

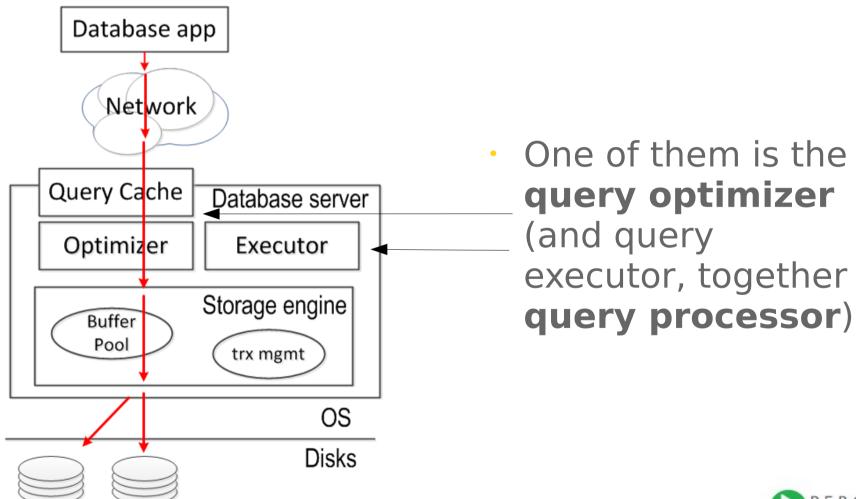
# Improving MySQL/MariaDB query performance through optimizer tuning

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# What is an optimizer and why do I need to tune it

#### Why does one need optimizer tuning

Database performance is affected by many factors



#### What is a query optimizer?

- Converts SQL into execution instructions
- Like GPS gives driving directions from A to B



		select_typ		•	table	-			possible_keys	key	1	k
	•	PRIMARY		•	part	•	ALL	•	PRIMARY	NULL	- <del>-</del>	]
1	ı	PRIMARY		ı	partsupp	Ι	ref	Ι	PRIMARY, i ps partkey, i ps suppkey	i ps partkey	1	
1	ı	PRIMARY		ı					PRIMARY, i s nationkey	PRIMARY	1	
1	ı	PRIMARY		1	nation	1	eq_ref	Τ	PRIMARY,i_n_regionkey	PRIMARY	1	
1	١	PRIMARY		1	region	1	ALL	1	PRIMARY	NULL	1	
2	١	DEPENDENT	SUBQUERY	1	partsupp	1	ref	1	PRIMARY, i_ps_partkey, i_ps_suppkey	i_ps_partkey	1	
2	١	DEPENDENT	SUBQUERY	1					PRIMARY, i_s_nationkey	PRIMARY	1	
2	ı	DEPENDENT	SUBQUERY	1	nation	1	eq_ref	1	PRIMARY,i_n_regionkey	PRIMARY	1	
2	1	DEPENDENT	SUBQUERY	1	region	1	eq ref	1	PRIMARY	PRIMARY	1	

#### Can the optimizer work poorly?

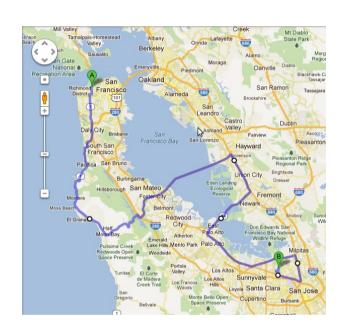
#### Yes, just like your GPS

It can produce a sub-optimal query plan (compare to GPS's directions)

Unneeded CPU/IO usage in the database is the same extra mileage

How can one tell if they have a problem in query optimizer?

- Like with the GPS:
  - Knowledge of possible routes
  - Common sense



This is what this tutorial is about



#### Start from a database performance problem

- Database response time is too high
  - or the server denies connections altogether
    - because running queries have occupied all @@max\_connections threads.
- Monitoring (cacti, etc) shows high CPU/ disk utilization
- Is this an optimizer problem?
  - Maybe.

#### Signs that it's caused by the optimizer

- Only certain kinds of queries are affected
  - Unlike concurrency/disk/binlog problems that cause slowdown for all queries
- These queries can be caught with
  - Slow query log
    - pt-query-digest
  - SHOW PROCESSLIST
  - SHOW PROFILE

#### Optimizer problems in Slow query log

Percona server/MariaDB:--log\_slow\_verbosity=query\_plan

```
# Thread_id: 1 Schema: dbt3sf10 QC_hit: No
# Query_time: 2.452373 Lock_time: 0.000113 Rows_sent: 0 Rows_examined: 1500000
# Full_scan: Yes Full_join: No Tmp_table: No Tmp_table_on_disk: No
# Filesort: No Filesort_on_disk: No Merge_passes: 0
SET timestamp=1333385770;
select * from customer where c_acctbal < -1000;
```

- Look for big Query\_time (>> Lock\_time)
- Look for Rows\_examined >> Rows\_sent
  - Although there are "legitimate" cases for this
- Query\_plan members (violet), if present, also show \*possible\* problem areas

#### **SHOW PROCESSLIST;**

```
MariaDB [dbt3sf10]> show processlist;
  Id | User | Host
                                     Command |
                                                Time |
                                                       State
                                                                       Info
                                                       Sending data
   1 | root | localhost | dbt3sf10 |
                                                 176
                                                                       select
                                     Query
   2 | root | localhost | dbt3sf10
                                     Query
                                                       NULL
                                                                       show pr
             localhost | dbt3sf10
                                                       freeing items
                                                                       insert
      root l
                                     Query
                                                  19
                                                  77
             localhost | dbt3sf10
                                      Sleep
                                                                       NULL
                                                      Updating
      root | localhost | dbt3sf10
                                     Query
                                                                       update
   6 | root | localhost | dbt3sf10 | Query
                                                       statistics
                                                                       select
```

- States to look for:
  - optimizing
  - executing
  - Statistics
  - Sending data

#### SHOW PROFILE

- Uses the same set of states
- Shows where exactly the query has spent time

mysql> show profile for	or query 1;	
Status	Duration	
starting   Opening tables	0.000097     0.000038	Phases that can occupy lots of time:
System lock   Table lock   init	0.000014     0.000019     0.000027	Storage engine: transaction start, etc
optimizing statistics	0.000027	Optimization (search for query plan) Optimization and sometimes InnoDB locking
preparing   executing   Sending data	0.000023     0.000014     18.501940	Execution (of the query plan)
sending data   end   query end	0.000015     0.000006	
freeing items   logging slow query	0.000029   0.000004	Storage engine: transaction commit, binary logging
logging slow query   cleaning up	0.000050   0.000006	Diriary logging  10  PERCONA  10

#### Global observations summary

- A query [pattern] that
  - Spends a lot of time in optimization/execution phase
  - Examines many more records than you think it should
  - Uses Full\_scan, Full\_join, etc
  - · .. or all of the above
- Is a candidate for further investigation

### Reasons why a query is "heavy"

 Some queries are inherently hard to compute



 Some get a plan which looks good but turns out bad due to circumstances that were not accounted for



And some are just bad query plan choices



#### **Know thy optimizer**

- How to tell which case you're dealing with?
  - Need to know "the distance"
    - The lower bound of amount of effort one needs to expend to arrive at the destination
    - Easy in geography, not so easy in SQL
  - Need to know "the roads"
    - Available table access strategies
    - How they are composed into query plans







### Optimizer troubleshooting

From simple to complex queries

#### Let's take a simple, slow query

select \* from orders where o\_orderDate='1992-06-06' and o\_clerk='Clerk#000009506';

```
# Thread id: 2 Schema: dbt3sf10 QC hit: No
```

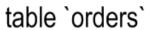
- # Query\_time: 41.129178 Lock\_time: 0.000174 Rows\_sent: 2 Rows\_examined: 15000000
- #Full scan: Yes Full join: No Tmp table: No Tmp table on disk: No
- # Filesort: No Filesort\_on\_disk: No Merge\_passes: 0

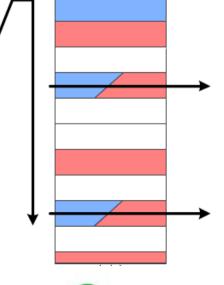
use dbt3sf10;

SET timestamp=1333432984;

select \* from orders where o\_orderDate='1992-06-06' and o\_clerk='Clerk#000009506';

- Full scan examined lots of rows
- But only a few (two) of them "have contributed to the result set"
  - It was unnecessary to read the rest
- Query plan is poor, it seems.





