FEAST Eigenvalue Solver v4.0 User Guide



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References

If you are using FEAST, please consider citing one or more publications below in your work.

Main reference

E. Polizzi, Density-Matrix-Based Algorithms for Solving Eigenvalue Problems, Phys. Rev. B. Vol. 79, 115112 (2009)

Math analysis

P. Tang, E. Polizzi, FEAST as a Subspace Iteration EigenSolver Accelerated by Approximate Spectral Projection; SIAM Journal on Matrix Analysis and Applications (SIMAX) 35(2), 354-390 - (2014)

Non-Hermitian solver

J. Kestyn, E. Polizzi, P. T. P. Tang, FEAST Eigensolver for Non-Hermitian Problems, SIAM Journal on Scientific Computing (SISC), 38-5, ppS772-S799 (2016);

Hermitian using Zolotarev quadrature

S. Güttel, E. Polizzi, P. T. P. Tang, G. Viaud, Optimized Quadrature Rules and Load Balancing for the FEAST Eigenvalue Solver,

SIAM Journal on Scientific Computing (SISC), 37 (4), pp2100-2122 (2015).

Eigenvalue count using stochastic estimates

E. Di Napoli, E. Polizzi, Y. Saad, Efficient Estimation of Eigenvalue Counts in an Interval, Numerical Linear Algebra with Applications, V23, I4, pp674-692,(2016).

Polynomial Non-linear eigenvalue problem – Residual Inverse Iterations

B. Gavin, A. Miedlar, E. Polizzi, FEAST Eigensolver for Nonlinear Eigenvalue Problems Journal of Computational Science, V. 27, 107, (2018)

IFEAST

B. Gavin, E. Polizzi, Krylov eigenvalue strategy using the FEAST algorithm with inexact system solves Numerical Linear Algebra with Applications, vol 25, number 5, 20 pages (2018).

PFEAST

J. Kesyn, V. Kalantzis, E. Polizzi, Y. Saad, PFEAST: A High Performance Sparse Eigenvalue Solver Using Distributed-Memory Linear Solvers

Proceedings of the International Conference for High Performance Computing, Networking, Storage and Analysis, ACM/IEEE Supercomputing Conference (SCâĂŹ16), pp 16:1-16:12, (2016).

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1 Background

"The solution of the algebraic eigenvalue problem has for long had a particular fascination for me because it illustrates so well the difference between what might be termed classical mathematics and practical numerical analysis. The eigenvalue problem has a deceptively simple formulation and the background theory has been known for many years; yet the determination of accurate solutions presents a wide variety of challenging problems."

J. H. Wilkinson- The Algebraic Eigenvalue Problem- 1965

The eigenvalue problem is ubiquitous in science and engineering applications. It can be encountered under different forms: Hermitian or non-Hermitian, and linear or non-linear. The eigenvalue problem has led to many challenging numerical questions and a central problem: how can we compute eigenvalues and eigenvectors in an efficient manner and how accurate are they?

The FEAST library package represents an unified framework for solving various family of eigenvalue problems and addressing the issues of numerical accuracy, robustness, performance and parallel scalability. Its originality lies with a new transformative numerical approach to the traditional eigenvalue algorithm design - the FEAST algorithm.

1.1 The FEAST Algorithm

The FEAST algorithm is a general purpose eigenvalue solver which takes its inspiration from the densitymatrix representation and contour integration technique in quantum mechanics¹. The algorithm gathers key elements from complex analysis, numerical linear algebra and approximation theory, to construct an optimal subspace iteration technique making use of approximate spectral projectors². FEAST can be applied for solving both standard and generalized forms of the Hermitian or non-Hermitian problems (linear or nonlinear), and it belongs to the family of contour integration eigensolvers. Once a given search interval is selected, FEAST's main computational task consists of a numerical quadrature computation that involves solving independent linear systems along a complex contour, each with multiple right hand sides. A Rayleigh-Ritz procedure is then used to generate a reduced dense eigenvalue problem orders of magnitude smaller than the original one (the size of this reduced problem is of the order of the number of eigenpairs inside the search interval/contour). FEAST offers a set of appealing features: (i) Remarkable robustness with well-defined convergence rate; (ii) All multiplicities naturally captured; (iii) No explicit orthogonalization procedure on long vectors required in practice; (iv) Reusable subspace as initial guess when solving a series of eigenvalue problems; and (v) Efficient use of both blocked BLAS-3 operations and parallel resources for solving the linear systems with multiple right hand sides. FEAST can exploit a key strength of modern computer architectures, namely, multiple levels of parallelism. Natural parallelism appears at three different levels (L1, L2 or L3): (L1) search contours can be treated separately (no overlap), (L2) linear systems can be solved independently across the quadrature nodes of the complex contour, and (L3) each complex linear system with multiple right-hand-sides can be solved in parallel. Parallel resources can be placed at all three levels simultaneously in order to achieve scalability and optimal use of the computing platform. Within a parallel environment, the main numerical task can be reduced to the solution of a single linear system using direct or iterative parallel solvers.

1.2 The FEAST Solver

FEAST release dates with main features are listed below:

```
v1.0 (Sep. 2009): Hermitian problem (standard/generalized)
```

v2.0 (Mar. 2012): SMP+MPI+RCI interfaces

v2.1 (Feb. 2013): Adoption by Intel-MKL

v3.0 (Jun. 2015): Support for non-Hermitian

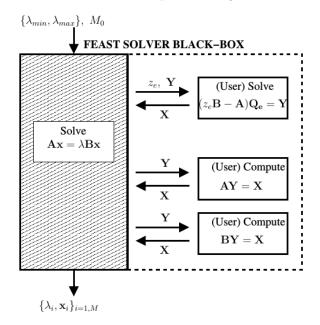
v4.0 (Feb. 2020): Residual inverse iterations - mixed precision - IFEAST (FEAST w/o factorization) - PFEAST (3 MPI levels) - Support for non-linear (polynomial) - Support for extreme eigenvalues (lowest/largest)

 $^{^{1}\}mathrm{E.}$ Polizzi, Phys. Rev. B. Vol. 79, 115112 (2009)

²P. Tang, E. Polizzi, SIMAX 35(2), $354 \hat{a} \hat{A} \hat{S} 390$ - (2014)

The FEAST package v2.1 has been featured as Intel-MKL's principal HPC eigensolver since 2013^3 . The current version of the FEAST package (v4.0) released in Feb. 2020 represents a significant upgrade since the entire FEAST package has been re-coded to perform residual inverse iterations⁴. As a result, v4.0 is in average much faster than v2.1 or v3.0 (×3 – 4 using new default optimization parameters), and it became possible to implement new important features such as IFEAST (using Inexact Iterative solver) and Non-linear polynomial FEAST. Furthermore, v4.0 features PFEAST with its 3-MPI levels of parallelism.

FEAST is a comprehensive numerical library offering both simplicity and flexibility, and packaged around a "black-box" interface as depicted in Figure 1 for the Hermitian problem.



"Black-box" interface for the Hermitian problem. In normal mode, FEAST requires a search interval and a search subspace size M_0 . It includes features such as reverse communication interfaces (RCI) that are matrix format independent, and linear system solver independent, as well as ready to use driver interfaces for dense, banded and sparse systems. For the driver interfaces the "black-box" region extends then to the right dashed box, and only the system matrices are required as inputs from the users. The RCI interfaces represent the kernel of FEAST which can be customized by the users to allow maximum flexibility for their specific applications. Users have then the possibility to integrate their own linear system solvers (direct or iterative - with or without preconditioner) and handle their own matrix-vector multiplication procedure.

The current main features of the FEAST v4.0 package include:

- Standard or Generalized Hermitian and non-Hermitian eigenvalue problems (left/right eigenvectors and bi-orthonormal basis);
- Polynomial eigenvalue problems such as quadratic, cubic, quartic, etc. (left/right eigenvectors);
- Finding eigenpair within a search contour (normal mode); Finding extreme eigenvalues (lowest/largest) for sparse Hermitian systems;
- Real/Complex arithmetic and mixed precision (single precision operations leading to double precision final results);
- Two libraries: **SMP version** (one node), and **MPI version** (multi-nodes);
- Reverse communication interfaces (RCI).
- Driver interfaces for dense (using LAPACK), banded (using SPIKE), and sparse-CSR formats (using MKL-PARDISO);
- IFEAST- FEAST w/o factorization for sparse-CSR drivers (using BiCGStab);
- PFEAST- FEAST using 3 levels of MPI parallelism for HPC (MPI solvers includes MKL-Cluster-PARDISO and PBiCGStab); Sparse and RCI interfaces compatible with local row-distributed data.
- A set of flexible and useful practical options (quadrature rules, contour shapes, stopping criteria, initial guess, fast stochastic estimates for eigenvalue counts, etc.)
- Portability: FEAST routines can be called from any Fortran or C codes.
- FEAST interfaces only require (any optimized) LAPACK and BLAS packages.
- Large number of driver examples, utility routines, and documentation.

 $^{^3} https://software.intel.com/en-us/articles/introduction-to-the-intel-mkl-extended-eigensolver.$

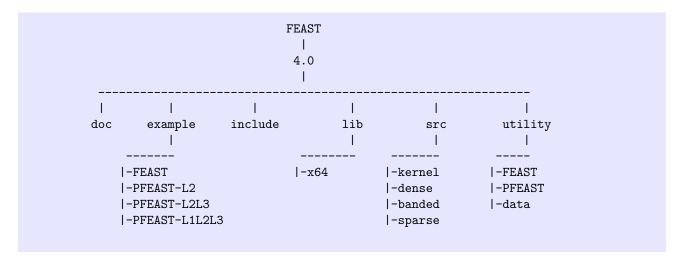
⁴B. Gavin, A. Miedlar, E. Polizzi, Journal of Computational Science, V. 27, 107, (2018); B. Gavin, E. Polizzi, in preparation (2020).

2 Installation and Setup: A Step by Step Procedure

2.1 Installation

Please follow the following steps (here for Linux/Unix systems):

- 1. Download the latest FEAST package version feast_4.0.tgz in http://www.feast-solver.org
- 2. Put the file in your preferred directory such as \$HOME directory or (for example) /opt/ directory if you have ROOT privilege.
- 3. Execute: tar -xzvf feast_4.0.tgz to create the following FEAST tree directory.



4. If <FEAST directory> denotes the package's main directory after installation, for example

~/home/FEAST/4.0 or /opt/FEAST/4.0,

it is not mandatory but recommended to define the Shell variable \$FEASTROOT, e.g.

respectively for the BASH or CSH shells. One of this command can be placed in the appropriate shell startup file in \$HOME (i.e .bashrc or .cshrc).

2.2 Compilation

Go to the directory \$FEASTROOT/src and execute make to see all available options. The same FEAST source code is used for compiling FEAST-SMP (libfeast) and/or FEAST-MPI (libpfeast). The command is:

where you can select the following options:

- <arch> : it is the name of the directory \$FEASTROOT/lib/<arch> where the FEAST libraries will be located once compiled (you can use the name of your architecture). Default is x64
 - <f90> : it is your own Fortran90 compiler (possible choices: ifort, gfortran, pgf90). Default is ifort
- <mpi>: (mandatory for compiling libpfeast only) it is your MPI library (possible choices: impi, mpich, openmpi). Defaults to impi (intel MPI)
- <mkl>: it enables Intel-MKL math library instructions (possible choices: yes, no). Default is yes.
 - if <mkl>=yes, at the linking stage, FEAST will have to be linked with Intel MKL.

• if <mkl>=no, at the linking stage, FEAST can be linked with any BLAS/LAPACK. Not using MKL will impact the behavior and performance of the FEAST sparse Driver interfaces: (i) it would not be possible to use MKL-PARDISO and cluster-MKL-PARDISO so the FEAST sparse interfaces will instead be calling IFEAST (BiCGStab); (ii) in-built sparse mat-vec routines (used by the IFEAST sparse interfaces) will be slower.

For example, if the above default options look fine with you, just use:

- make feast to compile the FEAST-SMP library and create the file libfeast.a in \$FEASTROOT/lib/<arch>
- make pfeast to compile the FEAST-MPI library and create the file libpfeast.a in \$FEASTROOT/lib/<arch>

Congratulations, FEAST is now successfully installed and compiled on your computer!!

