

# TEMPEST Attacks Against AES

Covertly stealing keys for \$200

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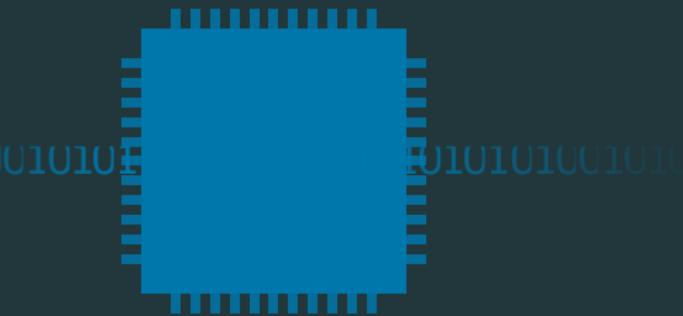
Craig Ramsay & Jasper Lohuis

September 22, 2017

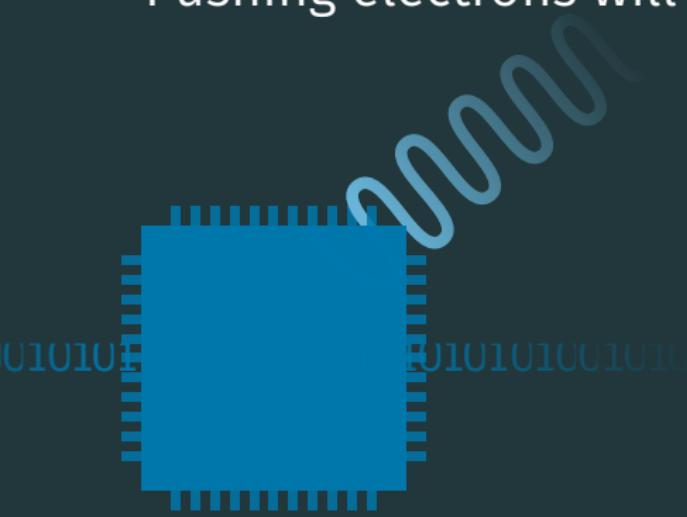
# Introduction

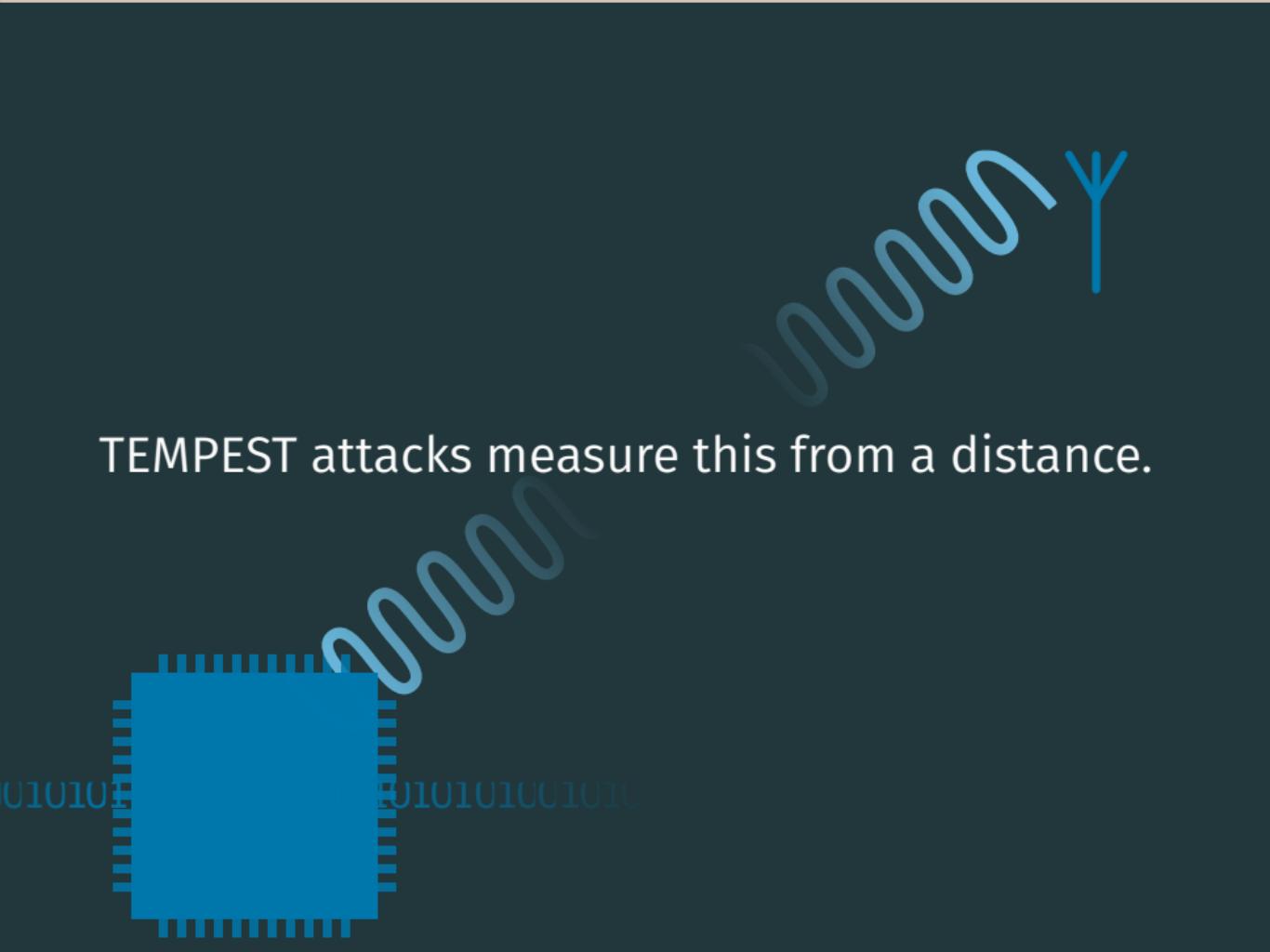
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Your code just pushes electrons around.



Pushing electrons will make magnetic fields.





TEMPEST attacks measure this from a distance.



TEMPEST attacks measure this from a distance.

# Project people

Duncan Lew

First intern. Close-by FPGA  
attacks



Freek van Tienen

We'll see!



Craig Ramsay

Radio-based workflow &  
attacking ARM

Jasper Lohuis

Cheap shielding, SDRs &  
antennas

Thanks for feeding us, folks

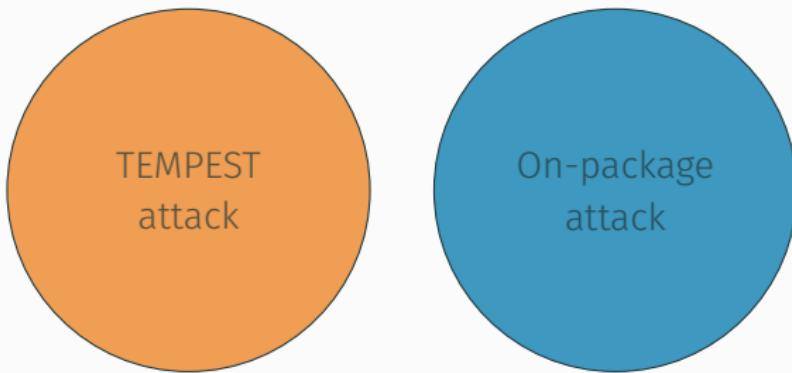


**FOX IT**

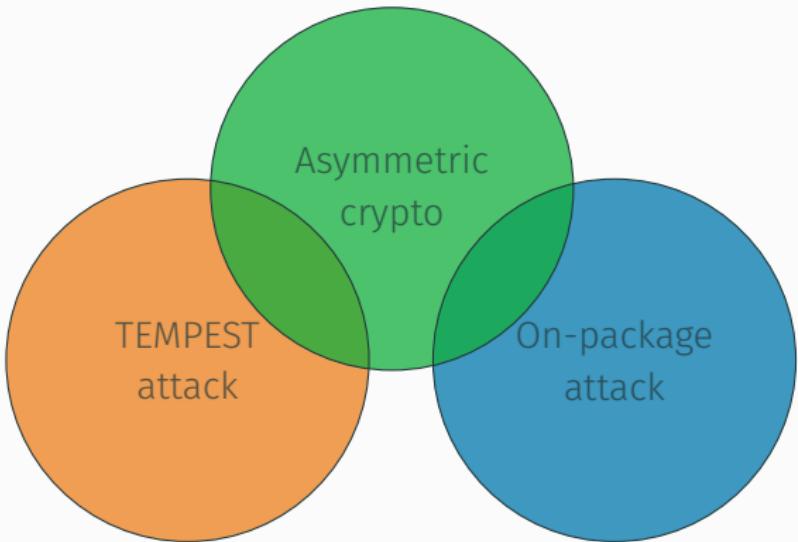
FOR A MORE SECURE SOCIETY

riscure

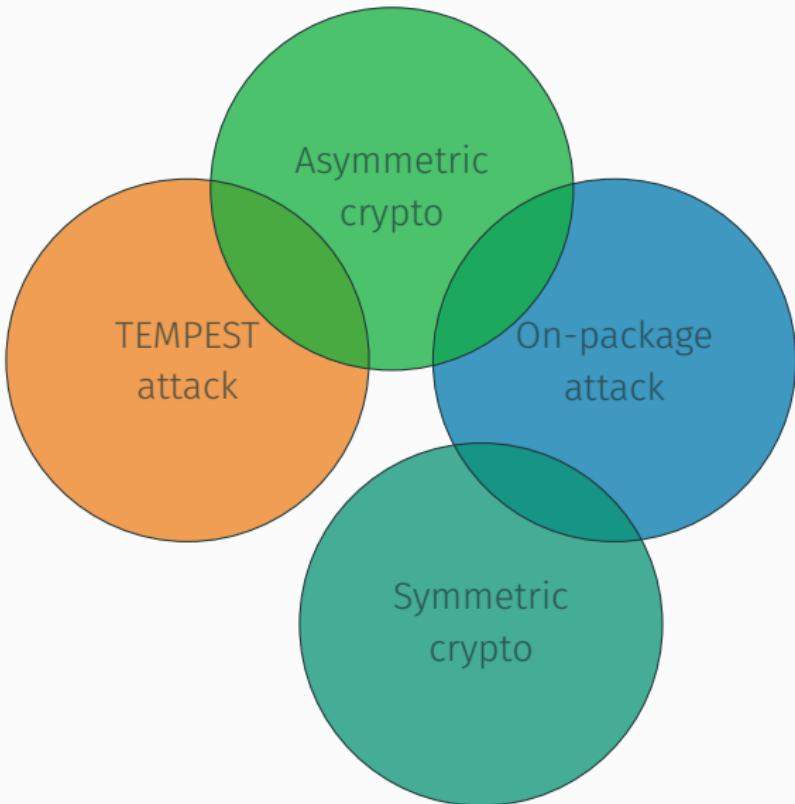
## Existing Work



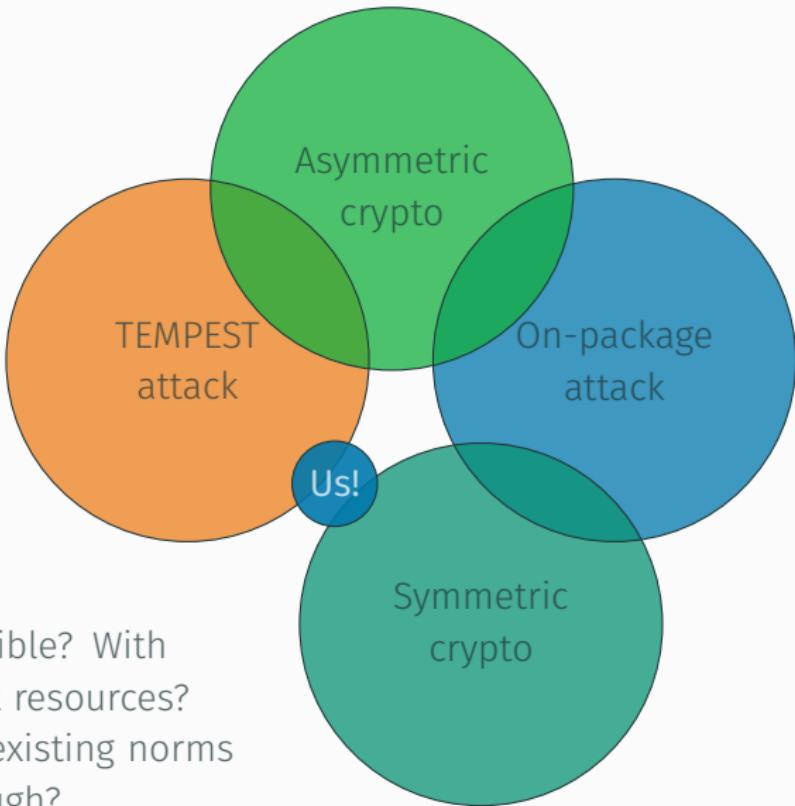
## Existing Work



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# Existing Work



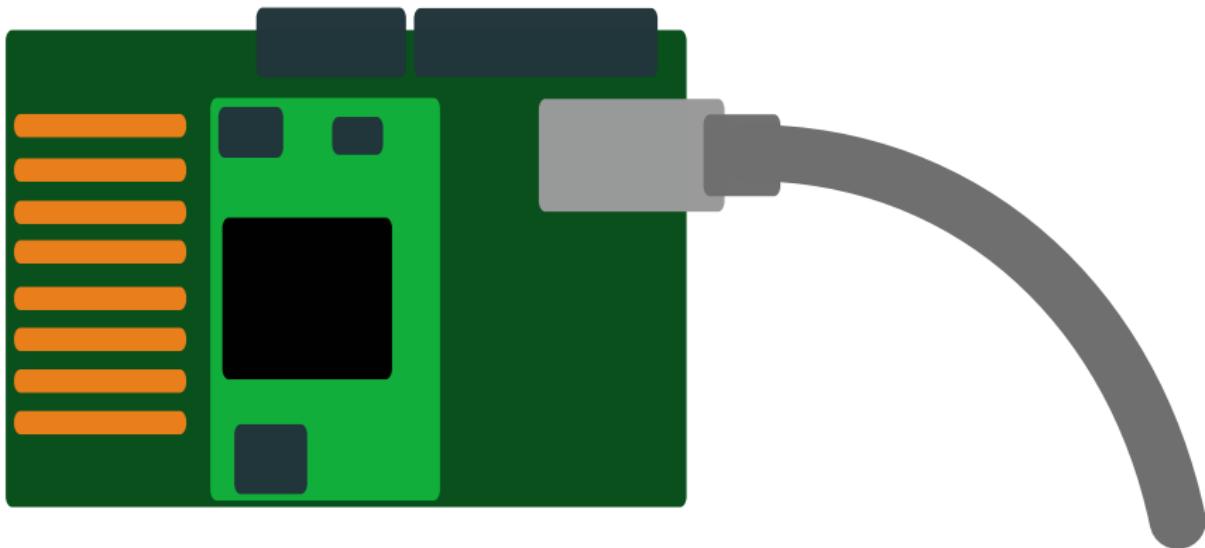
## Replicating On-Package Attack

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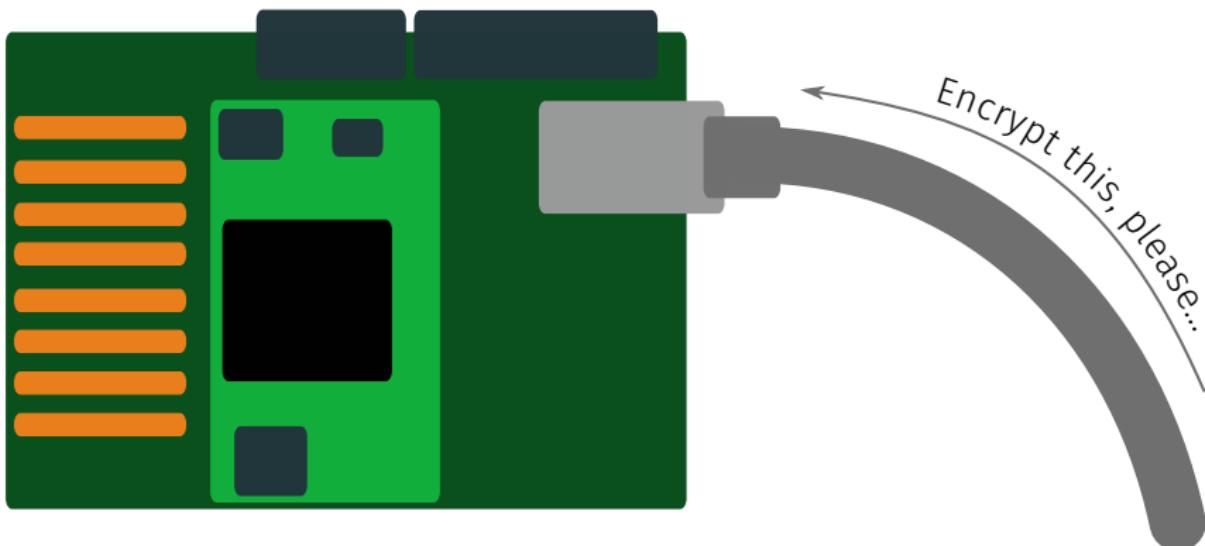
# Overview



