April 27, 2016

1 Effects

Fix some set of resources R. A resource is some language primitive that has the authority to directly perform I/O operations. Elements of the set R are denoted by r. M is another fixed set of methods on resources. An effect is a member of the set of pairs $R \times M$. A set of effects is denoted by ε. In this system we cannot dynamically create resources or resource-operations.

Throughout we refer to the notions of effects and captures. A piece of code C has the effect (r, m) if operation m is performed on resource r during execution of C. C captures the effect (r, m) if it has the authority to perform operation m on resource r at some point during its execution.

We use $r.m$ as syntactic sugar for the effect (r, m). For example, $FileIO.append$ instead of $(FileIO, append)$.

Non-resource types may be structural or primitive (is this right?). The set of primitive types P contains familiar types. For example, $Int \in P$. If a type is not a resource and not primitive, we call it composite. Structural types are sets of method declarations. An example is $\{\texttt{def}\ double(x : Int) : Int\}$. An interesting composite type is `Unit`, which is equivalent to the empty set $\varnothing$.

2 Fully-Annotated Programs

In this first system every method in the program is explicitly annotated with its set of effects.

2.1 Grammar

$$
\begin{array}{lll}
e & ::= x & \textit{expressions} \\
& \mid\ \ \texttt{new}\ x \Rightarrow \overline{\sigma = e} & \\
& \mid\ \ e.m(e) & \\
& \mid\ \ r & \\
\\
\tau & ::= \{\bar{\sigma}\} \mid \{\bar{r}\} & \textit{types} \\
\\
\sigma & ::= d\ \texttt{with}\ \varepsilon & \textit{labeled decls.} \\
\\
\Gamma & ::= \varnothing & \\
& \mid\ \ \Gamma,\ x : \tau &
\end{array}
$$

Notes:

– Declarations (σ-terms) are annotated by what effects they have.
– All methods take exactly one argument. If a method specifies no argument, then the argument is implicitly oaf type `Unit`.

2.2 Rules

$\boxed{\Gamma \vdash e : \tau\ \texttt{with}\ \varnothing}$

$$
\frac{}{\Gamma,\ x : \tau \vdash x : \tau\ \texttt{with}\ \varnothing}\ (\varepsilon\text{-}\textsc{Var})
\qquad
\frac{}{\Gamma,\ r : \{r\} \vdash r : \{r\}\ \texttt{with}\ \varnothing}\ (\varepsilon\text{-}\textsc{Resource})
$$

$$
\frac{\Gamma,\ x : \{\bar{\sigma}\} \vdash \overline{\sigma = e}\ \texttt{OK}}{\Gamma \vdash \texttt{new}\ x \Rightarrow \overline{\sigma = e} : \{\bar{\sigma}\}\ \texttt{with}\ \varnothing}\ (\varepsilon\text{-}\textsc{NewObj})
$$

$\boxed{\Gamma \vdash \sigma = e \ \texttt{OK}}$

$$\frac{\Gamma,\ x : \tau \vdash e : \tau' \ \texttt{with}\ \varepsilon \quad \sigma = \texttt{def}\ m(x : \tau) : \tau' \ \texttt{with}\ \varepsilon}{\Gamma \vdash \sigma = e \ \texttt{OK}} \ (\varepsilon\text{-}\textsc{ValidImpl}_\sigma)$$

$\boxed{\Gamma \vdash e_1.m(e_2) : \tau \ \texttt{with}\ \varepsilon}$

$$\frac{\Gamma \vdash e_1 : \{\bar{r}\} \ \texttt{with}\ \varepsilon_1 \quad \Gamma \vdash e_2 : \tau_2 \ \texttt{with}\ \varepsilon_2 \quad m \in M}{\Gamma \vdash e_1.m(e_2) : \{\bar{r}\} \ \texttt{with}\ \{\bar{r}, m\} \cup \varepsilon_1 \cup \varepsilon_2} \ (\varepsilon\text{-}\textsc{MethCallResource})$$

$$\frac{\Gamma \vdash e_1 : \{\bar{\sigma}\} \ \texttt{with}\ \varepsilon_1 \quad \Gamma \vdash e_2 : \tau_2 \ \texttt{with}\ \varepsilon_2 \quad \sigma_i = \texttt{def}\ m_i(y : \tau_2) : \tau \ \texttt{with}\ \varepsilon}{\Gamma \vdash e_1.m_i(e_2) : \tau \ \texttt{with}\ \varepsilon_1 \cup \varepsilon_2 \cup \varepsilon} \ (\varepsilon\text{-}\textsc{MethCallObj})$$

Notes:

- Every expression in the program must be explicitly annotated; either as σ-terms or by what they capture.
- The rules ε-Var, ε-Resource, and ε-NewObj have in their consequents an expression typed with no effect: merely having an object or resource is not an effect; you must do something with it, like a call a method on it, in order for it to be an effect.
- ε-ValidImpl says that the return type and effects of the body of a method must agree with what its signature says.
- In ε-MethCallResource, we may only call a method m on a resource r if m is a predefined operation in the set M. Invoking m returns the resource r you called it upon (which has potentially different state afterwards).