2.3 Linking FEAST

In order to use the FEAST library for your main application code, you will then need to add the following instructions in your Makefile:

- for the LIBRARY PATH: -L\$FEASTROOT/lib/<arch>
- ullet for the LIBRARY LINKS using FEAST-SMP: -lfeast using FEAST-MPI: -lpfeast
- for the INCLUDE PATH (mandatory only for C codes): -I\$(FEASTROOT)/include

Remarks

- 1- If FEAST was compiled with the option MKL=yes, your must also link with the MKL libraries. Otherwise, you can link with any BLAS, LAPACK libraries.
- 2- If you use the FEAST banded interfaces, you need to install the SPIKE solver www.spike-solver.org SPIKE must be compiled using the same Fortran compiler used for compiling FEAST.
- 3- For C codes, the user must include the following instructions in the header:

```
#include "feast.h"
#include "feast_dense.h" //for feast dense interfaces
#include "feast_banded.h" //for feast banded interfaces
#include "feast_sparse.h" //for feast sparse interfaces
```

Using PFEAST (MPI sparse linear system solver), you must use instead:

```
#include "pfeast.h"
#include "pfeast_sparse.h"
```

2.4 HelloWorld Example (F90, C, MPI-F90, MPI-C)

This example solves a 4-by-4 real symmetric standard eigenvalue system $\mathbf{A}\mathbf{x} = \lambda\mathbf{x}$ (using dense format) where

$$\mathbf{A} = \begin{pmatrix} 2 & -1 & -1 & 0 \\ -1 & 3 & -1 & -1 \\ -1 & -1 & 3 & -1 \\ 0 & -1 & -1 & 2 \end{pmatrix}. \tag{1}$$

The four eigenvalue solutions are $\lambda = \{0, 2, 4, 4\}$. Let us suppose that one can specify a search interval (such as [3, 5]), a single call to the **dfeast_syev** subroutine should return the solutions associated with $\{4, 4\}$. The FEAST parameters need first to be set to their default values by a call to the **feastinit** subroutine. Below, we provide examples written in F90, C, MPI-F90, and MPI-C.

F90

A Fortran90 source code of helloworld.f90 is provided below:

```
program helloworld
     implicit none
     !! 4x4 eigenvalue system
     integer, parameter :: N=4
     character(len=1) :: UPLO='F' ! 'L' or 'U' also fine
     double precision, dimension (N*N) :: A=(/ 2.0d0, -1.0d0, -1.0d0, 0.0d0,&
                                                      \& -1.0 \, \mathrm{d0} \,, \quad \  3.0 \, \mathrm{d0} \,, -1.0 \, \mathrm{d0} \,, -1.0 \, \mathrm{d0} \,, \&
                                                      \&-1.0d0, -1.0d0, 3.0d0, -1.0d0,&
                                                      & 0.0\,d0, -1.0\,d0, -1.0\,d0, 2.0\,d0/)
9
10
     !! input parameters for FEAST
     integer, dimension (64) :: fpm
11
     integer :: M0=3 ! search subspace dimension
12
     {\color{red} \textbf{double precision}} \ :: \ {\color{blue} \textbf{Emin=3.0d0}} \ , \ {\color{blue} \textbf{Emax=5.0d0}} \ ! \ {\color{blue} \textbf{search interval}}
13
     !! output variables for FEAST
14
     double precision, dimension(:), allocatable :: E, res
15
     double precision, dimension (:,:), allocatable :: X
16
17
     double precision :: epsout
     integer :: loop, info, M, i
19
20 !!! Allocate memory for eigenvalues.eigenvectors, residual
     allocate(E(M0), X(N, M0), res(M0))
21
22
23 !!!!!!!!! FEAST
     call feastinit (fpm)
24
     fpm(1)=1 !! change from default value (print info on screen)
25
     call dfeast_syev (UPLO, N, A, N, fpm, epsout, loop, Emin, Emax, MO, E, X, M, res, info)
26
27
28 !!!!!!!!! REPORT
29
     if (info==0) then
         print *, 'Solutions (Eigenvalues/Eigenvectors/Residuals)'
30
         do i = 1 M
31
            print *, 'E=', E(i), 'X=', X(:,i), 'Res=', res(i)
32
             print *,''
33
         enddo
34
35
     endif
37 end program helloworld
```

To create the executable, compile and link the source program with the feast library, one can use (for example):

• ifort -o helloworld helloworld.f90 -L\$FEASTROOT/lib/<arch> -lfeast -mkl if FEAST was compiled with ifort and MKL flag was set to 'yes'.

```
gfortran -o helloworld helloworld.f90 -L$FEASTROOT>/lib/<arch> -lfeast
• -Wl,-start-group -lmkl_gf_lp64 -lmkl_gnu_thread -lmkl_core -Wl,-end-group -lgomp
-lpthread -lm -ldl
```

if FEAST was compiled with gfortran and MKL flag was set to 'yes'.

Remarks:

- 1-Many other options are possible. For example, you can link with your own BLAS/LAPACK libraries, you can compile FEAST with ifort and compile the helloworld example using gfortran via extra flag options, or vice-versa, etc.
- 2- FEAST is using a linear solver that can be threaded (the LAPACK dense solver is used for the helloworld example). Using MKL, you can control the number of threads by setting up the value of MKL_NUM_THREADS. For example, in BASH shell:

```
export MKL_NUM_THREADS=<omp>
```

where $\langle omp \rangle$ represents the number of threads (cores).

A run of the resulting executable looks like:

./helloworld

and the output of the run should be:

```
****** FEAST v4.0 BEGIN **********
*****************
Routine DFEAST_SYEV
Solving AX=eX with A real symmetric
List of input parameters fpm(1:64)-- if different from default
  fpm(1)=
| FEAST data
| Emin
                  1 3.000000000000000E+00
                     5.00000000000000E+00
 #Contour nodes
                  | 8 (half-contour)
 Quadrature rule
                  | Gauss
 Ellipse ratio y/x |
                     0.30
 System solver
                     LAPACK dense
                     Single precision
 FEAST uses MKL?
                     Yes
 Fact. stored?
                     Yes
 Initial Guess
                     Random
 Size system
| Size subspace
                        3
| FEAST runs
#It | #Eig |
                              | Error-Trace
                                                       | Max-Residual
                      Trace
              7.9999999999999964E+00
7.9999999999999982E+00
                                        1.000000000000000E+00
                                                                  2.2017475048923482E-08
              7.99999999999982E+00
                                        3.5527136788005011E-16
                                                                  1.0614112804180586E-15
==>FEAST has successfully converged with Residual tolerance <1E-12
  # FEAST outside it.
  # Eigenvalue found
                           2 from 3.999999999999987E+00 to 4.00000000000000E+00
| FEAST-RCI timing |
| Fact. cases(10,20)|
                        0.0009
```

```
| Solve cases(11,12)|
                        0.0197
| A*x cases(30,31)|
                        0.0000
                        0.0000
       cases(40,41)|
 B*x
| Misc. time
                        0.0004
| Total time (s)
                        0.0210
*****************
****** FEAST- END**************
Solutions (Eigenvalues/Eigenvectors/Residuals)
E= 4.000000000000000
                         X= -0.409757405384329
                                                   4.544205370554897E-003
 0.814970605398106
                       -0.409757405384329
                                                 1.061411280418059E-015
                                             Res=
    4.000000000000000
                         X = -0.286529001556041
                                                   0.866013481533371
-0.292955478421287
                       -0.286529001556041
                                             Res= 6.481268641478987E-016
%\end{verbatim}
```

\mathbf{C}

Similarly to the F90 example, the corresponding C source code for the helloworld example (helloworld.c) is provided below:

```
1 #include <stdio.h>
2 #include <stdlib.h>
#include "feast.h"
#include "feast_dense.h"
5 int main() {
    /* 4x4 eigenvalue system */
    int
          N=4;
    double A[16] = \{2.0, -1.0, -1.0, 0.0, -1.0, 3.0, -1.0, -1.0, -1.0, -1.0, 3.0, -1.0, 0.0, -1.0, -1.0, 2.0\}
    /* input parameters for FEAST */
10
    int fpm[64];
11
    int
          M0=3; //search subspace dimension
12
    double Emin=3.0, Emax=5.0; // search interval
13
    /* output variables for FEAST */
14
    double *E, *res, *X;
    double epsout;
16
    int loop, info, M, i;
17
18
     /* Allocate memory for eigenvalues.eigenvectors/residual */
19
    E=calloc(M0, sizeof(double)); //eigenvalues
    res=calloc(MO, sizeof(double));//eigenvectors
21
    X=calloc(N*M0, sizeof(double));//residual
22
23
    /* !!!!!!!! FEAST !!!!!!!!*/
24
    feastinit (fpm);
25
    fpm[0]=1; /*change from default value */
26
27
    dfeast_syev(&UPLO,&N,A,&LDA,fpm,&epsout,&loop,&Emin,&Emax,&M0,E,X,&M,res,&info);
28
     /*!!!!!!!! REPORT !!!!!!!*/
29
    if (info==0) {
30
       printf("Solutions_(Eigenvalues/Eigenvectors/Residuals)\n");
31
       for (i=0; i \le M-1; i=i+1){
32
         printf("E=%.15e<sub>\u00b1</sub>X=%.15e<sub>\u00b1</sub>%.15e<sub>\u00b1</sub>%.15e<sub>\u00b1</sub>%.15e<sub>\u00b1</sub>Res=%.15e\n"
33
34
                  *(E+i), *(X+i*N), *(X+1+i*N), *(X+2+i*N), *(X+3+i*N), *(res+i));
35
    }
36
37
    return 0;
```

To create the executable, compile and link the source program with the feast library, one can use (for example):

```
icc -qopenmp -I$FEASTROOT/include -o helloworld helloworld.c -L$FEASTROOT/lib/x64
-lfeast -mkl -lirc -lifcore -lifcoremt
```

if FEAST was compiled with ifort and MKL flag was set to 'yes'.

```
gcc -fopenmp -I$FEASTROOT/include -o helloworld helloworld.c -L$FEASTROOT/lib/x64
-lfeast -Wl,-start-group -lmkl_gf_lp64 -lmkl_gnu_thread -lmkl_core -Wl,-end-group
-lgomp -lpthread -lm -ldl -lgfortran
```

if FEAST was compiled with gfortran and MKL flag was set to 'yes'.

MPI-F90

FEAST can be straightforwardly parallelized using MPI at level L2 (where the FEAST inner linear systems are automatically distributed among MPI processes). As a reminder level L1 corresponds to the parallelization of the search interval, and level L3 corresponds to the parallelization of each linear system using row-data-distribution and MPI solver. Examples using L1-L2-L3 (MPI-MPI-MPI) are discussed in the PFEAST Section 3.6).

You can create the file phelloworld.f90 by cc-paste the content of helloworld.f90 and by just adding a few lines at the beginning and at the end of the program, i.e.

```
!!!! add at the very beginning
include 'mpif.h'

!!!! add after variable declarations
integer :: code
call MPI_INIT(code)

!!!! add at the end
call MPI_FINALIZE(code)
```

Your program must be compiled using the same MPI implementation used to compile the FEAST-MPI library. Once compiled, your source program must now be linked with the pfeast library. You can use (for example):

• mpiifort -o phelloworld phelloworld.f90 -L\$FEASTROOT/lib/<arch> -lpfeast -mkl if FEAST was compiled with ifort, MKL flag was set to 'yes', and MPI was chosen to be 'impi' (intel mpi).

```
mpif90.mpich -fc=gfortran phelloworld.f90 -o phelloworld -L$FEASTROOT>/lib/<arch>
-lpfeast -Wl,-start-group -lmkl_gf_lp64 -lmkl_gnu_thread -lmkl_core -Wl,-end-group
-lgomp -lpthread -lm -ldl -lifcore
```

if FEAST was compiled with gfortran, MKL flag was set to 'yes', and MPI was chosen to be 'mpich'.

A run of the resulting executable looks like:

```
mpirun -ppn 1 -n <np>./phelloworld
```

where < np > represents the number of MPI processes (here we also choose 1 MPI process per compute node with the option -ppn 1)

Remarks:

1-Scalability performances will be optimal here when the number of MPI processes <np> reaches the number of contour points (provided by the default value fpm(2)=8 in this example).

2-Since FEAST is also threaded, make sure that your number of selected threads <code><omp></code> times the number of mpi processes <code><np></code> for a given compute node (i.e. <code><omp>*<np></code>) does not exceed the number of your physical cores.