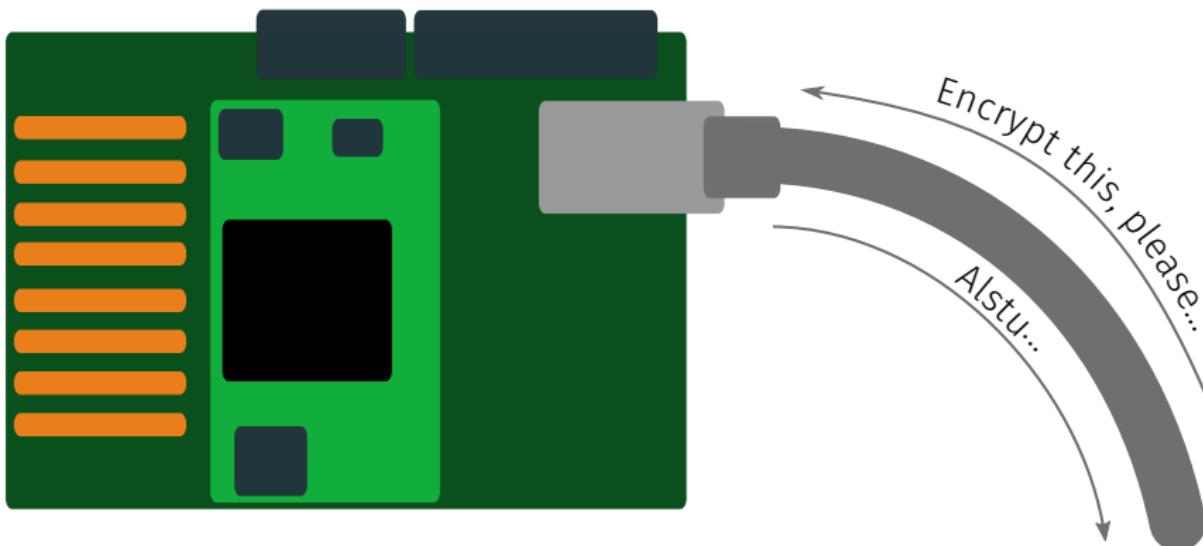
# Overview



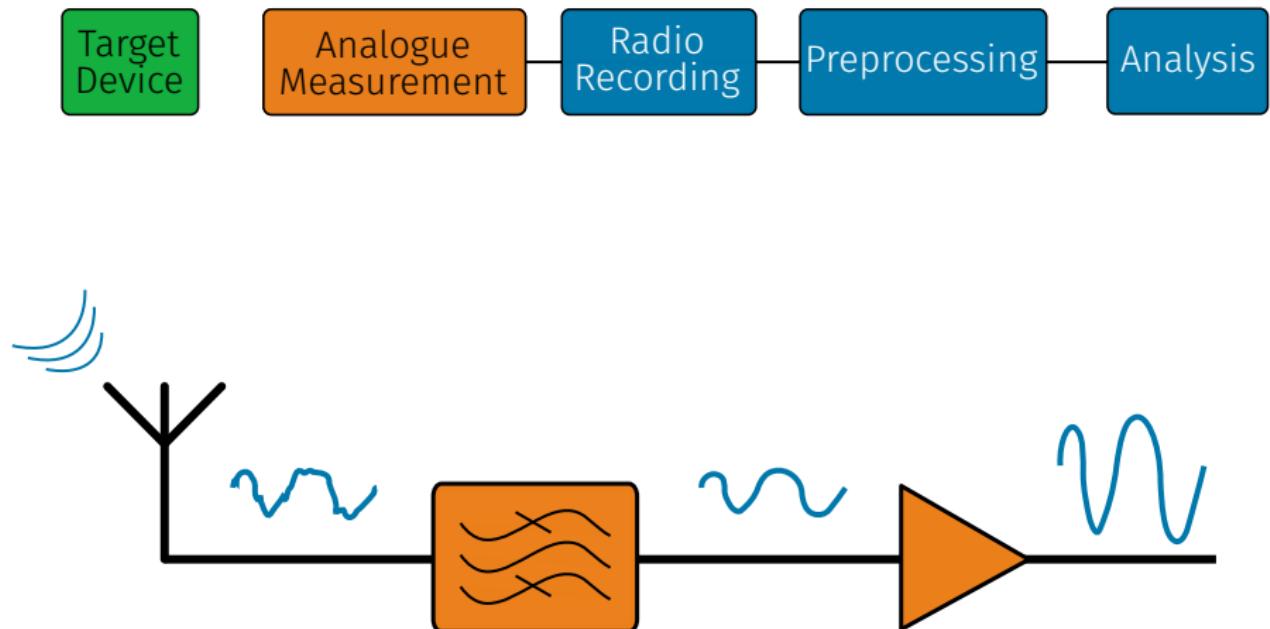
# Overview



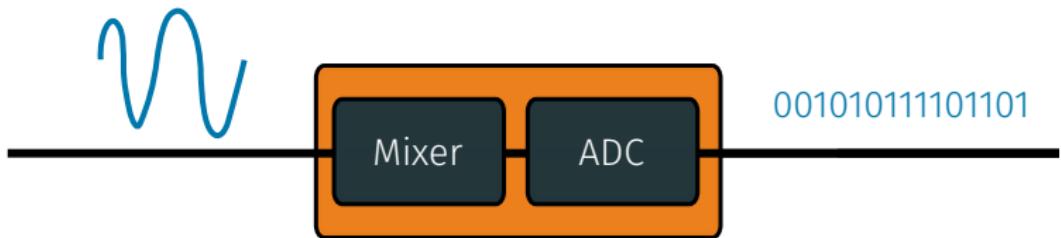
# Overview



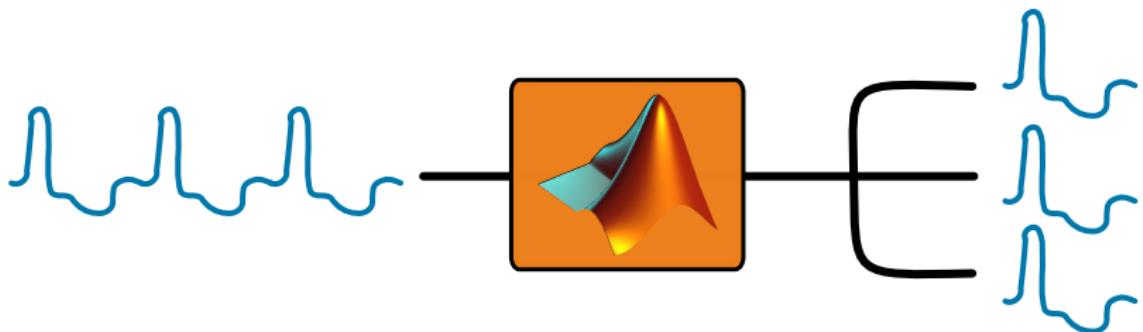
# Overview



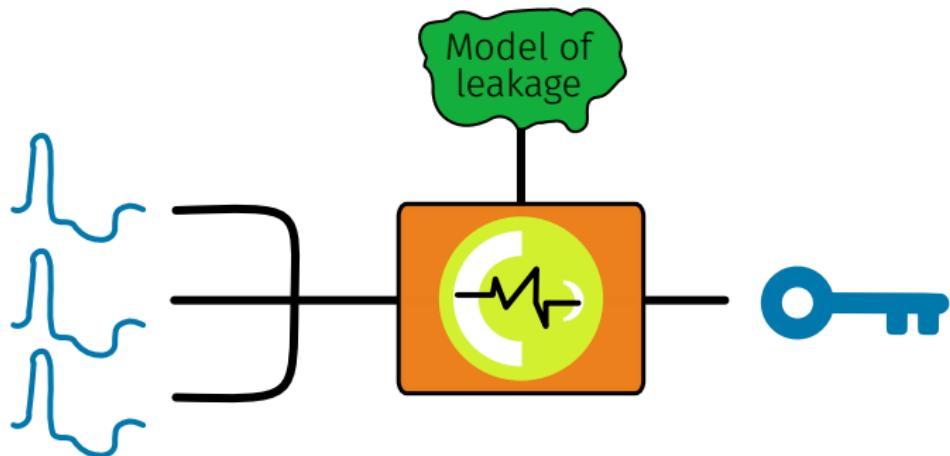
# Overview



# Overview



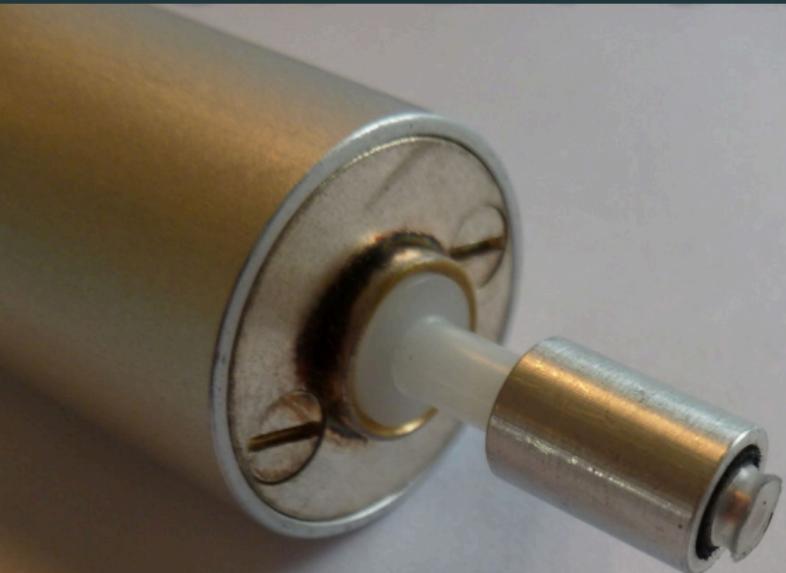
# Overview



## Measuring the field



## Measuring the field



## Recording —Low-end



## Recording —Low-end

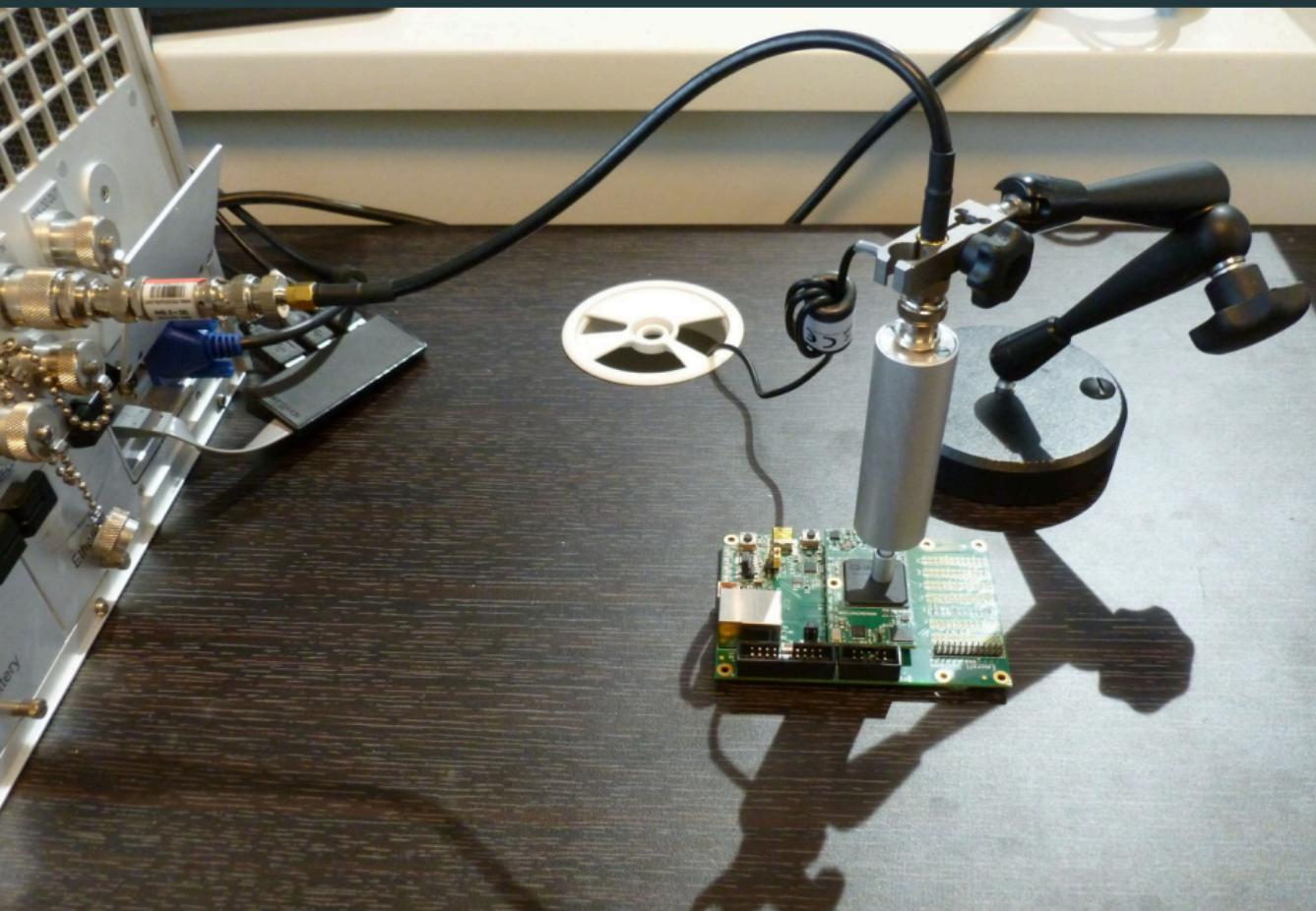


€20

2.4 MHz Bandwidth

5.2 MB/s

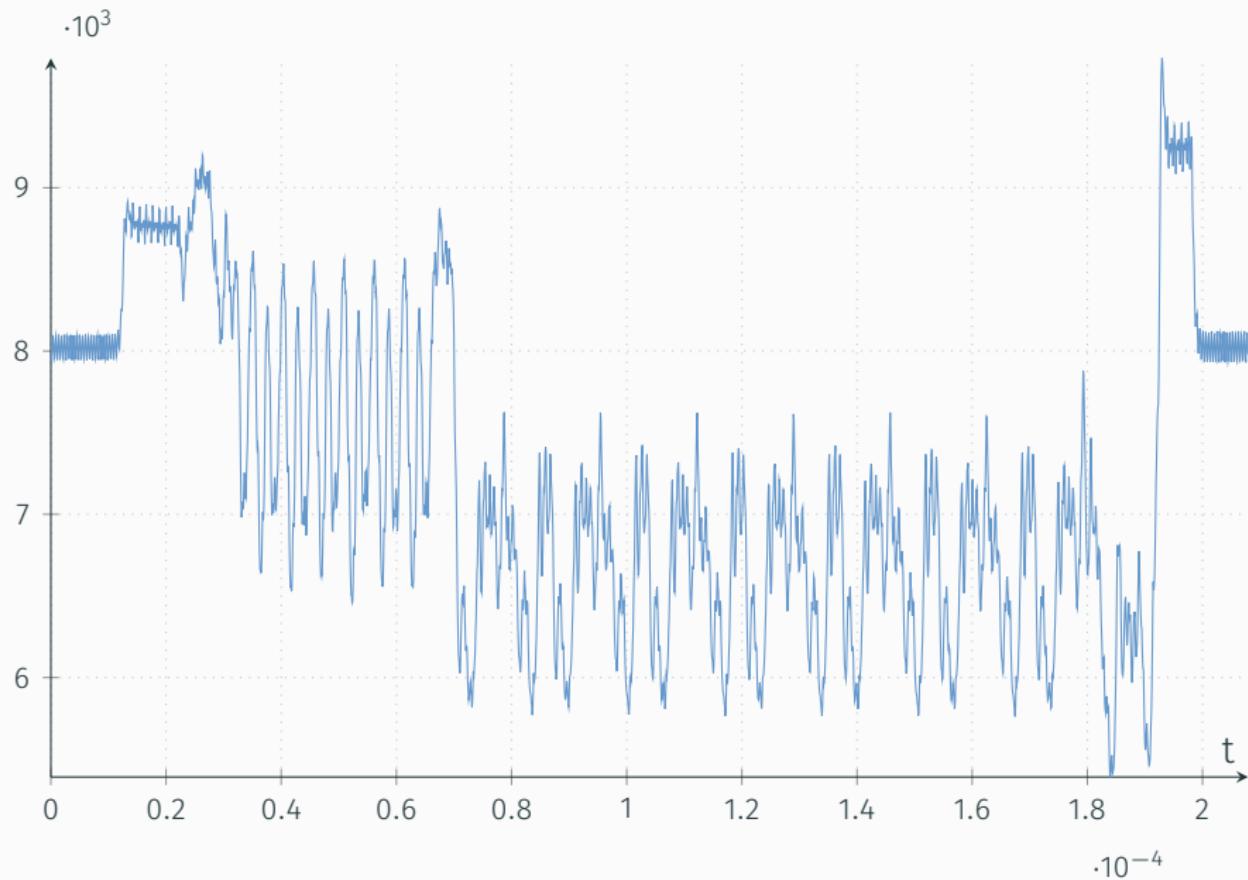
# Positioning



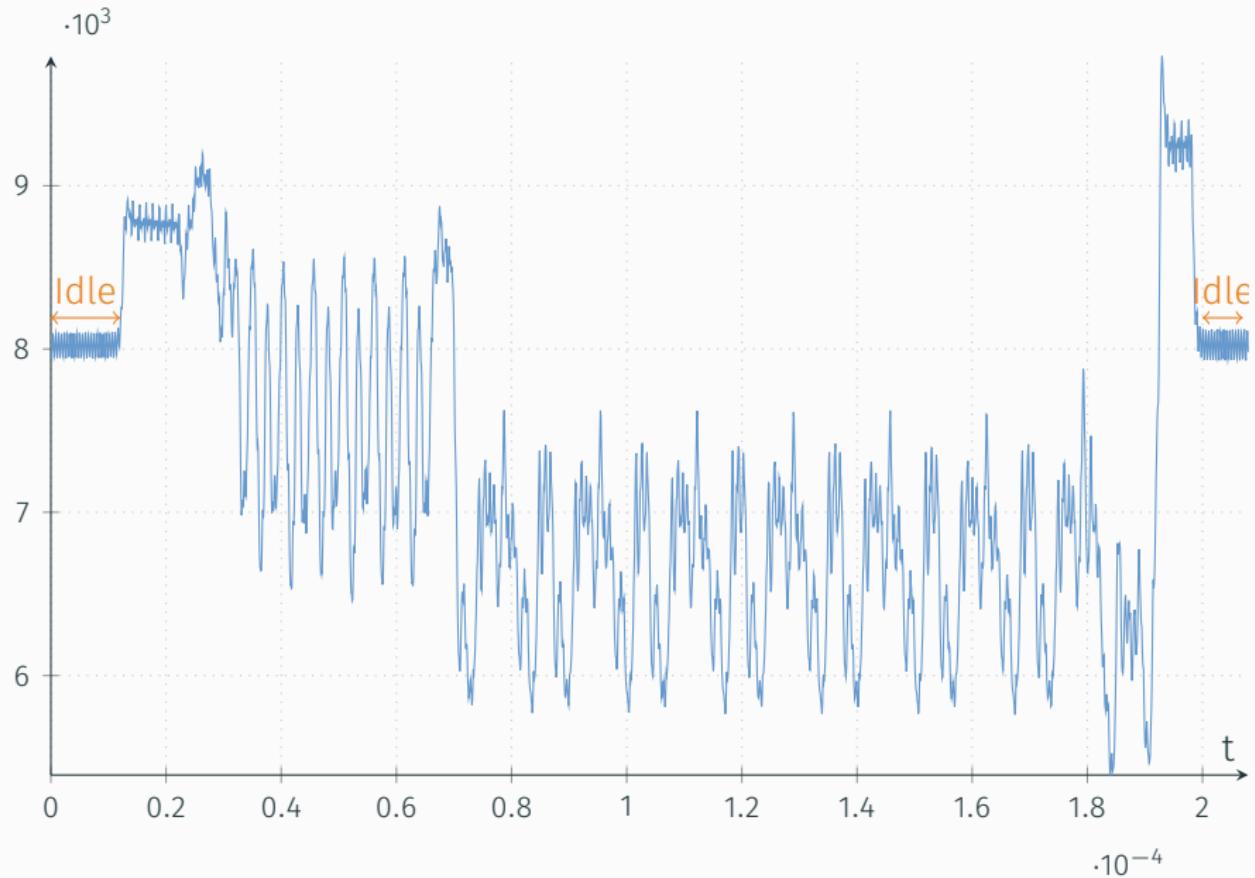
# Positioning



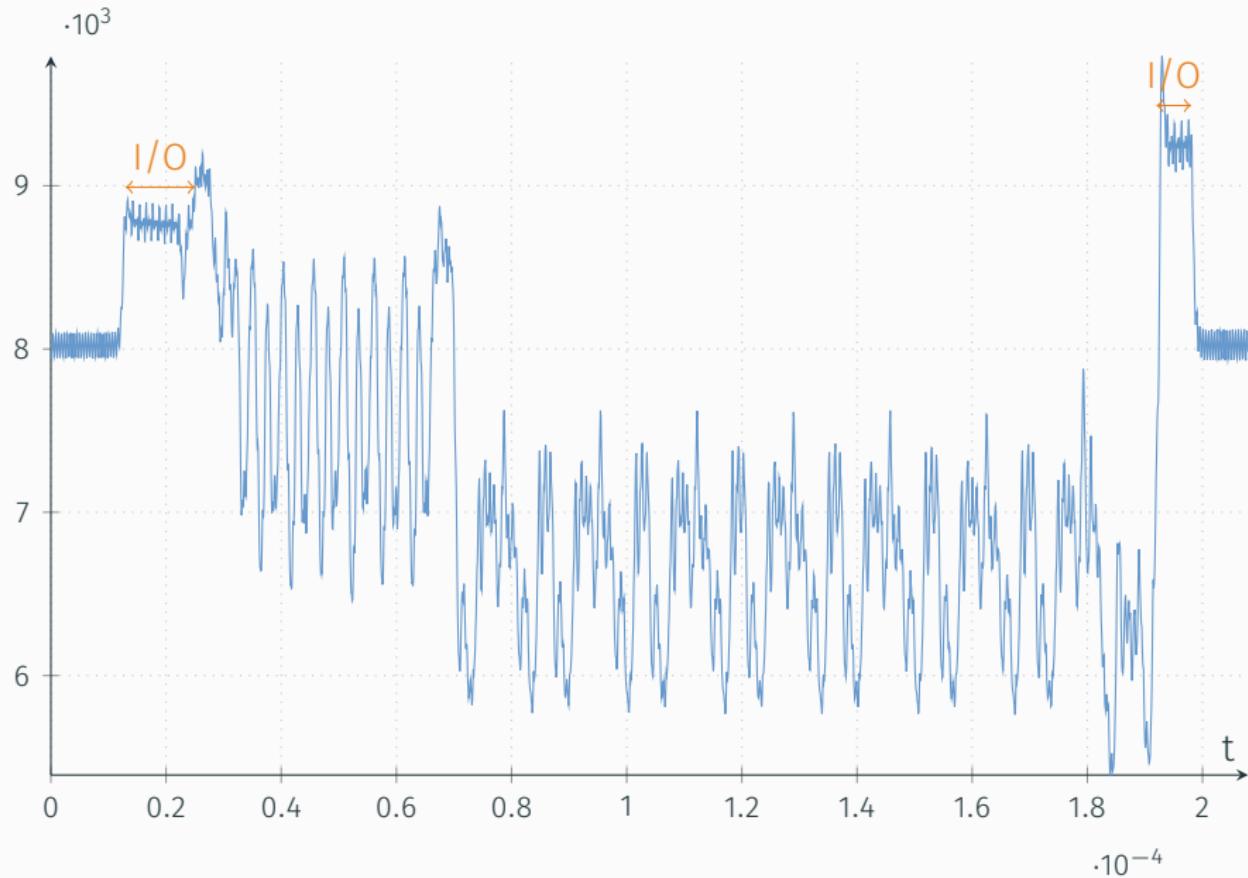
# ARM Software Trace



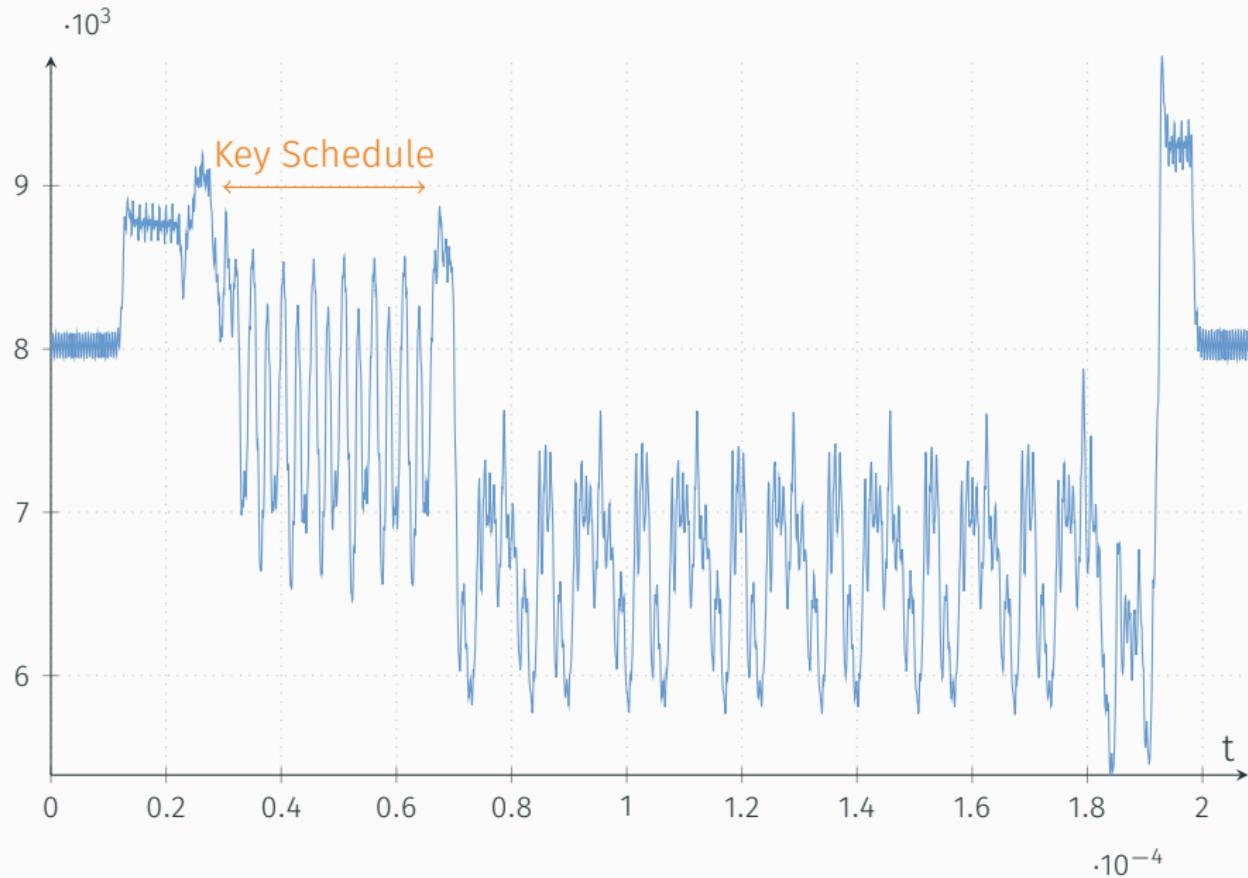
# ARM Software Trace



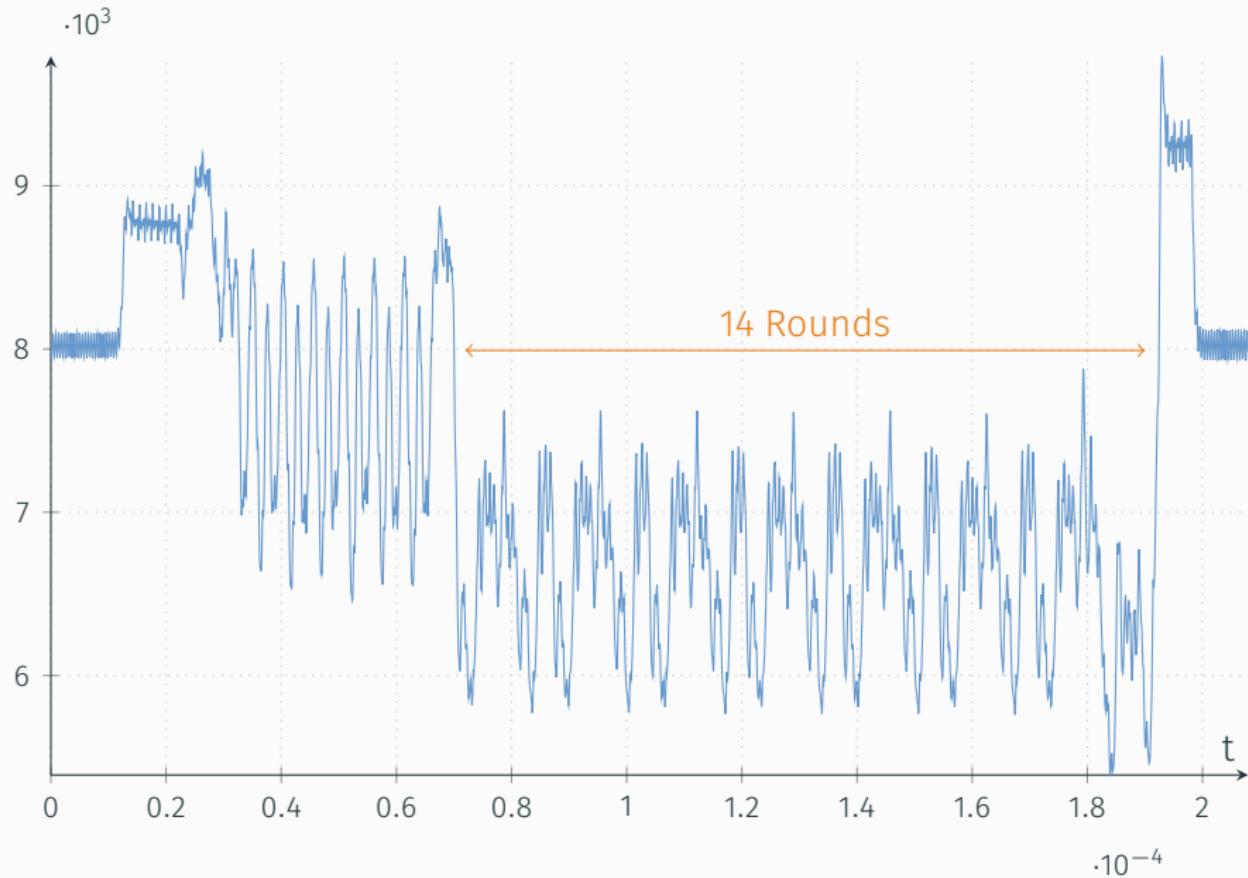
# ARM Software Trace



# ARM Software Trace



# ARM Software Trace



Nice... but still, how do you get a key?

(This part is just existing SCA techniques)

Our trace is related to “power consumption”.

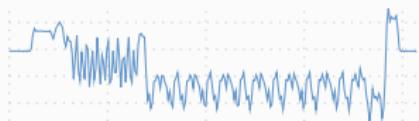
$$P \approx P_{static} + P_{noise} + P_{data} + P_{operation}$$

$$P \approx P_{\text{static}} + P_{\text{noise}} + P_{\text{data}} + P_{\text{operation}}$$

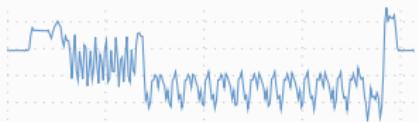
# Correlation Intro

Input Byte

