A peek under the Blue Coat

ProxySG internals

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- Introduction
- Storage: filesystems and registry
- Binaries
- Wernel and OS mechanisms
- Understanding internals
- 6 Security mechanisms
- Conclusion



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What? Why?

Blue Coat ProxySG?

- enterprise (Web) proxy
- one of the most deployed in big companies
- lots of complex features:
 - URL categorization (WebSense and others)
 - video streaming / instant messaging specific handling
 - MAPI and SMB proxy / cache / prefetcher
 - etc.
- runs proprietary SGOS

BlueCost

Why research ProxySG?

- widely used in Airbus Group
- interesting target for malicious actors: log bypass, Internet exposed, MITM, etc.
- no known previous research: unknown security level
 - security bulletins: mostly OpenSSL and Web administration interface bugs

Research

Study objectives:

- assess the global security level
- write recommendations for secure deployment
- be prepared for forensics in case of a compromised ProxySG

Why publish?

- first public info but surely not first research
- ullet foster research \Longrightarrow better security

Today's presentation:

- raw technical results, as a starting point for research
- goes from low level (FS) to high level, following our approach
- applies to all ProxySG models and 6.x versions up to Q1 2015



Getting started

Running ProxySG:

- hardware: commodity x86 CPUs, HDD, etc.
- VMware appliances

Common versions:

- 5.5: older version, EOL Aug 2014
- 6.2: previous *long term release*, EOL Oct 2015
- 6.5: latest long term release, recommended by BC

To get a first look, we need to access the filesystem:

- 6.? (≥ 6.4): small FAT32 partition containing proprietary BCFS image
- older versions: fully proprietary disk partitionning/data (no FAT32)



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On disk data: intro

Hardware

Basic architecture: 3 disks (or more)

- small CompactFlash or SSD for OS (FAT32)
- 2 or more drives for data (proprietary FS)

Filesystems

- static, read-only FS for OS (BCFS):
 - OS files
 - low level (static) configuration: kernel options, resource limits
- cache engine FS based on hash tables (CEFS) (Patent US7539818)
- registry in CEFS for settings

Remarks

- unknowns:
 - CEFS structures
 - log storage format
- on-disk partition structures are very complex
- today: only static FS (BCFS) for OS files



System disk organization (BIOS mode)

Files on FAT32 partition

```
/sgos/boot/systems/system1
/sgos/boot/cmpnts/starter.si
/sgos/boot/cmpnts/boot.exe
/sgos/boot/meta.txt
/sgos/fbr.con
```

Both starter.si and system1 use BCFS

bootloader: starter.si

- 6 MiB
- basic SGOS (UP kernel, drivers, no application)
- looks up available systems
- displays GRUB-like boot menu

Real OS: system1

- 210 MiB
- full blown OS:
 - SMP kernel
 - Web UI
 - actual applications
 - etc.

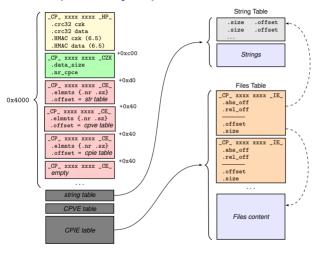


Boot sequence (BIOS)

- BIOS
- MBR
- boot sector of active partition
- boot.exe, found by hardcoded sector number
- kernel.exe, first file entry in starter.si FS
- 6 kernel starts sequencer.exe, second entry in starter.si
- sequencer.exe parses the main.cfg script and starts the necessary drivers
- main.cfg finally launches starter.exe which displays the boot menu
- starter.exe loads the selected system



BCFS (read-only FS) format



How to extract?

- read CPCE entries, note offsets for strings table and files table
- parse files table (CPIE) linearly
- get file name from strings table

How to modify?