#### Status variables

23 rows in set (0.01 sec)

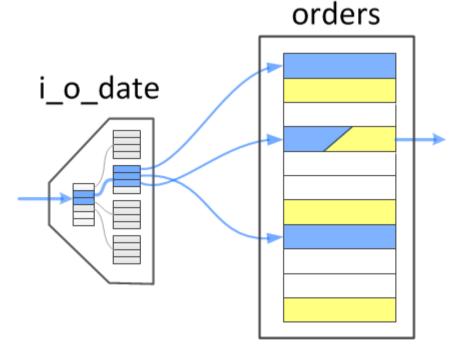
```
MariaDB [dbt3sf10]> flush status;
Query OK, 0 rows affected (0.00 sec)
MariaDB [dbt3sf10] > select * from orders where o orderDate='1992-06-06' and o clerk='Clerk#000009506';
MariaDB [dbt3sf10]> show status like 'Handler%';
| Variable name
| Handler commit
| Handler delete
| Handler discover
| Handler icp attempts
| Handler icp match
| Handler mrr init
| Handler mrr key refills
| Handler mrr rowid refills | 0
| Handler prepare
| Handler read first
| Handler read key
| Handler read next
 Handler read prev
| Handler read rnd
                                                 Indeed, 15M sequential reads
| Handler read rnd deleted
| Handler read rnd next
                            1 15000001
| Handler rollback
| Handler savepoint
| Handler savepoint rollback | 0
| Handler tmp update
| Handler tmp write
                            1 0
| Handler update
 Handler write
```

#### With index, we'll read fewer rows

```
alter table orders add key i_o_orderdate (o_orderdate); select * from orders where o_orderDate='1992-06-06' and o_clerk='Clerk#000009506';
```

id			possible_keys	<del>-</del>	key_len			
1 1	SIMPLE		i_o_orderdate		4	const	6303	Using where

- The index allows to only read records with i\_o\_orderdate='Clerk..'
  - (At the cost of having to read index entries also)
- Still not perfect: we're reading records that have the wrong o\_orderDate
- Let's check how it executes:



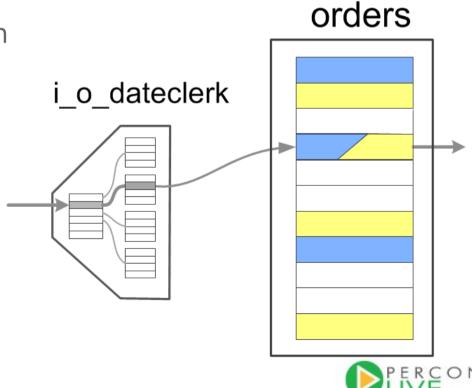
#### Status variables

```
MariaDB [dbt3sf10]> flush status;
MariaDB [dbt3sf10] > select * from orders where o orderDate='1992-06-06' and o clerk='Clerk#000009506';
... (0.39 sec)
MariaDB [dbt3sf10]> show status like 'Handler%';
| Variable_name
| Handler commit
| Handler delete
| Handler discover
| Handler icp attempts
| Handler icp match
| Handler mrr init
| Handler mrr key refills
| Handler mrr rowid refills | 0
| Handler prepare
| Handler read first
                                                  1 index lookup
 Handler read key
                                                  6K forward index reads
                            I 6122
 Handler read next
| Handler read prev
                            1 0
| Handler read rnd
| Handler read rnd deleted
| Handler read rnd next
| Handler rollback
| Handler savepoint
| Handler savepoint rollback | 0
| Handler tmp update
 Handler tmp write
 Handler update
 Handler write
23 rows in set (0.00 sec)
```

#### Can have an even better index

alter table orders add key i\_o\_date\_clerk (o\_orderdate, o\_clerk); select \* from orders where o\_orderDate='1992-06-06' and o\_clerk='Clerk#000009506';

- Added a wide index covering both columns
- Now, only records that match the whole WHERE will be read
  - Perfect!



#### Status variables

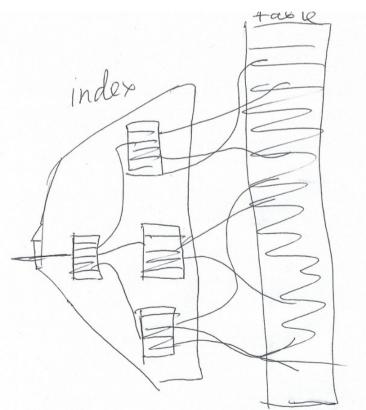
23 rows in set (0.00 sec)

```
MariaDB [dbt3sf10] > select * from orders where o orderDate='1992-06-06' and o clerk='Clerk#000009506';
... (0.00 sec)
MariaDB [dbt3sf10]> show status like 'Handler%';
 Variable name
 Handler commit
| Handler delete
| Handler discover
| Handler icp attempts
| Handler icp match
| Handler mrr init
                             1 0
| Handler mrr key refills
| Handler mrr rowid refills
| Handler prepare
| Handler read first
                                                   1 index lookup, and that's it
| Handler read key
 Handler read next
| Handler read prev
 Handler read rnd
| Handler read rnd deleted
| Handler read rnd next
| Handler rollback
| Handler savepoint
| Handler savepoint rollback | 0
| Handler tmp update
                                                         Let's try to generalize..
| Handler tmp write
| Handler update
 Handler write
```

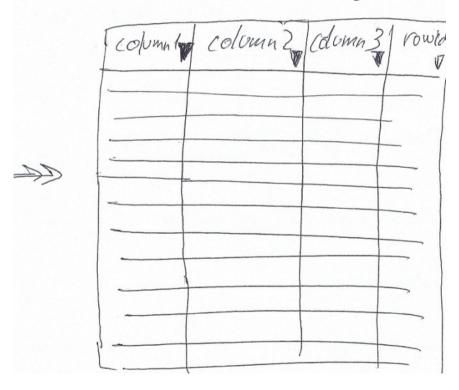
#### Index fundamentals

ALTER TABLE tbl ADD INDEX(column1, column2, column3)

On the physical level, it's a B-TREE



On a logical level: an ordered directory



## Querying data through index: 'ref'

alter table orders add index(o\_orderDATE, o\_shippriority, o\_custkey)

ref access can use any prefix:

o\_orderDATE='2011-11-02' AND o\_shippriority=2

0_	orderDATE='2011-11-08' AND
C	_shippriority=3 AND o_custkey= 4

Not allowed:

o\_shipprioirity=2

o\_shipprioirity=2

o_orderDATE	o_shipprio	o_custkey
2011-11-01	1	10
2011-11-02	1	10
2011-11-02	1	10
2011-11-02	1	11
2011-11-02	2	10
2011-11-03	2	10
2011-11-03	2	11
2011-11-04	3	9
2011-11-04	4	10
2011-11-05	3	11
2011-11-06	1	10
2011-11-08	2	4
2011-11-09	1	10

#### Ref(const) access summary

- Uses equalities to make index lookups
- #rows estimate is usually precise
- ANALYZE will not help
   EXPLAIN: type=ref, ref=const, key len

- Status variables:
  - 1 x Handler\_read\_key
  - N x Handler read next
- Estimates and control.. will talk later

#### Range access

#### Range access

- Can use equality and non-equality comparisons
- Think of index as a directed axis.
- and "ranges" on it where the WHERE condition is TRUE:

... WHERE o\_orderDATE > '2010-01-01'

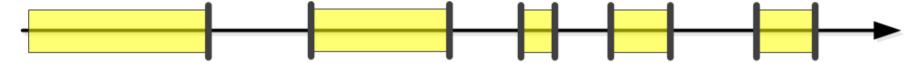


... WHERE o\_orderDATE > '2010-01-01' AND o\_orderDATE < '2010-01-07'



#### Range access (2)

- Arbitrarily deep AND/OR formulas are supported
- The optimizer will split out the "useful" parts of WHERE
- And produce a list of disjoint ranges to be scanned:



"Useful" conditions compare key with constant:

tbl.key>const

tbl.key=const

tbl.key LIKE 'foo%'

tbl.key IS NULL

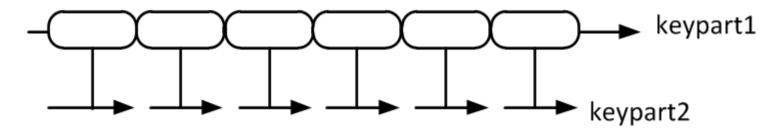
tbl.key BETWEEN c1 AND c2

tbl.key IN (const1, const2 ...)