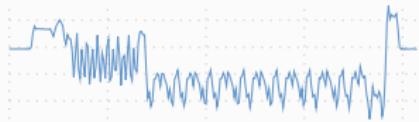
0011 0101



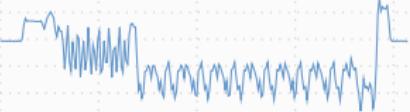
0100 0000



1101 1110



# Correlation Intro

| Input Byte | (# '1' bits) | Power estimate   |
|------------|--------------|--|
| 0011 0101  | 4            |  |
| 0100 0000  | 1            |  |
| 1101 1110  | 6            |  |

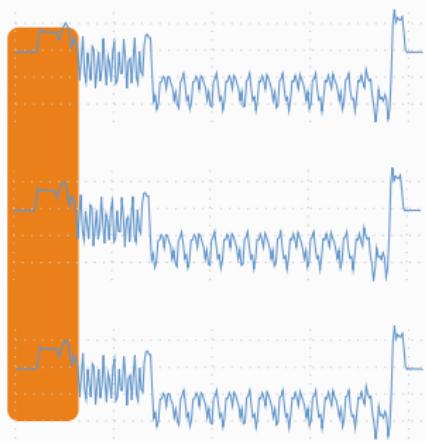
## Correlation Intro

## Input Byte

(# '1' bits)  
Power estimate

0011 0101

4



0100 0000

1

1101 1110

6

## Correlate!

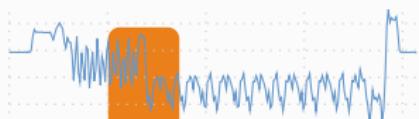
## Correlation Intro

## Input Byte

(# '1' bits)  
Power estimate

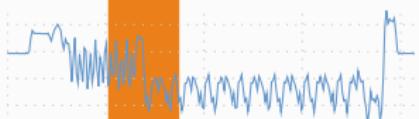
0011 0101

4



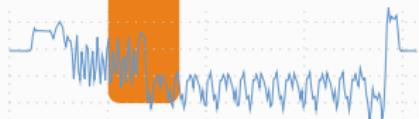
0100 0000

1



1101 1110

6



## Correlate!

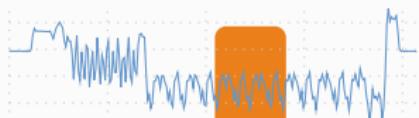
## Correlation Intro

## Input Byte

(# '1' bits)  
Power estimate

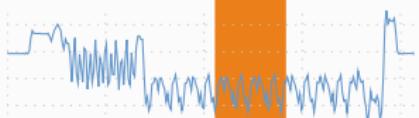
0011 0101

4



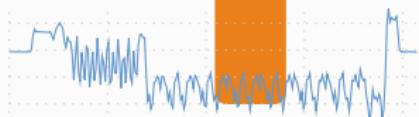
0100 0000

1



1101 1110

6



## Correlate!

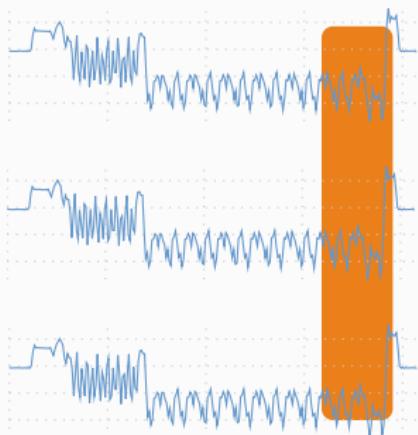
## Correlation Intro

## Input Byte

(# '1' bits)  
Power estimate

0011 0101

4



0100 0000

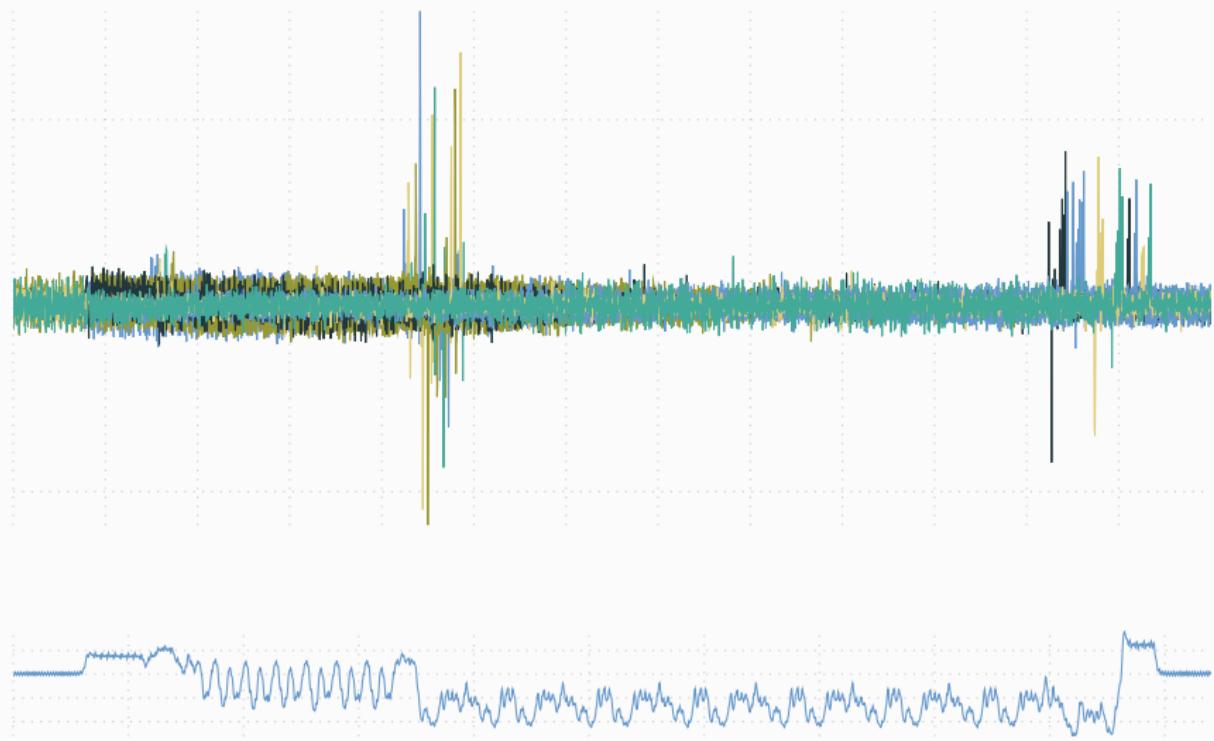
1

1101 1110

6

## Correlate!

# I/O Correlation



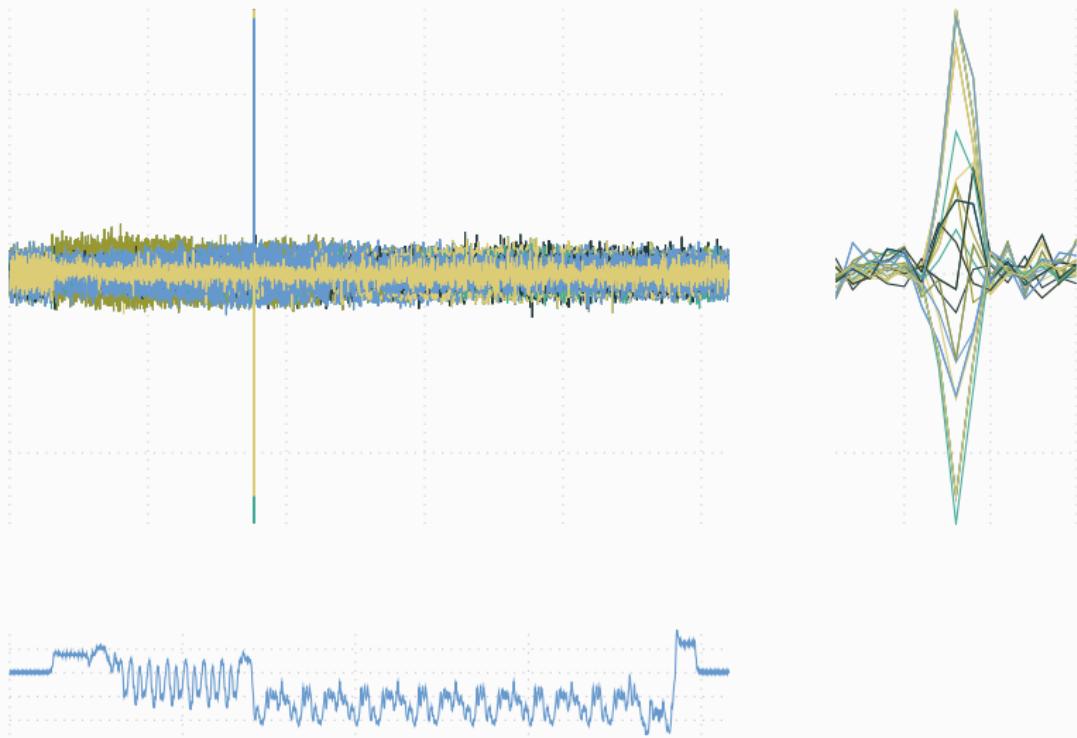
We can detect data!

