x} - \mathbf{b}\|_2^2$, we have:

$$\mathbb{E}[f(\mathbf{x}^{(t)}) - f(\mathbf{x}^*)] \le \left(1 - \frac{\lambda_{\min}(\mathbf{A}^\mathsf{T}\mathbf{A})}{\|\mathbf{A}\|_F^2}\right)^t [f(\mathbf{x}^{(0)}) - f(\mathbf{x}^*)]$$

Corollary: After $T = O(\frac{\|\mathbf{A}\|_F^2}{\lambda_{\min}(\mathbf{A}^T\mathbf{A})}\log \frac{1}{\epsilon})$ we obtain error $\epsilon \|\mathbf{b}\|_2^2$.

Is this more or less iterations than the $T = O(\frac{\lambda_{\max}(\mathbf{A}^T \mathbf{A})}{\lambda_{\min}(\mathbf{A}^T \mathbf{A})} \log \frac{1}{\epsilon})$ required for gradient descent to converge?

Recall useful linear algebraic fact:

$$\|\mathbf{A}\|_F^2 = \operatorname{tr}(\mathbf{A}^T\mathbf{A}) = \sum_{i=1}^d \lambda_i(\mathbf{A}^T\mathbf{A})$$

$$\lambda_{\max}(\mathbf{A}^T\mathbf{A}) \leq \|\mathbf{A}\|_F^2 \leq d \cdot \lambda_{\max}(\mathbf{A}^T\mathbf{A})$$

For solving $\|\mathbf{A}\mathbf{x} - \mathbf{b}\|_2^2$,

 $(\# GD | \text{Iterations}) \le (\# RCD | \text{Iterations}) \le d \cdot (\# GD | \text{Iterations})$

But RCD iterations are cheaper by a factor of *d*.

COMPARISON

When does
$$\|\mathbf{A}\|_F^2 = \operatorname{tr}(\mathbf{A}^T \mathbf{A}) = \mathbf{d} \cdot \lambda_{\max}(\mathbf{A}^T \mathbf{A})$$
?

When does
$$\|\mathbf{A}\|_F^2 = \text{tr}(\mathbf{A}^T \mathbf{A}) = \mathbf{1} \cdot \lambda_{\text{max}}(\mathbf{A}^T \mathbf{A})$$
?

COMPARISON

Roughly:

Stochastic Gradient Descent performs well when <u>data points</u> (rows) are repetitive.

Stochastic Coordinate Descent performs well when <u>data</u> <u>features</u> (columns) are repetitive.



STATIONARY POINTS

We understand much less about optimizing non-convex functions in comparison to convex functions, but not nothing. In many cases, we're still figuring out the right questions to ask

Definition (Stationary point)

For a differentiable function f, a stationary point is any \mathbf{x} with:

$$\nabla f(\mathbf{x}) = \mathbf{0}$$

local/global minima - local/global maxima - saddle points

STATIONARY POINTS

Reasonable goal: Find an approximate stationary point $\hat{\mathbf{x}}$ with

$$\|\nabla f(\hat{\mathbf{x}})\|_2^2 \le \epsilon.$$

SMOOTHNESS FOR NON-CONVEX FUNTIONS

Definition

A differentiable (potentially non-convex) function f is β smooth if <u>for all</u> \mathbf{x} , \mathbf{y} ,

$$\|\nabla f(\mathbf{x}) - \nabla f(\mathbf{y})\|_2 \le \beta \|\mathbf{x} - \mathbf{y}\|_2$$

Corollary: For all x, y

$$\left|\nabla f(\mathbf{x})^{\mathsf{T}}(\mathbf{x}-\mathbf{y})-[f(\mathbf{x})-f(\mathbf{y})]\right|\leq \frac{\beta}{2}\|\mathbf{x}-\mathbf{y}\|_{2}^{2}.$$

GRADIENT DESCENT FINDS APPROXIMATE STATIONARY POINTS

Theorem

If GD is run with step size $\eta = \frac{1}{\beta}$ on a differentiable function f with global minimum \mathbf{x}^* then after $T = O(\frac{\beta[f(\mathbf{x}^{(1)}) - f(\mathbf{x}^*)]}{\epsilon})$ we will find an ϵ -approximate stationary point $\hat{\mathbf{x}}$.

$$\cdot \nabla f(\mathbf{x}^{(t)})^{\mathsf{T}}(\mathbf{x}^{(t)} - \mathbf{x}^{(t+1)}) - f(\mathbf{x}^{(t)}) + f(\mathbf{x}^{(t+1)}) \le \frac{\beta}{2} \|\mathbf{x}^{(t)} - \mathbf{x}^{(t+1)}\|_{2}^{2}.$$

•
$$f(\mathbf{x}^{(t+1]}) - f(\mathbf{x}^{(t)}) \le \frac{\beta}{2} \eta^2 \|\nabla f(\mathbf{x}^{(t)})\|_2^2 - \eta \|\nabla f(\mathbf{x}^{(t)})\|_2^2$$

$$f(\mathbf{x}^{(t+1])}) - f(\mathbf{x}^{(t)}) \le \frac{-\eta}{2} \|\nabla f(\mathbf{x}^{(t)})\|_2^2$$

$$\cdot \frac{1}{T} \sum_{t=1}^{T} \frac{\eta}{2} ||f(\mathbf{x}^{(t)})||_{2}^{2} \leq \frac{1}{T} \sum_{t=1}^{T} f(\mathbf{x}^{(t)}) - f(\mathbf{x}^{(t+1)})$$

$$\cdot \frac{\eta}{2} \min_{t} \|f(\mathbf{x}^{(t)})\|_{2}^{2} \le \frac{1}{T} \left[f(\mathbf{x})^{(1)} - f(\mathbf{x})^{(T)} \right]$$

QUESTIONS IN NON-CONVEX OPTIMIZATION

If GD can find a stationary point, are there algorithms which find a stationary point faster using preconditioning, acceleration, stocastic methods, etc.?

QUESTIONS IN NON-CONVEX OPTIMIZATION

What if my function only has global minima and stationary points? Randomized methods (SGD, perturbed gradient methods, etc.) can "escape" stationary points under some minor assumptions.

Example: $min_x \frac{-x^T A^T A x}{x^T x}$

- Global minimum: Top eigenvector of A^TA (i.e., top principal component of A).
- · Stationary points: All other eigenvectors of A.

Useful for lots of other matrix factorization problems beyond vanilla PCA.