MPI-C

Similarly to the MPI-F90 example, You can create the file phelloworld.c by cc-paste the content of helloworld.c and by just adding a few lines at the beginning and at the end of the program, i.e.

```
!!!! add at the very beginning
#include <mpi.h>

!!!! change the argument list of the main function
int main(int argc, char **argv)

!!!! add after variable declarations
MPI_Init(&argc,&argv);

!!!! add at the end
MPI_Finalize();
```

Your program must be compiled using the same MPI implementation used to compile the FEAST-MPI library. Once compiled, your source program must now be linked with the pfeast library. You can use (for example):

```
mpiicc -qopenmp -I$FEASTROOT/include -o phelloworld phelloworld.c
-L$FEASTROOT/lib/x64 -lpfeast -mkl -lirc -lifcore -lifcoremt
if FEAST was compiled with ifort, MKL flag was set to 'yes', and MPI was chosen to be 'impi' (intel mpi).
```

```
mpicc -cc=gcc -fopenmp -I$FEASTROOT/include -o phelloworld phelloworld.c
-L$FEASTROOT/lib/x64 -lpfeast -Wl,-start-group -lmkl_gf_lp64 -lmkl_gnu_thread
-lmkl_core -Wl,-end-group -lgomp -lpthread -lm -ldl -lgfortran
```

if FEAST was compiled with gfortran, MKL flag was set to 'yes', and MPI was chosen to be 'mpich'.

3 FEAST Interfaces

3.1 At a Glance

There are the two different types of interfaces available in the FEAST library:

Driver interfaces

- Optimal drivers acting on commonly used matrix data storage (dense, banded, sparse-CSR, row-distributed CSR).
- Use predefined linear system solvers: LAPACK (for dense), SPIKE (for banded), MKL-PARDISO (for sparse-CSR and FEAST), BiCGStab (for sparse-CSR and IFEAST), MKL-CLUSTER-PARDISO (for row-distributed CSR and PFEAST), PBiCGStab (for row-distributed CSR and PIFEAST).

Reverse communication interfaces (RCI)

- Constitute the kernel of FEAST, independent of the matrix data formats, so users can easily customize FEAST with their own explicit or implicit data format (or row-distributed data format).
- Mat-vec routines and direct/iterative linear system solvers must also be provided by the users.

Here is the complete list of all FEAST v4.0 interfaces (186 in total).

Properties	RCI interfaces	Dense/Banded interfaces	Sparse interfaces
Linear $AX = BX\Lambda$			
Real Sym. $A = A^T, B \text{ spd}$	dfeast_srci{x}	dfeast_{sy,sb}{ev,gv}{x}	{p}d{i}feast_scsr{ev,gv}{x}
Complex Herm. $A = A^H, B \text{ hpd}$	zfeast_hrci{x}	zfeast_{he,hb}{ev,gv}{x}	{p}z{i}feast_hcsr{ev,gv}{x}
Complex Sym. $A = A^T, B = B^T$	zfeast_srci{x}	zfeast_{sy,sb}{ev,gv}{x}	{p}z{i}feast_scsr{ev,gv}{x}
Real General	zfeast_grci{x}	dfeast_{ge,gb}{ev,gv}{x}	{p}d{i}feast_gcsr{ev,gv}{x}
Complex General	zfeast_grci{x}	zfeast_{ge,gb}{ev,gv}{x}	{p}z{i}feast_gcsr{ev,gv}{x}
Polynomial $\sum_{i} A_i X \Lambda^i = 0$			
Real Sym. $A_i = A_i^T$	zfeast_srcipev{x}	dfeast_sypev{x}	{p}d{i}feast_scsrpev{x}
Complex Herm. $A_i = A_i^H$	zfeast_grcipev{x}	zfeast_hepev{x}	{p}z{i}feast_hcsrpev{x}
Complex Sym. $A_i = A_i^T$	zfeast_srcipev{x}	zfeast_sypev{x}	{p}z{i}feast_scsrpev{x}
Real General	zfeast_grcipev{x}	dfeast_gepev{x}	{p}d{i}feast_gcsrpev{x}
Complex General.	zfeast_grcipev{x}	zfeast_gepev{x}	{p}z{i}feast_gcsrpev{x}

where

- dfeast and zfeast stand for real double precision and complex double precision, respectively.
- {ev,gv} stands for either standard (i.e. B=I) or generalized eigenvalue problems.
- {x} is optional stands for the expert FEAST version which enables customized quadrature nodes/weights.
- {i} is optional stands for the IFEAST version of the sparse interfaces using inexact iterative solver.
- {p} is optional stands for the PFEAST version of the sparse interfaces using distributed MPI solvers (direct or iterative).

In addition, all the input parameters for the FEAST algorithm are contained into an integer array of size 64 named here fpm. Prior calling the FEAST interfaces, this array needs to be initialized. There exists two FEAST initialization routines:

```
!!!! initialization for FEAST
feastinit(fpm)
!!!! initialization for PFEAST (needed if the sparse interfaces use a MPI solver)
pfeastinit(fpm,L1_comm_world,nL3)
```

All input FEAST parameters are then set to their default values. The detailed list of the fpm parameters is given in Table 1. Users can modify their values accordingly before calling the FEAST interfaces.

fpm(i) F90	Description	Default values	
fpm[i-1] C			
: 4	Runtime and algorithm options		
i=1	Print runtime comments 0: Off 1: On screen	0	
	n<0: Write/Append comments in the file feast< n >.log		
i=2	#contour points for Hermitian FEAST (half-contour)	8 using FEAST	
1-2	if fpm(16)=0,2, values permitted (1 to 20, 24, 32, 40, 48, 56)	4 using IFEAST	
	if fpm(16)=1, all positive values permitted	3 using Stochastic fpm(14)=2	
i=3	Stopping convergence criteria in double precision (0 to 16)	12	
1-0	$\epsilon = 10^{-\text{fpm(3)}}$	12	
i=4	Maximum number of FEAST refinement loop allowed (≥ 0)	20 using FEAST	
		50 using IFEAST	
i=5	Provide initial guess subspace (0: No; 1: Yes)	0	
i=6	Convergence criteria (for solutions in the search contour)	1	
	0: Using relative error on the trace epsout i.e. epsout< ϵ		
	1: Using relative residual res i.e. $\max_i \operatorname{res}(i) < \epsilon$		
i=8	#contour points for non-Herm./poly. FEAST (full-contour)	16 using FEAST	
	if fpm(16)=0, values permitted (2 to 40, 48, 64, 80, 96, 112)	8 using IFEAST	
	if fpm(16)=1, all values permitted (>2)	6 using Stochastic fpm(14)=2	
i=9	L2 communicator for PFEAST	set by call to pfeastinit	
i=10	Store linear system factorizations (0: No; 1: Yes).	1 using all Driver interfaces	
		0 using RCI interfaces	
i=14	0: FEAST normal execution	0	
	1: Return subspace Q after 1 contour		
	2: Stochastic estimate of #eigenvalues inside search contour		
i=15	#Contours for non-Hermitian or polynomial FEAST.	0 using non-sym. drivers	
	0: two-sided contour (compute right/left eigenvectors)	1 using Stochastic fpm(14)=2	
	1: one-sided contour (compute only right eigenvectors)	2 using sym. drivers	
: 10	2: one sided contour (left=right* eigenvectors)	0 for Hermitian FEAST	
i=16	Integration type (0: Gauss 1: Trapezoidal; 2: Zolotarev)		
	Remark: option 2 only for Hermitian	1 for non-Herm./poly. FEAST 1 for IFEAST	
i=18	Ellipse contour ratio 'vertical axis'/'horizontal axis' (≥ 0)	30 for Hermitian FEAST	
1-10	fpm(18)/100 = ratio $\sqrt{100}$ ratio	100 for non-Herm./poly. FEAST	
	1pm(10)/100 = 1atto	100 for IFEAST	
i=19	Ellipse rotation angle in degree from vertical axis [-180:180]	0	
1 10	Remark: only for non-Hermitian		
i=49	L3 communicator for PFEAST	set by call to pfeastinit	
	Driver interface options	Set 25, can to P104201111	
i=40	Search interval option for sparse Hermitian drivers	0	
1 10	0: search interval provided by user		
	-1: search M0/2 lowest eigenvalues- return search interval		
	1: search M0/2 largest eigenvalues- return search interval		
i=41	Matrix scaling for sparse drivers (0: No; 1: Yes).	1	
i=42	Mixed Precision for all drivers	1	
	0: use double precision linear system solvers		
	1: use single precision linear system solvers		
i=43	Automatic switch from FEAST to IFEAST drivers	0	
	(0:feast,1:ifeast)		
i=45	Accuracy of BiCGStab in IFEAST $\mu = 10^{-\text{fpm}(45)}$	1	
i=46	Maximum #iterations for BiCGStab in IFEAST 40		
i=60	Output: returns the total number of BicGstab iterations	N/A	
1-00			

Table 1: List FEAST parameters with default input values.

Remark: Using the C language, the components of the fpm array starts at 0 and stops at 63. Therefore, the components fpm[j] in C (j=0-63) must correspond to the components fpm(i) in Fortran (i=1-64) specified above (i.e. fpm[i-1]=fpm(i)).

Errors and warnings encountered during a run of the FEAST package are stored in an integer variable, info. If the value of the output info parameter is different than "0", either an error or warning was encountered. The possible return values for the info parameter along with the error code descriptions, are given in Table 2.

info	Classification	Description	
202	Error	Problem with size of the system N	
201	Error	Problem with size of subspace MO	
200	Error	Problem with Emin, Emax or Emid, r	
(100 + i)	Error	Problem with i^{th} value of the input FEAST parameter (i.e fpm(i))	
7	Warning	The search for extreme eigenvalues has failed, search contour must be set by user	
6	Warning	FEAST converges but subspace is not bi-orthonormal	
5	Warning	Only stochastic estimation of #eigenvalues returned fpm(14)=2	
4	Warning	Only the subspace has been returned using fpm(14)=1	
3	Warning	Size of the subspace MO is too small (MO<=M)	
2	Warning	No Convergence (#iteration loops>fpm(4))	
1	Warning	No Eigenvalue found in the search interval	
0	Successful exit		
-1	Error	Internal error conversion single/double	
-2	Error	Internal error of the inner system solver in FEAST Driver interfaces	
-3	Error	Internal error of the reduced eigenvalue solver	
		Possible cause for Hermitian problem: matrix B may not be positive definite	
-(100+i)	Error	Problem with the i^{th} argument of the FEAST interface	

Table 2: Return code descriptions for the parameter info.

Quick Tutorial using FEAST Drivers

- 1. Identify your FEAST driver. Look at the corresponding section of this documentation to set up the argument lists.
- 2. Specify a search contour enclosing the wanted eigenvalues (normal FEAST mode). Alternatively, you can also use the Hermitian sparse drivers to search for the MO/2 lowest (fpm(40)=-1) or largest (fpm(40)=1) eigenpairs (here MO must be set to two times the number of wanted eigenvalues).
- 3. In normal FEAST mode, specify the search subspace size MO as an overestimation of your estimated #eigenvalues M within the contour (typically MO≥ 1.5M). If needed, user can take advantage of fast stochastic estimates for M within a particular contour using fpm(14)=2.
- 4. Change the fpm default options if needed and run the code.

Tips

- The FEAST convergence rate depends on the choice of the search subspace size M0, the number of contour points, and the nature of the quadrature. To improve the convergence rate, you have the possibility to:
 - keep on increasing the number of quadrature nodes fpm(2) (fpm(8) for non-Hermitian/Polynomial).
 - keep on increasing M0 for the Gauss-Legendre or Trapezoidal quadrature.
 - switch to Zolotarev quadrature for the Hermitian problem with fpm(16)=2 (Zolotarev is ideally suited to deal with continuum spectra without the need to increase the subspace size $M0\sim M$).
- Although FEAST could be used to seek 1000's of eigenpairs within a single search contour, the size of the search subspace M0 is supposed to be much smaller than the size of the eigenvalue problem N. As a result, the arithmetic complexity would mainly depend on the inner system solve (i.e. O(NM₀) for narrow banded or sparse system solvers). If you are looking for a very large number of eigenvalues, it is recommended to consider multiple search intervals to be solved in parallel using PFEAST.
- FEAST v4.0 is using an inverse residual iteration algorithm which enables the linear systems to be solved with very low accuracy with no impact on the FEAST double precision convergence rate (!).

Consequently, all FEAST linear systems are solved in single precision by default (fpm(42)=1). Using the RCI interfaces, users can then plug in their own low accuracy (single precision or less) direct or iterative solver. Additionally, all the linear system factorizations can be kept in memory by (using fpm(10)=1) which improves performance but use more memory.