#### Range access for multi-part keys

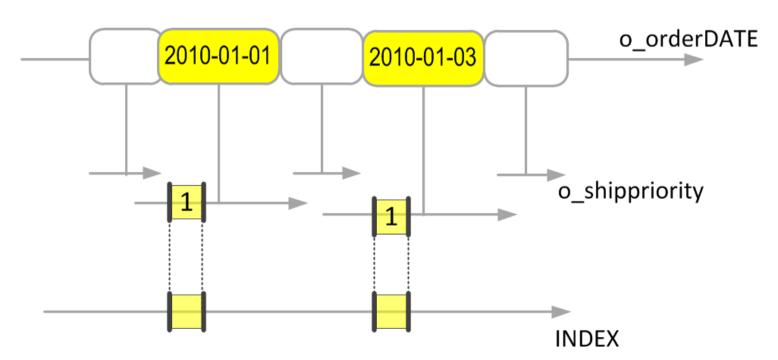
- Much more complex
- Ranges over (keypart1, keypart2, ...) tuple space
  - Think of it as axis with sub-axis at each point:



- Range building rules:
  - Number of ranges is a function of WHERE, not of data in the table
  - Do not use knowledge about domain density (e.g. no knowledge like "there is no integers between 1 and 2")

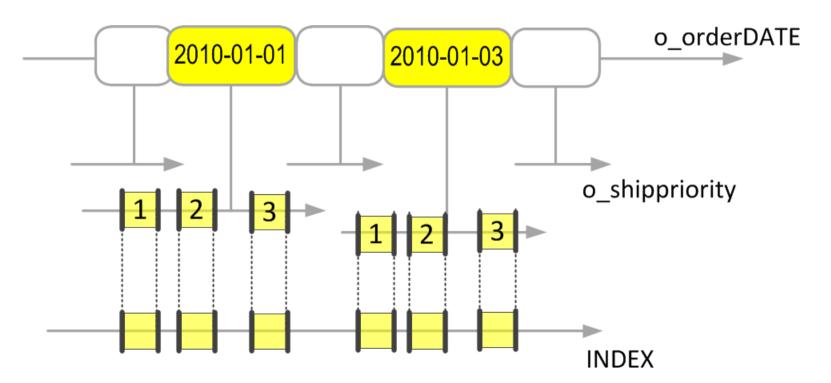
 Equality on 1<sup>st</sup> key part? Can put condition on 2<sup>nd</sup> key part into sub-axis

```
alter table orders add key i_o_date_clerk (o_orderdate, o_clerk); select * from orders where o_orderDATE IN ('2010-01-01', '2010-01-03') AND o_shippriority=1
```



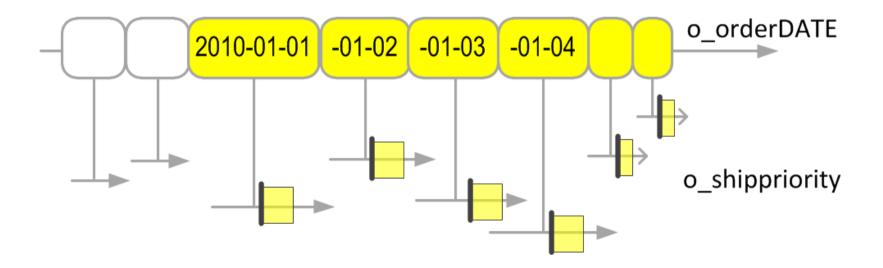
This can cause big number of ranges. That's ok

```
select * from orders
where
o_orderDATE IN ('2010-01-01', '2010-01-03') AND
(o_shippriority = 1 OR o_shippriority=2 OR o_shippriority=3)
```

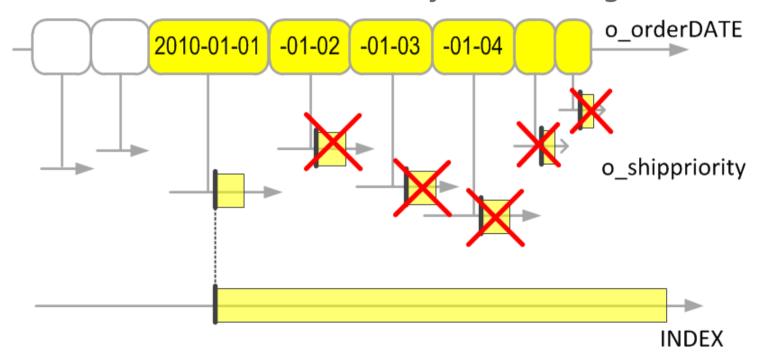


Non-equality on 1<sup>st</sup> keypart:

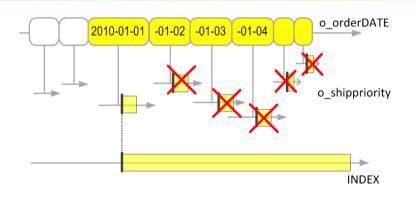
select \* from orders where o\_orderDATE >= '2010-01-01' AND o\_shippriority>3

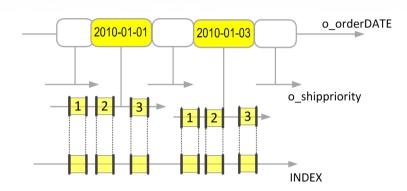


- Non-equality on 1<sup>st</sup> keypart:
- select \* from orders where o\_orderDATE >= '2010-01-01' AND o\_shippriority>3
- Can attach the second only to the edge "cell":



#### Multi-key-part range summary





- keyparts for which queries have equalities should go first
- multiple keyparts with non-equality ranges will not reduce the number of rows read
  - e.g. t.keypart1 BETWEEN ... AND t.keypart2 BETWEEN ... won't take advantage of condition on keypart2
- When keypartX is followed by other keyparts, then
   keypartX IN (1,... N) will work better than keypartX BETWEEN 1 AND N

#### Can I look at the ranges it has constructed?

- Normally, no. But there is a last-resort method:
  - Get a really small subset of your data onto separate server
  - Get a debug build there. Start it with -debug
  - Run your query with FORCE INDEX:

```
explain select * from orders force index(i_o_datepriokey)
where (o_orderDATE IN ('2010-02-01','2010-02-03') AND o_shippriority < 3) OR
(o_orderDATE >= '2010-03-10' AND o_shippriority < 4)
```

- Go into mysqld.trace
  - Grep for "query:" until you find the query
  - Grep for ">print\_quick"

#### Ranges in mysqld.trace

#### mysqld.trace

```
: | | query: explain select * from orders force index(i o datepriokey) where o orderDATE
IN ('2010-02-01','2010-02-03') AND o shippriority < 3 OR (o orderDATE > '2010-03-10' AND o
shippriority < 4)
T@10
      : | | >PROFILING::set query source
       : | | <PROFILING::set query source
T@10
       : | | >mysql parse
T@10
       : | | | | | | | >print quick
T@10
quick range select, key i o datepriokey, length: 9
 2010-02-01/\text{NULL} < X < 2010-02-01/3
 2010-02-03/NULL < X < 2010-02-03/3
 2010-03-10 < X
other keys: 0x0:
```