Let's find a value using **1 key byte**  
and correlate for all 256 possibilities

$32 \times 2^8$  guesses  
(instead of  $2^{256}$ )

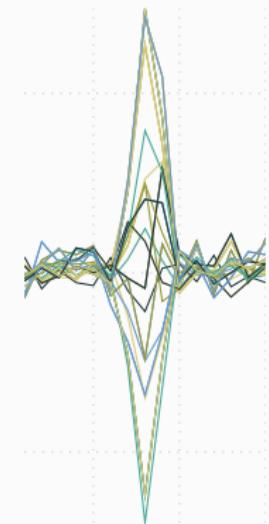
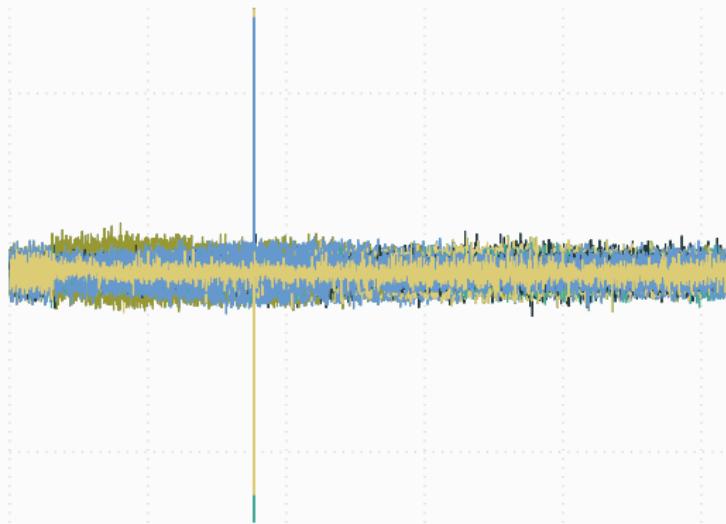
8192 guesses  
(instead of  $10^{77}$ )

# T Table Correlation



Known-key bitwise on T Table lookup

# T Table Correlation



Easy & works  
(but could do better)

Known-key bitwise on T Table lookup

*“You know can addresses leak too, right?”*

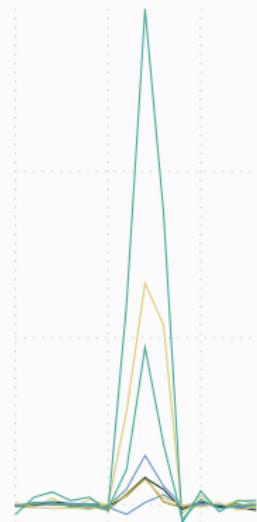
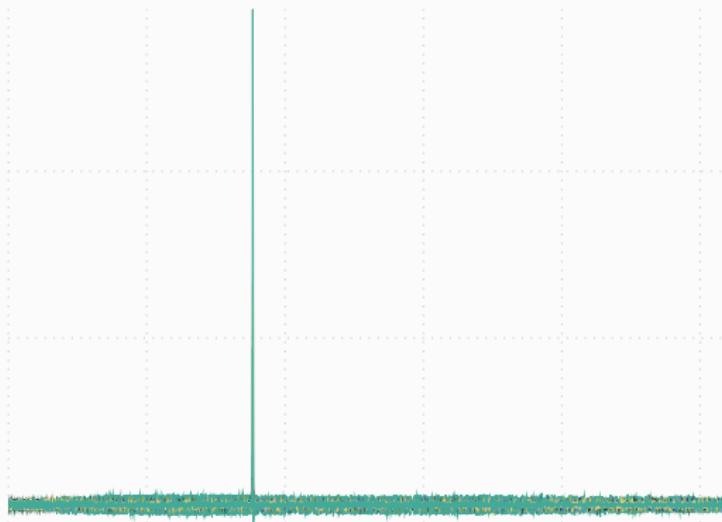
—Riscure, 2017

*“Oh... thanks.”* — Me, 2017

# ARM T Table Addresses

|                        |                                    |
|------------------------|------------------------------------|
| 1508: 4b4a             | ldr r3 , [pc, #296]                |
| 150a: 681b             | ldr r3 , [r3, #0]                  |
| 150c: 0e1b             | lsrs r3 , r3 , #24                 |
| 150e: 4a4d             | ldr r2 , [pc, \#308]               |
| <b>1510: f852 2023</b> | <b>ldr.w r2 , [r2, r3, lsl #2]</b> |
| 1514: 4b48             | ldr r3 , [pc, #288]                |
| 1516: 681b             | ldr r3 , [r3, #0]                  |
| <b>1518: 0c1b</b>      | lsrs r3 , r3 , #16                 |
| <b>151a: b2db</b>      | uxtb r3 , r3                       |
| <b>151c: 494a</b>      | ldr r1 , [pc, #296]                |

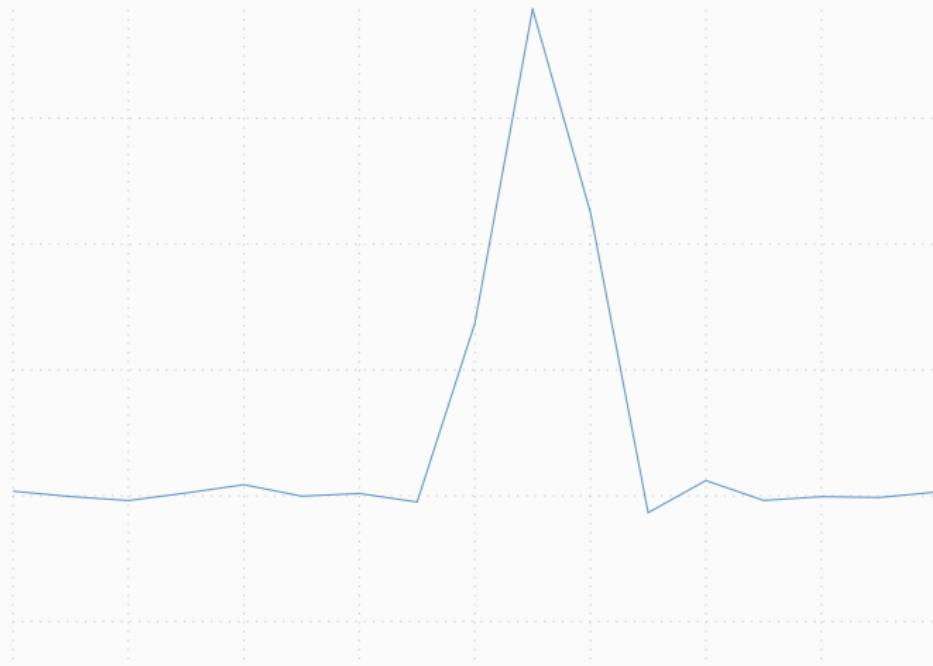
# T Table Address Correlation



Known-key bitwise on T Table lookup address  $\oplus$  previous address

If the correlation for the *correct* key byte is biggest, we have an attack.

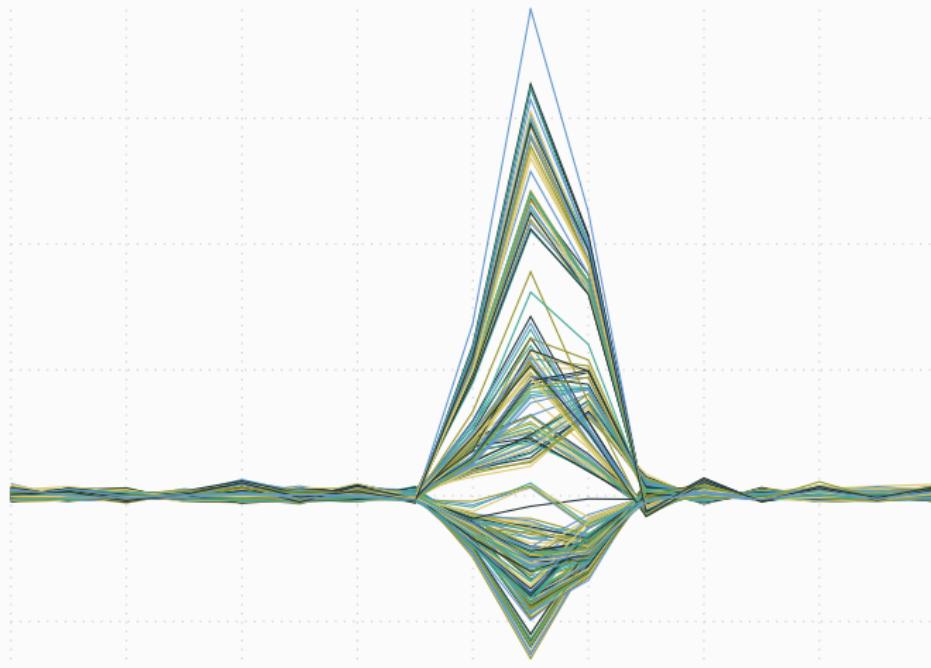
# T Table Attack



Correct key byte...

HD on T Table lookup address (real attack)

# T Table Attack



All key byte guesses. We win!

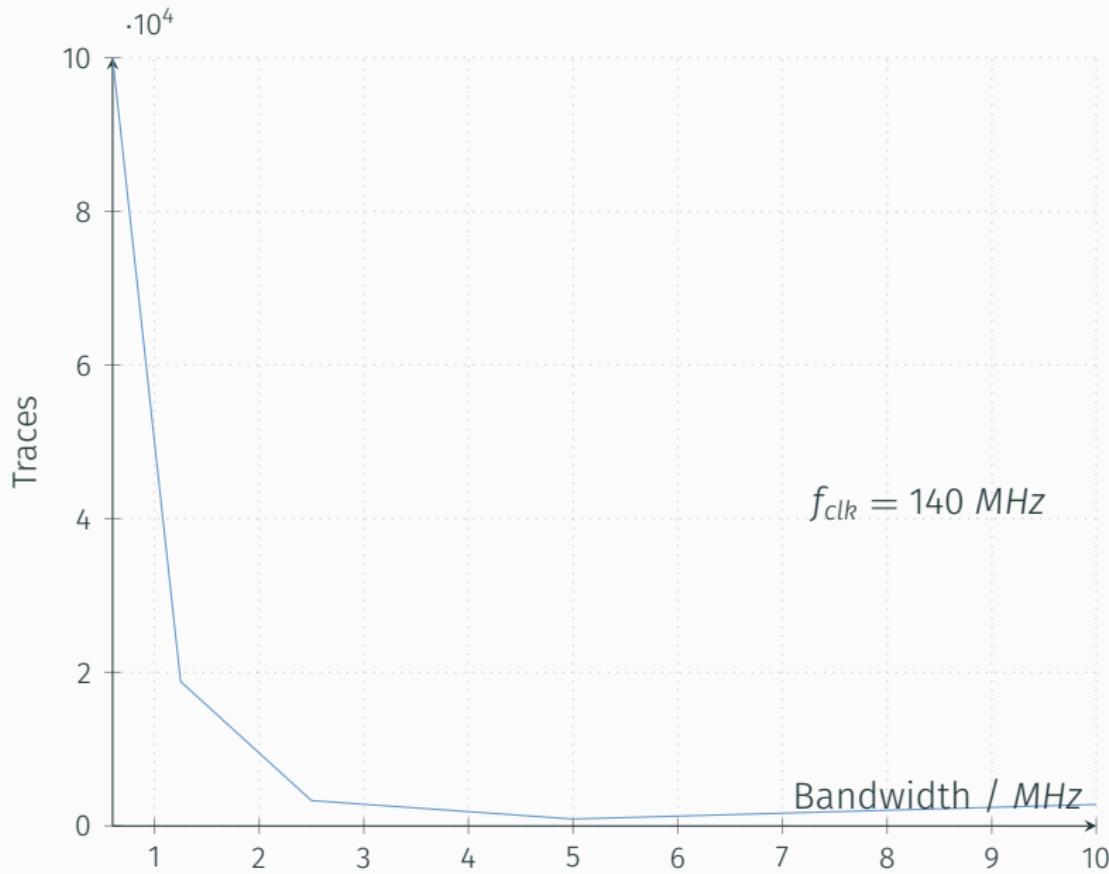
HD on T Table lookup address (real attack)

Repeat this for all 32 key bytes and we have the full key

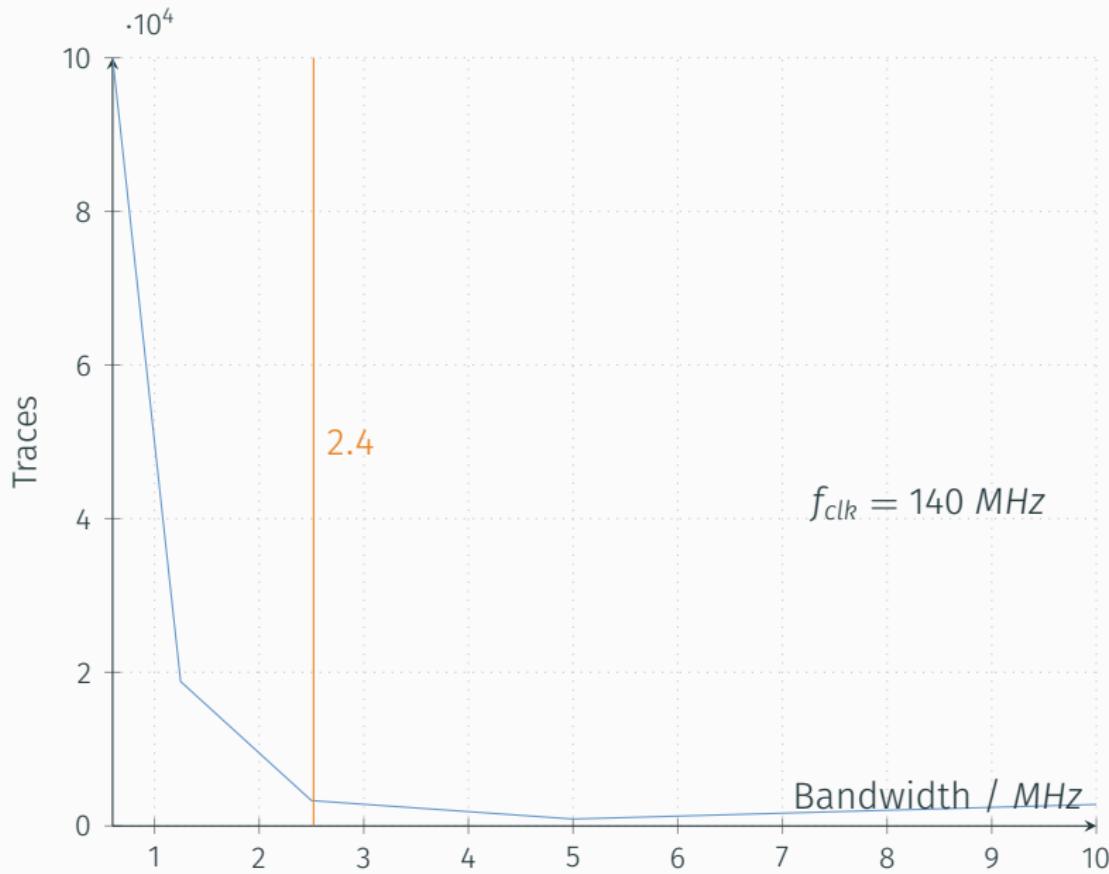
## On-package attack results

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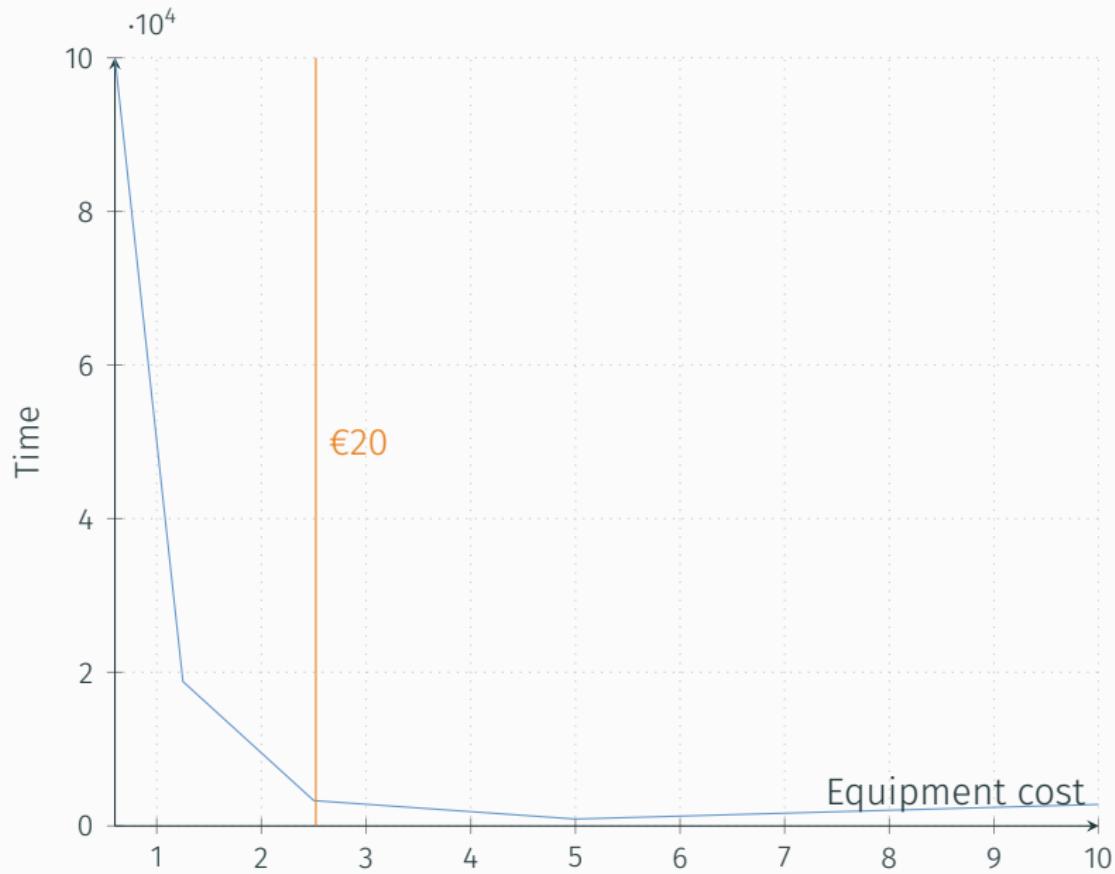
# Bandwidth vs # traces



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# Bandwidth vs # traces

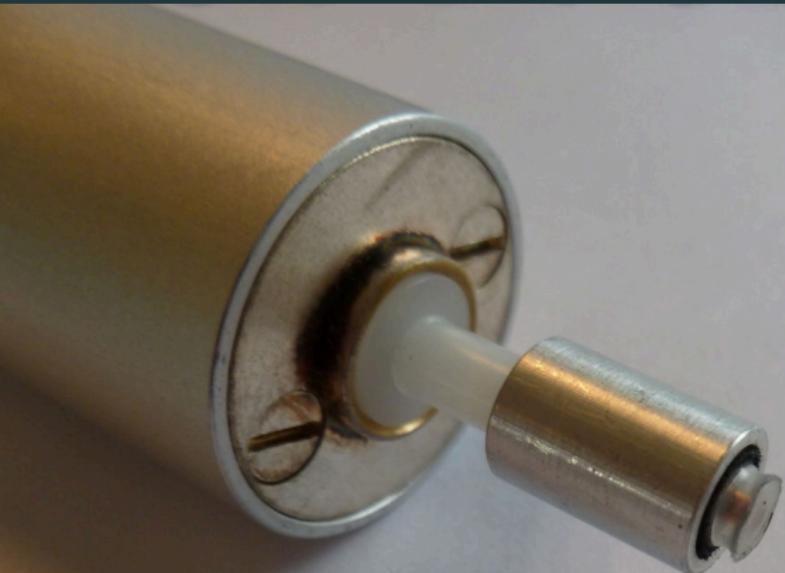


## Getting some distance

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Only need to improve analogue side.  
Analysis is the same.