- Using the FEAST-SMP library, parallelism at the third level L3 (linear system solves) can only be achieved using the threading capabilities of the linear system solver and via the shell variable MKL_NUM_THREADS if Intel-MKL is used or the shell variable OMP_NUM_THREADS if SPIKE is used for the banded interfaces.
- Using the FEAST-MPI library, you can trivially parallelize the second level L2 (contour points) and keep on using the same FEAST/IFEAST driver interfaces with shared memory solver at L3. Scalability performances will be optimal when the number of MPI processes reaches the number of contour points (either fpm(2) for FEAST Hermitian or fpm(8) for FEAST non-Hermitian and Polynomial).
- The FEAST-MPI library also offers the possibility to use a MPI solver at level L3. This scheme is called PFEAST and it is detailed in Section 3.6.

3.2 FEAST Hermitian

We note the following:

• The Table below details the series of arguments in each $\{List-A\}$, and $\{List-B\}$ that are specific to the type of matrix format represented above by F (as a placeholder).

	\mathbf{F}	List-A	List-B
Dense	{y,e}	{ UPLO, N, A, LDA }	{ B, LDB }
Banded	b	{ UPLO, N, ka, A, LDA }	{ kb, B, LDB }
Sparse	csr	{ UPLO, N, A, IA, JA }	{ B, IB, JB }

- Table 4 details the specific matrix-format arguments in {List-A} and {List-B}
- Table 3 details the common arguments in all the Hermitian FEAST interfaces above,
- Table 5 details the arguments for the Hermitian RCI interfaces (in red above).

	Type	I/O	Description
fpm	integer(64)	in/out	FEAST input parameters (see Table 1)
epsout	double real	out	Trace relative error $ trace_k - trace_{k-1} / \max(Emin , Emax)$
loop	integer	out	# of FEAST subspace iterations
Emin, Emax	double real	in/out	Lower and Upper bounds of search interval
			Remark: Output values if fpm(40)=+-1 for sparse drivers
MO	integer	in/out	Search subspace dimension
			On entry: initial guess MO≥M
			On exit: new suitable MO if guess too large
			Remark: M0=2*Wanted if fpm(40)=+-1 for sparse drivers
E	double real(MO)	in/out	Eigenvalues
			On entry: initial guess if fpm(5)=1 (previous FEAST run)
			On exit: Eigenvalues solutions E(1:M)
			Remark: the E(M+1:M0) values are outside [Emin, Emax]
X	double real(N,MO) using dfeast	in/out	Eigenvectors (N: size of the system)
	double complex(N,MO) using zfeast		On entry: initial guess if fpm(5)=1 (previous FEAST run)
			On exit: Eigenvectors solutions X(1:N,1:M)
			Remark: if fpm(14)=1, first Q subspace on exit
M	integer	out	#Eigenvalues found in [Emin, Emax]
			#Estimated eigenvalues if fpm(14)=2
res	double real(MO)	out	Relative residual $ \mathbf{A}\mathbf{x_i} - \lambda_i \mathbf{B}\mathbf{x_i} _2 / \alpha \mathbf{B}\mathbf{x_i} _2$
			$ \text{ with } \alpha = \max(\texttt{Emin} , \texttt{Emax})$
info	integer	out	Error handling (see Table 2 for all INFO codes)
Zne,Wne	double complex(fpm(2))	in	Custom integration nodes and weights- Expert mode

Table 3: List of common arguments for the FEAST Hermitian Driver interfaces.

	Type	I/O	Description
Comn	non		
UPLO	character(len=1)	in	Matrix Storage ('F', 'L', 'U') 'F': Full; 'L': Lower; 'U': Upper
N	integer	in	Size of the system
Dense	9		
A	double real(LDA,N) using dfeast double complex(LDA,N) using zfeast	in	Eigenvalue system (Stiffness) matrix
LDA	integer	in	Leading dimension of A LDA>=N
В	double real(LDB,N) using dfeast double complex(LDA,N) using zfeast	in	Eigenvalue system (Mass) matrix
LDB	integer	in	Leading dimension of B LDB>=N
Bande	ed		
ka	integer	in	The number of sub or super-diagonals within the band of A.
A	double real(LDA,N) using dfeast double complex(LDA,N) using zfeast	in	Eigenvalue system (Stiffness) matrix
LDA	integer	in	Leading dimension of A LDA>=2*ka+1 if UPLO='F' LDA>=ka+1 if UPLO='L' or 'U'
kb	integer	in	The number of sub or super-diagonals within the band of B.
В	double real(LDB,N)using dfeastdouble complex(LDB,N)using zfeast	in	Eigenvalue system (Mass) matrix
LDB	integer	in	Leading dimension of B LDB>=2*kb+1 if UPLO='F' LDB>=kb+1 if UPLO='L' or 'U'
Spars	e-csr		
A	double real(IA(N+1)-1) using dfeast double complex(IA(N+1)-1) using zfeast	in	Eigenvalue system (Stiffness) matrix - CSR values
IA	integer(N+1)	in	Sparse CSR Row array of A.
JA	integer(IA(N+1)-1)	in	Sparse CSR Column array of A.
В	double real(IB(N+1)-1) using dfeast double complex(IB(N+1)-1) using zfeast	in	Eigenvalue system (Mass) matrix - CSR values
IB	integer(N+1)	in	Sparse CSR Row array of B.
JB	integer(IB(N+1)-1)	in	Sparse CSR Column array of B.

Table 4: List of arguments that are matrix-format specific for the FEAST Driver interfaces. Applicable to Hermitian and Non-Hermitian Drivers.

	Type		I/O	Description
ijob	integer		in/out	On entry: ijob=-1 (initialization)
				On exit: ID of the FEAST_RCI operation
N	integer		in	Size of the system
Ze	double complex		out	Coordinate along the complex contour
work1	double real(N,MO)	using dfeast	in/out	Workspace
	double $complex(N,MO)$	using zfeast		
work2	double complex(N,MO)		in/out	Workspace
Aq, Bq	double real(MO,MO)	using dfeast	in/out	Workspace for the reduced eigenvalue problem
	$\operatorname{double\ complex}(\texttt{MO,MO})$	using zfeast		

Table 5: List of arguments for the FEAST Hermitian RCI interfaces.

Hermitian Driver Interfaces: Examples

Let us consider the following systems:

System1 a "real symmetric" generalized eigenvalue problem $\mathbf{A}\mathbf{x} = \lambda \mathbf{B}\mathbf{x}$, where \mathbf{A} is real symmetric and \mathbf{B} is symmetric positive definite. \mathbf{A} and \mathbf{B} are of the size N=1671 and have the same sparsity pattern with number of non-zero elements NNZ=11435.

System2 a "complex Hermitian" standard eigenvalue problem $\mathbf{A}\mathbf{x} = \lambda\mathbf{x}$, where \mathbf{A} is complex Hermitian. \mathbf{A} is of size N = 600 with number of non-zero elements NNZ = 2988.

The \$FEASTROOT/example/FEAST directory provides Fortran and C implementation of these systems using both dense, banded and sparse-CSR storage. Here, the complete list of routines:

	System1	System2
dense		
	{F90,C}dense_dfeast_sygv	{F90,C}dense_zfeast_heev
banded		
	{F90,C}dense_dfeast_sbgv	{F90,C}dense_zfeast_hbev
sparse		
	{F90,C}dense_dfeast_scsrgv	{F90,C}dense_zfeast_hcsrev
	{F90,C}dense_dfeast_scsrgv_lowest*	
	(*using dfeast_scsrgv to compute few lowest eig.)	

The \$FEASTROOT/example/PFEAST-L2 directory provides the parallel implementation of all these routines using FEAST-MPI. The routine names are preceded by the letter P. The MPI parallelization operates only at the second level L2 where all the (fpm(2)) linear systems are distributed among the MPI processes.

Hermitian RCI Interfaces

Using the FEAST_RCI interfaces, the ijob parameter must first be initialized with the value -1. Once the RCI interface is called, the value of the ijob output parameter, if different than 0, is used to identify the FEAST operation that needs to be completed by the user. Users have then the possibility to customize their own matrix direct or iterative factorization and linear solve techniques as well as their own matrix multiplication routine.

Here is a general (F90) template example of RCI for solving real symmetric problem:

```
1 ijob=-1! initialization
_2 do while (ijob/=0)
3 call dfeast_srci(ijob, N, Ze, work1, work2, Aq, Bq, fpm, epsout, loop, Emin, Emax, M0, E, X, M, res, info)
   select case(ijob)
  case (10) !!Factorize the complex matrix Az <=(ZeB-A) - or factorize a preconditioner of ZeB-A
            !!REMARK: Az can be formed and factorized using single precision arithmetic
  ..... <<< user entry
   case (11) !! Solve the linear system with fpm(23) rhs; Az * Qz=work2(1:N,1:fpm(23))
            !!Result (in place) in work2 \leq Qz(1:N,1:fpm(23))
            !!REMARKS: -Solve can be performed in single precision
10
            1.1
                        -Low accuracy iterative solver are ok
11
  ..... <<< user entry
12
   case (30) !! Perform multiplication A * X(1:N,i:j) result in work1(1:N,i:j)
13
            11
                      where i=fpm(24) and j=fpm(24)+fpm(25)-1
14
15 . . . . . .
            ..... <<< user entry
  case (40) !! Perform multiplication B * X(1:N,i:j) result in work1(1:N,i:j)
                       where i=fpm(24) and j=fpm(24)+fpm(25)-1
17
            !!REMARK: user must set work1(1:N,i:j)=X(1:N,i:j) if B=I
18
19 ..... <<< user entry
20 end select
21 end do
22
```

Here is a general (F90) template example of RCI for solving complex Hermitian problem:

```
ı ijob=-1! initialization
_2 do while (ijob/=0)
3 call zfeast hrci(ijob, N, Ze, work1, work2, Aq, Bq, fpm, epsout, loop, Emin, Emax, MO, E, X, M, res, info)
   select case(ijob)
   case (10) !! Factorize the complex matrix Az <=(ZeB-A) - or factorize a preconditioner of ZeB-A
            !!REMARK: Az can be formed and factorized using single precision arithmetic
          ..... <<< user entry
  case(11) !!Solve the linear system with fpm(23) rhs; Az * Qz=work2(1:N,1:fpm(23))
            !! Result (in place) in work2 \le Qz(1:N,1:fpm(23))
            !!REMARKS: -Solve can be performed in single precision
10
            1.1
                        -Low accuracy iterative solver are ok
11
12 ..... <<< user entry
   case (20) !![Optional: *only if * needed by case (21)]
13
            !!Factorize the complex matrix Az^H
14
            !!REMARKS: -The matrix Az from case(10) cannot be overwritten
15
            1.1
                       -case (20) becomes obsolete if the solve in case (21) can be performed
            1.1
                         by reusing the factorization in case (10)
17
         .... <<< user entry
18
  case(21) !!Solve the linear system with fpm(23) rhs; Az^H * Qz=work2(1:N,1:fpm(23))
19
            !! Result (in place) in work2 \le Qz(1:N,1:fpm(23))
20
           ..... <<< user entry
  case (30) !! Perform multiplication A * X(1:N,i:j) result in work1(1:N,i:j)
22
23
                       where i=fpm(24) and j=fpm(24)+fpm(25)-1
           ..... <<< user entry
24
  case (40) !!Perform multiplication B * X(1:N,i:j) result in work1(1:N,i:j)
25
            !! where i=fpm(24) and j=fpm(24)+fpm(25)-1
            !!REMARK: user must set work1(1:N,i:j)=X(1:N,i:j) if B=I
27
  ..... <<< user entry
28
29 end select
30 end do
31
```

3.3 FEAST Non-Hermitian

We note the following:

• The Table below details the series of arguments in each {List-A}, and {List-B} that are specific to the type of matrix format represented above by F (as a placeholder).

	\mathbf{F}	List-A	List-B
Dense			
Symmetric	У	{UPLO, N, A, LDA}	{B, LDB}
General	е	{N, A, LDA}	{B, LDB}
Banded			
Symmetric	b	{UPLO, N, ka, A, LDA}	{kb, B, LDB}
General	Ъ	{N, kla, kua, A, LDA}	{klb, kub, B, LDB}
Sparse			
Symmetric	csr	{UPLO, N, A, IA, JA}	{B, IB, JB}
General	csr	{N, A, IA, JA}	{B, IB, JB}

- Similarly to the Hermitian case, Table 4 details the specific matrix-format arguments in {List-A} and {List-B}. For the banded drivers and the real/complex general cases, kla (resp. klb) represents the number of sub-diagonals for matrix A (resp. matrix B), and kua (resp. kub) the number of super-diagonals for matrix A (resp. matrix B).
- Table 7 details the common arguments in all the non-Hermitian FEAST interfaces above.
- Table 6 details the arguments for the non-Hermitian RCI interfaces (in red above).