EXPLAIN only shows key\_len which is max # keyparts used

+	+	table	type	possible_keys	key	key_len	+   ref	rows	+   E
1	SIMPLE	orders	range	i_o_datepriokey	i_o_datepriokey	9	NULL	] 3	U

#### Range access estimates

+	id	select_type	table	type	possible_keys	   key	+   key_len	+   ref	rows	+   E
	1	SIMPLE	orders	range	i_o_datepriokey	i_o_datepriokey	9	NULL	3	ָּע ע

- Are obtained by doing dives into the index
  - Dive at the start of the range
  - Dive at the end
  - Estimate the difference
- Usually are precise (not more than 2x wrong)
- Are not affected/do not need ANALYZE TABLE
- Cost you disk IO!
  - Indexes that are never worth using will still be probed!
  - Check your buffer pool + EXPLAINs to see if this happens!
    - Quick fix: IGNORE INDEX (unusable\_index)



#### Range access: Status variables

Variable_name	Value
Handler commit	+   1
Handler delete	1 0
Handler_delete   Handler discover	1 0
Handler_icp_attempts	1 9
Handler icp match	19
Handler mrr init	1 0
Handler mrr key refills	1 0
Handler mrr rowid refills	1 0
Handler prepare	1 0
Handler_prepare   Handler read first	1 0
Handler read key	0   9
Handler_read_key   Handler read next	1 9 1 9
Handler read prev	l 0
Handler_read_prev	1 0
Handler_read_rnd deleted	1 0
Handler_read_rnd_dereced	1 0
Handler_read_rnd_next   Handler rollback	1 0
·	1 0
Handler_savepoint	1 0
Handler_savepoint_rollback	
Handler_tmp_update	0
Handler_tmp_write	0
Handler_update	0
Handler_write	] 0
23 (0.05)	+

#### Increments

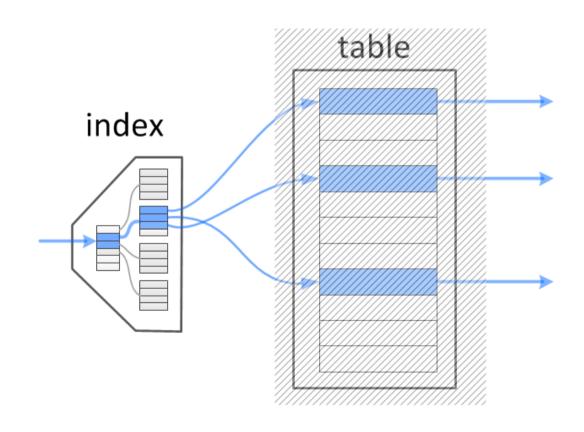
- index lookup counter
- walk-forward-in-index counter



## Index-only scans

#### **Index-only scans**

- Besides lookups, indexes have another use
- If the index has all columns you need, it is possible to avoid accessing table rows altogether
- This is called "index only" scan
- Can be applied to ref, and range accesses.



## Index-only read example

```
explain select sum(o totalprice)
from orders
where o orderdate between '1995-01-01' and '1995-01-07';
   ------
|id|select_type|table |type |possible_keys|key |key_len|ref |rows|Extra |
| 1|SIMPLE | orders|range|i o orderdate|i o orderdate|4 | NULL|4358|Using index condition|
alter table orders add key o date price (o orderDATE, o totalprice);
# re-run the EXPLAIN:
alter table orders add key o date price (o orderDATE, o totalprice);
explain select sum(o totalprice) from orders where o orderdate between '1995-01-01' and '1995-02-01';
    ------
|id|select_type|table |type |possible_keys|key |key_len|ref |rows |Extra |
 | 1|SIMPLE |orders|range|... |o_date_price|4 |NULL|39424|Using where Using index|
```

scale=1, in-mem: from 0.2 to 0.1 sec

#### Index-only reads and status variables

#### Surprisingly, no difference

```
select sum(o totalprice)
from orders use index(o date price)
where o orderdate between '1995-01-01' and
                        '1995-02-01';
Handler commit
 Handler delete
 Handler discover
 Handler icp attempts
 Handler icp match
 Handler mrr init
 Handler mrr key refills
 Handler mrr rowid refills | 0
 Handler prepare
 Handler read first
 Handler read key
 Handler read next
                           I 20137
 Handler read prev
 Handler read rnd
 Handler read rnd deleted
 Handler read rnd next
 Handler rollback
 Handler savepoint
 Handler savepoint rollback | 0
 Handler tmp update
 Handler tmp write
 Handler update
 Handler write
```

```
select sum(o totalprice)
from orders use index(i o orderdate)
where o orderdate between '1995-01-01' and
                        '1995-02-01';
 Variable name
 Handler commit
 Handler delete
 Handler discover
 Handler icp attempts
                           | 20137
 Handler icp match
                           | 20137
| Handler mrr init
 Handler mrr key refills
 Handler mrr rowid refills | 0
 Handler prepare
 Handler read first
 Handler read key
 Handler read next
                           I 20137
| Handler read prev
                           1 0
 Handler read rnd
 Handler read rnd deleted
 Handler read rnd next
 Handler rollback
 Handler savepoint
 Handler savepoint rollback | 0
 Handler tmp update
 Handler tmp write
 Handler update
 Handler write
```

#### Index-only reads and status variables

Let's look at InnoDB "lower-level" counters

show status like 'Innodb%';				
+	İ	Value	    -+	
<pre>  Innodb_buffer_pool_read_requests   Innodb_buffer_pool_reads   Innodb_buffer_pool_wait_free</pre>	     	422426 10034 0	   	+2566
Innodb_rows_inserted   Innodb_rows_read   Innodb_rows_updated		0 241651 0		+20138

- These are global non FLUSHable counters, so we manually calculate increment and show it on the right
- \_rows\_read are the same, too
- \_buffer\_pool\_reads are different
- A bit sense-less counting
  - Have plans to fix in MariaDB.

## Index-only reads summary

- Can be used with any index-based access methods
- Allows to skip accessing table records
  - Some savings for CPU-bound loads
  - Huge savings for IO-bound loads
  - Extra columns in the index a cost, though
    - Index is bigger
    - [Potentially] more frequently updated
- EXPLAIN shows "Using index"
- Counting in for Handler\_xxx and Innodb\_row\_reads status variables
  - Watch lower-level counters.

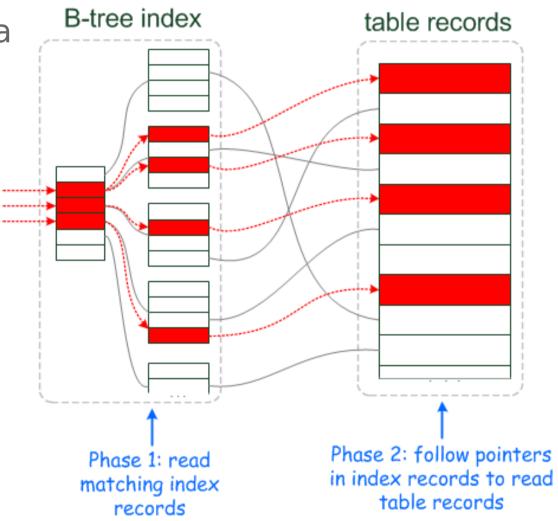
#### Index Condition Pushdown

MySQL 5.6+, MariaDB 5.3+

## Index Condition Pushdown idea

A non-index only read is a two-step process:

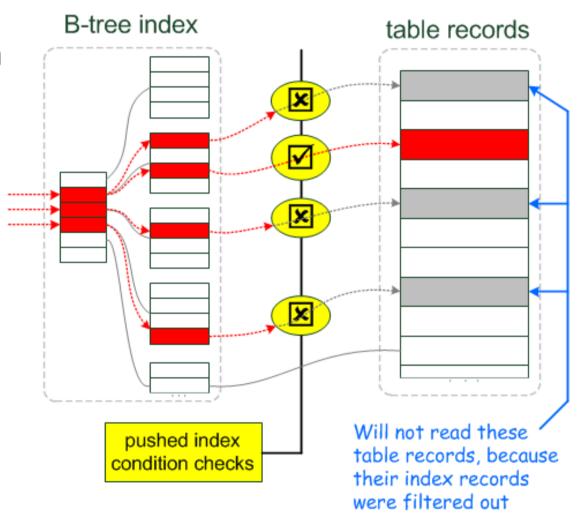
- 1. Read index
- 2. Read record
- 3. Check the WHERE condition.



