phull: p-hull in R

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WZUR, Warsaw, October 4, 2009



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Outline

p-hull and its properties

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Given an arbitrary $0 , <math>x_0, y_0 \in \mathbb{R}$, $a \ge 0$ and $b \ge 0$, let

$$E_{p,a,b}^{(x_0,y_0)} = \left\{ (x,y) \in \mathbb{R}^2 : \left| \frac{y - y_0}{b} \right|^p + \left| \frac{x - x_0}{a} \right|^p \le 1 \right\}. \tag{1}$$

Moreover, for $p = \infty$ we have

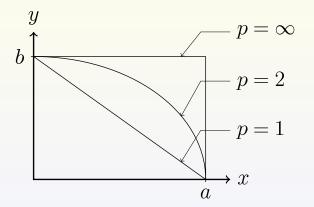
$$E_{p,a,b}^{(x_0,y_0)} = \left\{ (x,y) \in \mathbb{R}^2 : \max\left\{ \left| \frac{y - y_0}{b} \right|, \left| \frac{x - x_0}{a} \right| \right\} \le 1 \right\}, \quad (2)$$

and for p=0

$$E_{p,a,b}^{(x_0,y_0)} = \left\{ (x,y) \in \mathbb{R}^2 : \begin{array}{c} x \in [x_0 - a, x_0 + a] & \wedge & y = y_0 \\ y \in [y_0 - b, y_0 + b] & \wedge & x = x_0 \end{array} \right\}.$$
(3)

We call $E_{p,a,b}^{(x_0,y_0)}$ the *p*-ellipse of size (a,b) centered at (x_0,y_0) .

Illustration: $\partial E_{p,a,b}^{(0,0)} \cap \mathbb{R}_0^+ \times \mathbb{R}_0^+$.



We are given a finite planar set $Q = \{q_1, q_2, \dots, q_n\}$, such that $q_i = (x_i, y_i) \in \mathbb{R}^2$, $i = 1, \dots, n$ $(n \ge 4)$.

Let

$$x_1 = \min_{p_i \in P} x_i,$$

$$x_r = \max_{p_i \in P} x_i,$$

$$y_b = \min_{p_i \in P} y_i,$$

$$y_t = \max_{p_i \in P} y_i.$$

Then $B(Q) = [x_{\rm l}, x_{\rm r}] \times [y_{\rm t}, y_{\rm b}]$ is the minimal bounding rectangle of Q.

For a fixed $p \ge 0$ let

$$\begin{array}{lcl} C_p^{\mathrm{bl}}(Q) & = & \bigcup_{a,b: \ Q\not\in \mathrm{int} \ E_{p,a,b}^{(x_1,y_b)} \ E_{p,a,b}^{(x_1,y_b)}, \\ C_p^{\mathrm{br}}(Q) & = & \bigcup_{a,b: \ Q\not\in \mathrm{int} \ E_{p,a,b}^{(x_r,y_b)} \ E_{p,a,b}^{(x_r,y_b)}, \\ C_p^{\mathrm{tr}}(Q) & = & \bigcup_{a,b: \ Q\not\in \mathrm{int} \ E_{p,a,b}^{(x_r,y_t)} \ E_{p,a,b}^{(x_r,y_t)}, \\ C_p^{\mathrm{tl}}(Q) & = & \bigcup_{a,b: \ Q\not\in \mathrm{int} \ E_{p,a,b}^{(x_1,y_t)} \ E_{p,a,b}^{(x_1,y_t)}. \end{array}$$

We further on assume int $C_p^{\mathrm{bl}}(Q)$, int $C_p^{\mathrm{br}}(Q)$, int $C_p^{\mathrm{tr}}(Q)$, int $C_p^{\mathrm{tr}}(Q)$ are mutually exclusive.

p-hull

Definition

Let $Q=\{q_1,q_2,\ldots,q_n\}\subset\mathbb{R}^2$ and $p\geq 0$. The p-hull of Q, denoted by $H_p(Q)$, is defined by

$$H_p(Q) = \partial \left(B(Q) \setminus C_p^{\text{bl}}(Q) \setminus C_p^{\text{br}}(Q) \setminus C_p^{\text{tr}}(Q) \setminus C_p^{\text{tl}}(Q) \right). \tag{4}$$

Properties of a p-hull

Proposition

Let $Q = \{q_1, q_2, \dots, q_n\} \subset \mathbb{R}^2$ and $p \geq 0$. Then we have the following.