3 Partially-Annotated Programs

In this second system methods may either be fully labeled with their effects or have no labels. When they have no labels a conservative effect inference is performed using rules which provide an upper-bound (not necessarily tight) on the effects of the code when executed.

3.1 Grammar

$$
\begin{array}{rcll}
e & ::= & x & \textit{expressions} \\
 & | & \texttt{new}_\sigma\ x \Rightarrow \overline{\sigma = e} & \\
 & | & \texttt{new}_d\ x \Rightarrow \overline{d = e} & \\
 & | & e.m(e) & \\
 & | & r & \\
 & & & \\
\tau & ::= & \{\bar{\sigma}\} & \textit{types} \\
 & | & \{\bar{r}\} & \\
 & | & \{\bar{d}\} & \\
 & | & \{\bar{d}\ \texttt{captures}\ \varepsilon\} & \\
 & & & \\
\sigma & ::= & d\ \texttt{with}\ \varepsilon & \textit{labeled decls.} \\
 & & & \\
d & ::= & \texttt{def}\ m(x : \tau) : \tau & \textit{unlabeled decls.}
\end{array}
$$

Notes:

- σ denotes a declaration with effect labels. d denotes a declaration without effect labels.
- There are two new expressions: $\texttt{new}_\sigma$ for objects whose declarations are annotated; $\texttt{new}_d$ for objects whose declarations aren't.
- $\{\bar{d}\ \texttt{captures}\ \varepsilon\}$ is a special kind of type that doesn't appear in the source program, but may be assigned as a consequence of the capture rules.

3.2 Rules

In addition to the rules from the previous system, the partially-annotated system has the following rules.

$\boxed{\Gamma \vdash e : \tau}$

$$\frac{}{\Gamma,\ x : \tau \vdash x : \tau}\ (\text{T-Var}) \qquad \frac{}{\Gamma,\ r : \{\bar{r}\} \vdash r : \{\bar{r}\}}\ (\text{T-Resource})$$

$$\frac{\Gamma \vdash r : \{\bar{r}\} \quad \Gamma \vdash e : \tau \quad m \in M}{\Gamma \vdash r.m(e_1) : \{\bar{r}\}}\ (\text{T-MethCall}_r)$$

$$\frac{\Gamma \vdash e_1 : \{\bar{\sigma}\},\ \texttt{def}\ m(x : \tau_1) : \tau_2\ \texttt{with}\ \varepsilon \in \{\bar{\sigma}\} \quad \Gamma \vdash e_2 : \tau_1}{\Gamma \vdash e_1.m(e_2) : \tau_2}\ (\text{T-MethCall}_\sigma)$$

$$\frac{\Gamma \vdash e_1 : \{\bar{d}\},\ \texttt{def}\ m(x : \tau_1) : \tau_2 \in \{\bar{d}\} \quad \Gamma \vdash e_2 : \tau_1}{\Gamma \vdash e_1.m(e_2) : \tau_2}\ (\text{T-MethCall}_d)$$

$$\frac{\Gamma \vdash \sigma_i = e_i\ \texttt{OK}}{\Gamma \vdash\ \texttt{new}_\sigma\ x \Rightarrow \overline{\sigma = e} : \{\bar{\sigma}\}}\ (\text{T-New}_\sigma) \qquad \frac{\Gamma \vdash d_i = e_i\ \texttt{OK}}{\Gamma \vdash\ \texttt{new}_d\ x \Rightarrow \overline{d = e} : \{\bar{d}\}}\ (\text{T-New}_d)$$

$\boxed{\Gamma \vdash d = e \ \texttt{OK}}$

$$\frac{d = \texttt{def}\ m(x : \tau_1) : \tau_2 \quad \Gamma \vdash e : \tau_2}{\Gamma \vdash d = e\ \texttt{OK}}\ (\varepsilon\text{-}\textsc{ValidImpl}_d)$$