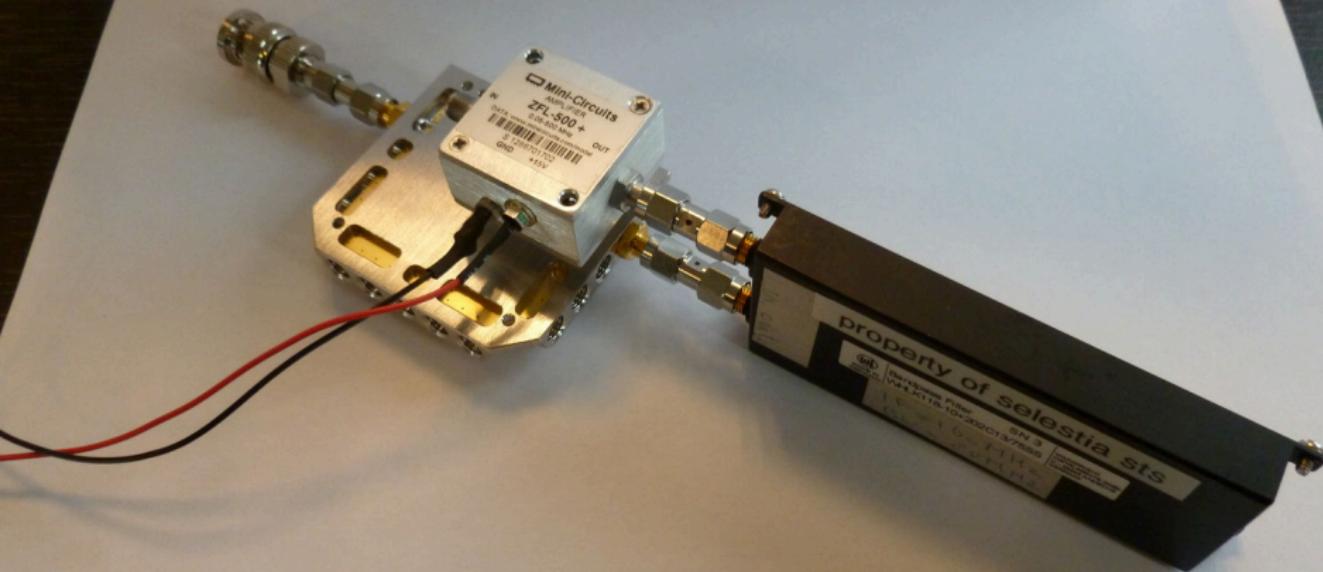
Loop size



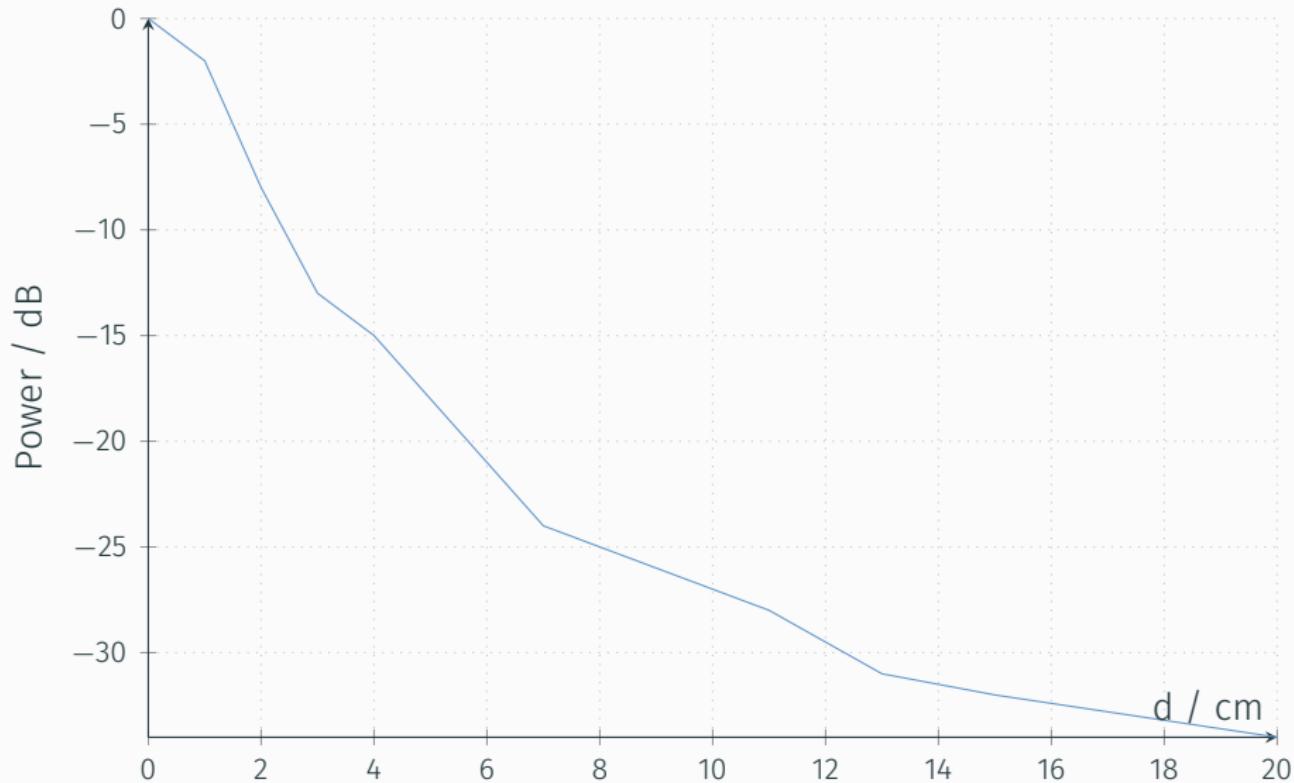
Loop size

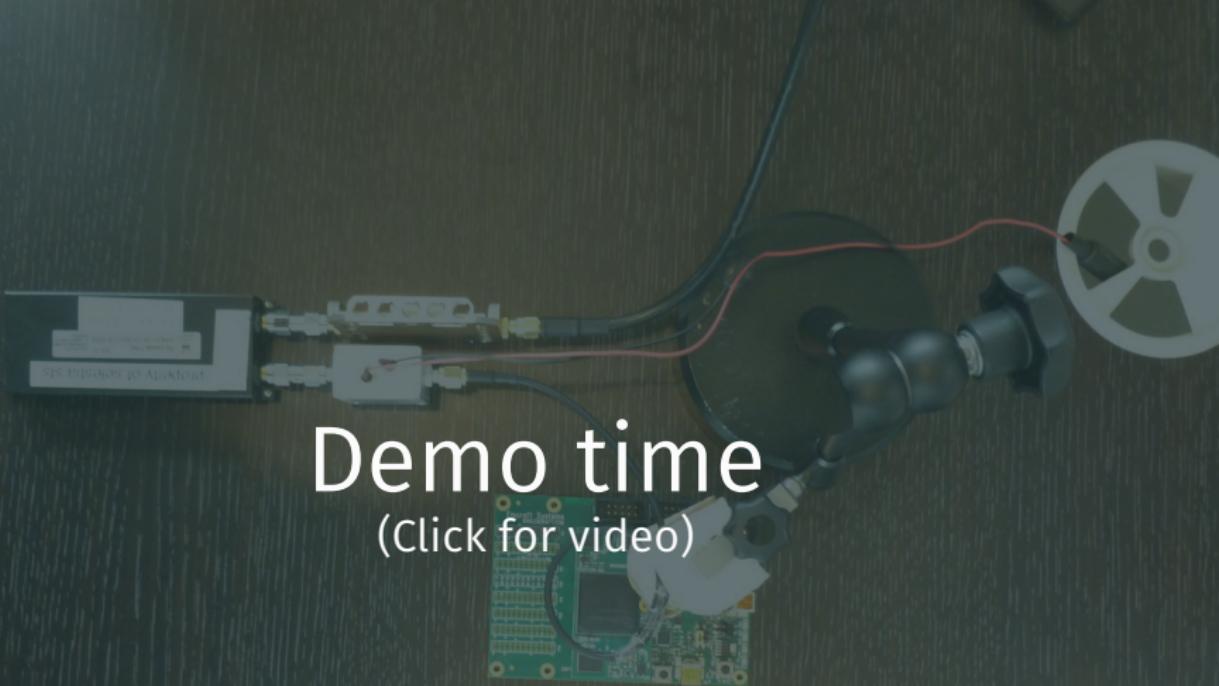


# Amplification and filtering



# Small loop distance

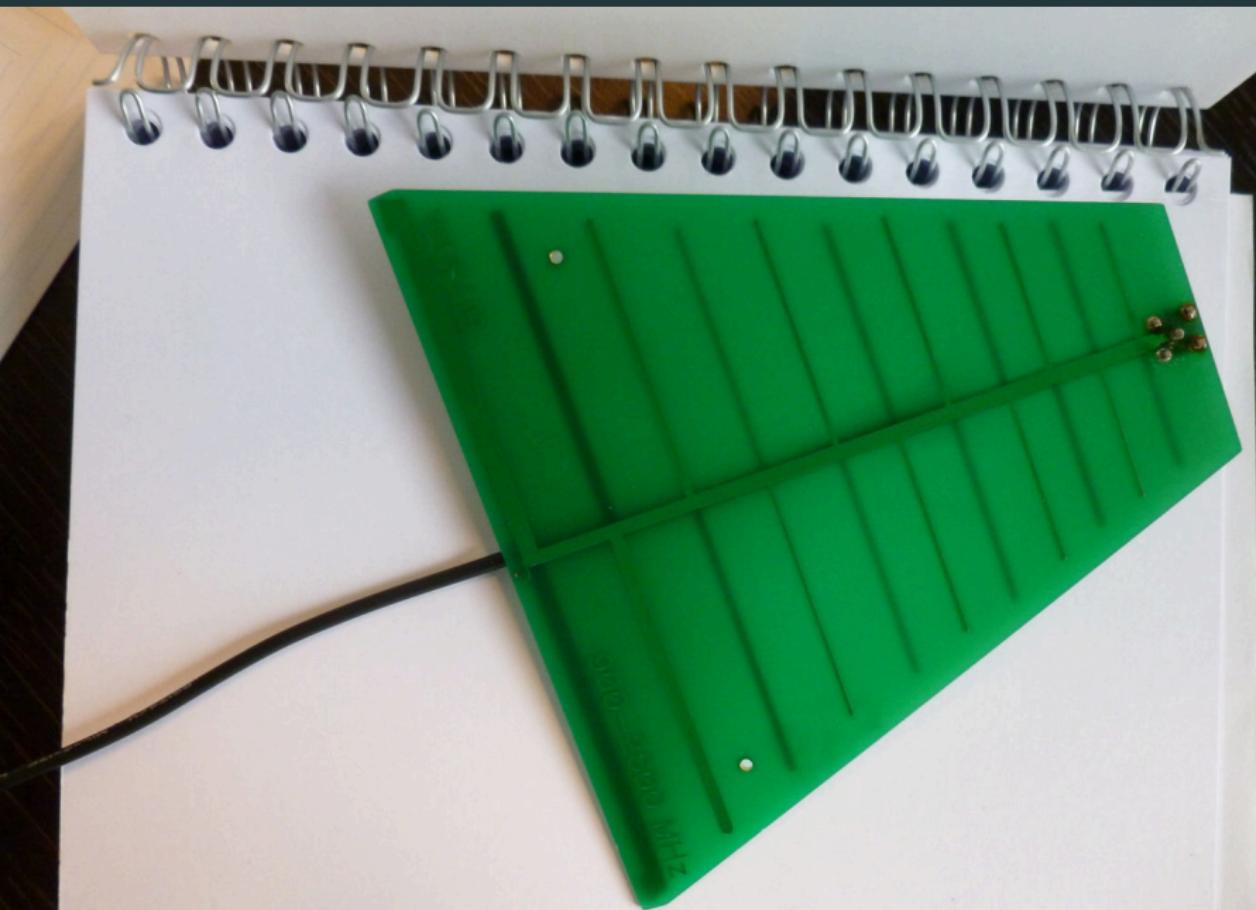




**Demo time**  
(click for video)

Small loops are amazing for under  $\approx 5\text{ cm}$ .  
Won't get us to  $1\text{ m}$  though.