	Type	I/O	Description
ijob	integer	in/out	On entry: ijob=-1 (initialization)
			On exit: ID of the FEAST_RCI operation
N	integer	in	Size of the system
Ze	double complex	out	Coordinate along the complex contour
work1	double $complex(N,MO)$	in/out	Workspace
	or		
	double $complex(N,2*M0)$		
	(if left vector calculated for non-sym.		
	<pre>interfaces and fpm(15)=0)</pre>		
work2	double complex(N,MO)	in/out	Workspace
Aq, Bq	double $complex(MO,MO)$	in/out	Workspace for the reduced eigenvalue problem

Table 6: List of arguments for the FEAST RCI interfaces. Applicable to Non-Hermitian and Polynomial Drivers.

	Туре	I/O	Description
fpm	integer(64)	in/out	FEAST input parameters (see Table 1)
epsout	double real	out	Trace relative error $ trace_k - trace_{k-1} / \max(Emid + r)$
loop	integer	out	# of FEAST subspace iterations
Emid	double complex	in	Coordinate center of the contour ellipse
r	double real	in	Horizontal radius of the contour ellipse
MO	integer	in/out	Search subspace dimension
			On entry: initial guess MO≥M
			On exit: new suitable MO if guess too large
E	double complex(MO)	in/out	Eigenvalues
			On entry: initial guess if fpm(5)=1 (previous FEAST run)
			On exit: Eigenvalues solutions E(1:M)
			Remark: the E(M+1:M0) values are outside the contour
X	double $complex(N,MO)$	in/out	Eigenvectors (N: size of the system)
	or		On entry: initial guess if fpm(5)=1 (previous FEAST run)
	double complex(N,2*MO)		On exit: (right) Eigenvectors solutions X(1:N,1:M)
	(if left vectors calculated for		Remarks: -left vectors (if calculated) in X(1:N,M0+1:M0+M)
	non-sym. drivers and fpm(15)=0)		-if fpm(14)=1, first Q subspace on exit
M	integer	out	#Eigenvalues found inside contour
			#Estimated eigenvalues if fpm(14)=2
res	double complex(MO)	out	Relative residual res(1:M) (right); res(MO+1,MO+M) (left)
	or		(right) $ \mathbf{A}\mathbf{x_i} - \lambda_i \mathbf{B}\mathbf{x_i} _2 / \alpha \mathbf{B}\mathbf{x_i} _2$
	double complex(2*M0)		(left) $ \mathbf{A}^{\mathbf{H}}\mathbf{x}_{i} - \lambda_{i}^{*}\mathbf{B}^{\mathbf{H}}\mathbf{x}_{i} _{2}/ \alpha\mathbf{B}^{\mathbf{H}}\mathbf{x}_{i} _{2}$
	(if left vectors calculated)		with $\alpha = \max(\texttt{Emid} + \mathbf{r})$
info	integer	out	Error handling (see Table 2 for all INFO codes)
Zne, Wne	double complex(fpm(8))	in	Custom integration nodes and weights- Expert mode

Table 7: List of common arguments for the FEAST Driver interfaces. Applicable to Non-Hermitian and Polynomial Drivers.

non-Hermitian Driver Interfaces: Examples

Let us consider the following systems:

System3 a "real non-symmetric" generalized eigenvalue problem $\mathbf{A}\mathbf{x} = \lambda \mathbf{B}\mathbf{x}$, where \mathbf{A} and \mathbf{B} are real non-symmetric. \mathbf{A} and \mathbf{B} are of the size N=1671 and have the same sparsity pattern with number of non-zero elements NNZ = 13011.

System4 a "complex symmetric" standard eigenvalue problem $\mathbf{A}\mathbf{x} = \lambda\mathbf{x}$, where \mathbf{A} is complex symmetric. \mathbf{A} is of size N = 801 with number of non-zero elements NNZ = 24591.

The \$FEASTROOT/example/FEAST directory provides Fortran and C implementation of these systems using both dense, banded and sparse-CSR storage. Here, the complete list of routines (System4 includes examples for using expert routines as well):

	System3	System4			
dense					
	{F90,C}dense_dfeast_gegv	{F90,C}dense_zfeast_syev{x}			
banded					
	{F90,C}dense_dfeast_gbgv	{F90,C}dense_zfeast_sbev{x}			
sparse					
	{F90,C}dense_dfeast_gcsrgv	{F90,C}dense_zfeast_scsrev{x}			

The \$FEASTROOT/example/PFEAST-L2 directory provides the parallel implementation of all these routines using FEAST-MPI. The routine names are preceded by the letter P. The MPI parallelization operates only at the second level L2 where all the (fpm(2)) linear systems are distributed among the MPI processes.

non-Hermitian RCI Interfaces

The RCI template for solving the complex symmetric problem is the same than the one used for solving the real symmetric case in Section 3.2 (just replace dfeast_srci{x} by zfeast_srci{x}).

Here is a general (F90) template example of RCI for solving real/complex general problem:

```
1 ijob=-1! initialization
_2 do while (ijob/=0)
3 call zfeast_grci(ijob, N, Ze, work1, work2, Aq, Bq, fpm, epsout, loop, Emid, r, M0, E, X, M, res, info)
  select case(ijob)
   case (10) !! Factorize the complex matrix Az <=(ZeB-A) - or factorize a preconditioner of ZeB-A
             !!REMARK: Az can be formed and factorized using single precision arithmetic
            ..... <<< user entry
   case (11) !! Solve the linear system with fpm(23) rhs; Az * Qz=work2(1:N,1:fpm(23))
             !!Result (in place) in work2 <= Qz(1:N,1:fpm(23))
             !!REMARKS: -Solve can be performed in single precision
10
             1.1
                         -Low accuracy iterative solver are ok
11
12 ..... <<< user entry
   case (20) !![Optional: *only if * needed by case (21)]
             !!Factorize the complex matrix Az^H
14
             !!REMARKS: -The matrix Az from case (10) cannot be overwritten
15
             1.1
                       -case (20) becomes obsolete if the solve in case (21) can be performed
16
             1.1
                          by reusing the factorization in case (10)
17
18
          .... <<< user entry
   case(21) !!Solve the linear system with fpm(23) rhs; Az^H * Qz=work2(1:N,1:fpm(23))
19
            !! Result (in place) in work2 \le Qz(1:N,1:fpm(23))
20
            \ldots \ldots <\!\!<\!\!< user\ entry
21
   case (30) !! Perform multiplication A * X(1:N,i:j) result in work1(1:N,i:j)
22
            1.1
                        where i=fpm(24) and j=fpm(24)+fpm(25)-1
23
          .... <<< user entry
24
   case (31) !! Perform multiplication A^H * X(1:N, i:j) result in work1(1:N, i:j)
            1.1
                       where i=fpm(34) and j=fpm(34)+fpm(35)-1
26
  ..... <<< user entry
27
   case (40) !! Perform multiplication B * X(1:N, i:j) result in work1(1:N, i:j)
28
                     where i=fpm(24) and j=fpm(24)+fpm(25)-1
29
             ! ! REMARK: user must set work1 (1:N, i:j) = X (1:N, i:j) if B = I
30
            ..... <<< user entry
31
   case (41) !! Perform multiplication B^H * X(1:N, i:j) result in work1(1:N, i:j)
32
             !! where i = fpm(34) and j = fpm(34) + fpm(35) - 1
33
             ! ! REMARK: user must set work1(1:N, i:j) = X(1:N, i:j) if B=I
34
           ..... <<< user entry
36 end select
37 end do
38
```

3.4 FEAST Polynomial (quadratic, cubic, quartic, etc.)

```
Solving \sum_{i=0}^{p} \lambda^i A_i x = 0 !!!!  \{ \text{Ai} \} \text{ Real-Sym., Complex Herm., Complex Sym., Real General and Complex General dfeast_s } Fpev \{x \} (\{ \text{List-A} \}, \text{fpm,epsout,loop,Emid,r,M0,E,X,M,res,info,} \{\text{Zne,Wne} \}) z \text{feast\_h } Fpev \{x \} (\{ \text{List-A} \}, \text{fpm,epsout,loop,Emid,r,M0,E,X,M,res,info,} \{\text{Zne,Wne} \}) z \text{feast\_s } Fpev \{x \} (\{ \text{List-A} \}, \text{fpm,epsout,loop,Emid,r,M0,E,X,M,res,info,} \{\text{Zne,Wne} \}) d \text{feast\_g } Fpev \{x \} (\{ \text{List-A} \}, \text{fpm,epsout,loop,Emid,r,M0,E,X,M,res,info,} \{\text{Zne,Wne} \}) z \text{feast\_g } Fpev \{x \} (\{ \text{List-A} \}, \text{fpm,epsout,loop,Emid,r,M0,E,X,M,res,info,} \{\text{Zne,Wne} \}) !!!! \quad RCI \quad (\text{format independent}) \quad - \text{Real/Complex Symmetric and Hermitian/Real/Complex General} z \text{feast\_srcipev} \{x \} (\text{ijob,p,N,Ze,work1,work2,Aq,Bq,fpm,epsout,loop,Emid,r,M0,E,X,M,res,info,} \{\text{Zne,Wne} \}) z \text{feast\_grcipev} \{x \} (\text{ijob,p,N,Ze,work1,work2,Aq,Bq,fpm,epsout,loop,Emid,r,M0,E,X,M,res,info,} \{\text{Zne,Wne} \})
```

We note the following:

• The Table below details the series of arguments in $\{List-A\}$ that are specific to the type of matrix format represented above by F (as a placeholder). Remark: the banded format is not supported.

	F	List-A			
Dense					
Symmetric/Hermitian	{y,e}	{UPLO, p, N, A, LDA}			
General	е	{p, N, A, LDA}			
Sparse					
Symmetric/Hermitian	csr	{UPLO, p, N, A, IA, JA}			
General	csr	{p, N, A, IA, JA}			

- Table 8 details the specific matrix-format arguments in {List-A}.
- The common arguments in all the polynomial FEAST interfaces above are identical to the ones given for the non-Hermitian case in Table 7.
- The argument for the RCI interfaces (in red above) are also identical to the ones given for the non-Hermitian case in Table 7. We note the addition of the argument integer p which stands for the degree of the polynomial (e.g. p=2 fro quadratic, p=3 for cubic etc.)

	Type	I/O	Description					
Comm	Common							
UPLO	character(len=1)	in	Matrix Storage ('F', 'L', 'U')					
			'F': Full; 'L': Lower; 'U': Upper					
N	integer	in	Size of the system					
p	integer	in	Degree of the polynomial					
Dense								
A	double real(LDA,N,p+1) using dfeast	in	All system matrices A(:,i)					
	double complex(LDA,N,p+1) using zfeast							
LDA	integer	in	1st Leading dimension of A LDA>=N					
Sparse	e-csr - where nnz_max stands for the max of nor	-zero e	elements among all A sparse matrices					
A	double real(nnz_max,p+1) using dfeast	in	All system matrices A(:,i) - CSR values					
	double complex(nnz_max,p+1) using zfeast							
IA	integer(N+1,p+1)	in	All sparse CSR Row array of A(:,i).					
JA	<pre>integer(nnz_max,p+1)</pre>	in	All sparse CSR Column array of A(:,i).					

Table 8: List of arguments that are matrix-format specific for the FEAST Polynomial Driver interfaces. Remark: $A(:,i)==A_{i-1}$ in Fortran.

Polynomial Driver Interfaces: Examples

Let us consider the following system:

System5 a quadratic eigenvalue problem $(\mathbf{A_2}\lambda^2 + \mathbf{A_1}\lambda + \mathbf{A_0})\mathbf{x} = \mathbf{0}$, where $\mathbf{A_2}, \mathbf{A_1}, \mathbf{A_0}$ are real symmetric. The size of the system is N = 1000 with 2998 non-zero elements for A_0 and A_1 , and 1000 for A_2 .

The \$FEASTROOT/example/FEAST directory provides Fortran and C implementation of this system using both dense and sparse-CSR storage. Here, the complete list of routines:

	System5
dense	
	{F90,C}dense_dfeast_sypev
sparse	
	{F90,C}dense_dfeast_scsrpev

The \$FEASTROOT/example/PFEAST-L2 directory provides the parallel implementation of these routines using FEAST-MPI. The routine names are preceded by the letter P. The MPI parallelization operates only at the second level L2 where all the (fpm(2)) linear systems are distributed among the MPI processes.

