High-Performance Decimal Floating-Point Arithmetic: Algorithms and Implementation in C++

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This article presents a comprehensive, portable C++ implementation of decimal floating-point arithmetic conforming to IEEE 754-2019 standards[8]. Our system offers three IEEE 754-compliant types, and three additional types that prioritize performance over strict standard adherence, providing flexible options for various computational needs. We describe in detail the Decimal system architecture, its standard library, and usage guidelines. The implementation incorporates novel algorithms for key operations, significantly improving performance in common use cases. Rigorous testing results demonstrate the system's correctness and IEEE 754 compliance. Performance benchmarks show competitive or superior results compared to existing implementations. This work addresses the growing demand for precise decimal arithmetic in financial, scientific, and engineering applications, offering a robust, efficient solution for C++ developers.

CCS Concepts: • General and reference \rightarrow Design; Performance; • Mathematics of computing \rightarrow Mathematical software performance; • Software and its engineering \rightarrow Software architectures.

Additional Key Words and Phrases: C++, Decimal Floating Point, Object Oriented, Special Functions, System Architecture

ACM Reference Format:

1 INTRODUCTION

Decimal floating-point arithmetic is crucial for many applications[4], particularly in financial and scientific computing. While several C++ decimal floating-point packages exist, they often lack IEEE 754[8] conformance, interoperability with the C++ Standard Template Library (STL), or both, and have limited portability. This paper presents a novel decimal system that addresses these limitations and advances decimal floating-point technology. Our system is standalone, relies on a minimal subset of the C++ STL, and has been tested on a wide range of devices, from S390X mainframes to AVR boards. It provides seamless interoperability with existing C++ standard types and full IEEE 754 conformance, features not collectively offered by any known package. This work significantly enhances the toolset available for high-precision decimal computations across diverse computing environments.

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2 THE DECIMAL SYSTEM

2.1 Background

Decimal floating point is a method of representing floating point numbers using base-10 instead of base-2 like most are familiar with. Decimal floating point is not a new concept per se. Mechanical devices such as the abacus and slide rule use base-10 numbers. Early computers such as the IBM 650 used decimal floating point numbers prior to IEEE standardization[2]. Some modern architectures support decimal floating point in hardware such as IBM System Z[15]. The advantage of decimal floating point is that it can represent human readable floating point numbers exactly unlike binary floating point.

The major difference between the layout of binary and decimal floating-point numbers lies in their internal structure and how they represent the same numerical value. Figure 1 illustrates this difference by showing the bit layout of the same number in both formats.

Binary floating-point numbers consist of three parts:

- The sign bit (shown in blue)
- The exponent (shown in green)
- The significand (shown in red)

Decimal floating-point numbers, on the other hand, are composed of four distinct parts:

- The sign bit (shown in blue)
- The combination field (shown in orange)
- The exponent continuation (shown in green)
- The coefficient continuation (shown in red)

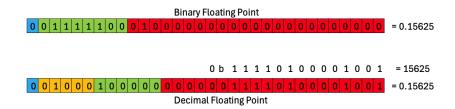


Fig. 1. Bit layouts of Binary and Decimal Floating Point Numbers

Another idiosyncrasy that is found in decimal floating point types are called cohorts[8]. A cohort is all the different ways to represent the same value. For example one can represent the number 10 as: 0.1×10^2 , 1×10^1 , 10×10^0 , 100×10^{-1} , and so on. The ramification of this is we must normalize the significand and exponents of a number before we can attempt doing comparisons or other basic mathematical operations.

2.2 Provided Types

The decimal system provides 6 total types: decimal32, decimal64, decimal128, decimal32_fast, decimal64_fast, and decimal128_fast. The first three types are conformant to IEEE-754 standards while the latter three provided identical numerical results without the constraints of IEEE-754.