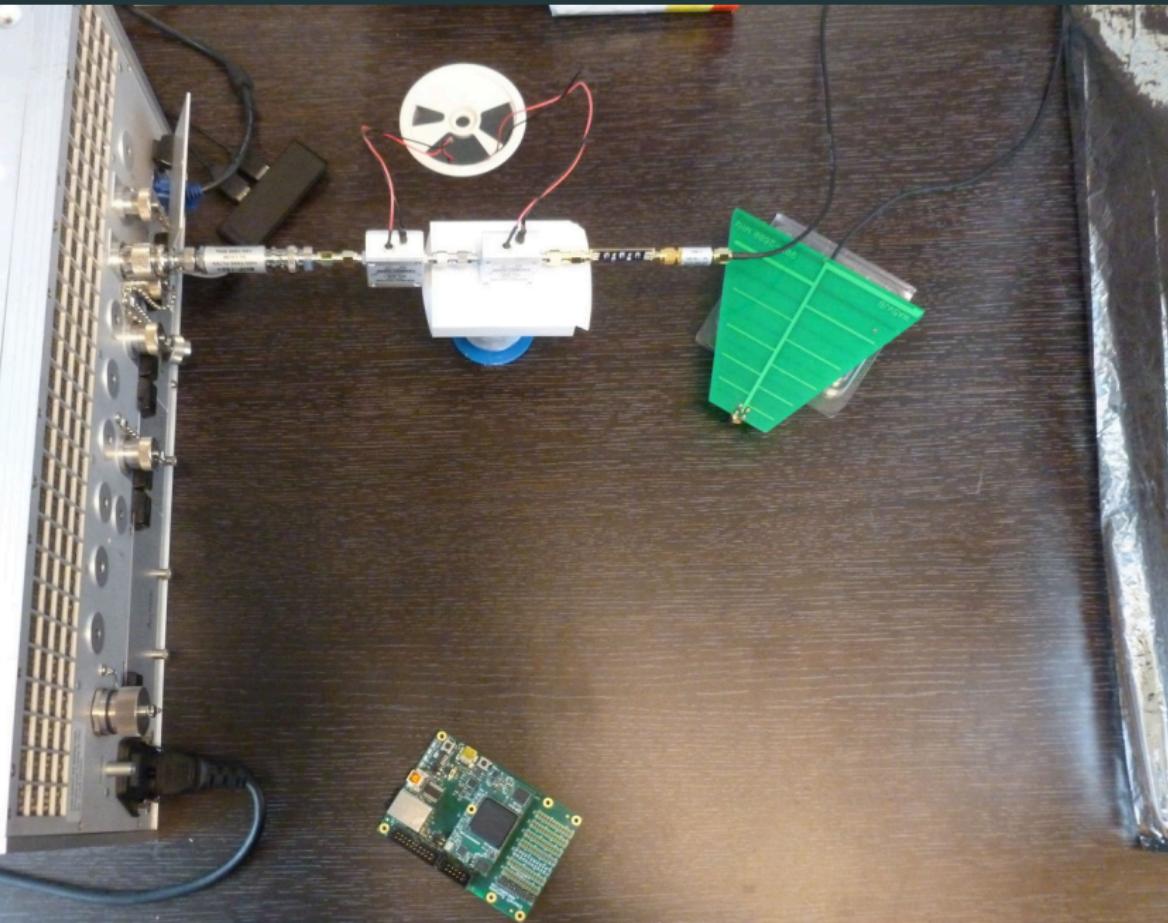
# Log-periodic antenna



## Log-periodic distance



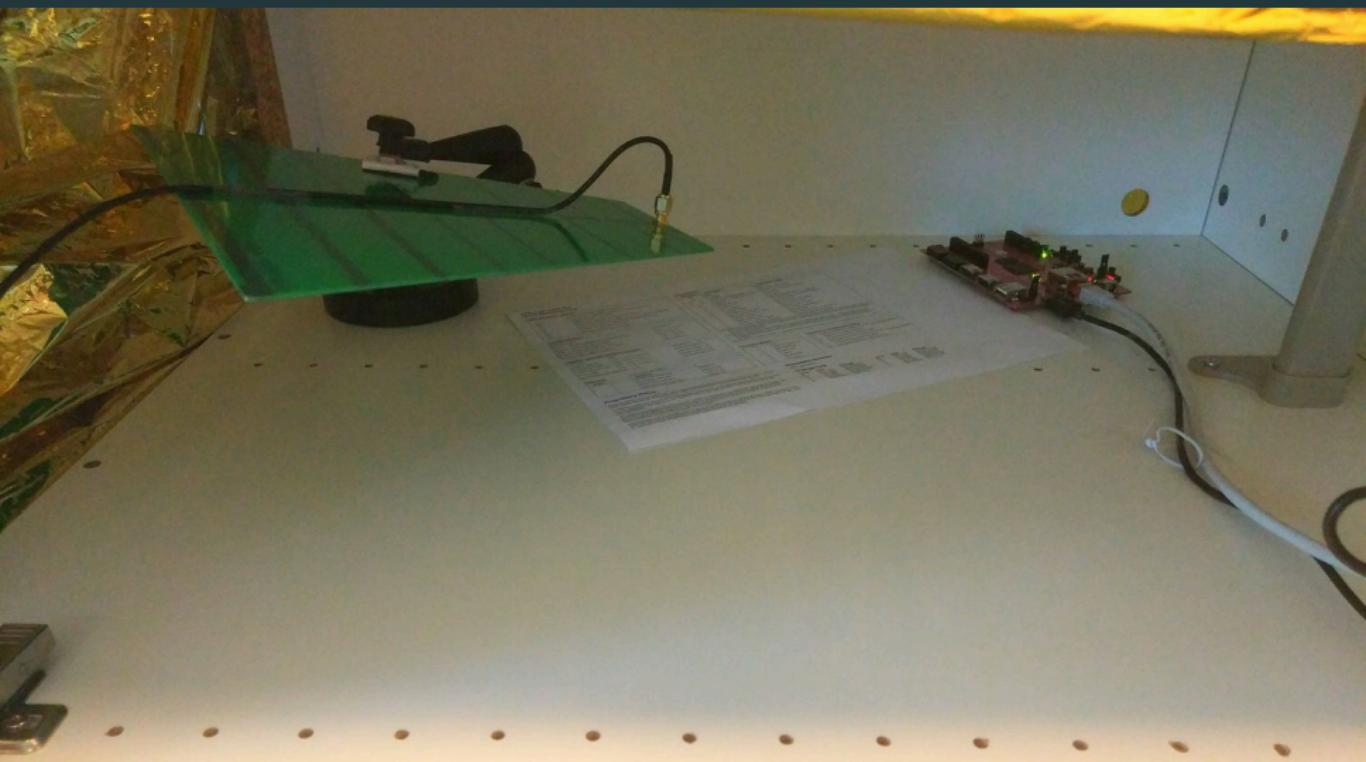
# Example setup

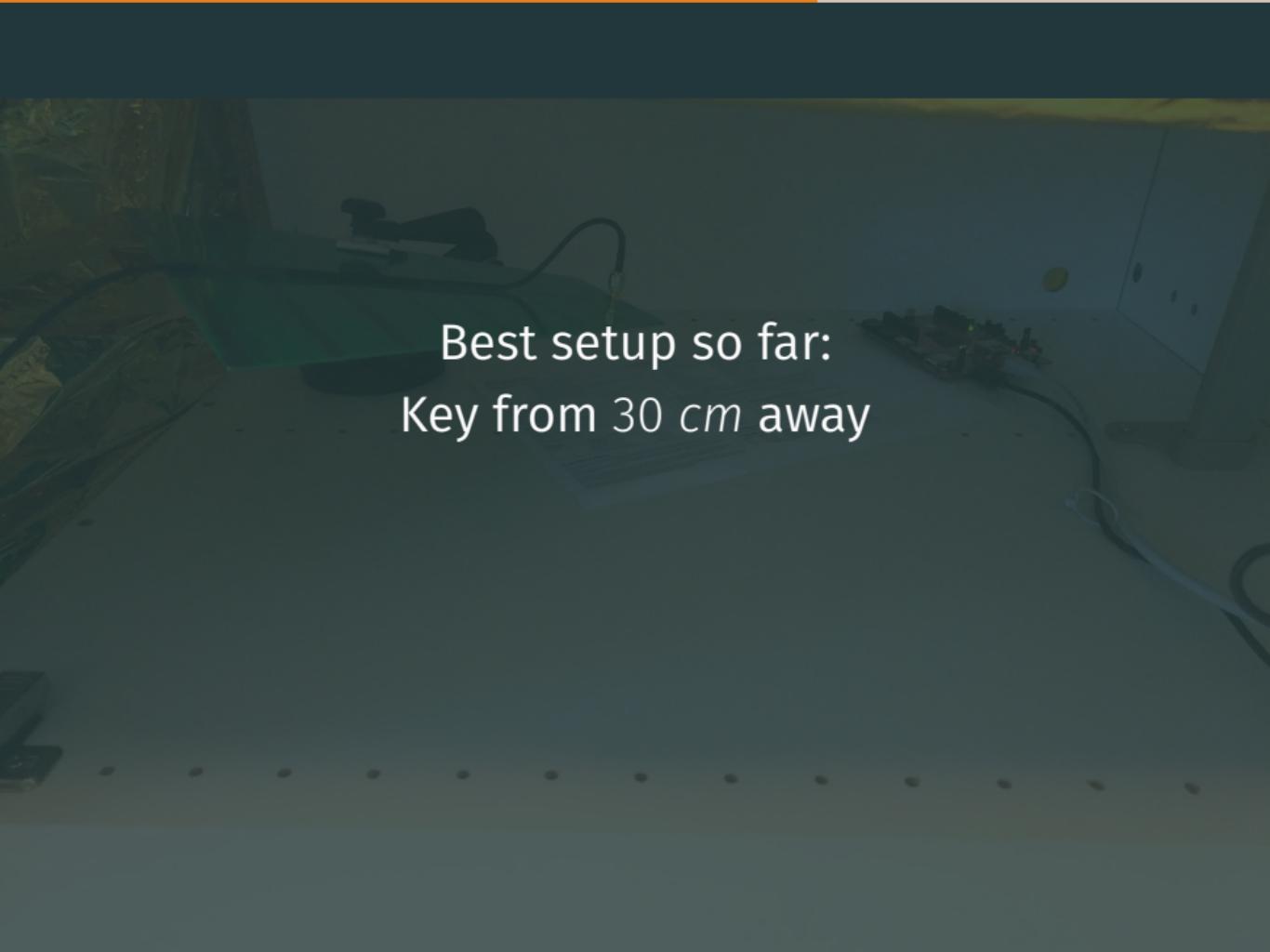


# DIY shielding

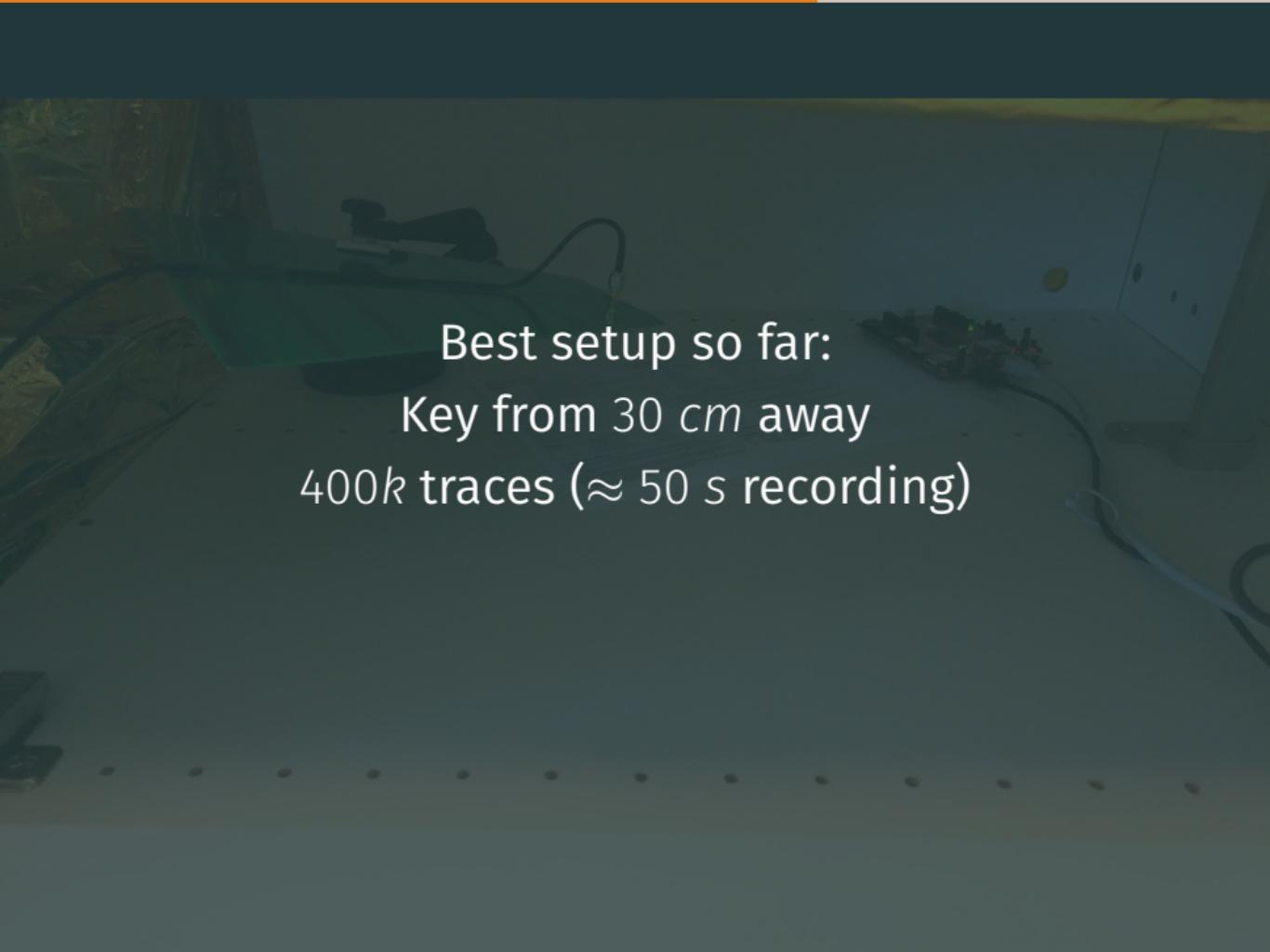


# Real setup

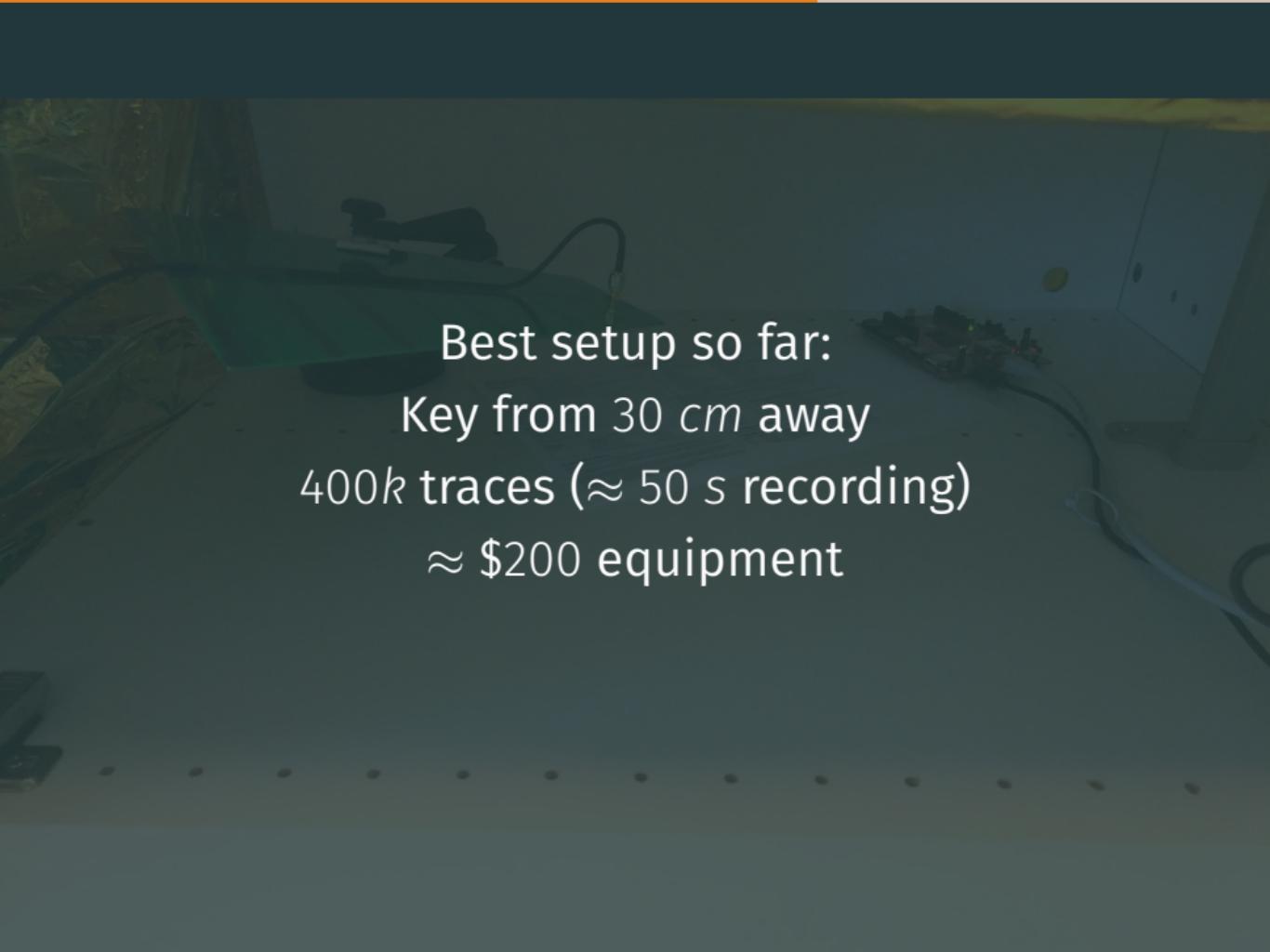




Best setup so far:  
Key from 30 cm away

A dark, low-light photograph of a computer keyboard. A small electronic device with a red LED and several wires is connected to one of the keyboard's ports. The background is mostly black.

Best setup so far:  
Key from 30 cm away  
 $400k$  traces ( $\approx 50$  s recording)

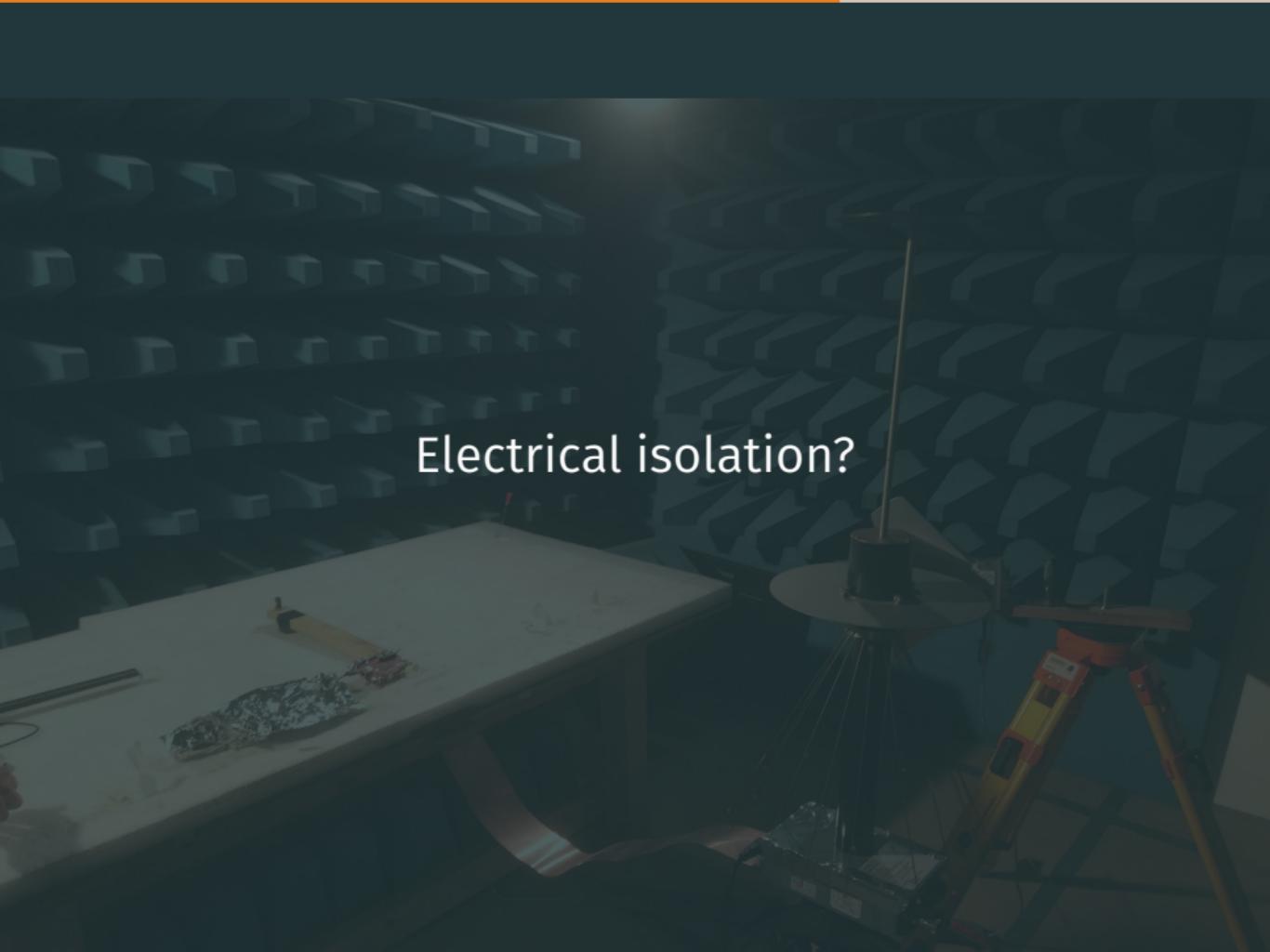
A dark, low-light photograph showing a person lying on their stomach on a bed. A laptop is open and positioned in front of them. To the right, there's some electronic equipment with cables. The overall scene is dimly lit.

Best setup so far:  
Key from 30 cm away  
400k traces ( $\approx$  50 s recording)  
 $\approx$  \$200 equipment

...and in ideal conditions?  
(Thanks, OSPL)

# Anechoic Chamber





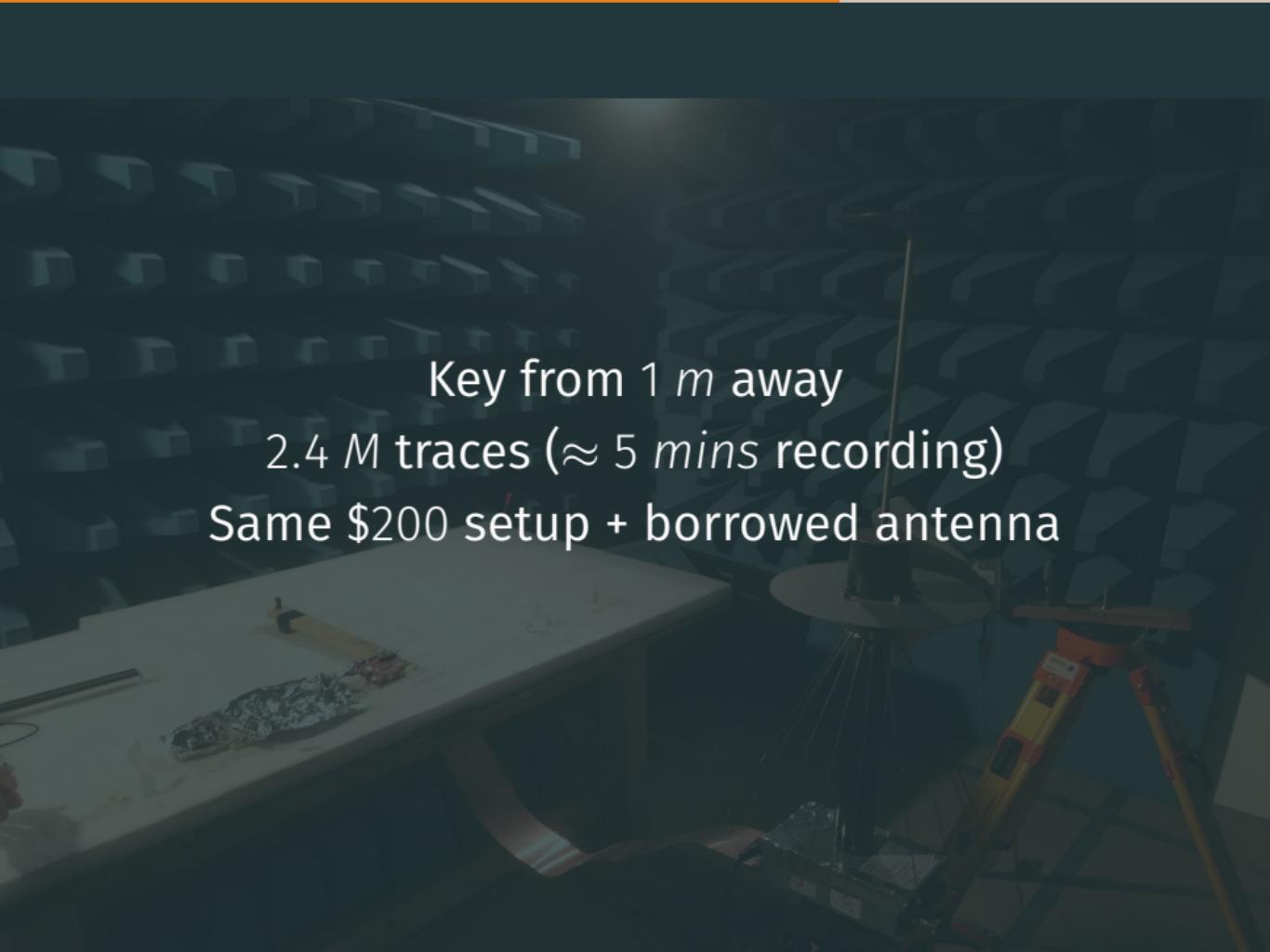
Electrical isolation?

Key from 1 m away



Key from 1 m away  
2.4 M traces ( $\approx$  5 mins recording)





Key from 1 m away  
2.4 M traces ( $\approx$  5 mins recording)  
Same \$200 setup + borrowed antenna

## Conclusion

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## Conclusion

- Break OpenSSL's AES with a wire and a \$20 dongle
- Radio hardware → *really* speeds up attack
- Increase attack distance with new analogue front-ends
  - First known demonstration
- 1 m works in 5 minutes...

Thanks! Questions?



