Insertion Sort in C



What is Insertion Sort?

- A simple sorting algorithm
- Works like sorting playing cards in your hand
- Builds the final sorted array one item at a time



Algorithm Idea

- 1. Start from the **second element**
- 2. Compare with elements before it
- 3. Shift larger elements to the right
- 4. Insert the element into the correct position
- 5. Repeat until array is sorted



C Implementation

```
#include <stdio.h>
void insertionSort(int arr[], int n) {
    for (int i = 1; i < n; i++) {
        int key = arr[i];
        int j = i - 1;
        while (j >= 0 && arr[j] > key) {
            arr[j + 1] = arr[j];
            j--;
        arr[j + 1] = key;
```



What is Binary Search?

- Efficient algorithm for finding an element in a sorted array
- Works by repeatedly dividing the search interval in half
- Compare target with the middle element:
 - o If equal → found
 - If smaller → search left half
 - If larger → search right half



Algorithm Idea

- 1. Start with low = 0, high = n-1
- 2. Compute mid = (low + high) / 2
- 3. Compare arr[mid] with key
 - o If arr[mid] == key → return mid
 - o If arr[mid] > key → search left half
 - Else → search right half
- 4. Repeat until low > high



Recursive Implementation (C)

```
int binarySearchRec(int arr[], int low, int high, int key) {
    if (low > high) return −1;
    int mid = (low + high) / 2;
    if (arr[mid] == key)
        return mid;
    else if (arr[mid] > key)
        return binarySearchRec(arr, low, mid - 1, key);
    else
        return binarySearchRec(arr, mid + 1, high, key);
```



Example Dry Run

Array: {1, 3, 5, 7, 9, 11, 13}

Search for 7

low=0, high=6 \rightarrow mid=3 \rightarrow arr[3]=7 \rightarrow found at index 3

Search for 6

low=0, high= $6 \rightarrow \text{mid}=3 \rightarrow \text{arr}[3]=7 \text{ (too big)}$

low=0, high= $2 \rightarrow \text{mid}=1 \rightarrow \text{arr}[1]=3$ (too small)

low=2, high= $2 \rightarrow \text{mid}=2 \rightarrow \text{arr}[2]=5$ (too small)

low=3, high= $2 \rightarrow \text{stop} \rightarrow \text{not found}$



What is the Fibonacci Sequence?

- Each number is the sum of the two preceding numbers
- Defined as:
 - \circ F(0) = 0
 - \circ F(1) = 1
 - \circ F(n) = F(n-1) + F(n-2), for n ≥ 2

Example sequence:



Recursive Definition in C

```
int fibonacci(int n) {
   if (n == 0) return 0; // base case
   if (n == 1) return 1; // base case
   return fibonacci(n-1) + fibonacci(n-2); // recursion
}
```



Example Trace: F(5)

```
fibonacci(5)
= fibonacci(4) + fibonacci(3)
= (fibonacci(3) + fibonacci(2)) + (fibonacci(2) + fibonacci(1))
= ...
= 5
```



Tree of calls (partial):

```
F(5)

/ \

F(4) F(3)

/ \ / \

F(3) F(2) F(2) F(1)
```



Towers of Hanoi

- You have 3 pegs: Source (A), Auxiliary (B), Destination (C)
- n disks of different sizes are stacked on Source peg
- Goal: Move all disks to Destination peg
- Rules:
 - i. Only one disk moved at a time
 - ii. A disk can only be placed on an empty peg or on a larger disk
 - iii. Disks must maintain order



Recursive Idea

To move n disks from $A \rightarrow C$ using B:

- 1. Move n-1 disks from $A \rightarrow B$ (using C)
- 2. Move the largest disk from $A \rightarrow C$
- 3. Move n-1 disks from $B \rightarrow C$ (using A)



Example for 3 Disks

- 1. Move 2 disks from $A \rightarrow B$
- 2. Move disk 3 from $A \rightarrow C$
- 3. Move 2 disks from $B \rightarrow C$

Moves:

- $A \rightarrow C$
- $A \rightarrow B$
- $C \rightarrow B$
- $A \rightarrow C$
- $B \rightarrow A$
- $B \rightarrow C$



C Implementation

```
#include <stdio.h>
void hanoi(int n, char from, char to, char aux) {
    if (n == 1) {
        printf("Move disk 1 from %c to %c\n", from, to);
        return;
    hanoi(n - 1, from, aux, to);
    printf("Move disk %d from %c to %c\n", n, from, to);
    hanoi(n - 1, aux, to, from);
int main() {
    int n = 3;
    hanoi(n, 'A', 'C', 'B');
    return 0;
```

What is Merge Sort?

- A **Divide and Conquer** sorting algorithm
- Steps:
 - i. Divide array into two halves
 - ii. Recursively sort both halves
 - iii. Merge the two sorted halves
- Always O(n log n) time complexity



Algorithm Outline

- 1. If array has 0 or 1 elements → already sorted
- 2. Otherwise:
 - Divide array into left and right halves
 - Recursively sort each half
 - Merge two sorted halves into one sorted array



Merge Function

```
void merge(int arr[], int l, int m, int r) {
    int n1 = m - l + 1;
    int n2 = r - m;
    int L[n1], R[n2];
    for (int i = 0; i < n1; i++) L[i] = arr[l + i];</pre>
    for (int j = 0; j < n2; j++) R[j] = arr[m + 1 + j];
    int i = 0, j = 0, k = 1;
    while (i < n1 \&\& j < n2) {
        if (L[i] \le R[i]) arr[k++] = L[i++];
        else arr[k++] = R[j++];
    while (i < n1) arr[k++] = L[i++];
    while (j < n2) arr[k++] = R[j++];
```

Merge Sort Function

```
void mergeSort(int arr[], int l, int r) {
    if (l < r) {
        int m = l + (r - l) / 2;
        // Recursively sort left and right halves
        mergeSort(arr, l, m);
        mergeSort(arr, m + 1, r);
        // Merge sorted halves
        merge(arr, l, m, r);
```



Example Dry Run

Array = $\{12, 11, 13, 5, 6, 7\}$

Divide \rightarrow {12, 11, 13} and {5, 6, 7}

Recursively sort:

 $\{12, 11, 13\} \rightarrow \{11, 12, 13\}$

 $\{5, 6, 7\} \rightarrow \{5, 6, 7\}$

Merge:

 $\{11, 12, 13\}$ and $\{5, 6, 7\} \rightarrow \{5, 6, 7, 11, 12, 13\}$



Visualization of Recursive Calls

```
mergeSort([12,11,13,5,6,7])
  → mergeSort([12,11,13])
  → mergeSort([12,11])
  → mergeSort([12]) & mergeSort([11])
  → merge([12], [11]) → [11,12]
  → mergeSort([13]) → [13]
  → merge([11,12], [13]) → [11,12,13]
  → mergeSort([5,6,7]) → [5,6,7]
  → merge([11,12,13], [5,6,7]) → [5,6,7,11,12,13]
```