## Index Condition Pushdown idea

#### Index Condition Pushdown

- 1. Read index
- 2. Check condition on index columns
- 3. Read record
- 4. Check the WHERE condition



#### Index condition pushdown example

#### Start without ICP

```
alter table orders add key i o clerk(o clerk);
   select o clerk, o orderkey, o shippriority
   from orders
   where o clerk IN ('Clerk#00000001', 'Clerk#00000201', 'Clerk#000003001') and
      dayofweek(o orderDATE)=1:
# MariaDB 5.2.x/ MySQL 5.5.x
 |id|select type|table |type |possible keys|key |key len|ref |rows|Extra
+--+------
# MariaDB 5.3+, MySQL 5.6+
SIMPLE |orders|range|i_o_clerk |i_o_clerk|16 |NULL|2978|<mark>Using index condition;</mark> Using where|
```

#### **Counters**

alter table orders add key i o clerk(o clerk);

```
select o clerk, o orderkey, o shippriority
   from orders
   where o clerk IN ('Clerk#00000001', 'Clerk#00000201', 'Clerk#000003001') and
         dayofweek(o orderDATE)=1;
# MariaDB 5.3+, MySQL 5.6+
  |id|select_type|table |type |possible_keys|key |key_len|ref |rows|Extra
          |orders|range|i o clerk |i o clerk|16 |NULL|2978|Using index condition; Using where|
                                    MariaDB [(none)]> show status like 'Innodb%';
  Handler commit
                                     Handler delete
 Handler discover
                                      Innodb buffer pool read requests | 32103 | + 24416
  Handler icp attempts
                       1 2979
                                      Innodb rows read | 3406 | + 2979
  Handler icp match
                       1 2979
  Handler mrr init
  Handler mrr key refills
  Handler mrr rowid refills
  Handler prepare
  Handler read first
  Handler read key
  Handler read next
  Handler read prev
  Handler read rnd
```

### Index condition pushdown example

Now, add an index that ICP could use:

```
alter table orders add key i_o_clerk_date(o_clerk, o_orderDATE); select o_clerk, o_orderkey, o_shippriority from orders where o_clerk IN ('Clerk#000000001', 'Clerk#000000201', 'Clerk#0000003001') and dayofweek(o_orderDATE)=1;
```

#### **Counters**

```
alter table orders add key i_o_clerk_date(o_clerk, o_orderDATE); select o_clerk, o_orderkey, o_shippriority from orders where o_clerk IN ('Clerk#000000001', 'Clerk#000000201', 'Clerk#0000003001') and dayofweek(o_orderDATE)=1;
```

```
|orders|range|i_o_clerk_date,i_o_clerk|i_o_clerk_date|
                                                      |NULL|2978|Using index condition|
MariaDB [(none)]> show status like 'Innodb%';
Handler commit
                                         Handler delete
Handler discover
                                         | Innodb buffer pool read requests | 7687
                                                                           1 +2585
Handler icp attempts
                                          Innodb rows read
                      1 2979
                                                                     1 427
                                                                           1 + 427
Handler icp match
                      1 427
                            | = 1/7 th
Handler mrr init
Handler mrr key refills
Handler mrr rowid refills | 0
Handler prepare
Handler read first
Handler read key
Handler read next
                      1 427
                             = 1/7th
Handler read prev
Handler read rnd
```

# Index Condition Pushdown optimization ... or lack thereof

```
alter table orders add key i_o_clerk_date(o_clerk, o_orderDATE); select o_clerk, o_orderkey, o_shippriority from orders where o_clerk IN ('Clerk#000000001', 'Clerk#000000201', 'Clerk#0000003001') and dayofweek(o orderDATE)=1;
```

#### INDEX (o\_clerk)

INDEX (o\_clerk, o\_orderDATE);

- Range selectivity on 1<sup>st</sup> keypart is (or, should be) the same for both
- Index entries are smaller
- ICP is applicable
  - But is condition dayofweek(...)=1 selective?

#### Current optimizer actions:

- #1. Choose an index to use
- #2. If possible, use IndexConditionPushdown with it
  - => Possibility of ICP doesn't affect index choice.



# Index Condition Pushdown optimization ... or lack thereof

#### What to do, then?

- USE/IGNORE INDEX hints
- Create only indexes such that ICP will be used
  - Include approriate extra columns at the end
  - In our example:
    - Create INDEX (o clerk, o orderDATE)
    - Don't create INDEX (o\_clerk)

## Index Merge

### **Index Merge Union**

- ORed conditions on different columns cannot be resolved with any single index
- Example: airline on-time performance data, 15M flights:

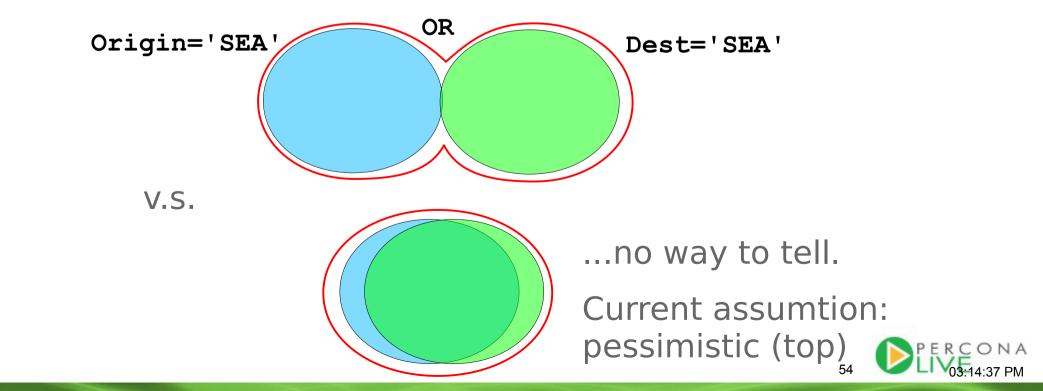
select \* from ontime where (Origin='SEA' or Dest='SEA'); INDEX(Origin) Ontime table Union op Origin='SEA' INDEX(Dest) Dest='SEA'



## Index Merge union

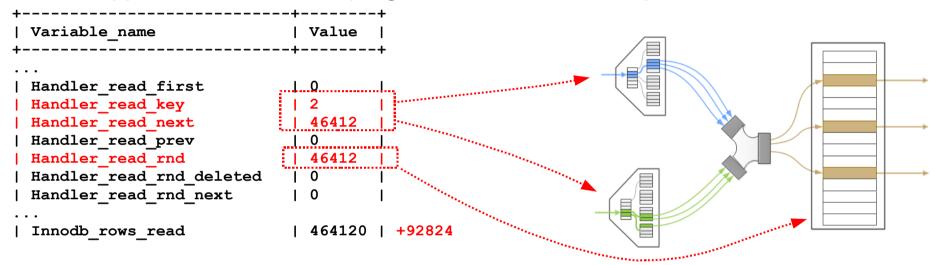
select \* from ontime where (Origin='SEA' or Dest='SEA');

Optimization is challenging:

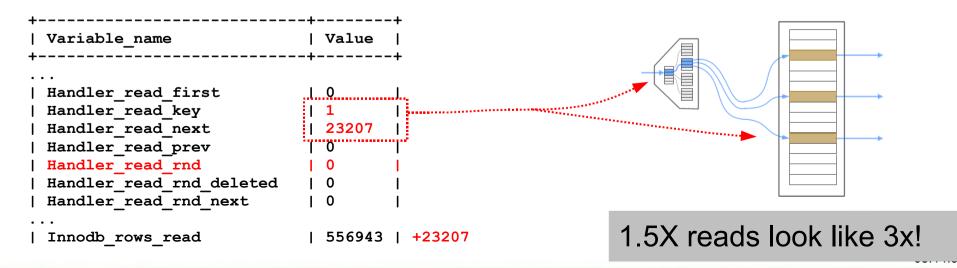


#### Index merge: statistics counters

select count(\*) from ontime where (Origin='SEA' or Dest='SEA'); -- 46412 rows



select count(\*) from ontime where Origin='SEA' -- 23207 rows



### Index merge properties

[MySQL, Percona, MariaDB 5.2x] index\_merge plans are removed from consideration if a range access is possible:

#### Solutions:

- Upgrade to MariaDB 5.3+ :-)
- IGNORE INDEX(\$index\_used\_for\_range)
  - there is no "FORCE INDEX MERGE" hint



## Index Merge/Union summary

- Used for ORs spanning multiple indexes
  - Each part of OR must have an index it could use
- Counter increments not quite comparable with singleindex access methods
- Can be turned off globally with set optimizer\_switch='index\_merge=off'
- Can be turned off locally with IGNORE INDEX(some merged index)
- Can be blocked by potential ref/range accesses
  - Fix1: Upgrade to MariaDB 5.3
  - Fix2: IGNORE INDEX(other\_range\_indexes)
     USE INDEX (index\_merge\_indexes)



### Index Merge/Intersection

- Like Index Merge/union, but for ANDs
- Aimed at cases when there is no composite index

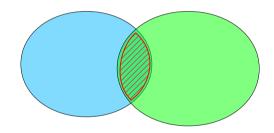
select avg(arrdelay) from ontime where <a href="depdel15=1">depdel15=1</a> and <a href="OriginState">OriginState</a> = 'CA';

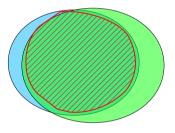
- MySQL, MariaDB 5.2: only support equality conditions
- MariaDB 5.3: any range condition supported
  - Must be manually enabled: set optimizer\_switch='index\_merge\_sort\_intersection=on'
  - EXPLAIN: Using sort-intersect

#### Index Merge/intersect optimization

Optimization challenges are similar to UNION

select avg(arrdelay) from ontime where depdel15=1 and OriginState ='CA';





- There is no way to force index\_merge/intersect
- To disable
  - set optimizer\_switch='index\_merge\_intersection=off
  - IGNORE INDEX(one\_of\_used\_indexes)
- Tip: showeling through lots of data?
   set optimizer\_switch='index\_merge\_sort\_intersection=on's set sort\_buffer\_size=...

# JOINs

## Consider a join

```
select count(*)
 from
  customer, orders
 where
  c custkey=o custkey and c acctbal < -500 and o orderpriority='1-URGENT';
   |customer[range|... |c acctbal |9 | NULL
| 1|SIMPLE
                                       | 6802|Using where; Using in
for each record R1 in customer
       for each matching record R2 in orders
                              EXPLAIN shows the loops,
         pass (R1, R2) into join output
                               from outer to inner
```

```
select count(*)
 from
   customer, orders
 where
   c custkey=o custkey and c acctbal < -500 and o orderpriority='1-URGENT';
|id|select type|table |type |possible keys|key |key len|ref
   1 | SIMPLE
          |customer|range|...
                           |c acctbal |9
                                               INULL
                                                             |6802|Using where; Using in
          |orders |ref |i o custkey |i o custkey|5
| 1|SIMPLE
                                              |customer.c custkey| 7|Using where
       for each record R1 in customer
         if (where [customer] is satisfied)
           for each matching record R2 in orders
                                                 "Using where" means
            if (where[orders] is satisfied)
                                                 checking a part of WHERE
              pass (R1, R2) into join output
                                                 Which part?
```

```
select count(*)
 from
   customer, orders
 where
   c custkey=o custkey and c acctbal < -500 and o orderpriority='1-URGENT';
|id|select_type|table |type|possible_keys|key |key|len|ref
   -----
1|SIMPLE
          |customer|range|...
                           |c acctbal |9
                                               INULL
                                                              |6802|Using where; Using in
          orders | ref | i o custkey | i o custkey | 5
1|SIMPLE
                                               |customer.c custkey| 7|Using where
       for each record R1 in customer
         if (where [customer] is satisfied)
           for each matching record R2 in orders
                                                   The part that refers to the
                                                   tables for which we know
            if (where[orders] is satisfied)
              pass (R1, R2) into join output
                                                   R{n}, the "current record".
```

```
select count(*)
 from
   customer, orders
 where
   c custkey=o custkey and c acctbal < -500 and o orderpriority='1-URGENT';
|id|select type|table |type |possible keys|key |key len|ref
   -----
1|SIMPLE
          |customer|range|... |c acctbal |9
                                                INULL
                                                               |6802|Using where; Using in
| 1|SIMPLE | orders | ref | i o custkey | i o custkey | 5 | customer.c custkey | 7 | Using where
  for each record R1 in customer // 6802 loops 4
   if (where[customer] is satisfied)
     for each matching record R2 in orders
                                         // 7 loops
       if (where[orders] is satisfied)
                                         "rows" shows how many rows are
         pass (R1, R2) into join output
                                          read from each table
                                         For `orders`, 7 rows will be read
                                          6802 times
```

```
select count(*)
  from
    customer, orders
  where
    c custkey=o custkey and c acctbal < -500 and o orderpriority='1-URGENT';
      . - - - - - + - - - - - + - - - - + - - - - - - + - - - - - + - - - - - + - - - - - - - - - - - - -
|id|select type|table |type|possible keys|key |key len|ref
                                                   NULL
                                                                   |6802|Using where; Using in
 1|SIMPLE
        |customer|range|...
                             |c acctbal |9
          |orders |ref |i o custkey |i o custkey|5
                                                 |customer.c custkey| 7 |Using where
1|SIMPLE
  for each record R1 in customer // 6802 loops
    if (where[customer] is satisfied)
      Wait, what about
        if (where[orders] is satisfied)
         pass (R1, R2) into join output
                                                            this 'if'?
                                                           It may reject some of
                                                            the 6802 values of R1
```

```
explain extended
  select * from customer, orders
  where
    c custkey=o custkey and c acctbal < 1000 and o orderpriority='1-URGENT';
                 |type|possible keys|key |key len|ref
                                                                 | 132987 | 39.44 | Using where |
          |orders |ref |i o custkey |i o custkey|5
                                                  |customer.c custkey|
                                                                       7 100.00 Using where
1|SIMPLE
 for each record R1 in customer // 6802 loops
   if (where[customer] is satisfied) 
     for each matching record R2 in orders // 7 loops
       if (where[orders] is satisfied) ←
                                                 "filtered" column shows the
        pass (R1, R2) into join output
                                                  percentage of records that are
                                                 expected to pass the "if".
                                                 100% is the most frequent
                                                  (optimizer puts 100% when there
                                                 are no reliable estimates for filtering)
```

#### Ref access

- `orders` is accessed with ref access: index lookups on index i\_o\_custkey using value of customer.c\_custkey
- This is similar to key=const ref access we've discussed earlier
  - But there is no "const", the value of customer.c\_custkey is different for each lookup.
- Each lookup is expected to produce 7 matching rows.
  - This information is from Index Statistics. We'll cover it later.

