## CSC 226 - Assignment 4 - Solutions

I. I first did this iteratively, using X t y directly. So, I check if X[i] = y[j] at any [i,j] in the array to determine if les [i,j] came from les [i-1,j-1] or not. (equiv. to a "x" in the path [i,j] porition).

If X[i] = y[j] then I stere X[i] (n y[j]) in s[ix], the substinct to be printed. Otherwise, we need to cleternihe if les [i,j] came from les [i-1,j) of les [i,j-1] (the bigger of the two). If they are liqual it comes from les [i,j-1] by viritue of the Les algorithm on slide 98 of lecture 17.

Algorithm printles (Array 11cs (DE), String X, Stringy)
Input: Array of lengths of LCS of strings X+y.
Ouput: String s, the los of x+y.

n = x. length()

m = y, length()

New string s of length [Ics[m][n]]

k = s. length() - 1

i = n

j = m

while (i >0 && j >0) do

if x[i] = y[j] then

s[k] = x[i]

i+i-1; j+j-1; k+k-1

else if ||cs[i-1,j] > ||cs[i,j-1] + hen i \(i-1;\)
else j \(i-1;\)
else |
else

Note: 1. It can also be dere recussively, in which ease you can print the strong out directly as you see it. 2. It seems you can also do it without having to look at the styrings it ty by just of looking at the lengths.

You you scan left in a new, as soon as the value decreases by I you have found a character in the substray. 2. Let N be a flow network and let f be a flow for N. Furthermore, let X be an st-cut
of N. Let state ses + t & T, and let F =  $\sum_{v \in S} \left( \sum_{e \in E(v)} f(e) - \sum_{e \in E(v)} f(e) \right)$ the sum of the potal reflew through all vertices in S, some of which include edges across the cut. First, note that thutses, & fie) - & fie) =0 by the conservation rule. Thus,  $F = \mathcal{E} f(e) - \mathcal{E} f(e) = |f|$ .  $e \in E(s)$   $e \in E(s)$ Furthermore, note that for every CFE that is not a crossing edge one of two cases apply: (1) Both f(e) a -f(e) are part of sum F for (2) Both f(e) & -f(e) an not part of sum Frako by net value of O.

This leaves only crossing edges contributing to F, files for those going from S to T and -files for those going from T to S. Thus, it is also true that

and trerefore |f| = f(x) |f| = F = f(x)

prove tris. I suspect many will do induction or contradiction proof.

3.	V	5	A	B	C	D	E	F	
	CVID	0	+00	+00	400	+00	400	+00	
	pass 1	0	2	9	3	+00	400	+00	
	pass 2	0	2	9	3	6	7	5	
	Dass 3		2		3	6	7	5	
×	pass 4	0	2	6	3	6	7	5	

No more iterations are necessary as there were no changes from iteration to 3 to iteration 4.