

Assume we are at state s and string so far is aw . Let the set of states from s and symbol a be $\{s_1, \dots, s_k\}$ by the transition relation Δ , i.e., $\Delta(s, a) = \{s_1, \dots, s_k\}$. The current configuration is (s, aw) , and we have k next configurations:

$$\begin{aligned} (s, aw) &\rightarrow (s_1, w) \\ &\quad \downarrow \\ &\quad (s_2, w) \\ &\quad \vdots \\ &\quad (s_k, w) \end{aligned}$$

We can push them into a stack for future processing



Then we pop the stack to get the configuration (s_k, w) and repeat the process above

pseudo code sketch

execute:

input: w : a string

output: result : true if accepted by NFA
false otherwise

$\text{result} := \text{false}$; push((q, w))

while stack \neq empty

 pop() \rightarrow Conf

 let conf be (s, w_1) .

 if $w_1 = \epsilon$ & s is final

$\text{result} := \text{true}$

 break

 else if $w_1 = \epsilon$

 then continue

 else if $w_1 = \epsilon$

 then continue

 else let $w_1 = bw_2$

 let $\Delta(s, b) = \{s_1, \dots, s_n\}$

 push((s_1, w_2))

 push((s_n, w_2))

Algorithm ideas for NFA