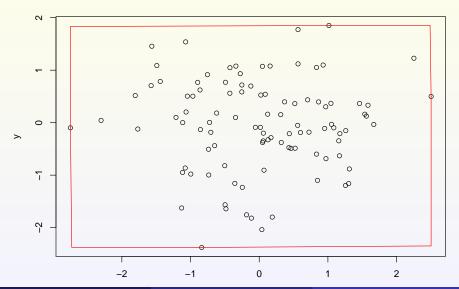
- If p=1 then $H_p(Q)$ is the convex hull of Q.
- ② If $p = \infty$ then $H_p(Q)$ is the X-Y hull of Q (see Nicholl et al, 1983).
- $If p = 0 then H_p(Q) = \partial B(Q).$

Properties of a p-hull

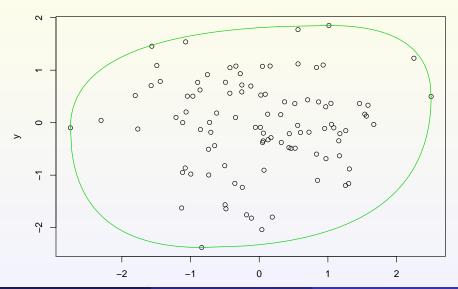
Other properties:

- $H_p(Q)$ is translation- and scale-invariant for any $p \ge 0$, i.e. $\circledast H_p(Q) = H_p(\circledast Q)$.
- ② $H_p(Q)$ is not rotation-invariant (thus it is orientation-dependent) for $p \neq 1$.
- **3** $H_p(Q)$ is convex for $p \leq 1$.

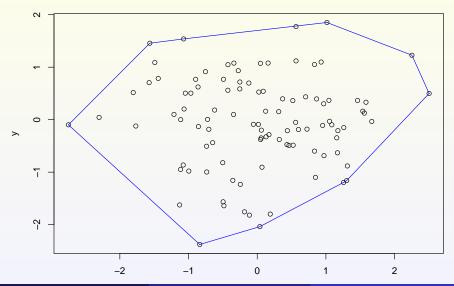
Example: p = 0.1



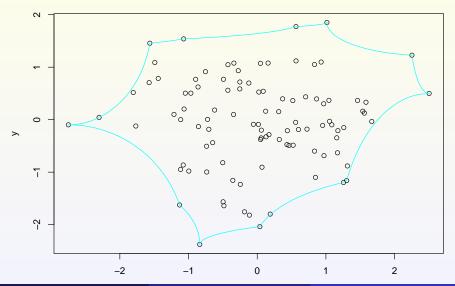
Example: p = 0.5



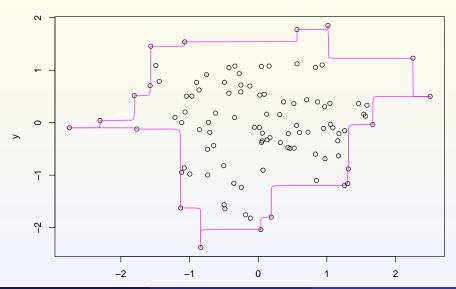
Example: p = 1.0



Example: p = 2.0



Example: p = 50



Applications

Among possible applications are:

- Scientometrics: calculation of so-called L_p indices (Gągolewski, Grzegorzewski, 2009),
- ② p.d.f. support estimation: given a random sample of points from distribution given by unknown p.d.f. f(x,y), estimate $\mathrm{supp}\, f$,
- **3** ...