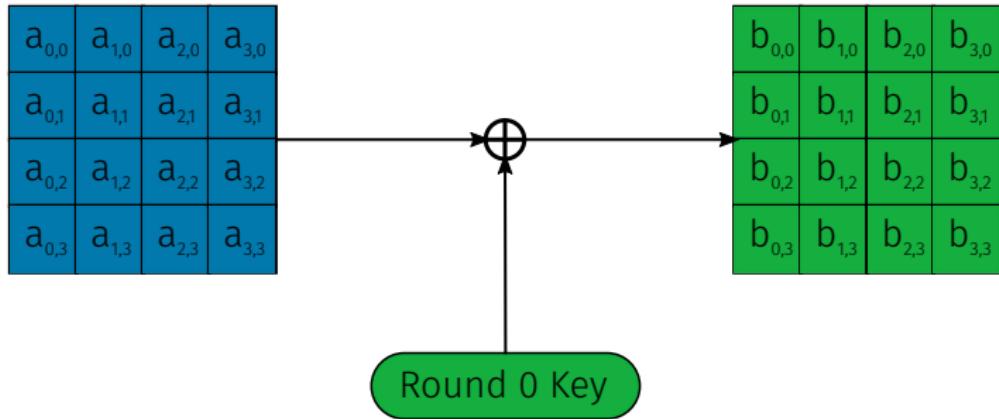
## Backup slides

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## Selecting an intermediate

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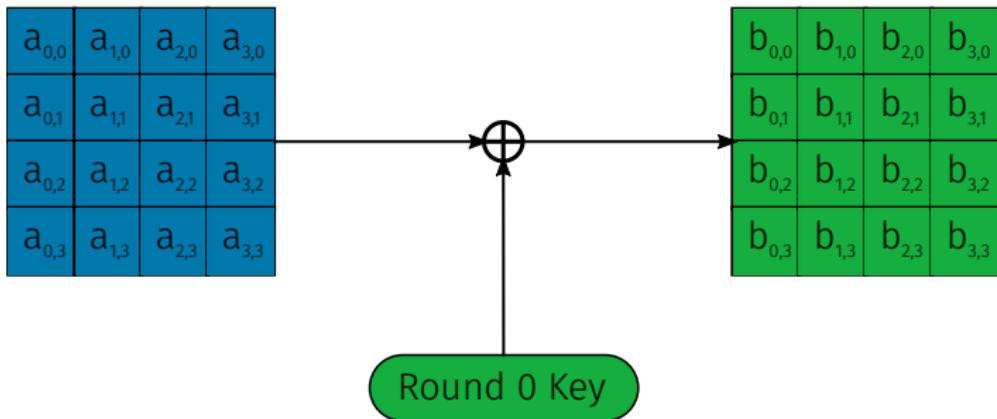
# OpenSSL AES Round 0



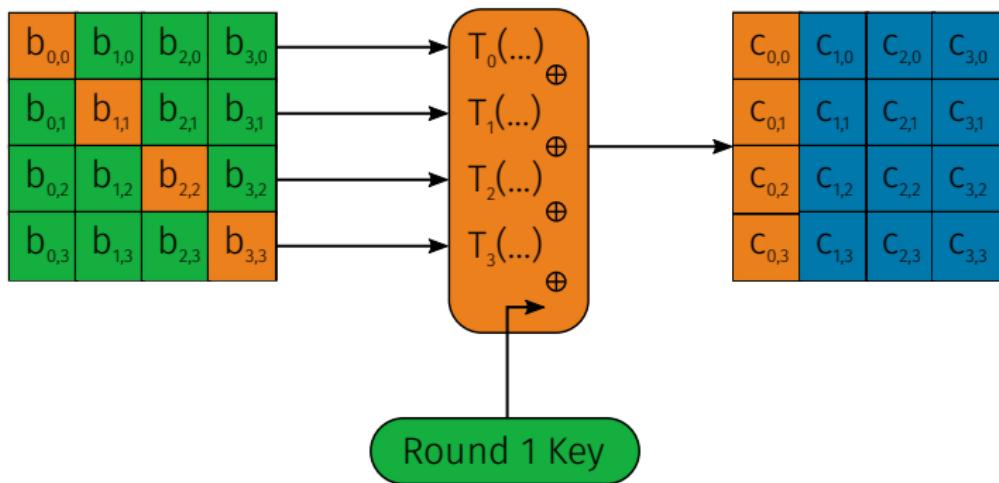
# OpenSSL AES Round 0

Attack XOR with key?

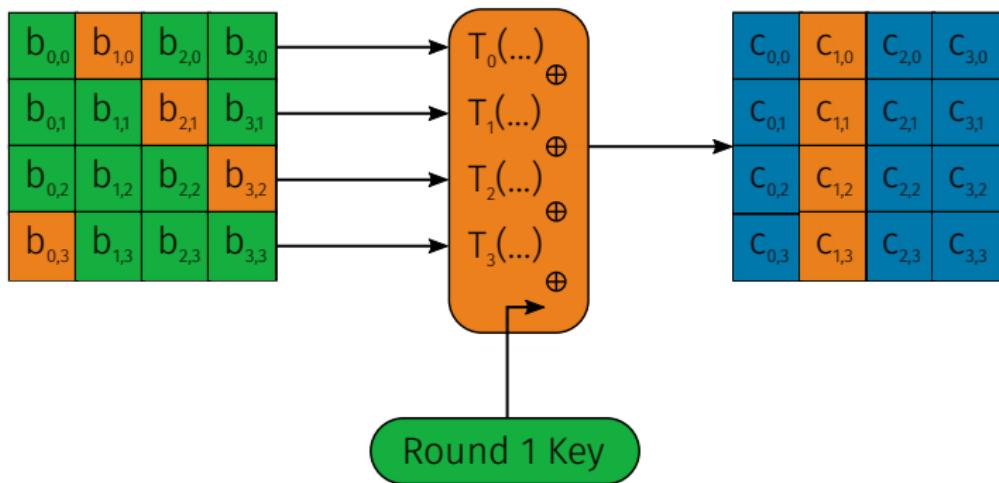
...Can do better!



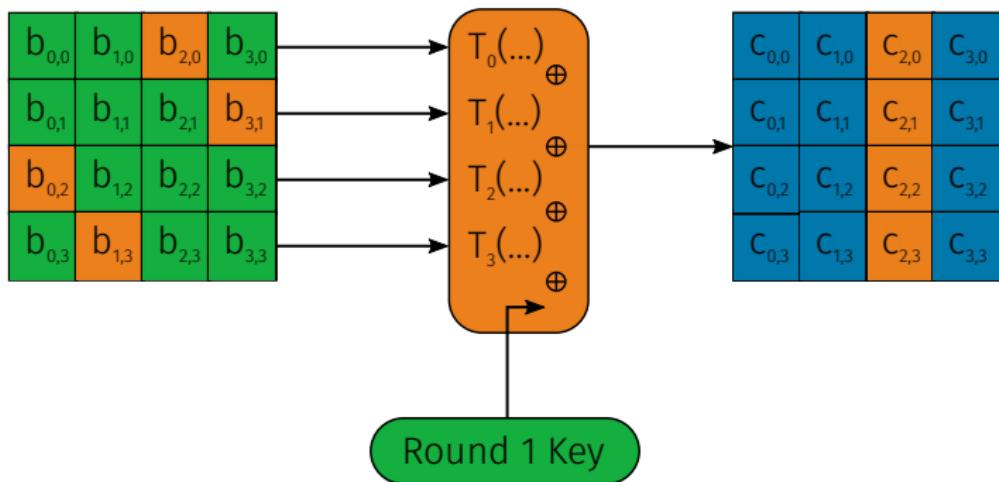
# OpenSSL AES Round 1



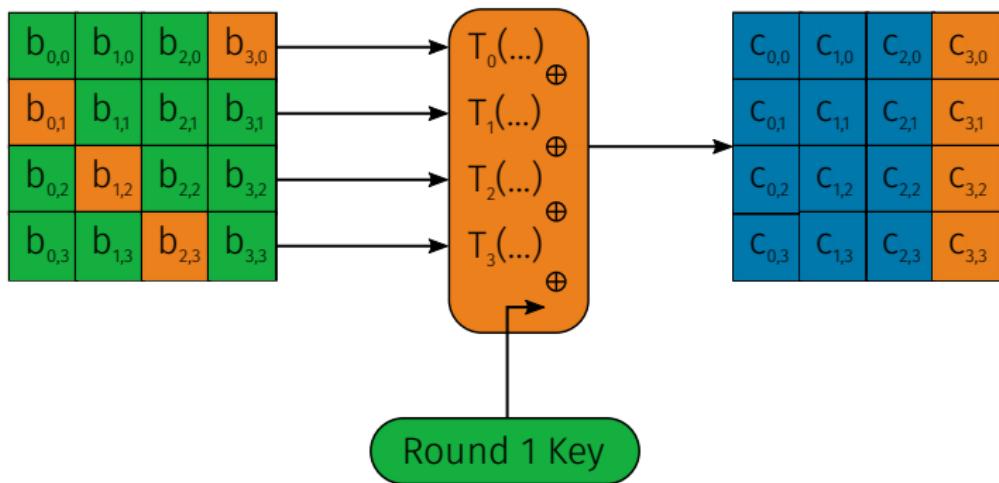
# OpenSSL AES Round 1



# OpenSSL AES Round 1



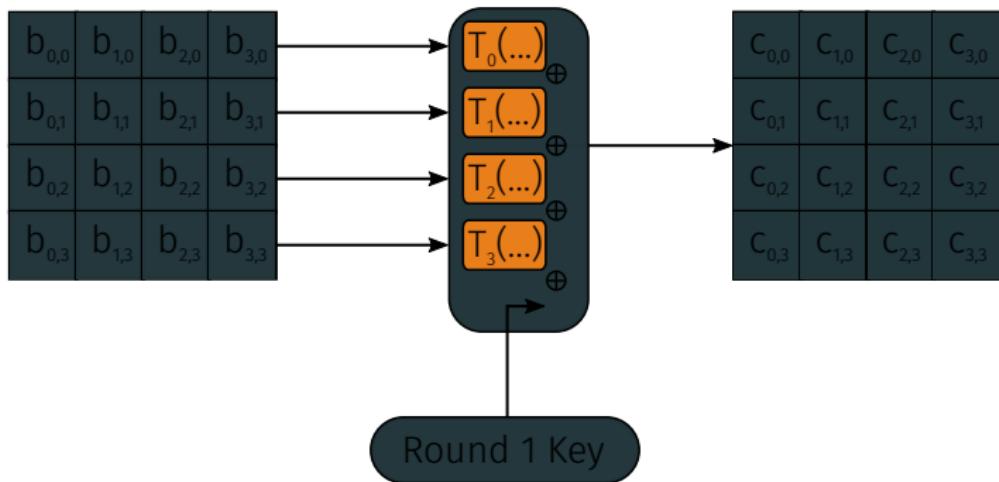
# OpenSSL AES Round 1



# OpenSSL AES Round 1

Attack these lookups.

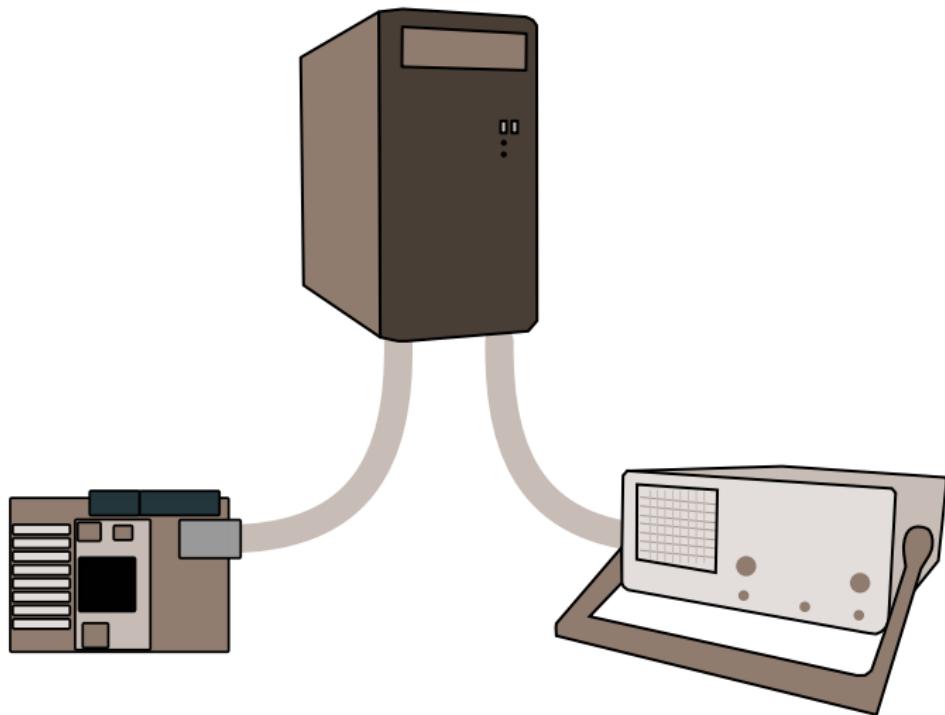
The non-linearity is useful.



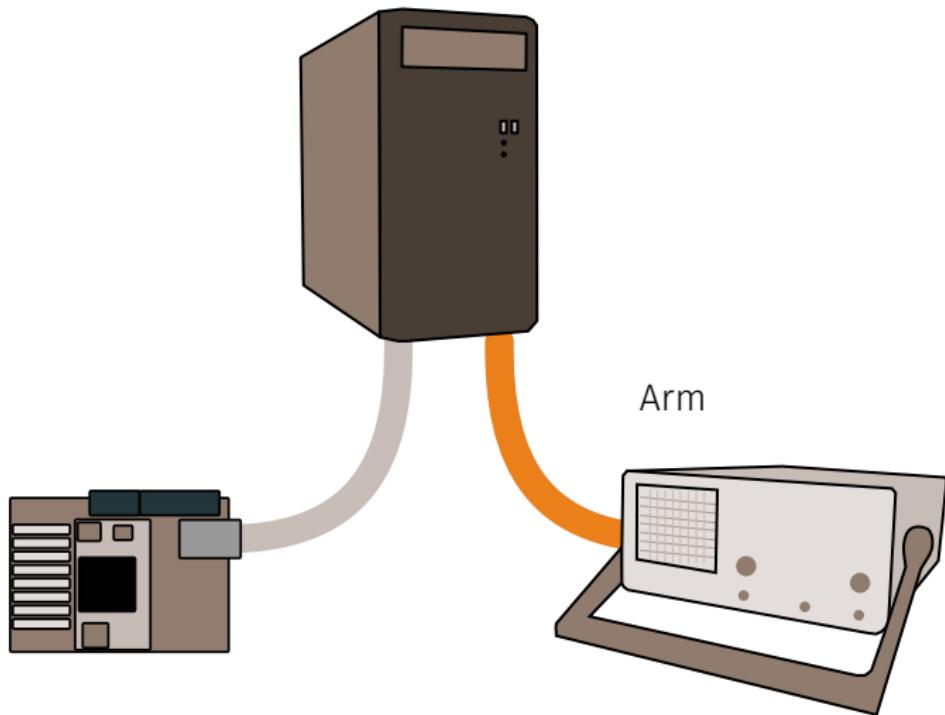
## Our setup vs traditional setup

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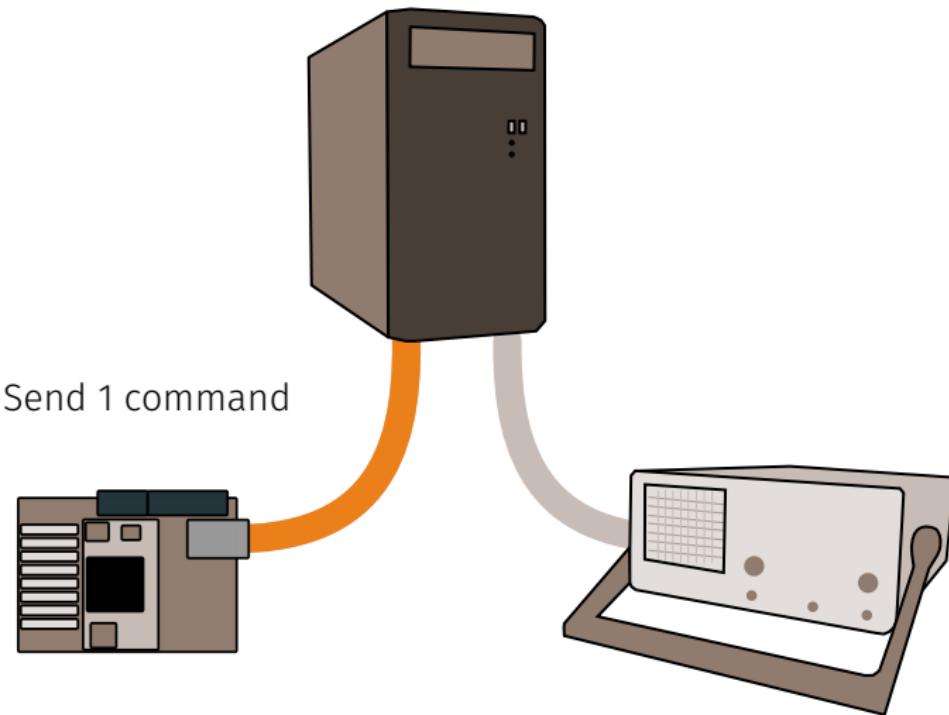
# Recording comparison – 'Scope