The three conformant types is the design specifications provided in IEEE-754 namely their properties:

| Parameter | decimal32 | decimal64 | decimal128 |
|--------------------------------|-----------|-----------|------------|
| Storage Width | 32 | 64 | 128 |
| Precision (decimal digits) | 7 | 16 | 34 |
| Max Exponent | 96 | 384 | 6144 |
| Exponent Bias | 101 | 398 | 6176 |
| Sign Width | 1 | 1 | 1 |
| Combination Field Width | 11 | 13 | 17 |
| Significand Continuation Width | 20 | 50 | 110 |

Table 1. Decimal Floating-Point Format Parameters

The three so-called fast types (i.e. decimalXX_fast) have the same values of precision, range, and exponent as their analogous conformant type, but require more space.

2.3 System Architecture

The decimal system architecture is robust and flexible. Each type is implemented completely independently of each other, but they share many of the same function implementations from the STL.

2.4 Using the System

Every effort was made during design and implementation to ensure that using the decimal system was straightforward and intuitive. The library contains zero dependencies, and uses only a small subsection of the C++ STL. The required language standard for the library is C++14.

```
1 #include <boost/decimal.hpp>
2 #include <iostream>
3 #include <iomanip>
5 int main()
6
  {
      using namespace boost::decimal;
7
8
g
      constexpr decimal32 val_1 {100};
                                                 // Construction from an
      integer
      constexpr decimal32 val_2 {10, 1};
                                                // Construction from an
10
      integer and exponent
      constexpr decimal32 val_3 {1, 2, false}; // Construction from an
11
      integer, exponent, and sign
12
       std::cout << "Val_1: " << val_1 << '\n'
13
                 << "Val_2: " << val_2 << '\n'
14
```

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```
15
                  << "Val_3: " << val_3 << '\n';
16
17
       if (val_1 == val_2 && val_2 == val_3 && val_1 == val_3)
18
           std::cout << "All equal values" << std::endl;</pre>
19
20
21
22
       constexpr decimal64 val_4 {decimal64{2, -1} + decimal64{1, -1}};
23
       constexpr double float_val_4 {0.2 + 0.1};
       const decimal64 val_5 { float_val_4 }; // Explicit Conversion from
24
       double
25
26
       std::cout << std::setprecision(17) << "Val_4: " << val_4 << '\n'
                      "Float: " << float_val_4 << '\n'
27
                     "Val_5: " << val_5 << '\n';
28
29
30
       if (val_4 == val_5)
31
32
           std::cout << "Floats are equal" << std::endl;</pre>
33
34
       else
35
36
           std::cout << "Floats are not equal" << std::endl;</pre>
37
38
39
       return 0;
40 }
```

Listing 1. Basic Usage

In Listing 1 we show how to construct the type and that it works just like a built-in flowing point type. The following is an example that leverages one of decimal-floating points main strengths and one that we expect many people will use.

```
1 #include <boost/decimal.hpp>
2 #include <iostream>
3
  #include <cassert>
4
5 int main()
6
  {
7
      using namespace boost::decimal;
8
      decimal64 val {0.25}; // Construction from a double (not recommended
      but explicit construction is allowed)
10
      char buffer[256];
11
       auto r_to = to_chars(buffer, buffer + sizeof(buffer) - 1, val);
12
      assert(r_to); // checks std::errc()
13
      *r_to.ptr = '\0';
14
15
```

```
16
       decimal64 return_value;
       auto r_from = from_chars(buffer, buffer + std::strlen(buffer),
17
      return_value);
       assert(r_from);
18
19
20
       assert(val == return_value);
21
22
       std::cout << " Initial Value: " << val << '\n'
                 << "Returned Value: " << return_value << std::endl;
23
24
25
       return 0;
26 }
```

Listing 2. Basic Usage

In Listing 2 we show how to serialize and parse numbers. These techniques and functions are an extension of Boost.Charconv[3].

3 IMPLEMENTATION DETAILS

3.1 IEEE 754 Conformant Decimal Types

The decimal system provides three IEEE-754 compliant types: decimal 32, decimal 64, and decimal 128. Each of these are constructed using a single unsigned integer of the same width to hold the bits. For example decimal 32 consists of a single std::uint32_t.decimal 128 uses a custom implementation of a 128-bit unsigned integer for portability reasons rather than relying on the existence of unsigned __int128. The discussion of implementing big integers is outside the scope of the paper, but papers and implementations can be found in numerous places [3][12][10].

The decoding step of decimal32 is somewhat involved because the values of the significand and exponent depend on the bits in the combination field. Our first case is where the bits in the combination field are in the form 00XXX, 01XXX, and 10XXX.