Computation

Let

$$\begin{array}{lll} q_{\rm bl_1} & = & \displaystyle \mathop{\arg\min}_{q_i \in Q: \ x_i = x_1} y_i, & q_{\rm bl_2} & = & \displaystyle \mathop{\arg\min}_{q_i \in Q: \ y_i = y_{\rm b}} x_i, \\ q_{\rm br_1} & = & \displaystyle \mathop{\arg\max}_{q_i \in Q: \ y_i = y_{\rm b}} x_i, & q_{\rm br_2} & = & \displaystyle \mathop{\arg\min}_{q_i \in Q: \ x_i = x_{\rm r}} y_i, \\ q_{\rm tr_1} & = & \displaystyle \mathop{\arg\max}_{q_i \in Q: \ x_i = x_{\rm r}} x_i, & q_{\rm tr_2} & = & \displaystyle \mathop{\arg\max}_{q_i \in Q: \ y_i = y_{\rm t}} x_i, \\ q_{\rm tl_1} & = & \displaystyle \mathop{\arg\min}_{q_i \in Q: \ y_i = y_{\rm t}} x_i, & q_{\rm tl_2} & = & \displaystyle \mathop{\arg\max}_{q_i \in Q: \ x_i = x_{\rm l}} y_i. \end{array}$$

Note that all the points $\in \partial B(Q)$.

Decomposition:

$$H_{p}(Q) = \partial \left(B(Q) \setminus C_{p}^{\text{bl}}(Q) \setminus C_{p}^{\text{br}}(Q) \setminus C_{p}^{\text{tr}}(Q) \setminus C_{p}^{\text{tl}}(Q) \right)$$

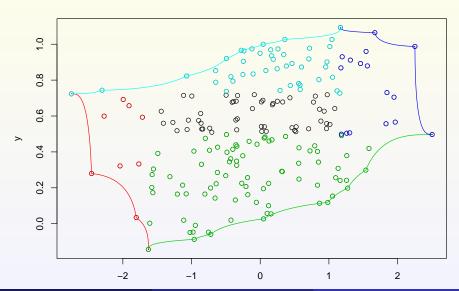
$$= \left(\partial C_{p}^{\text{bl}}(Q) \cup \partial C_{p}^{\text{br}}(Q) \cup \partial C_{p}^{\text{tr}}(Q) \cup \partial C_{p}^{\text{tl}}(Q) \right) \cap B(Q)$$

$$\cup \overline{q_{\text{bl}_{2}}q_{\text{br}_{1}}} \cup \overline{q_{\text{br}_{2}}q_{\text{tr}_{1}}} \cup \overline{q_{\text{tr}_{2}}q_{\text{tl}_{1}}} \cup \overline{q_{\text{tl}_{2}}q_{\text{bl}_{1}}}.$$

$$(5)$$

Moreover:

$$\partial C_{p}^{\text{bl}}(Q) = \partial C_{p}^{\text{bl}}(\{q_{i} \in Q : x_{i} \leq x_{\text{bl}_{2}} \land y_{i} \leq y_{\text{bl}_{1}}\}),
\partial C_{p}^{\text{br}}(Q) = \partial C_{p}^{\text{br}}(\{q_{i} \in Q : x_{i} \geq x_{\text{br}_{1}} \land y_{i} \leq y_{\text{br}_{2}}\}),
\partial C_{p}^{\text{tr}}(Q) = \partial C_{p}^{\text{tr}}(\{q_{i} \in Q : x_{i} \geq x_{\text{tr}_{2}} \land y_{i} \geq y_{\text{tr}_{1}}\}),
\partial C_{p}^{\text{tl}}(Q) = \partial C_{p}^{\text{tl}}(\{q_{i} \in Q : x_{i} \leq x_{\text{tl}_{1}} \land y_{i} \geq y_{\text{tl}_{2}}\}).$$
(6)