How good a join can/should be

#### Let's collect some data

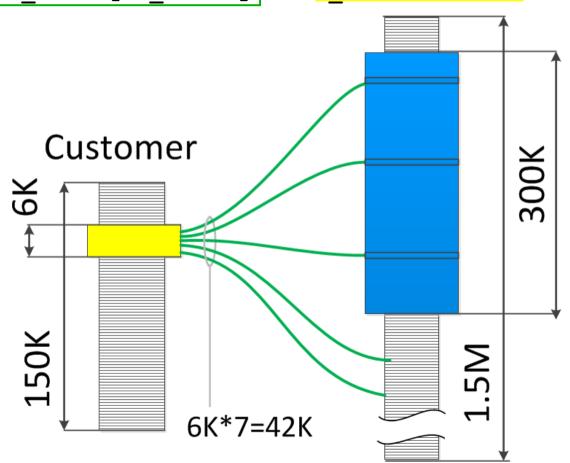
```
select count(*)
  from
   customer, orders
  where
   c custkey=o custkey and c acctbal < -500 and o orderpriority='1-URGENT';
|id|select_type|table |type|possible_keys|key |key_len|ref |rows|E
| 1|SIMPLE | customer|range|... | c_acctbal | 9 | NULL | 6802|Using where; Using in
| 1|SIMPLE | orders | ref | i_o_custkey | i_o_custkey | 5 | customer.c_custkey | 7 | Using where
 [explain] select count(*) from customer; -- 150K
 [explain] select count(*) from customer where c acctbal < -500 -- 6802
 [explain] select count(*) from orders -- 1.5M
 [explain] select count(*) from orders where o orderpriority='1-URGENT' - 300K
 One customer (c custkey) - 7 orders
```

And use it to analyze the join:

### **Analyzing the join**

select count(\*) from customer, orders
where

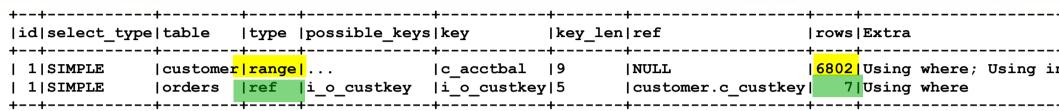
c custkey=o custkey and c acctbal < 1000 and o orderpriority='1-URGENT'

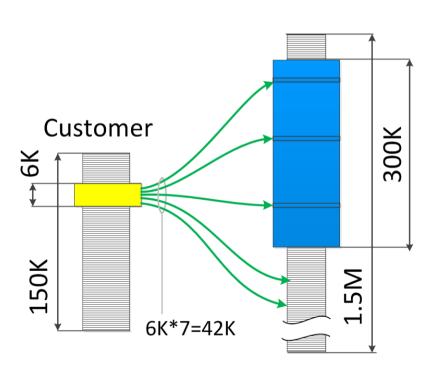


There are three ways to compute the join:

- Yellow → Green
- Blue → Green
- · Yellow, Blue

## Join plan analysis





- We're using the Yellow->Green read
- The yellow part is read efficiently
  - Range access is used, it scans
     6K rows
- The green part read efficiently
  - ref access, 6K lookups, each producing 7 rows on average
- This is probably better than the two other variants

### Tracking query plan execution

- EXPLAINs are what the optimizer \*expects \* to happen
- What actually happens?
- MySQL has no EXPLAIN ANALYZE
- Possible solutions
  - Status variables
  - "userstat"
  - performance\_schema?

#### Check status variables

```
MariaDB [dbt3sf1]> show status like 'Handler%';
  Variable name
  Handler commit
  Handler delete
 Handler discover
 Handler icp attempts
 Handler icp match
 Handler mrr init
  Handler mrr key refills
  Handler mrr rowid refills
  Handler prepare
  Handler read first
  Handler read key
                              1 6805
  Handler read next
                              I 75951
  Handler read prev
                               0
  Handler read rnd
  Handler read rnd deleted
  Handler read rnd next
                              1 0
  Handler rollback
 Handler savepoint
```

- Accesses to all tables are summed together!
- It is still possible to do analysis:
  - Run a "sub-join" with the 1<sup>st</sup> table;
     note the counters
  - Run a "sub-join" with the 1<sup>st</sup> and 2<sup>nd</sup> tables from the join order; note the counters; substract counters from #1
  - . . . .
- This is slow and painful.

#### More powerful: userstatv2

Percona Server and MariaDB:

```
MariaDB [dbt3sf1]> show table statistics;
 Table_schema | Table name | Rows read | Rows_changed | Rows changed x #indexes
             | customer | 6805
 dbt3sf1
              orders
                             69147 🔏
 dbt3sf1
2 rows in set (0.00 sec)
MariaDB [dbt3sf1] > show index statistics;
 Table schema | Table name | Index name | Rows read
 dbt3sf1 | customer
                          l c acctbal
                                            6805
 dbt3sf1
                          i o custkey |
              | orders
2 rows in set (0.00 sec)
```

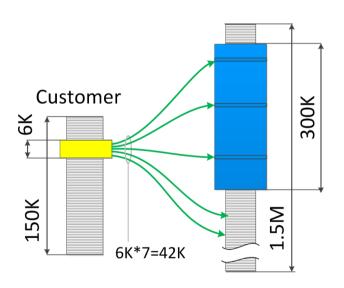
Much better but self-joins are still painful

#### Recap

- Ok we now know
  - Single table access methods
    - ref(const)
    - Range
    - index\_merge
  - Nested loop join algorithm
  - EXPLAIN
  - How to check if execution follows EXPLAIN

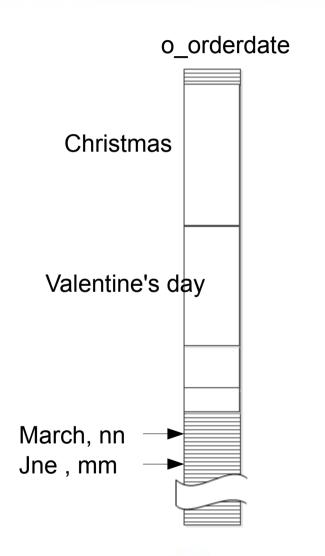
#### Typical reasons of poor join performance

- Poor access strategy for the 1<sup>st</sup> table
  - Highly-selective non-indexed conditions? Create an index!
- Joins on non-equality conditions
  - e.g. "same day" joins on DATETIME columns
  - Tip: Convert to equality join at all costs
- No index to do ref access for join
  - Cross-product joins inherently slow
- Correlated data (...)
- Wrong index statistics (...)



#### Non-uniform distributions

- Some data is highly nonuniform
- Performance depends on hitting the popular values
- No one-for-all solution



# Index and Table Statistics

# Database statistics and plan stability (InnoDB/XtraDB)

- Sample [8] random index leaf pages
- Table statistics (stored)
  - rows estimated number of rows in a table
  - clust\_size cluster index (table/primary key) size in number of pages
  - other\_size other index (non primary key) size in number of pages
- Index statistics (stored)
  - fields #fields in the index
  - rows\_per\_key rows per 1 key value, per prefix fields ([1 column value], [2 columns value], [3 columns value], ...)
  - index\_total\_pages number of index pages
  - index\_leaf\_pages number of leaf pages



### Table statistics updates

- Statistics updated when:
  - ANALYZE TABLE tbl name [, tbl name] ...
  - SHOW TABLE STATUS, SHOW INDEX
  - Access to INFORMATION\_SCHEMA.[TABLES| STATISTICS]
  - A table is opened for the first time (that is, after server restart)
  - A table has changed a lot (1/16th of the table updated/deleted/inserted)
  - When InnoDB Monitor is turned ON
  - Others (?)