Polynomial RCI Interfaces

Here is a general (F90) template example of RCI for solving real/complex symmetric problem:

```
1 !!! Here your polynomial matrices are in stored A[i] (i=1,...,p+1) (p polynomial degree)
_{2} !!! All the matrices are real or complex symmetric
3 ijob=-1! initialization
4 do while (ijob/=0)
5 call zfeast_srcipev(ijob,p,N,Ze,work1,work2,Aq,Bq,fpm,epsout,loop,Emid,r,M0,E,X,M,res,info)
6 select case(ijob)
   case (10) !!Form and Factorize P(Ze)
            !!Example for the quadratic problem: P(Ze)=A[3]*Ze**2+A[2]*Ze+A[1]
            !!REMARK: P(Ze) can be formed and factorized using single precision arithmetic
  ..... <<< user entry
10
  case(11) !!Solve the linear system with fpm(23) rhs; P(Ze)* Qz=work2(1:N,1:fpm(23))
            !!Result (in place) in work2 \ll Qz(1:N,1:fpm(23))
12
            ! !REMARKS:
                       -Solve can be performed in single precision
13
            1.1
                        -Low accuracy iterative solver are ok
14
         ..... <<< user entry
15
   case (30) !!Perform multiplication A[fpm(57)] * X(1:N,i:j) result in work1(1:N,i:j)
                       where i=fpm(24) and j=fpm(24)+fpm(25)-1; fpm(57) take the values 1...p+1
17
      ..... <<< user entry
18
   end select
19
20 end do
21
```

Here is a general (F90) template example of RCI for solving Hermitian or real/complex general problem:

```
1!!! Here your polynomial matrices are in stored A[i] (i=1,..,p+1) (p polynomial degree)
2 !!! At least one matrix is not real/complex symmetric
3 ijob=-1! initialization
4 do while (ijob/=0)
 \verb| call zfeast\_grcipev(ijob,p,N,Ze,work1,work2,Aq,Bq,fpm,epsout,loop,Emid,r,M0,E,X,M,res,info)| \\
  select case(ijob)
   case(10) !!Form and Factorize P(Ze)
            !!Example for the quadratic problem: P(Ze)=A[3]*Ze**2+A[2]*Ze+A[1]
            !!REMARK: P(Ze) can be formed and factorized using single precision arithmetic
            ..... <<< user entry
case(11) !! Solve the linear system with fpm(23) rhs; P(Ze)* Qz=work2(1:N,1:fpm(23))
11
            !!Result (in place) in work2 \ll Qz(1:N,1:fpm(23))
12
                       -Solve can be performed in single precision
13
            1.1
14
                        -Low accuracy iterative solver are ok
15 ..... <<< user entry
```

```
case(20) !![Optional: *only if * needed by case(21)]
            !! Factorize the complex matrix P(Ze)^H
17
            !!REMARKS: -The factorization P(Ze) from case(10) cannot be overwritten
18
                       -case(20) becomes obsolete if the solve in case(21) can be performed
            1.1
19
                         by reusing the factorization in case (10)
            1.1
20
         ..... <<< user entry
    \textbf{case} \ (21) \ \texttt{!!Solve the linear system with } fpm(23) \ rhs; \ P(Ze)^H * Qz=work2(1:N,1:fpm(23)) 
22
23
            !! Result (in place) in work2 \le Qz(1:N,1:fpm(23))
           ..... <<< user entry
24
   case (30) !! Perform multiplication A[fpm(57)] * X(1:N,i:j) result in work1(1:N,i:j)
25
          1.1
                       where i=fpm(24) and j=fpm(24)+fpm(25)-1; fpm(57) will the values 1...p+1
26
   ..... <<< user entry
27
   case(31) !!Perform multiplication A^H[fpm(57)] * X(1:N,i:j) result in work1(1:N,i:j)
28
    !! where i = fpm(34) and j = fpm(34) + fpm(35) - 1; fpm(57) will the values 1 \dots p+1
29
30 ..... <<< user entry
31 end select
32 end do
33
```

3.5 IFEAST (FEAST w/o Factorization)

IFEAST stands for using FEAST using inexact iterative solver for solving the linear systems (instead of using a direct solver). IFEAST only supports the sparse driver interfaces where the MKL-PARDISO solver is then replaced by a built-in BiCGstab solver. IFEAST is particularly effective if the sparse system matrix is very large (typically >1M) and/or the direct factorization becomes too expensive (both in time and memory). Two options are possible for calling IFEAST:

Option1 In the naming convention of all FEAST sparse drivers, replace feast by ifeast.

Option2 Keep the name of your FEAST sparse driver unchanged but use the new flag value fpm(43)=1.

As an example, let us re-work the helloworld example presented in Section 2.4. Here are the few lines that need to be changed in the code in order to make use of IFEAST:

```
!!!! the matrix A needs first to be defined in sparse CSR format;
  2 !!!! add these lines in the variable declaration section of the program
  _3 integer, parameter :: NNZ
  4 double precision, dimension (NNZ) :: sA = (/2.0d0, -1.0d0, -1.0d0, -1.0d0, 3.0d0, -1.0d0, 
                                                                                                                          \&-1.0d0, -1.0d0, 3.0d0, -1.0d0, -1.0d0, -1.0d0, 2.0d0/)
 6 integer, dimension (N+1) :: IA = (/1, 4, 8, 12, 15/)
       integer, dimension (NNZ) :: JA = (/1, 2, 3, 1, 2, 3, 4, 1, 2, 3, 4, 2, 3, 4/)
10 !!!! The direct call to IFEAST can be done as follow (option 1)
11 call feastinit (fpm)
12 fpm(1)=1 !! change from default value (print info on screen)
13 call difeast_scsrev(UPLO, N, sA, IA, JA, fpm, epsout, loop, Emin, Emax, M0, E, X, M, res, info)
14
15 !!!! Alternatively, an indirect call to IFEAST is also possible (option 2)
16 call feastinit (fpm)
17 fpm(1)=1 !! change from default value (print info on screen)
18 fpm(43)=1 !! switch solver in FEAST sparse driver from MKL-PARDISO to BicGStab
19 call dfeast_scsrev (UPLO, N, sA, IA, JA, fpm, epsout, loop, Emin, Emax, MO, E, X, M, res, info)
```

Here is the output of the run:

```
****** FEAST v4.0 BEGIN **********
*************
Routine DIFEAST_SCSREV
Solving AX=eX with A real symmetric
List of input parameters fpm(1:64)-- if different from default
  fpm(1)=
| FEAST data
                     3.000000000000000E+00
| Emin
 Emax
                     5.00000000000000E+00
 #Contour nodes
                     4 (half-contour)
 Quadrature rule
                     Trapezoidal
 Ellipse ratio y/x |
                     1.00
                     BiCGstab
 System solver
                     eps=1E-1; maxit= 40
                     Single precision
                     Matrix scaled
 FEAST uses MKL?
                     Yes
 Fact. stored?
                     Yes
 Initial Guess
                     Random
 Size system
                        4
| Size subspace
                         3
| FEAST runs
                  -1
```

```
#It |
      #Eig |
                       Trace
                                              Error-Trace
                                                                         Max-Residual
                       1.0501874385226984E-05; res max=
 #it
                                                          3.1005099299363792E-04
         4; res min=
                       1.8933478742837906E-02; res max=
 #it
         3; res min=
                                                          9.8125219345092773E-02
 #it
         2; res min=
                       2.9427373781800270E-02; res max=
                                                          8.7035216391086578E-02
 #it
                       1.8594998866319656E-02; res max=
                                                          4.0714181959629059E-02
         2; res min=
        2
              7.9995148090616173E+00
                                           1.00000000000000E+00
                                                                       8.7043125405406700E-03
 0
 #it
         2; res min=
                       2.8199069201946259E-02; res max=
                                                          5.2359659224748611E-02
 #it
         2: res min=
                       6.4112566411495209E-02; res max=
                                                          8.7011836469173431E-02
 #it
         2; res min=
                       3.0766267329454422E-02; res max=
                                                          5.8814667165279388E-02
 #it
                       1.4842565171420574E-02; res max=
                                                          1.5974638983607292E-02
         2; res min=
 1
               7.9999999809553799E+00
                                           9.7034378752525186E-05
                                                                       5.5062636929826442E-05
 #it
         1; res min=
                       5.3539618849754333E-02; res max=
                                                          5.3563684225082397E-02
                                                          8.7261073291301727E-02
 #it
         1; res min=
                       8.7251774966716766E-02; res max=
                       7.2112381458282471E-02; res max=
                                                          7.2119705379009247E-02
 #it
         1; res min=
 #it
         1; res min=
                       6.5362729132175446E-02; res max=
                                                          6.5378636121749878E-02
               7.99999999999591E+00
                                           3.8089158493903597E-09
                                                                       7.6857681986636782E-08
 #it
                       3.4910961985588074E-02; res max=
                                                          3.4911289811134338E-02
         1: res min=
 #it
                       8.4501713514328003E-02; res max=
                                                          8.4501914680004120E-02
         1; res min=
 #it
                       6.9185368716716766E-02: res max=
                                                          6.9185607135295868E-02
         1: res min=
                                                          6.2304504215717316E-02
 #it
         1; res min=
                       6.2304046005010605E-02; res max=
              8.000000000000036E+00
                                           8.8817841970012523E-15
                                                                       2.5198258072819870E-12
 3
                       3.4680232405662537E-02; res max=
 #it
         1; res min=
                                                          3.4876041114330292E-02
 #it
         1: res min=
                       8.4406383335590363E-02; res max=
                                                          8.4497839212417603E-02
                       6.9181196391582489E-02; res max=
 #it
         1; res min=
                                                          6.9283291697502136E-02
         1; res min=
 #it
                       6.2299706041812897E-02; res max=
                                                          6.2503643333911896E-02
 4
               8.00000000000000E+00
                                           7.1054273576010023E-16
                                                                       3.0767402982137245E-16
==>FEAST has successfully converged with Residual tolerance <1E-12
  # FEAST outside it.
                             4
  # Inner BiCGstab it.
                            31
  # Eigenvalue found
                             2 from
                                      3.99999999999991E+00 to
                                                                4.0000000000000009E+00
| FEAST-RCI timing |
| Fact. cases(10.20)|
                          0.0000
 Solve cases(11,12)|
                          0.0189
       cases(30,31)|
                          0.0000
I B*x
       cases(40,41)|
                          0.0001
 Misc. time
                          0.0005
| Total time (s)
                          0.0195
****** FEAST- END**************
*****************
Solutions (Eigenvalues/Eigenvectors/Residuals)
     4.000000000000000
                                                      -0.853553390593266
                          X= 0.353553390593291
E=
                         0.353553390593291
                                                Res= 3.076740298213724E-016
 0.146446609406686
     4.00000000000000
                           X= 0.353553390593257
                                                       0.146446609406767
 -0.853553390593281
                         0.353553390593257
                                                Res= 2.361858070262224E-016
```

Some Remarks:

- IFEAST is using different default fpm parameters than FEAST. In particular, the trapezoidal rule is used along a half-circle with 4 contour integration points, as well as the BiCGStab iterative solver.
- Each FEAST loop reports the total number of BiCGstab iterations for each linear system solve needed to reach the accuracy defined in fpm(45). The total number of BiCGtab iterations will be reported in fpm(60) (here 31, of course IFEAST is rather ineffective for such small system).
- You may have noticed that the values of the output eigenvectors for this example are different than the ones reported in Section 2.4. As a reminder, eigenvectors are not unique, both solutions are here orthonormal and correct (they span the same eigenvector subspace).

3.6 PFEAST and PIFEAST (MPI-solver)

In FEAST v4.0, the FEAST-MPI library enables the use of MPI linear system solvers at L3. As a result, the three level of parallelisms of FEAST (L1-L2-L3) can all support MPI (FEAST is internally using three MPI communicators). This MPI-MPI-MPI programming model is named PFEAST. PFEAST currently supports all the sparse drivers (for Hermitian/non-Hermitian/Polynomial problems) as well as all the RCI interfaces. Using RCI, an expert developer could straightforwardly customize PFEAST using highly-efficient application-specific MPI solvers such as: domain decompositions, or iterative/hybrid solvers with/without preconditioners.

Here some information about the use of PFEAST:

Initialization The feastinit(fpm) routine must be replaced by pfeastinit(fpm, L1_comm_world, nL3) which, in addition of setting up default fpm values, is going to initialize all MPI communicators.