```
1 // Comb. Exponent Significand
2 // s 00 TTT (00)eeeeee (0TTT)[sssssssssss][sssssssss]
3 // s 01 TTT (01)eeeeee (0TTT)[sssssssssss][sssssssss]
4 // s 10 TTT (10)eeeeee (0TTT)[sssssssssss][sssssssss]
```

Listing 3. Combination Field Pattern 1

In Listing 3 we can see that with these combination field bit patterns the first two bits of the combination field are appended to the front of the exponent, the next three bits are appended to the front of the significand. The second case is where the bits in the combination field are in the form 1100X, 1101X, and 1110X.

```
1 // Comb. Exponent Significand
2 // s 1100 T (00)eeeeee (100T)[sssssssssss][sssssssss]
3 // s 1101 T (01)eeeeee (100T)[sssssssssss][sssssssss]
4 // s 1110 T (10)eeeeee (100T)[sssssssssss][sssssssss]
```

Listing 4. Combination Field Pattern 2

In Listing 4 we can see that with these combination field bit patterns there is now an implied leading bit to the significand, followed by 00T bits. Bits 2 and 3 are now what represents the leading bits of the exponent as opposed to 0 and one like in Listing 3.

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The final case is where the bits in the combination field are in the form 1111X which is used for non-finite numbers:

```
1 // Comb. Exponent Significand
2 // s 1111 0 eeeeee [ssssssssss][sssssssss] = inf
3 // s 1111 1 eeeeee [ssssssssss][sssssssss] = qnan
4 // s 1111 1 1eeeee [sssssssssss][sssssssss] = snan
```

Listing 5. Combination Field Pattern 3

Since this combination field pattern is only for non-finite numbers it makes the implementation of isnan, isinf, etc. straightforward. We can also see the signbit remains unused so we can represent signed infinities. One can also use the remaining exponent and significand bits to provide a payload to NaN analogous to the use of std::nan[6].

Now that we have seen what each bit pattern of the combination field corresponds to; the full implementation of decoding the significand then becomes:

```
1 constexpr auto decimal32::full_significand() const noexcept ->
      significand_type
2
  {
3
      significand_type significand {};
5
      if ((bits_ & detail::d32_comb_11_mask) == detail::d32_comb_11_mask)
6
           // Only need the one bit of T because the other 3 are implied
7
           significand = (bits_ & detail::d32_comb_11_significand_bits) ==
      detail::d32_comb_11_significand_bits ?
           static_cast < std::uint32_t > (0b1001'0000000000'000000000) :
9
           static_cast < std::uint32_t > (0b1000'000000000'0000000000);
10
11
12
      else
13
           // Last three bits in the combination field, so we need to shift
14
      past the exp field
           // which is next
15
16
           significand |= (bits_ & detail::
      d32_comb_00_01_10_significand_bits) >> detail::d32_exponent_bits;
17
18
19
       significand |= (bits_ & detail::d32_significand_mask);
20
       return significand;
21
22 }
```

Listing 6. Decoding decimal 32 significand

Encoding the value is a similar to the decoding process found in Listing 6, but more complex process in that we need to work backwards through the steps to figure out what the value of the combination field needs to be, and then we use a series of masks to write the value.

3.2 IEEE 754 Fast Types

Now that we have discussed the three IEEE-754 conformant types we will turn our attention to much more performant, but non-compliant types. The library offers: decimal32_fast, decimal64_fast, and decimal128_fast. These types are similar to how instead of std::uint32_t you can use std::uint_fast32_t which on x86_64 platforms generally aliases to std::uint64_t. This is the classic optimization of trading space for time. Again we will focus on the 32-bit type for clarity and simplicity. Unlike decimal32 consisting of a data std::uint32_t, decimal32_fast actually contains its internal state in a structure.

