

算法 2 数据结构的高效实现

```
1: procedure INITIALIZE-TREE
2:   INITIALIZE-TREE-SIMPLE
3:   build heavy-light decomposition
4:   calculate  $g_u$  ( $\forall u$ )
5:   for each chain  $c$  do
6:     let  $c.segment-tree$  be a new segment tree of  $g_u$ 
7:     initialize  $c.segment-tree$ 
8:   end for
9: end procedure
10:
11: procedure MODIFY-CHILD-VALUE( $u, v, b$ )
12:   assume  $v_1, \dots, v_k$  are children of  $u$ 
13:   assume  $v = v_i$ , and  $v_j$  is the heavy-child of  $u$  ( $i \neq j$ )
14:    $u.processor.MODIFY(i, b)$ 
15:    $u.processor.MODIFY(j, 0)$ 
16:    $g_u(0) \leftarrow u.processor.QUERY$ 
17:    $u.processor.MODIFY(j, 1)$ 
18:    $g_u(1) \leftarrow u.processor.QUERY$ 
19: end procedure
20:
21: procedure RECALCULATE( $u$ )
22:   assume  $u$  is the  $i^{\text{th}}$  node of chain  $c$ 
23:    $r \leftarrow c.top$ 
24:    $c.segment-tree.MODIFY(i, g_u)$ 
25:    $g^* \leftarrow c.segment-tree.QUERY$ 
26:    $W(r) \leftarrow g^*(0)$   $\triangleright$  此时  $g^*(0) = g^*(1)$ 
27:   if  $r \neq \text{root}$  then
28:      $p \leftarrow r.parent$ 
29:     MODIFY-CHILD-VALUE( $p, r, W(r)$ )
30:     RECALCULATE( $p$ )
31:   end if
32: end procedure
33:
34: procedure MODIFY-TREE( $u, b$ )
35:    $g_u(0) \leftarrow b, g_u(1) \leftarrow b$ 
36:   RECALCULATE( $u$ )
37: end procedure
38:
39: function QUERY-TREE
40:   return  $W(\text{root})$ 
41: end function
```