- L1_comm_world represents your own defined MPI communicator for a given search contour (containing nL1 total MPI processes), if only a single contour is used it must take the value MPI_COMM_WORLD.
- nL3 is an (in/out) integer input that indicates the number of MPI processes you wish to use at level L3 (MPI system solver). The value of nL3 will be reassigned to nL1 if nL1 is not a multiple of nL3. In addition, if nL1< nL3 then nL3= nL1 on exit. Ideally, nL1/nL3 should be a multiple of a number of contour points (if it is equal to the number of contour points, then L2 is optimally used). Furthermore, L3 can also be threaded (MPI calling OpenMP on each local distributed system), make sure that your number of selected threads <omp> times nL1 does not exceed the number of available physical cores of your cluster.
- This initialization routine is setting up the L3 communicator fpm(49) (n particular) that may be needed to distribute your matrix.

Sparse Drivers In the naming convention of the FEAST sparse drivers, all routine names must be preceded by the letter p. In particular, you must replace z{i}feast by pz{i}feast, or d{i}feast by pd{i}feast. We actually use the names PFEAST and PIFEAST to indicate the MPI version of the FEAST and IFEAST interfaces, respectively. PFEAST is using MKL-Cluster-PARDISO and PIFEAST is using a built-in MPI-BiCGstab iterative solver (PBiCGStab). PIFEAST is also using its own built-in highly efficient MPI sparse mat-vec library.

PFEAST/PIFEAST drivers allow two options for the row distribution of matrices and solution vectors:

- global and common to all L3-MPI processes (The row-distribution will then take place internally). The argument list for all interfaces stay unchanged.
- locally row distributed among all L3-MPI processes. The user is responsible for distributing the data using the fpm(49) L3 communicator. The argument list for all interfaces stay mostly unchanged but the size of matrix/vectors N which must now be local.

Furthermore, PFEAST/PIFEAST will automatically detect which option above you are using!

RCI Interfaces The names of the RCI interfaces do not change (the letter p is not needed). In the argument list, only the size N must be changed to its local value (i.e. local number of rows for the row-distributed vector and work arrays).

PFEAST with global L3 distribution: Examples

The \$FEASTROOT/example/PFEAST-L2L3 directory provides Fortran and C implementation of System1 to System5 examples (discussed previously). Here is the complete list of the PFEAST routines that are all using global sparse CSR-storage.

	Type	Routine
System1	Real Sym. Generalized	P{F90,C}sparse_pdfeast_scsrgv
System2	Complex Herm. Standard	P{F90,C}sparse_pzfeast_hcsrev
System3	Real non-Sym. Generalized	P{F90,C}sparse_pdfeast_gcsrgv
System4	Complex Sym. Standard	P{F90,C}sparse_pzfeast_scsrev
System5	Real Sym. Quadratic	P{F90,C}sparse_pdfeast_scsrpev

PFEAST with local L3 distribution: Example

As an example, let us re-work the helloworld example presented in Section 2.4 using 2 MPI processes to distributed the 4×4 matrix. We obtain (for example):

$$\mathbf{A} = \begin{pmatrix} 2 & -1 & -1 & 0 \\ -1 & 3 & -1 & -1 \\ -1 & -1 & 3 & -1 \\ 0 & -1 & -1 & 2 \end{pmatrix} = \begin{pmatrix} \mathbf{A_1} \\ \mathbf{A_2} \end{pmatrix}$$

A Fortran90 source code example is provided below:

```
program helloworld_pfeast_local
     implicit none
     include 'mpif.h'
     !! 4x4 global eigenvalue system == two 2x4 local matrices
     integer, parameter :: Nloc=2, NNZloc=7
     character(len=1) :: UPLO='F
     double precision, dimension(NNZloc) ::
     integer, dimension(Nloc+1) :: IAloc
    integer, dimension (NNZloc) :: JAloc
    !! input parameters for FEAST
10
     integer, dimension (64) :: fpm
11
     integer :: M0=3 ! search subspace dimension
12
     double precision :: Emin=3.0d0, Emax=5.0d0 ! search interval
13
     !! output variables for FEAST
14
     double precision, dimension(:), allocatable :: E, res
1.5
     double precision, dimension (:,:), allocatable :: X
     double precision :: epsout
17
     integer :: nL3, rank3, loop, info, M, i
19 !!! MPI
     integer :: code
20
     call MPI_INIT(code)
21
23 !!! Allocate memory for eigenvalues.eigenvectors, residual
     allocate(E(M0), X(Nloc, M0), res(M0))
24
25
26 !!!!!!!!! INITIALIZE PFEAST and DISTRIBUTE MATRIX
27
     {\tt call pfeastinit} \ ({\tt fpm} \ , \! M\! P\! I\! \_\! C\! O\! M\! M\! \_\! W\! O\! R\! L\! D\! , nL3)
28
29
     call MPI_COMM_RANK(fpm(49), rank3, code) !! find rank of new L3 communicator
30
     if (rank3==0) then
31
       Aloc = (/2.0d0, -1.0d0, -1.0d0, -1.0d0, 3.0d0, -1.0d0, -1.0d0/)
32
33
       IAloc = (/1, 4, 8/)
       JAloc = (/1, 2, 3, 1, 2, 3, 4/)
34
     elseif (rank3==1) then
35
       Aloc = (/-1.0d0, -1.0d0, 3.0d0, -1.0d0, -1.0d0, -1.0d0, 2.0d0/)
36
       IAloc = (/1, 5, 8/)
37
       JAloc = (/1, 2, 3, 4, 2, 3, 4/)
38
     endif
39
41 !!!!!!!!! PFEAST
     fpm(1)=1 !! change from default value (print info on screen)
42
     call pdfeast_scsrev (UPLO, Nloc, Aloc, IAloc, JAloc, fpm, epsout, loop, Emin, Emax, M0, E, X, M, res, info)
43
44
45 !!!!!!!!! REPORT
   if (info==0) then
46
47
        print *, 'Solutions (Eigenvalues/Eigenvectors/Residuals) at rank L3', rank3
        do i = 1,M
48
           print *,'Eigenvalue',i
print *,'E=',E(i),'X=',X(:,i),'Res=',res(i)
49
50
        enddo
51
     endif
52
54 end program helloworld_pfeast_local
```

Your program must be compiled using the same MPI implementation used to compile the FEAST-MPI library. Once compiled, your source program must now be linked with the pfeast library. You can use (for example):

```
mpiifort -o helloworld_pfeast_local helloworld_pfeast_local.f90

-L$FEASTROOT>/lib/<arch> -lpfeast -mkl=parallel -lmkl_blacs_intelmpi_lp64 -liomp5
-lpthread -lm -ldl
```

if FEAST was compiled with ifort, MKL flag was set to 'yes', and MPI was chosen to be 'impi' (intel mpi).

```
mpif90.mpich -fc=gfortran -o helloworld_pfeast_local helloworld_pfeast_local.f90
-L$FEASTROOT>/lib/<arch> -lpfeast -Wl,-no-as-needed -lmkl_gf_lp64 -lmkl_gnu_thread
-lmkl_core -lmkl_blacs_intelmpi_lp64 -lgomp -lpthread -lm -ldl
if FEAST was compiled with gfortran, MKL flag was set to 'yes', and MPI was chosen to be 'mpich'.
```

A run of the resulting executable looks like:

and the output of the run should be:

```
**************
****** FEAST v4.0 BEGIN ***********
**************
Routine PDFEAST_SCSREV
Solving AX=eX with A real symmetric
#MPI (total=L2*L3) 2= 1* 2
List of input parameters fpm(1:64)-- if different from default
  fpm(1)=
| FEAST data |
          | 3.00000000000000E+00
| Emin
 Emax
                | 5.00000000000000E+00
| #Contour nodes | 8 (half-contour)
| Quadrature rule | Gauss
 Ellipse ratio y/x | 0.30
                | MKL-Cluster-Pardiso
 System solver
                 | Single precision
                   Matrix scaled
 FEAST uses MKL?
                  Yes
 Fact. stored?
                   Yes
| Initial Guess
                   Random
 Size system
                       4
| Size subspace
                       3
| FEAST runs
                                                         | Max-Residual
#It | #Eig |
                                  | Error-Trace
                    Trace
 0
       2
             7.9999999999947E+00 1.0000000000000E+00 1.2829791863202860E-08
             8.00000000000000E+00
                                  1.0658141036401502E-15 6.6022321739723205E-16
==>FEAST has successfully converged with Residual tolerance <1E-12
  # FEAST outside it.
                         1
                         2 from 3.99999999999996E+00 to 4.0000000000000E+00
  # Eigenvalue found
| FEAST-RCI timing |
| Fact. cases(10,20)| 0.0058
| Solve cases(11,12)|
                      0.0025
| A*x cases(30,31)|
                       0.0000
```

```
0.0000
       cases(40,41)|
 B*x
 Misc. time
                          0.0005
| Total time (s)
                          0.0088
****** FEAST- END***************
Solutions (Eigenvalues/Eigenvectors/Residuals) at rank L3
Eigenvalue
                     1
                                                      -0.851189974005894
     4.000000000000000
                           X= 0.358971274554087
                                                                                    2.784560286672102E-016
Eigenvalue
                     2
     4.00000000000000
                           X= -0.348051180209196
                                                      -0.159610864767551
                                                                                    6.602232173972321E-016
Solutions (Eigenvalues/Eigenvectors/Residuals) at rank L3
Eigenvalue
                     1
                           X= 0.133247424897719
                                                       0.358971274554087
                                                                                    2.784560286672102E-016
     4.00000000000000
                                                                              Res=
Eigenvalue
                     2
     4.000000000000000
                           X= 0.855713225185942
                                                      -0.348051180209196
                                                                                   6.602232173972321E-016
```

PFEAST using 3 levels of parallelism: Example

Three levels of parallelism means that L1 is active and multiple search contours can be used simultaneously. FEAST v4.0 does not offer automatic partitioning of the overall eigenvalue spectrum, it is then up to the users to guess it. Users could take advantage of fast stochastic estimates with the flag fpm(14)=2. Once the eigenvalue spectrum partitioned, a single call to PFEAST will account for all L1-L2-L3 MPI parallelism.

The \$FEASTROOT/example/PFEAST-L1L2L3 directory provides Fortran and C implementation of the System2 example (discussed previously). It uses two search intervals. The name of the routines are:

System2					
sparse	3P{F90,C}dense_pzfeast_hcsrev				

Remark: Since multiple search intervals are involved, the option fpm(1)=1 (printing FEAST info on screen), may provide a bit confusing results to read. You can easily change this flag value using fpm(1)=-i with i is associated with the rank i-1 of the L1 MPI Communicator. Each search contour will then print all its FEAST results into separate files named feast{i}.log.

4 Complement

4.1 Matrix storage

Let us consider the following matrix A (as an example):

$$\mathbf{A} = \begin{pmatrix} a_{11} & a_{12} & 0 & 0 \\ a_{21} & a_{22} & a_{23} & 0 \\ 0 & a_{32} & a_{33} & a_{34} \\ 0 & 0 & a_{43} & a_{44} \end{pmatrix} \tag{2}$$

If the matrix presents some particular properties such as Hermitian $(a_{ij} = a_{ji}^*)$ for $i \neq j$ or symmetric $(a_{ij} = a_{ji})$, only half of the matrix elements need to be defined. Using the FEAST Driver interfaces, this matrix could be stored in dense, banded or sparse-CSR format as follows:

Dense The matrix is stored in a two dimensional array in a straightforward fashion. Using the options UPLO='L' or UPLO='U', the lower triangular or upper triangular part respectively, do not need to be referenced.

Banded The matrix is also stored in a two dimensional array following the banded LAPACK-type storage:

$$\mathbf{A} = \begin{pmatrix} * & a_{12} & a_{23} & a_{34} \\ a_{11} & a_{22} & a_{33} & a_{44} \\ a_{21} & a_{32} & a_{43} & * \end{pmatrix}$$

In contrast to LAPACK, no extra-storage space is necessary since LDA>=kl+ku+1 if UPLO='F' (LAPACK banded storage would require LDA>=2*kl+ku+1). For this example, the number of subdiagonals kl and superdiagonals is ku are both equal to 1. Using the option UPLO='L' or UPLO='U', the rows respectively above or below the diagonal elements row, do not need to be referenced (or stored).