### Displaying statistics

- MySQL 5.5, MariaDB 5.3, and older (GA versions)
  - Issue SQL statements to count rows/keys
  - Indirectly, look at EXPLAIN for simple queries
- MariaDB 5.5, Percona Server 5.5 (using XtraDB) (GA versions)
  - information\_schema.[innodb\_index\_stats, innodb\_table\_stats]
  - Read-only, always visible
- MySQL 5.6 (development)
  - mysql.[innodb\_index\_stats, innodb\_table\_stats]
  - User updatetable
  - Only available if innodb\_analyze\_is\_persistent=ON
- MariaDB \$NEXT (development)
  - MyISAM persistent updateable tables
  - + current MySQL and Percona mechanisms



### Plan [in]stability

Statistics may vary a lot (orders)

```
MariaDB [dbt3]> select * from information_schema.innodb_index_stats;
  | PRIMARY | 3, 1 | 4, 1 | 25% | 1_ps_partkey | 3, 0 | => | 4, 1 | | 25% | (4) | 1_ps_suppkey | 64, 0 | 91, 1 | 30% | (80) | 1_o_orderdate | 9597, 1 | 1660956, 0 | 99% | (6234) | (15)
  partsupp | PRIMARY | 3, 1
  partsupp
 partsupp | i_ps_suppkey | 64, 0 orders | i_o_orderdate | 9597, 1
                                                | 15, 0 | 0% (15)
  orders | i_o_custkey | 15, 1
  lineitem | i l receiptdate | 7425, 1, 1 | | 6665850, 1, 1 | 99.9% (23477)
MariaDB [dbt3]> select * from information_schema.innodb_table_stats;
  table name | rows
                                    lrows
 partsupp | 6524766 | 9101065 | 28% (8000000) orders | 15039855 | ==> | 14948612 | 0.6% (15000000)
  lineitem | 60062904 | | 59992655 | 0.1% (59986052)
```

### Controlling statistics (GA versions)

- MySQL 5.5, MariaDB 5.3, and older
  - Manual tuning: optimizer hints, system variables
- MySQL 5.5, MariaDB 5.5 with InnoDB plugin
  - innodb\_stats\_on\_metadata = OFF update only on restart, ANALYZE
  - innodb\_stats\_sample\_pages = 8 is default increase precision
  - No way to "freeze" db statistics in all cases
- MariaDB 5.5, Percona Server 5.5 (using XtraDB)
  - Can "freeze" the current InnoDB statistics
  - innodb\_use\_sys\_stats\_table=ON use I\_S.INNODB\_SYS\_STATS
  - innodb\_stats\_auto\_update=OFF recalculate except for "first open" and "ANALYZE TABLE"
  - No manual control over statistics, only refresh by random sampling

### Plan [in]stability

Same query, different statistics

=>

- Different access methods and/or
- Different JOIN orders
- => different query plans
- Query performance may change a lot when statistics changes
- BEWARE WHEN BENCHMARKING THE OPTIMIZER

# Controlling statistics (MySQL dev)

MySQL 5.6 (development version, public code)

- Persistent and user-updatetable InnoDB statistics
  - innodb analyze is persistent = ON,
  - updated only on ANALYZE TABLE
- Control the precision of sampling [default 8]
  - innodb\_stats\_persistent\_sample\_pages,
  - innodb\_stats\_persistent\_sample\_pages
- No new statistics compared to older versions

# Controlling statistics (MariaDB dev)

MariaDB \$NEXT (development version, public code)

- Current MySQL and Percona InnoDB statistics
- Engine-independent, persistent, user-updateable statistics
- Precise
- Additional statistics per column (even when there is no index):
  - min\_value, max\_value: minimum/maximum value per column
  - nulls\_ratio: fraction of null values in a column
  - avg\_length: average size of values in a column
  - avg\_frequency: average number of rows with the same value
- => better query plans

Code: bzr branch lp:~maria-captains/maria/maria-5.5-mwl248

### Resources on InnoDB statistics

#### From the InnoDB/XtraDB developers:

https://mysqlperformanceblog.com/doc/percona-server/5.5/diagnostics/innodb stats.html

http://dev.mysql.com/doc/refman/5.6/en/innodb-other-changes-statistics-estimation.html

http://dev.mysql.com/doc/refman/5.6/en/innodb-parameters.html

http://dev.mysql.com/doc/refman/5.6/en/innodb-performance.html#innodb-persistent-stats

http://blogs.innodb.com/wp/2011/04/innodb-persistent-statistics-at-last/

http://dev.mysql.com/doc/refman/5.6/en/analyze-table.html

# Optimizer differences between MySQL branches

### Optimizer in different MySQL branches

- Base version: MySQL 5.5
  - Optimizer is not different from MySQL 5.1
- Percona Server X.Y
  - Closely follows MySQL X.Y
  - Some improvements in diagnostics (eg. userstats)
- MariaDB 5.3
  - Lots of new features
- MariaDB 5.5
  - For the most part, is a merge of Maria-5.3 with MySQL-5.5
  - But has a couple of small improvements
- MySQL 5.6
  - Lots of new features, intersection with MariaDB 5.3



# Optimizer feature comparison (1)

Feature	MariaDB 5.3/5.5	MySQL 5.5	MySQL (5.6 dev)
Disk access optimizations			
Index Condition Pushdown (ICP)	YES	-	YES
Disk-sweep Multi-range read (DS-MRR)	YES	-	YES
DS-MRR with Key-ordered retrieval	YES	-	-
Index_merge / Sort_intersection	YES	-	-
Cost-based choice of range vs. index_merge	YES	-	-
ORDER BY LIMIT <small_limit></small_limit>	-	-	YES
Use extended (hidden) primary keys for innodb/xtradb	YES (5.5)	-	-
Join optimizations			
Batched key access (BKA)	YES	-	YES
Block hash join	YES	-	-
User-set memory limits on all join buffers	YES	-	-
Apply early outer table ON conditions	YES	-	-
Null-rejecting conditions tested early for NULLs	YES		

# Optimizer feature comparison (2)

Feature	MariaDB 5.5	MySQL 5.5	MySQL (5.6 dev)
Subquery optimizations			
In-to-exists	YES	YES	YES
Semi-join	YES	-	-YES
Materialization	YES	-	YES
NULL-aware Materialization	YES	-	-
Cost choice of materialization vs in-to-exists	YES	-	-
Subquery cache	YES	-	-
Fast explain with subqueries	YES		-
Optimizations for derived tables / views			
Delayed materialization of derived tables / materialized views	YES	-	YES
Instant EXPLAIN for derived tables	YES	-	YES
Derived Table with Keys optimization	YES	-	YES
Fields of merge-able views and derived tables used in equality optimizations	YES	- 92	- - - - - - - - - - - - - - - - - - -

# **Optimizer feature comparison (3)**

Feature	MariaDB 5.3/5.5	MySQL 5.5	MySQL (5.6 dev)
Execution control			
LIMIT ROWS EXAMINED rows_limit	YES (5.5)	-	-
Optimizer control (optimizer switch)			
Systematic control of all optimizer strategies	YES	-	partial
EXPLAIN improvements			
Explain for DELETE, INSERT, REPLACE, and UPDATE	-	-	YES
EXPLAIN in JSON format	-	-	YES
More detailed and consistent EXPLAIN for subqueries	YES	-	-

# Optimizer switch (2)

#### MariaDB 5.5

#### **Index merge:**

index\_merge=on,
index\_merge\_union=on,
index\_merge\_sort\_union=on,
index\_merge\_intersection=on,
index\_merge\_sort\_intersection=off

#### **Condition pushdown:**

engine\_condition\_pushdown=off,
index condition pushdown=on,

mrr=off, mrr\_cost\_based=off, mrr\_sort\_keys=off

#### MySQL 5.6

#### **Index merge:**

index\_merge=on,
index\_merge\_union=on,
index\_merge\_sort\_union=on,
index\_merge\_intersection=on

#### **Condition pushdown:**

engine\_condition\_pushdown=on,
index condition pushdown=on,

mrr=on, mrr\_cost\_based=on

# Optimizer switch (2)

#### MariaDB 5.5

#### **Subqueries:**

materialization=on,
semijoin=on,
loosescan=on,
firstmatch=on,
in\_to\_exists=on,
partial\_match\_rowid\_merge=on,
partial\_match\_table\_scan=on,
subquery\_cache=on,

Derived tables: derived\_merge=on, derived\_with\_keys=on

#### MySQL 5.6

#### **Subqueries:**

materialization=on, semijoin=on, loosescan=on, firstmatch=on

# Optimizer switch (3)

#### MariaDB 5.5

joins
join\_cache\_bka=on,
join\_cache\_incremental=on,
join\_cache\_hashed=on,
outer\_join\_with\_cache=on,
semijoin\_with\_cache=on,
optimize\_join\_buffer\_size=off
+ other system variables

Other features table\_elimination=on, extended\_keys=off

MySQL 5.6

#### Joins:

batched\_key\_access=off
block\_nested\_loop=on,

Thank You!

Q & A