```
1 class decimal32 fast final
2 {
3 public:
      using significand_type = std::uint_fast32_t;
      using exponent_type = std::uint_fast8_t;
5
      using biased_exponent_type = std::int_fast32_t;
6
8
  private:
      // In regular decimal32 we have to decode the 24 bits of the
      significand and the 8 bits of the exp
10
      // Here we just use them directly at the cost of at least 2 extra
      bytes of internal state
      // since the fast integer types will be at least 32 and 8 bits
11
      respectively
12
       significand_type significand_ {};
       exponent_type exponent_ {};
14
      bool sign_ {};
16
17
       // Continued
18 };
```

Listing 7. Internal state of decimal32_fast

In Listing 7 we see that now instead of having to encode and decode the number every time we perform an operation we can now just access the value directly:

```
1 constexpr auto full_significand() const noexcept -> significand_type
2 {
3    return significand_;
4 }
```

Listing 8. Decoding decimal32_fast significand

Now compare the decoding process found in the conformant type in Listing 6 to that in Listing 8. Encoding the value is similarly trivial. Each of the fast types use the same approach, but with different types used for the internal state to ensure that the range and precision is equal to that of the IEEE 754 conformant types.

3.3 Standard Library

3.3.1 Special Functions. The implementation of transcendental, and C++17 mathematical special functions[5][9] extensively uses Páde Approximations and Remez Polynomials.

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4 RESULTS

4.1 Testing and Precision

4.2 Performance

As this system is implemented in software it will never be as performant as binary floating point is on computing systems with floating-point hardware. For example the table below shows the runtime difference for comparison operations between float, double, and the types provided by this system:

| Type | Runtime (µs) | Ratio to double |
|----------------|--------------|-----------------|
| float | 8,587 | 1.376 |
| double | 6,240 | 1.000 |
| decimal32 | 275,597 | 44.166 |
| decimal64 | 296,929 | 47.587 |
| decimal32_fast | 99,664 | 15.972 |
| decimal64_fast | 102,132 | 16.367 |

Table 2. Comparison of runtime and ratio for comparison operations

The gap in performance increases even further for the multiplication operation which is typically only a few cycles in hardware[1]:

| Type | Runtime (µs) | Ratio to double |
|----------------|--------------|-----------------|
| float | 1,646 | 0.957 |
| double | 1,720 | 1.000 |
| decimal32 | 313,219 | 182.104 |
| decimal64 | 583,818 | 339.429 |
| decimal32_fast | 86,093 | 50.054 |
| decimal64_fast | 333,582 | 193.943 |

Table 3. Comparison of runtime and ratio for multiplication operations

For floating-point operations that are implemented in software like parsing and serializing numbers the performance gap is quite smaller:

| Type | Runtime (μs) | Ratio to double |
|-----------|--------------|-----------------|
| float | 235,816 | 0.953 |
| double | 247,307 | 1.000 |
| decimal32 | 366,682 | 1.483 |
| decimal64 | 485,965 | 1.965 |

Table 4. Comparison of runtime and ratio for from_chars

5 CONCLUSION AND OUTLOOK

The portable C++ decimal system has been presented. The system allows for users to easily employ decimal-floating point numbers in their existing code bases. For the first time a complete and interoperable system has been fully provided.

| Type | Runtime (μs) | Ratio to double |
|-----------|--------------|-----------------|
| float | 316,300 | 1.040 |
| double | 304,272 | 1.000 |
| decimal32 | 406,053 | 1.335 |
| decimal64 | 678,451 | 2.230 |

Table 5. Comparison of runtime and ratio for to_chars

During the course of our research and implementation we have limited ourselves to 32, 64, and 128-bit types. Further research can be conducted into making higher precision types as the properties of such numbers are specified in IEEE 754. We could also expand our range of compatibility to allow the types to be massively parallelized such as providing support for CUDA devices.

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