

CSC 226 - ASSIGNMENT 3 - SOLUTIONS

1. Suppose G is a connected graph with distinct, positive integer edge weights. Prove that G has exactly one MST T .

Proof: Assume for purposes of contradiction that there exists two MSTs for G , call them T and T' .

- Since T + T' are distinct MSTs each contains at least one edge that the other does not.
- Let $e \in T$ + $e' \in T'$ be the edges with minimum weight that ~~do~~ not exist in the other tree, respectively.
i.e. $e \in T$ but $e \notin T'$ + $e' \in T'$ but $e' \notin T$.
- Since all the edge weights are distinct, assume without loss of generality that $w(e) < w(e')$.
- Now, place edge e in tree T' creating cycle C . Furthermore, cycle C must contain at least one edge $f \in T'$ not in T otherwise C would be a cycle in T + T would not be a tree.
- $w(f)$ must be greater than $w(e)$ since $w(e) < w(e') \leq w(f)$ by assumption above.
- Remove f from T' and we have a new spanning tree T^* with $w(T^*) < w(T)$.
- This is a contradiction since T was a minimum spanning tree.
- Therefore, G has exactly one MST.

2. Let V' be any subset of V in $G = (V, E)$ and let e be a minimum weight edge that has exactly one endpoint in V' . Then, there exists a minimum spanning tree T in G that contains e .

Proof: By contradiction, assume that no MST exists that contains e as described above.

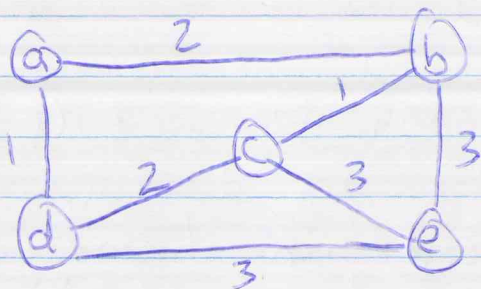
→ Let T be a MST of G and add e to T , creating a cycle C .

→ Since e connects V' + $V - V'$ then there must exist some other edge f that connects V' and $V - V'$ or e would have had to be part of T , such that $f \in C$ too.

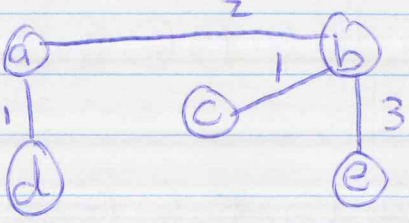
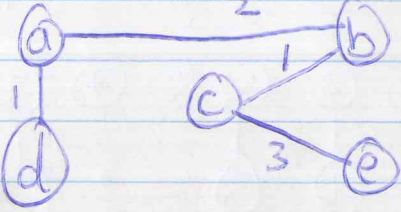
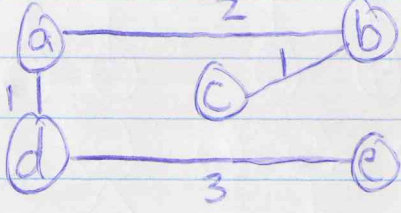
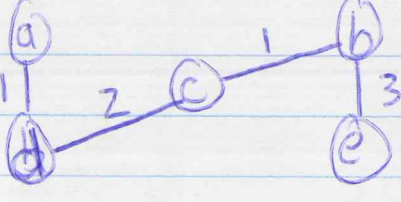
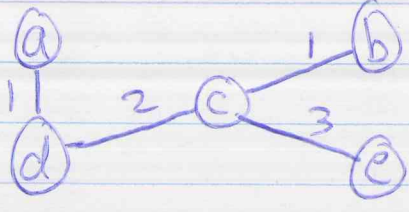
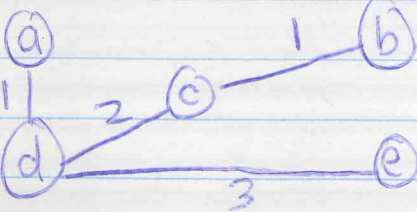
→ $w(f) \geq w(e)$ by definition of e above and thus, if we remove f from C we get a new spanning tree T' with weight $w(T') \leq w(T)$, implying that T' is a MST containing e , a contradiction.

→ Therefore, there must exist at least one MST of G containing e .