- cannot increase file size
- 2 fix CRC and HMAC



System image configuration variables (CPVE)

- offset and size specified by 3rd _CP_ _CE_ entry
- modifying the variable implies fixing CRC/HMAC and reboot
- variable names can be found in sequencer.exe

Structure

```
struct cpve_entry {
    uint32_t magic1; /* _CP_ */
    uint64_t unk;
    uint32_t magic2; /* _VE_ */
    uint16_t number;
    uint16_t section;
    uint32_t unk2;
    uint64_t value; }
```

Known variables (section, number: description)

Section 4, kernel:

- 4,0: flags:
 - 0x8: GDB monitor enabled
 - 0x200: int3 at OS startup
 - 0x400: kernel debug logs enabled
- 4,1: arch_flags
 - 1: activate Write Protect in cr0
- 4,3: console_speed (in bauds)



Cache Engine FS (CEFS): writable storage

- hash-table object storage with disk backend
- mostly used for cache data:
 - web content
 - CIFS files
 - MAPI mails
 - etc.
- regular files are also supported, with prefix /legacy/cache_engine/

Some files (paths straight from the code, no typo)

- .../persistent/replicated/authorized_keys
- .../persistent/replicated/volatile//config/v9/registry/registry.xml
- .../transient//snmp.log
- .../persistent/replicated/licensing_certificate



Registry: settings storage

- tree structure used for all settings
- entries are referenced by strings like "config:Authenticator:local users"
- on-disk storage: xml file on writable CEFS

URLs (admin rights needed)

```
/registry/show
/registry/registry.html
/registry/registry.xml
/registry/debug
```

Interesting CLI extensions (cf slide 24)

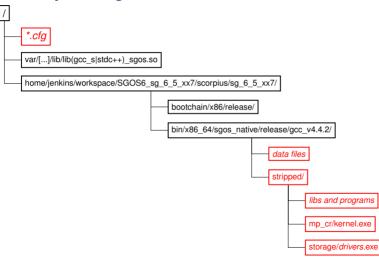
```
reg-set
reg-delete
reg-list
reg-trace
```



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OS Filesystem organization





ELF files: kernel, libs, programs

Everything interesting is located in .../stripped/:

- .exe, .exe.so and .so extensions (version 5 was using PE files)
- 32 or 64 bits ELF files, depending on model (RAM size?)
- everything in C++, compiled with g++ with custom sgos target
- lots of unit tests
- more than 2600 source files referenced
- everything is stripped, but lots of external symbols
- heavy template use: AMI::Config_Data::Config_Data::Config_Data::Storage_Class, AMI::Storage_PircaMI::Storage_PircaMI::Shared_PircaMI::Sha

"custom" ABI in 32 bits (probably gcc called with -mregparm):

EAX, EDX, ECX, stack

in 64 bits, standard SysV ABI:

RDI, RSI, RDX, RCX, R8, R9, stack



Known code?

Interesting open source libraries (version numbers from 6.5 release, Aug 2014):

- BGET: memory allocator (first dev in 1972!)
- NET-SNMP 5.4.2.1 (2008-10-31)
- newlib: libc
- expat 1.95.2: XML parser (2001!)
- libxml2 2.7.7-82143f4 (2010-11-04)
- OpenSSH 6.3 (2013-09-13)
- OpenSSL 1.0.1e (2013-02-11)
- zlib 1.2.3 (2005-07-18)

Blue Coat states that they backport fixes regularly (without necessarily changing the version string).



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Kernel

The kernel in practice

- kernel access partially abstracted in libknl_api.so
- small (~800 KiB), basic primitives:
 - interrupt/exception handling
 - semaphores/locks
 - message passing
 - drivers
- ds:1014h points to a "TEB"-like structure

Some syscalls

Nop

Suicide

Enable_event_logging

Register_worker_address

Symbol_address

Processor_voltage

Semaphore_signal_all

Grow_stack



Kernel: syscall

32 bits

- call dword ptr ds:1018h
- parameters in structure pointed by eax

kernel_req	struc		
field_0	dd	?	
return_code	dd	?	
return_code2	2 dd	?	
arg0	dd	?	
arg1	dd	?	
arg2	dd	?	
arg3	dd	?	
sys_num	dd	?	
kernel_req	ends		

64 bits

- call [ds:0FFFFFF8000000020h]
- parameters in structure pointed by rdi

knl_req64	struc
field_0	dq ?
retcode	dq ?
arg0	dq ?
arg1	dq ?
arg2	dq ?
arg3	dq ?
sysnum	dq ?
field_38	dq ?
knl_req64	ends