Sparse-CSR The non-zero elements of the matrix are stored using a set of one dimensional arrays (A,IA,JA) following the definition of the CSR (Compressed Sparse Row) format

$$\mathbf{A} = (a_{11}, a_{12}, a_{21}, a_{22}, a_{23}, a_{32}, a_{33}, a_{34}, a_{43}, a_{44})$$

$$\mathbf{IA} = (1, 3, 6, 9, 11)$$

$$\mathbf{JA} = (1, 2, 1, 2, 3, 2, 3, 4, 3, 4)$$

Using the option UPLO='L' or UPLO='U', one would get respectively

$$\mathbf{A} = (a_{11}, a_{21}, a_{22}, a_{32}, a_{33}, a_{43}, a_{44})$$
 $\mathbf{IA} = (1, 2, 4, 6, 8)$
 $\mathbf{JA} = (1, 1, 2, 2, 3, 3, 4)$
 $\mathbf{A} = (a_{11}, a_{12}, a_{22}, a_{23}, a_{33}, a_{34}, a_{44})$
 $\mathbf{IA} = (1, 3, 5, 7, 8)$
 $\mathbf{JA} = (1, 2, 2, 3, 3, 4, 4)$

4.2 Search contour

Figure 2 summarizes the different search contour options possible for both the Hermitian and non-Hermitian (including Polynomial) FEAST algorithms.

For the Hermitian case, the user must then specify a 1-dimensional real-valued search interval $[E_{min}, E_{max}]$. These two points are used to define a circular or ellipsoid contour \mathcal{C} centered on the real axis, and along which the complex integration nodes are generated. The choice of a particular quadrature rule will lead to a different set of relative positions for the nodes and associated quadrature weights. Since the eigenvalues are real, it is convenient to select a symmetric contour with the real axis $(\mathcal{C} = \mathcal{C}^*)$ since it only requires to operate the quadrature on the half-contour (e.g. upper half).

With a non-Hermitian/Polynomial problem, it is necessary to specify a 2-dimensional search interval that surrounds the wanted complex eigenvalues. Circular or ellipsoid contours can also be used and they can be generated using standard options included into FEAST v4.0. These are defined by a complex midpoint E_{mid}

and a radius r for a circle (for an ellipse the ratio between the horizontal axis 2r and vertical axis can also be specified, as well as an angle of rotation). A "Custom Contour" feature is also supported that can use arbitrary quadrature nodes and weights (provided by the users).

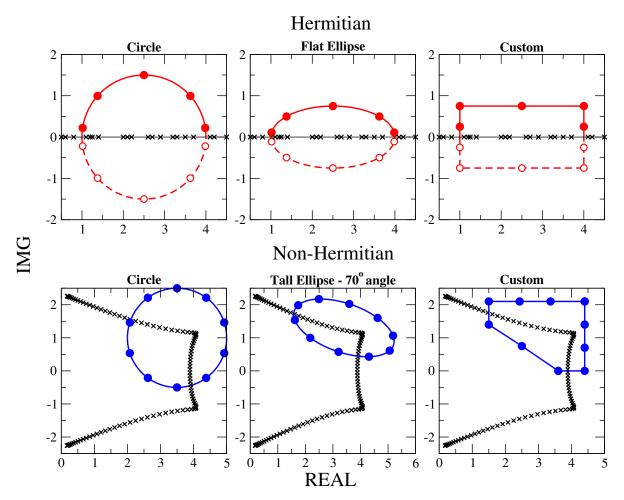


Figure 2: Various search contour examples for the Hermitian and the non-Hermitian/Polynomial FEAST algorithms. Both algorithms feature standard ellipsoid contour options and the possibility to define custom arbitrary shapes. In the Hermitian case, the contour is symmetric with the real axis and only the nodes in the upper-half may be generated. In the non-Hermitian/Polynomial case, a full contour is needed to enclose the wanted complex eigenvalues. Some data used to generate these plots:

Hermitian case: fpm(2)=5 for all, $[E_{min}, E_{max}] = [1, 4]$, r = 1.5 for all; fpm(18)=50 for the flat ellipse; and expert routine for the custom contour

Non-Hermitian/Polynomial case: fpm(8)=10 for all; $E_{mid} = 3.5 + i$ and r = 1.5 for circle; $E_{mid} = 3.4 + 1.3i$, r = 0.75, fpm(18)=200, fpm(19)=70 for tall rotated ellipse; and expert routine for the custom contour

The FEAST package provides a couple of utility routines that can return the integration nodes and weight used by the FEAST interfaces:

zfeast_contour(Emin,Emax,fpm2,fpm16,fpm18,Zne,Wne)

Returns FEAST integration nodes and weights for a half-contour contour defined by Emin and Emax. To be used with Hermitian FEAST interfaces.

zfeast_gcontour(Emid,r,fpm8,fpm16,fpm18,fpm19,Zne,Wne)

- Returns FEAST integration nodes and weights for a full contour defined by Emid and r. To be used with non-Hermitian and polynomial FEAST interfaces.

The description of the arguments list for these routines is given in Table 9 and Table 10.

	Type	I/O	Description
Emin, Emax	double real	in	Lower and Upper bounds of search interval
fpm2	integer	in	Value of fpm(2)- #contour point (half-contour)
fpm16	integer	in	Value of fpm(16)- Integration type
fpm18	integer	in	Value of fpm(18)- Ellipse definition
Zne	double complex	out	Integration nodes
Wne	double complex	out	Integration weights

Table 9: List of arguments for zfeast_contour.

	Type	I/O	Description
Emid	double complex	in	Coordinate center of the contour ellipse
r	double real	in	Horizontal radius of the contour ellipse
fpm8	integer	in	Value of fpm(8)- #contour point (full-contour)
fpm16	integer	in	Value of fpm(16)- Integration type
fpm18	integer	in	Value of fpm(18)- Ellipse definition
fpm19	integer	in	Value of fpm(19)- Ellipse rotation angle
Zne	double complex	out	Integration nodes
Wne	double complex	out	Integration weights

Table 10: List of arguments for zfeast_gcontour.

4.3 Contour Customization

The Custom Contour feature grants the flexibility to target specific eigenvalues in a complex plane. This feature must be used with "Expert" routines that take two additional arguments containing the complex integration nodes and weights. Custom contours can be employed by following three simple steps:

- 1. Define a contour (half-contour that encloses $[\lambda_{min}, \lambda_{max}]$ for the Hermitian problem, or full contour for the non-Hermitian/Polynomial problem),
- 2. Calculate corresponding integration nodes and weights, and
- 3. Call "Expert" FEAST routine (either Driver or RCI interfaces by adding a x at the end of the routine name).

Furthermore, the FEAST package provides a utility routine zfeast_customcontour that can assist the user to extract nodes and weights from a custom design arbitrary geometry in the complex plane (full-contour). Users must only define the geometry of their contour. The contour can be comprised of line segments and half ellipses. Two important points to note: (i) the actual contour will end up being a polygon defined by the integration points along the path, and (ii) only convex contours may be used. A geometry that contains P contour parts/pieces is defined using three arrays Zedge, Tedge, and Nedge. The interface is defined below and the description of the arguments list is given in Table 11.

zfeast_customcontour(Nc,P,Nedge,Tedge,Zedge,Zne,Wne)

As an example, the following code will generate the corresponding complex contour.

	Type	I/O	Description		
Nc	integer	in	The total number of integration nodes, should be equal to		
			SUM(Nedge(1:P))		
P	integer	in	Number of contour parts/pieces that make up the contour		
Zedge	integer(P)	in	Complex endpoints of each contour piece		
			Remark: * endpoints positioned in clockwise direction		
			* the k^{th} piece is $[{ t Zedge}(k)$, ${ t Zedge}(k+1)]$		
			* last piece is [Zedge(P), Zedge(1)]		
Tedge	integer(P)	in	The type of each contour piece:		
			*If Tedge(k)=0, k^{th} piece is a line		
			*If Tedge(k)>0, k^{th} piece is a (convex) half-ellipse		
			with Tedge(k)/100 = ratio a/b and a primary radius from the endpoints		
			Remark: 100 is a half-circle		
Nedge	integer(P)	in	#integration intervals to consider for each piece		
			define the accuracy of the trapezoidal rule by piece for FEAST		
Zne	double complex	out	Custom integration nodes for FEAST		
Wne	double complex	out	Custom integration weights for FEAST		

Table 11: List of arguments for zfeast_customcontour.

The \$FEASTROOT/example/FEAST directory provides Fortran and C implementation of the expert FEAST routines using a custom contour. It is applied on the System4 example using both dense, banded and sparse-CSR storage. Here, the complete list of routines:

	System4
dense	
	{F90,C}dense_zfeast_syevx
banded	
	{F90,C}dense_zfeast_sbevx
sparse	
	{F90,C}dense_zfeast_scsrevx

4.4 FEAST utility sparse drivers

If a sparse matrix can be provided by the user in coordinate/matrix market format, the \$FEASTROOT/utility directory offers a quick way to test all the FEAST parameter options and the efficiency/reliability/timing of the FEAST SPARSE driver interfaces. Two general drivers are provided for FEAST/IFEAST and PFEAST/PIFEAST, named driver_feast_sparse or driver_pfeast_sparse in their respective subdirectories. The command ">make all" should compile the drivers.

If we denote mytest a generic name for the user's eigenvalue system test $\mathbf{A}\mathbf{x} = \lambda \mathbf{x}$ or $\mathbf{A}\mathbf{x} = \lambda \mathbf{B}\mathbf{x}$. You will need to create the following three files:

• mytest.mtx should contain the matrix A in coordinate format; As a reminder, the coordinate format is defined row by row as

```
N N NNZ
: : : : :
i j real(valj) img(valj)
: : : : :
: : : : :
i iNNZ jNNZ real(valNNZ) img(valNNZ)
```

with N: size of matrix, and NNZ: number of non-zero elements.

- mytestB.mtx should contain the matrix B (if any) in coordinate format;
- mytest.in should contain the search interval, selected FEAST parameters, etc. The following .in file is given as a template example (here for solving a standard eigenvalue problem in double precision):

```
! s: symmetric, h: hermitian, g: general
s
        ! e=standard or g=generalized eigenvalue problem
g
          (d,z) precision i.e (double real, double complex)
d
        ! UPLO (L: lower, U: upper, F: full) for the coordinate format of matrices
F
0.18d0
        ! Emin
1.00d0
         Emax
25
        ! MO search subspace (MO>=M)
        !!!!!!!!! How many changes from default fpm(1,64) (use 1-64 indexing)
1 1
        !fpm(1)=1 !example comments on/off (0,1)
        !fpm(2)=4 !number of contour points
```

You may change any of the above options to fit your needs. For example, you could add as many fpm FEAST parameters as you wish. You can also use the flag fpm(43)=1 to switch to IFEAST. In addition, the L or U options for UPLO give you the possibility to provide only the lower or upper triangular part of the matrices mytest.mtx and mytestB.mtx in coordinate format.

Finally results and timing can be obtained by running the FEAST sparse driver:

```
./driver_feast_sparse <PATH_TO_MYTEST>/mytest
```

In addition, all the eigenvalue solutions will be locally saved in the file eig.out.

For the PFEAST sparse drivers, a run would look like (other options could be applied):

where < nL1 > represents the total number of MPI processes to use for a single contour, and nL3 is the number of MPI processes used for solving the linear systems. As a reminder, L3 uses MKL-Cluster-PARDISO with PFEAST, and PBicGStab with PIFEAST. The level L3 can also be threaded by setting the shell variable MKL_NUM_THREADS equal to the desired number of threads. Make sure that <omp>*<nl1> does not exceed the number of physical cores. Several combinations of <nl1>, <nl3> and <omp> are possible depending also on the value of the -ppn directive.

In order to illustrate a direct use of the utility drivers, several examples are provided in the directory \$FEASTROOT/utility/data summarized in Table 12.

	Real	Complex	Symmetric	Hermitian	General	Standard	Generalized
helloworld	X		X			X	
system1	X		X				X
system2		X		X		X	
system3	X				X		X
system4		X	X			X	
cnt	X		X				X
со	X		X				X
c6h6	X		X				X
Na5	X		X			X	
grcar	X				X	X	
qc324		X	X			X	
bcsstk11	X		X				X

Table 12: List of system matrices provided in the \$FEASTROOT/utility/data directory. System 1 to 4 corresponds to the matrices used in the example directory.

To run a specific test, you can execute using the FEAST driver (for example):

or using the PFEAST driver (for example):

Here we use 8 total compute cores and nL1=4 MPI, nL2=nL1/nL3=2 MPI, nL3=2 MPI, omp=2 threads.