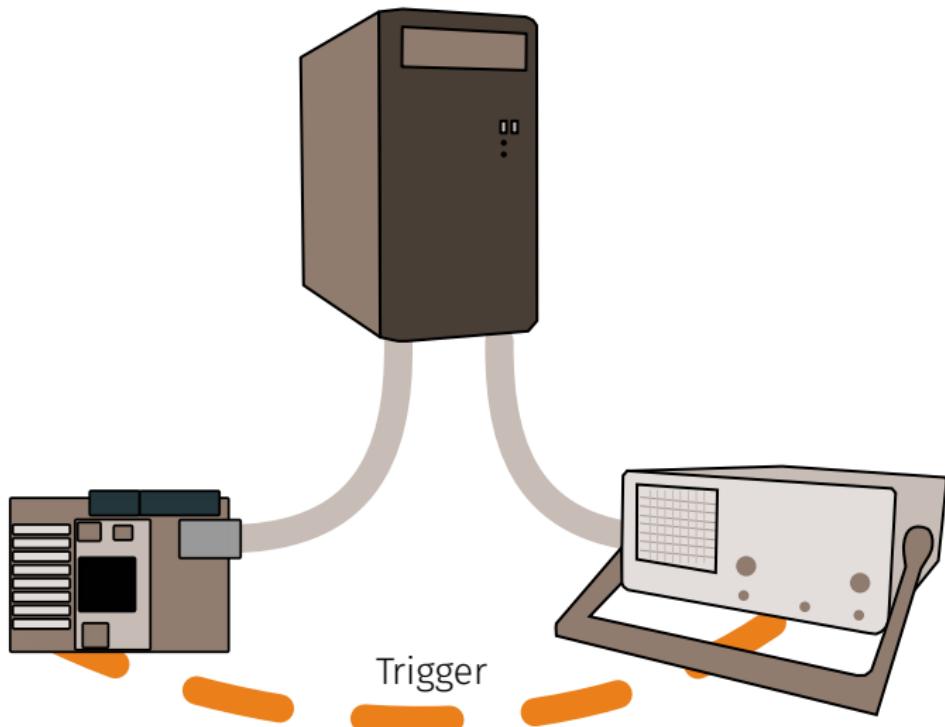
## Recording comparison – 'Scope



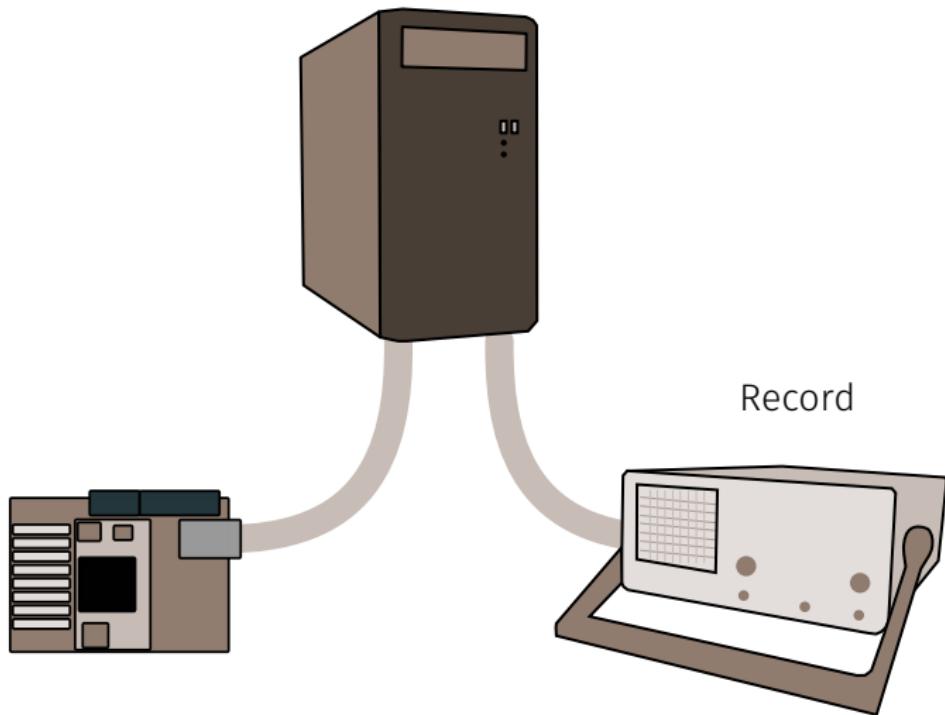
## Recording comparison – 'Scope



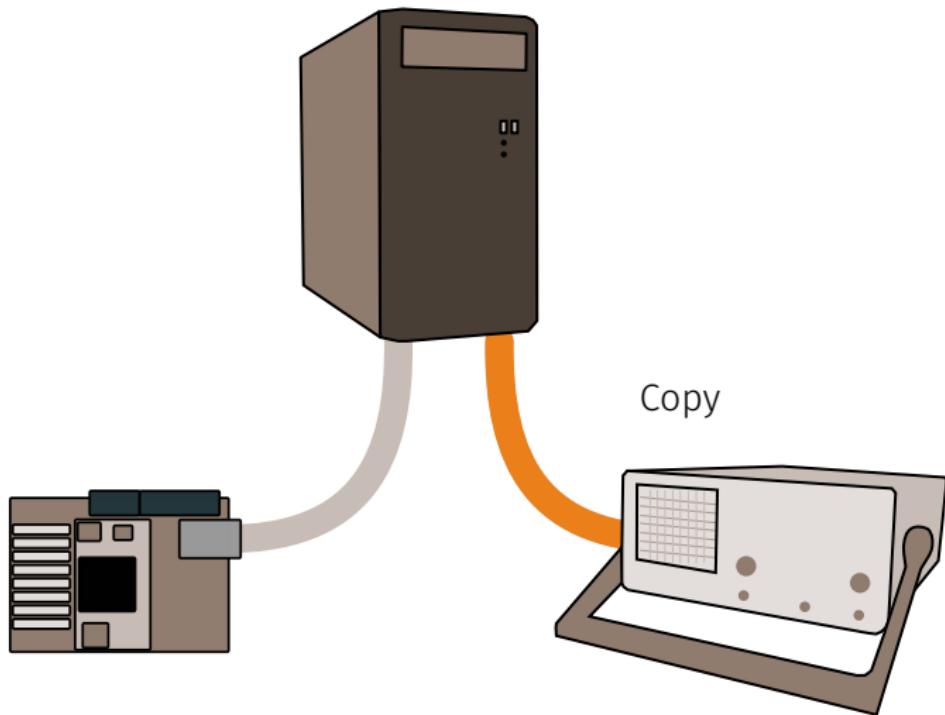
## Recording comparison – 'Scope



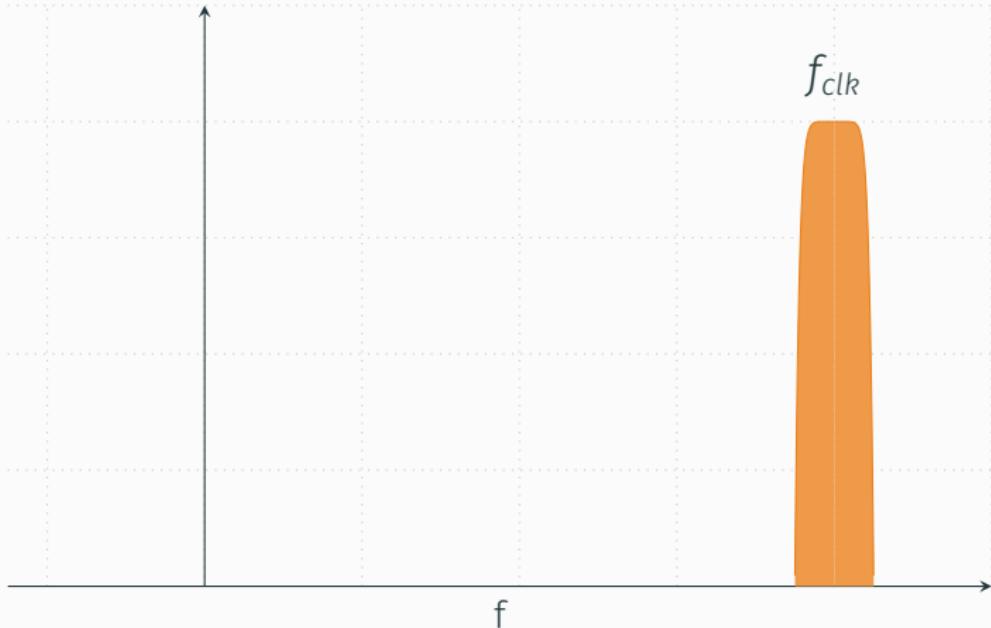
# Recording comparison – 'Scope



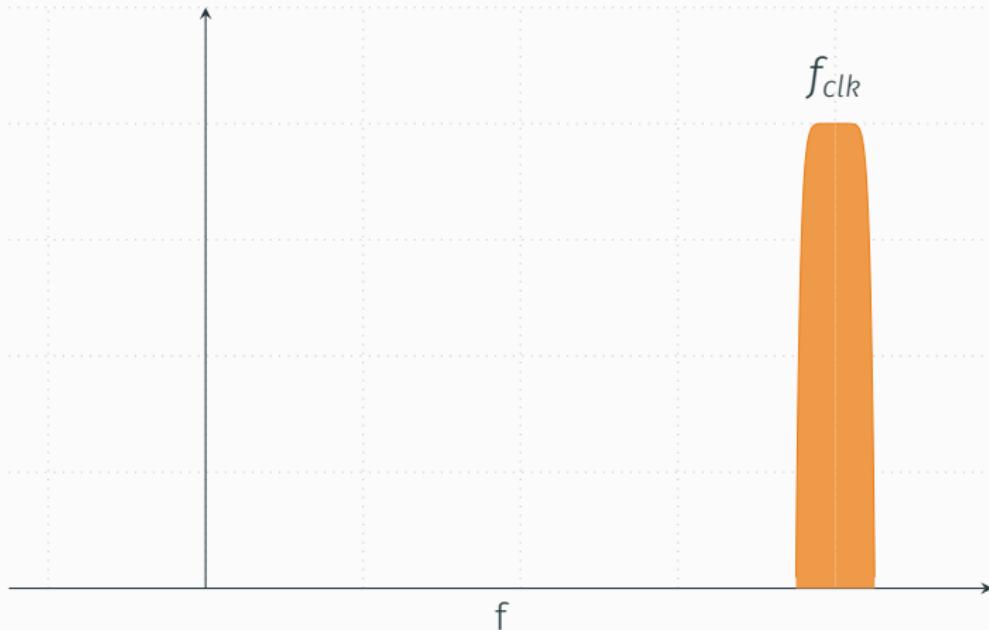
## Recording comparison – 'Scope



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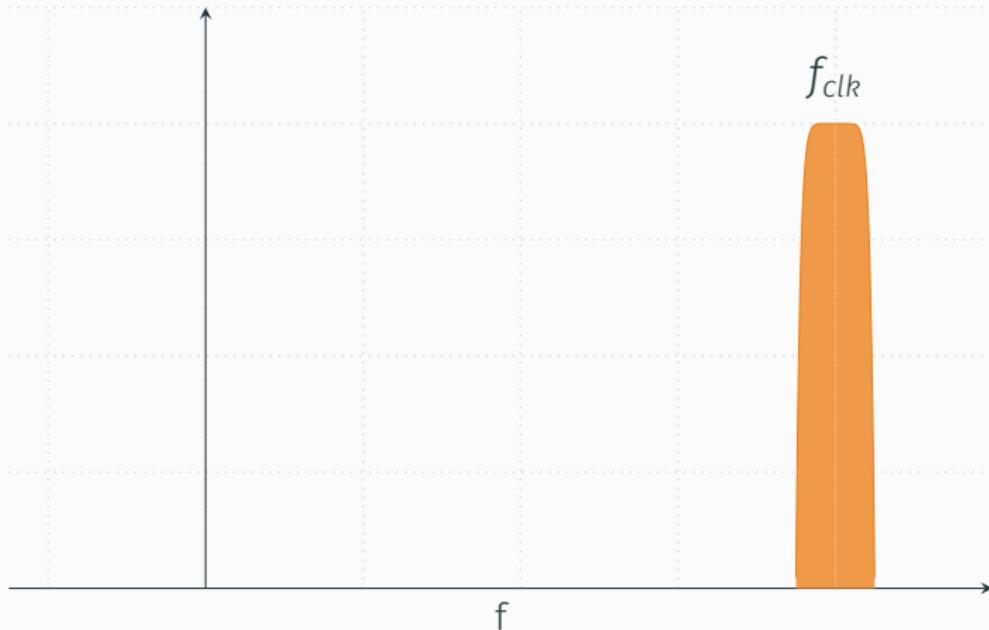


## Recording comparison – 'Scope

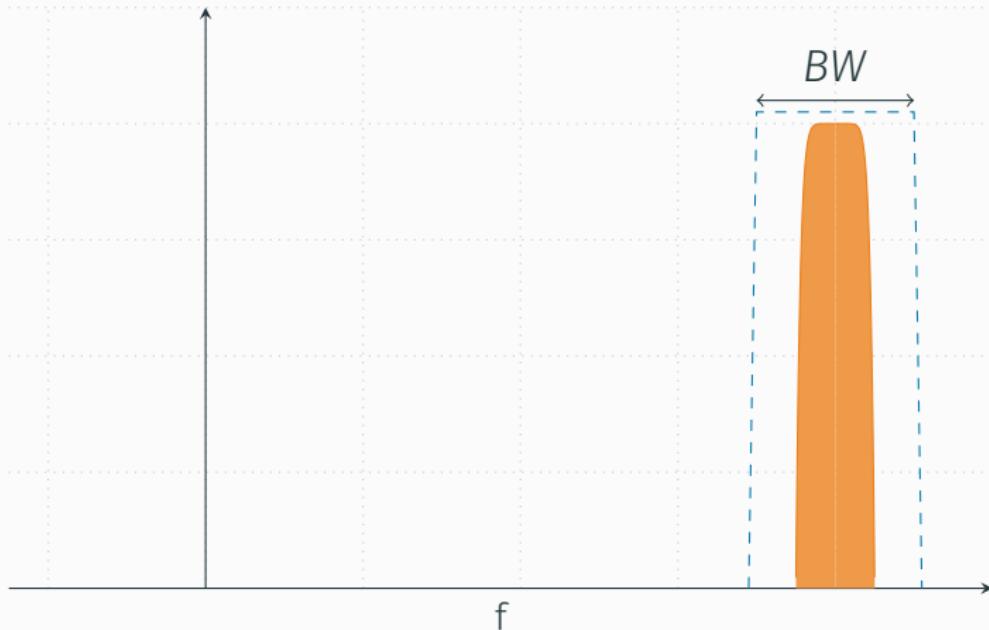


$$f_s > 2f_{clk}$$

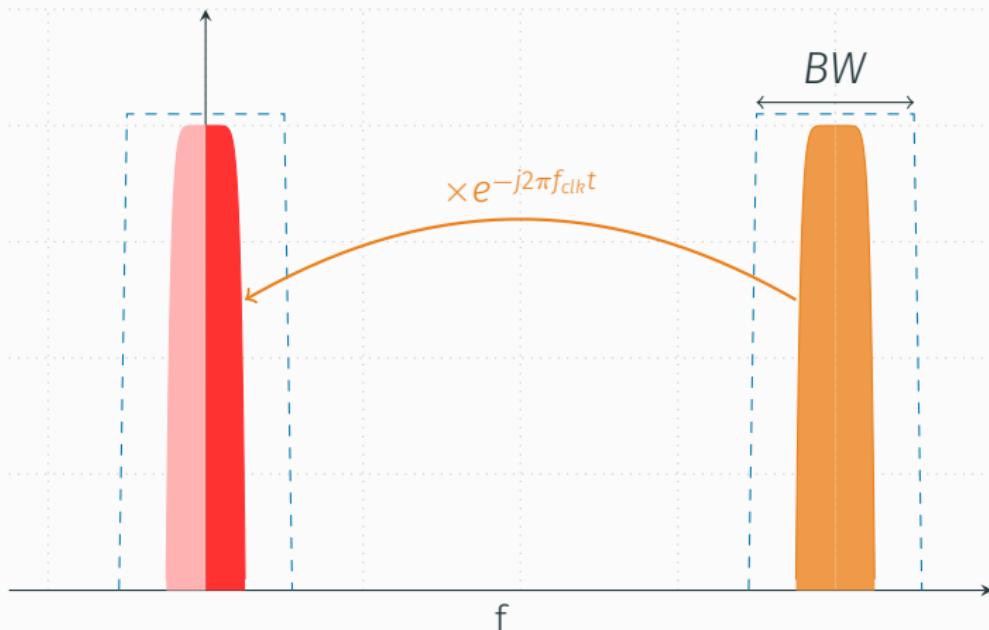
## Recording comparison – Radio



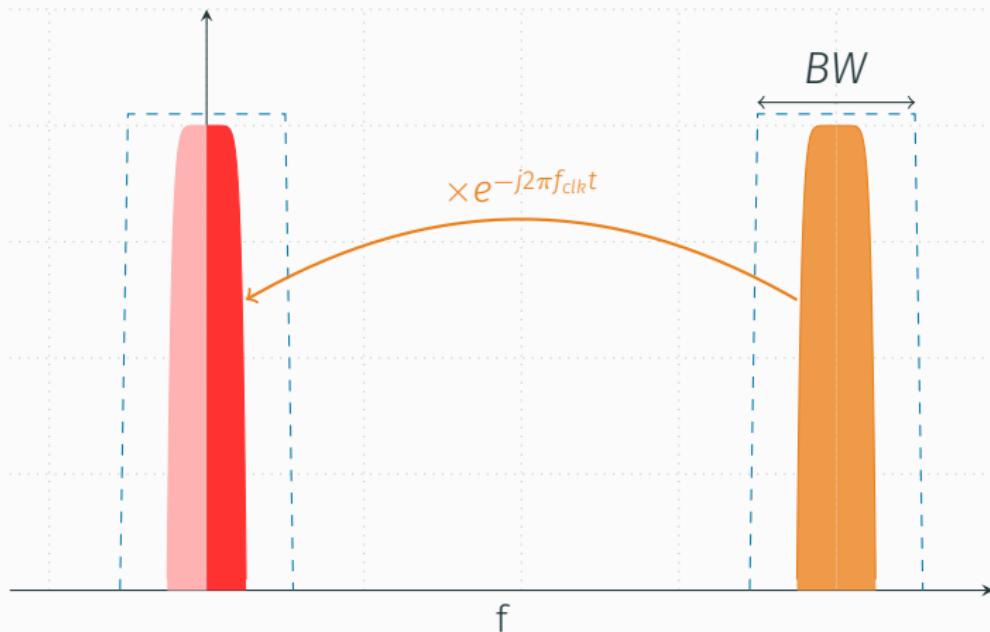
# Recording comparison – Radio



# Recording comparison — Radio

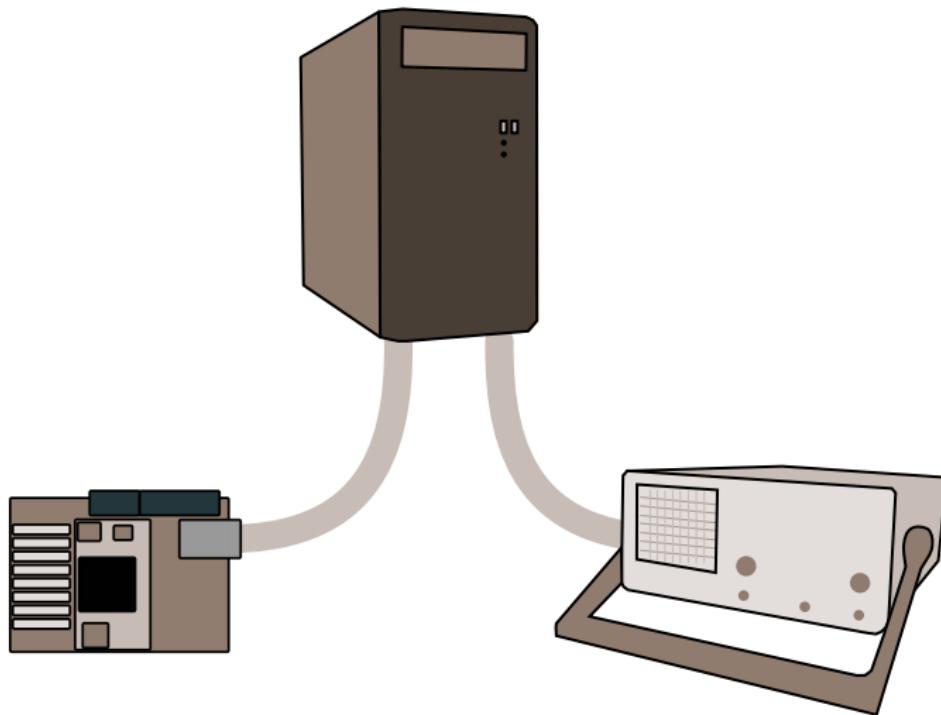


# Recording comparison – Radio

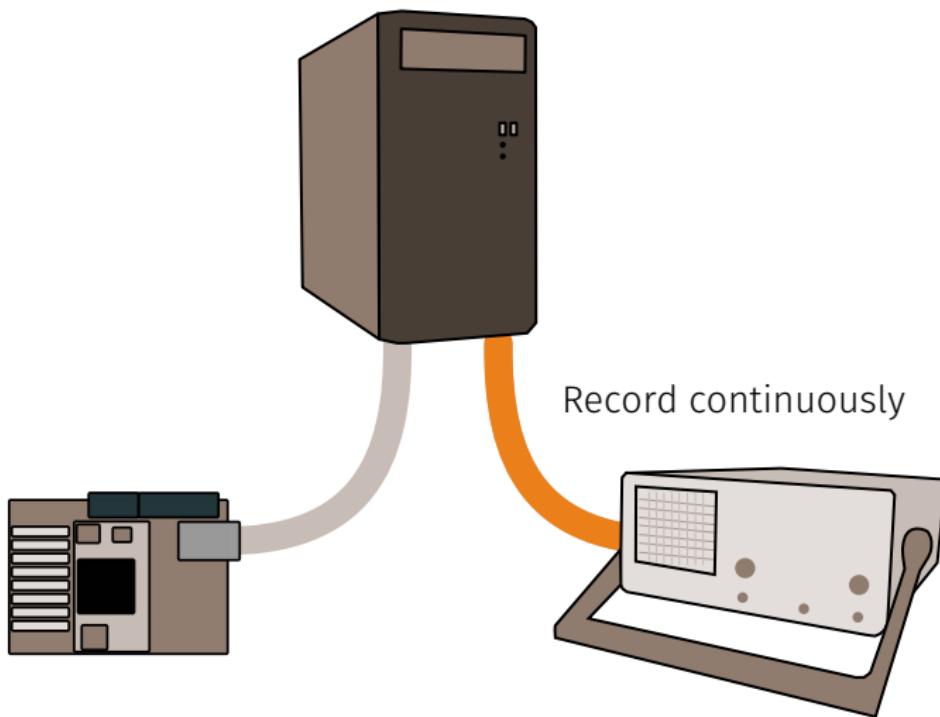


$$f_s > 2 \times BW$$

# Recording comparison — Radio



## Recording comparison – Radio



# Recording comparison – Radio