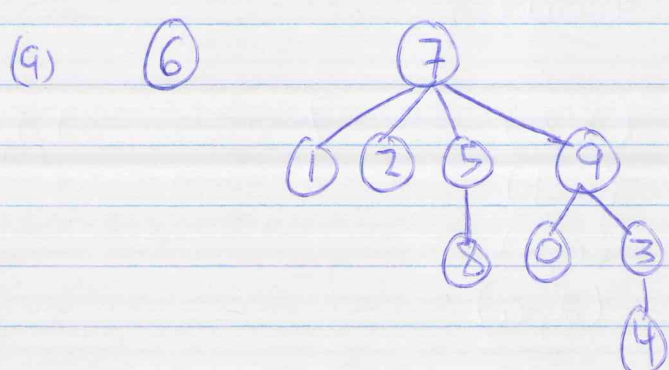
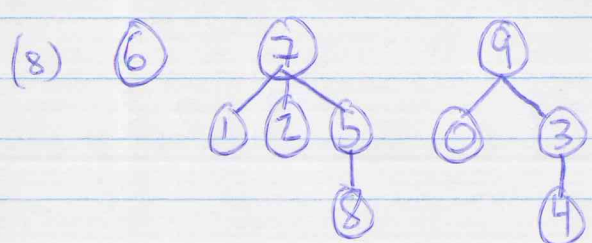
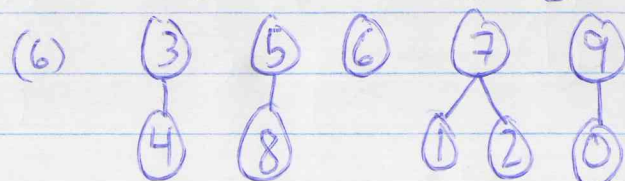
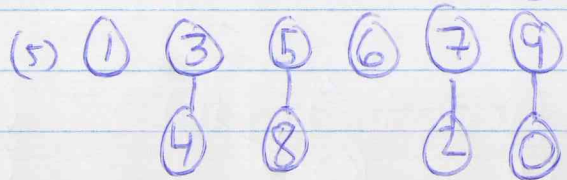
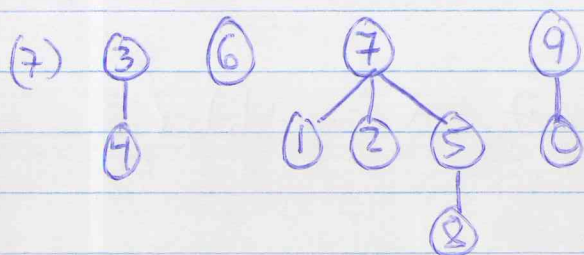
3. $G =$



QST: Conjecture: All MSTs T of a graph G have the same profile.

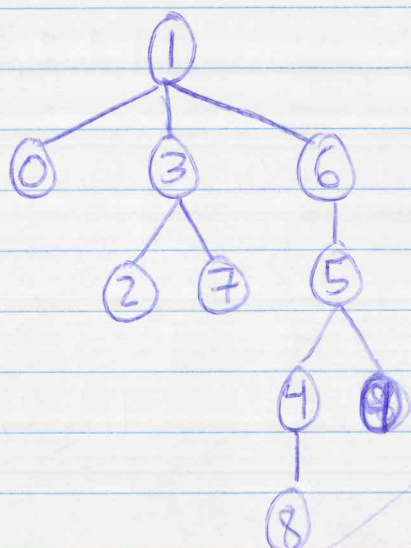
mst	edges	weight	profile
(1.) 	(a,b) (b,c) (b,e) (d,a)	7	1, 1, 2, 3
(2.) 	(a,b) (b,c) (c,e) (d,a)	7	1, 1, 2, 3
(3.) 	(a,b) (b,c) (d,a) (d,e)	7	1, 1, 2, 3
(4.) 	(b,c) (b,e) (c,d) (d,a)	7	1, 1, 2, 3
(5.) 	(b,c) (c,d) (c,e) (d,a)	7	1, 1, 2, 3
(6.) 	(b,c) (c,d) (d,a) (d,e)	7	1, 1, 2, 3

4.	i	0	1	2	3	4	5	6	7	8	9	(1)	0	1	2	3	4	5	6	7	8	9
(1)	del	0	1	2	3	4	5	6	7	8	9	(2)	1	2	3	4	5	6	7	8	9	
(2)	9-0	9	1	2	3	4	5	6	7	8	9	(3)	1	2	3	4	5	6	7	8	9	0
(3)	3-4	9	1	2	3	3	5	6	7	8	9	(4)	1	2	3	5	6	7	8	9		
(4)	5-8	9	1	2	3	3	5	6	7	5	9	(5)	1	2	3	5	6	7	8	9		
(5)	7-2	9	1	7	3	3	5	6	7	5	9	(6)	1	2	3	5	6	7	8	9		
(6)	2-1	9	7	7	3	3	5	6	7	5	9	(7)	1	2	3	5	6	7	8	9		
(7)	5-7	9	7	7	3	3	7	6	7	5	9	(8)	1	2	3	5	6	7	8	9		
(8)	0-3	9	7	7	9	3	7	6	7	5	9	(9)	1	2	3	5	6	7	8	9		
(9)	4-2	9	7	7	9	3	7	6	7	5	7											



5.

i	0	1	2	3	4	5	6	7	8	9
$id[i]$	1	1	3	1	5	6	1	3	4	5



→ At some point Node 5 was merged to node 6.

→ At that time size of 5 ~~was~~ had to be less than or equal to size of 6.

→ subsequently, any more ~~the~~ unions with this tree would append the ~~the~~ other tree directly to 6 as a child.

→ this could only happen when 6 + 5 were singletons.

Notes:

1. The proofs in question 1 + 2 will not necessarily be the same as mine. Try to follow their logic and mark accordingly.
2. Question 4: When the trees are the same size they can arbitrarily add one to the other. So, their final forest may look slightly different.
→ also, they may union-by-height instead of size (number of nodes) and that is okay too.