Memory organization

Back to the 90s

- protected mode
- everything in ring 0 (mentioned in US7539818 patent ;)

ELF mapping: at boot, once and for all

```
Unpacking executables...

Unpacking sequencer.exe elapsed time: 0s, 0ms, 326us

Unpacking ata.exe elapsed time: 0s, 0ms, 413us

[...]

Relocating executables...

Relocating sequencer.exe elapsed time: 0s, 2ms, 356us

Relocating ata.exe elapsed time: 0s, 0ms, 559us

10 executables relocated; total unpack and reloc time 0s, 20ms, 550us
```

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Making things easier: our tools

IDA plugins:

- automatically comment function with source filename (from debug logs)
- automatically rename functions from debug log strings
- automated syscall recognition (with syscall name, parameters)
- CLI structures dumper to list all CLI commands

BCFS: FUSE tool to mount system images:

- file access: read/write (without size change)
- CPVE access: read/write
- automatic CRC/HMAC calculation

Tools are internal PoCs and are not going to be released.



Getting more info: useful tricks

Enable debug info, by modifying BCFS (physical access or RCE needed):

- kernel "*printk*": CPVE 4,0 |= 0x400
- debug mode: set customer_release to 0 in main_cr.cfg

230+ CLI extensions in debug mode:

- list with ".", access with ".extension"
- examples: cfg, policy, cag, mgmt, etc.

Example commands

- .mgmt show-adv-urls
- .svc ashowstate
- .<ext> logaddmask all then .<ext> logshow
- .policy dbgtraceon



CLI extension example

```
ProxySG VA 1818181818>.cag
cag extension usage
COMMands:
  logshow : display contents of the debug log
  logaddmask : add a mask to the debug log
  logsubmask : remove a mask from the debug log
 log2console : toggling logging to the console logreset : reset the CAG debug log gzip-allow : determine if gzip allowed in responses
ProxvSG VA 1818181818>enable
Enable Password:
ProxySG VA 1818181818#_
```



GDB

Kernel includes GDB stub! But finding how to activate it took me weeks :(

- CPVE 4,0 |= 0x8
- multiplexed on COM1 with console
- send 0x18, 0x14 on COM port to activate
- (non-standard) text paging is handled server-side, patch client or use monitor util height 1000000

GDB monitor extensions (kernel side)

Current debug extensions:
name knl, Function 0x1261500
name util, Function 0x1028786E0
name scorpius, Function 0x1028487E0

Some knl extensions

- processes: display all active processes.
- pd: display the contents of a process descriptor.
- images: display details of loaded ELF files.



Practical understanding: HTTP parsing

Goal: find function for HTTP response parsing

- activate HTTP debug mode at https://x.x.x.x:8082/HTTP/debug
- make request through proxy
- get log
- read interesting function name
- look for function in libhttp.exe.so

Example log (simplified, most recent first):

```
HTTP CW 95B72F20: Parse_request called. beg=57DE3000 end=57DE30DC length=220
```

HTTP CW 95B72F20: Parse_request

HTTP CW 95B72F20: Should_tunnel_on_error

HTTP CW 95B72F20: Read_request

HTTP CW 95B72F20 POLICY: Evaluating PE_POLICY_CHECKPOINT_NEW_CONNECTION

HTTP CW 95B72F20: Transaction_startup

HTTP CW 95B72F20: Init_state

Going deeper: Hell

Locating the code is the easy part. Problems:

- HUGE functions (16 KiB!, see CFG)
- C++ everywhere
- IDA struggles with calling convention
- threads, everywhere!

Dynamic debugging howto:

- find image base using monitor knl image libhttp.exe.so in GDB
- relocate binary in IDA
- set breakpoint in Proxy SG CLI: conf t; debug; breakpoint-set 0 B X <ADDR>
- break and connect!