- Naïve algorithm: $O(n^3)$ time. :-(
- 2 Algorithm proposed by Gągolewski, Nowakiewicz, Dębski (2009) generalizes Graham's scan (Graham, 1972): $O(n \log n)$ time, O(n) memory.
- $\textbf{ 0} \textbf{ Output: } P \cap \partial C^{\mathrm{bl}}_p(Q) \textbf{ (without loss of generality)}.$
- Input: $p \ge 0$, $W = \{q_i \in Q: x_i \le x_{\text{bl}_2} \land y_i \le y_{\text{bl}_1}\}$ as an array sorted by x coordinate $w[1], \ldots, w[m]$.
- **5** Denotation: by $E_{p,q_i,q_j}^{(x_0,y_0)}$ we mean an p-ellipse centered at (x_0,y_0) interpolating $q_i \neq q_j$.

```
Create an empty stack S;
   Push w[1] into S;
2
3
   i := 2:
    while (i < n) and (w[i]_u \ge w[1]_u) do
4
5
        i := i + 1:
    Push w[i] into S;
6
    for j = i + 1, i + 2, ..., n do
7
        if (S[\#S]_{u} < w[j]_{u}) then {
8
            while (\# S \ge 2) and (S[\# S - 1] \in E_{n,S[\# S],w[i]}^{(x_{bl},y_{bl})}) do
9
                Pop from S:
10
            Push w[j] into S;
11
12
13
     return S:
```

Implementation: phull 0.1-2 — package available on CRAN. (http://cran.r-project.org/web/packages/phull/index.html)

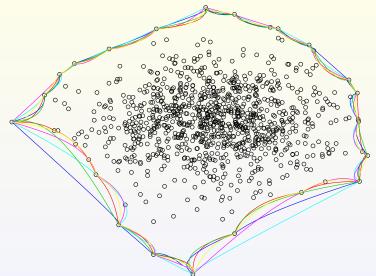
Example: axes rotation.

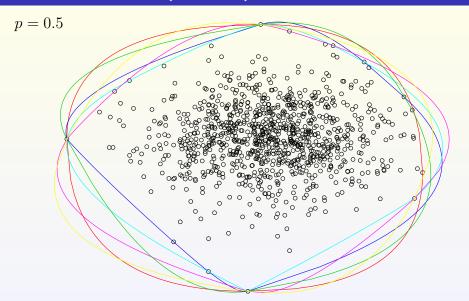
```
library(phull); # load the library
translateAndRotate <- function(data, x0, y0, angle)
{ ... }
rotateAndTranslate <- function(data, x0, y0, angle)
{ ... }</pre>
```

```
set.seed(98765); n <- 1000; p <- 3.0;
data <- matrix(c(rnorm(n), rt(n, 10)), ncol=2); # input data
nres <- 50; # "resolution"</pre>
ptest <- phull(data, p=p);</pre>
                                              # compute the p-hull
discr_0 <- as.matrix(ptest, nres=nres);</pre>
                                              # sample
print(ptest)
        p-hull, p=3
data: data
1000 points, bounding rectangle: (...)
```

```
data2 <- translateAndRotate(data, angle=-pi/6</pre>
   -ptest$xrange[1], -ptest$yrange[1]);
ptest2 <- phull(data2, p=p);</pre>
                                              # compute the p-hull
discr_30 <- as.matrix(ptest2, nres=nres); # sample</pre>
discr_30 <- rotateAndTranslate(discr_30, angle=pi/6,
   ptest$xrange[1], ptest$yrange[1]);
plot(data, type="p", pch=1);
lines(discr_0, col=2);
lines(discr_30, col=4);
... and so on...
```

p = 3





p = 20

Related packages

alphahull (Pateiro-Lopez, Rodriguez-Casal, 2009): α -shapes (Edelsbrunner et al, 1983).



References



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- M. Gągolewski, M. Nowakiewicz, M. Dębski (2009). Efficient algorithms for computing "geometric" scientific impact indices. Submitted for publication.
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- T. M. Nicholl, D. T. Lee, Y. Z. Liao, C. K. Wong (1983). On the X-Y convex hull of a set of X-Y polygons. *BIT* 23, 456–471.
 - B. Pateiro-Lopez, A. Rodriguez-Casal (2009). alphahull: Generalization of the convex hull of a sample of points in the plane, see alphahull @ CRAN.

Thank you for your attention.