$\boxed{\Gamma \vdash e_1.m(e_2) : \tau\ \texttt{with}\ \varepsilon}$

$$\frac{\varepsilon = effects(\Gamma') \quad \Gamma' \subseteq \Gamma \quad \Gamma', x : \{\bar{d}\ \texttt{captures}\ \varepsilon\} \vdash \overline{d = e}\ \texttt{OK}}{\Gamma \vdash\ \texttt{new}_d\ x \Rightarrow \overline{d = e} : \{\bar{d}\ \texttt{captures}\ \varepsilon\}\ \texttt{with}\ \varnothing}\ (\text{C-}\textsc{NewObj})$$

$$\frac{\Gamma \vdash e_1 : \{\bar{d}\ \texttt{captures}\ \varepsilon\}\ \texttt{with}\ \varepsilon_1 \quad \Gamma \vdash e_2 : \tau_2\ \texttt{with}\ \varepsilon_2 \quad d_i := \ \texttt{def}\ m_i(y : \tau_2) : \tau}{\Gamma \vdash e_1.m_i(e_2) : \tau\ \texttt{with}\ \varepsilon_1 \cup \varepsilon_2 \cup effects(\tau_2) \cup \varepsilon}\ (\text{C-}\textsc{MethCall})$$

Notes:

- The ε judgements are to be applied to annotated parts of the program; the C rules for unannotated parts.
- The rules ε-Var, ε-Resource, and ε-NewObj have in their antecedents an expression typed with no effect. Merely having an object or resource is not an effect; you must do something with it, like a call a method on it, in order for your program to have effects.
- The T judgements before standard typechecking, but they operate on annotated terms. They are needed to apply the ε-ValidImpl$_d$) rule.
- In applying C-NewObj the variable Γ is the current context. The variable Γ' is some sub-context. A good choice of sub-context is Γ restricted to the free variables in the method-body being typechecked. This means we only consider the effects used in the method-body and gives a better approximation of its effects.
- When an unannotated d-declaration is encountered it is first assigned a γ-type by C-NewObj. This annotates it as capturing a certain set of effects. C-MethCall can then conclude its effects to be what it captures.

3.3 Effects Function

The `effects` function returns the set of effects in a particular typing context.

A method m can return a resource r (or an object that returns r, and so on). Returning a resource isn't an effect but it means any unannotated program using m also captures r. To account for this, `effects` also uses `escapes`.

- $\texttt{effects}(\cdot) = \varnothing$
- $\texttt{effects}(\{\bar{r}\}) = \{(r, m) \mid r \in \bar{r}, m \in M\}$
- $\texttt{effects}(\{\bar{d}\ \texttt{captures}\ \varepsilon_1\}\ \texttt{with}\ \varepsilon_2) = \varepsilon_1 \cup \varepsilon_2 \cup \texttt{escapes}(\bar{d})$
- $\texttt{effects}(d\ \texttt{with}\ \varepsilon) = \varepsilon \cup \texttt{escapes}(d\ \texttt{with}\ \varepsilon)$
- $\texttt{effects}(\{\bar{\sigma}\}) = \bigcup_{\sigma \in \bar{\sigma}}\ \texttt{effects}(\sigma)$

3.4 Escapes Function

$\boxed{\texttt{escapes}(\tau)}$

$$\frac{}{\texttt{escapes}(\texttt{def}\ m(x:\tau):\{\bar{r}\}) = \{(r,m) \mid r \in \bar{r}, m \in M\}}\ (\textsc{Escapes-Resource})$$

$$\frac{\tau_2 \notin P \quad \tau_2 \notin R}{\texttt{escapes}(\texttt{def}\ m(x:\tau_1):\tau_2) = escapes(\tau_2)}\ (\textsc{Escapes-Structural})$$

$$\frac{\tau_2 \in P}{\texttt{escapes}(\texttt{def}\ m(x:\tau_1):\tau_2) = \varnothing}\ (\textsc{Escapes-Primitive})$$

$$\frac{}{\texttt{escapes}(\texttt{def}\ m(x:\tau_1):\tau_2\ with\ \varepsilon) = \texttt{escapes}(\texttt{def}\ m(x:\tau_1):\tau_2)}\ (\textsc{Escapes}_{\sigma})$$

$$\frac{}{\texttt{escapes}(\{\bar{d}\}) = \bigcup_{d \in \bar{d}}\ escapes(d)}\ (\textsc{Escapes}_{\bar{d}}) \qquad \frac{}{\texttt{escapes}(\{\bar{\sigma}\}) = \bigcup_{\sigma \in \bar{\sigma}}\ escapes(\sigma)}\ (\textsc{Escapes}_{\bar{\sigma}})$$