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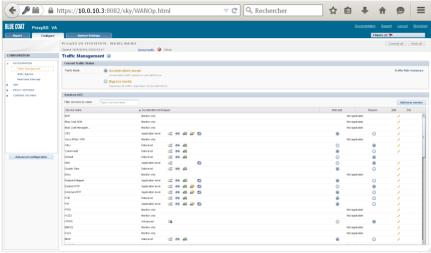


Application security

- authentication:
 - local passwords are hashed with FreeBSD MD5 crypt (\$1\$), Blowfish supported
 - dozens of schemes supported: LDAP, AD, etc.
- default protocols: only HTTPS and SSH
- · read-only or admin accounts
- OS trust:
 - PKCS7 signed updates (SHA-512/RSA-2048)
 - local images:
 - < 6.5: CRC only
 - ≥ 6.5: HMAC SHA-1
- crypto:
 - openssl
 - critical random data is generated securely



Administration interface (Flash)



Administration interface (Java)



Administration interface

- actually POSTs CLI commands, in an enable shell
- restricted commands for read-only users
- Java interface specifics:
 - also uses a kind for RPC mechanism (/Secure/Local/console/pod)
 - also implements its own HTTPS "client"

Request (simplified)

show version

----7d518638300904-

Response data

ProxySG VA 1818181818#(config)show version Version: SGOS 6.4.1.2 MACH5 Edition

Release id: 90192

UI Version: 6.4.1.2 Build: 90192 Serial number: XXXXXXXXXX

NIC 0 MAC: 000FF9B6006F

There were 0 errors and 0 warnings $\,$



System-level security

BAD

- no stack canaries
- no ASLR
- everything in ring0
- kernel callgate at a fixed address

GOOD

- NX enabled on most platforms since 6.2. 300/600 support added in 6.5.7.1 and 6.2.16.3
- BGET heap: asserts check for meta-data coherence (unlink attacks impossible)
- read-only FS for binaries makes it (way) harder to backdoor OS
- physical access (or code exec) is needed to change system image as updates are signed



Exploitability

Facilitating exploits

- previous slide :)
- vtables everywhere
- only C++ code => more memory corruption bugs (vs script/safer languages)

Hurdles

- no second chance (ring0)
- no ASLR but mapping different for each version
- custom payload needed
- guard pages



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Conclusion

Findings

- unusual, entirely proprietary OS design
- no user/kernel isolation or exploit hardening (historical for performance?)
- no vulnerabilities found (I didn't look for them!)...
- ...but Blue Coat release notes document plenty of fixes for "software restarts"

Recommendations

- use a dedicated (V)LAN for administration
- monitor the event log
- investigate reboots
- physically protect appliances
- use secure passwords (of course!)



Evolutions in ProxySG

Security enhancements in recent versions

- NX support for 300/600 added in 6.5.7.1 and 6.2.16.3
- bootchain and system image validation (hashes published by Blue Coat)
- Secure boot in pre-release, available in a future release
- debug (GDB, CLI extensions) support removed

We are currently discussing further security enhancements, such as user/supervisor separation, with Blue Coat. Release timing and platform support are still under discussion.



End

Questions?

Thanks!

- Stéphane D. for his work on BCFS and the tikz figures :)
- Stéphane L. and AGI for giving me the opportunity to work on Blue Coat



Backup slides



System disk organization (UEFI mode)

Files on FAT32 partition

/sgos/boot/systems/diag.si /sgos/boot/systems/system.si /sgos/boot/meta.txt /sgos/fbr.con /EFI/BOOT/BOOTx64.EFI /EFI/BOOT/osloader.si

.si files use BCFS

New: UEFI

- BOOTx64.EFI replaces starter.si
- osloader.si contains a copy of BOOTx64.EFI

New: diag

Linux diagnostic system:

- check hardware health
- interesting cli binary, with symbols:)



Boot sequence (UEFI)

- UEFI
- BOOTx64.EFI
- desired system is selected
- prekernel.exe is started, first file entry in system.si FS
- prekernel.exe setups GDT, IDT, etc. and starts kernel.exe (2nd entry)
- kernel starts sequencer.exe, (3rd entry)
- sequencer.exe parses the main_cr.cfg script
- main_cr.cfg includes main_common.cfg which starts everything

Way simpler than BIOS boot.

