# Chapter 1 - Introduction to Computer Networking

ex. on message transmission, on message reciept, on message error, etc. Sunday, January 17, 2016 PROTOCOL - desires the format and order of messages exchanged and actions to be laken MOSTS/END-SYSTEMS - interconnected devices (smartphones, computers, servers) running network applications COMMUNICATION LINK - fiber, copper wire, radio, satellite PACKET - block of data rosted between an origin and a destination ROUTERS/SWITCHES - devices that forward packets toward destination via links 15P - Internet Serice Provider, how end-systems access the internet INTERNET PROTOCOL STACK - the protocols of the various layers of the internet 5. APPLICATION (message) name of packets processed at this layer -> supporting retwork applications (ex. browsers) -> FTP, SMTP, HTTP 4. TRANSPORT -> EXISTS ON MOSTS ( segment) -> creater logical communication between two end-systems -> TCP, UDP 3. NETWORK (datagram) -> routes messages from source to destination -> IP, routing protocols 2. LINK EXISTS ON LINKS (frame) -> data transfer between neighbouring network elements -> Ethernet, 802.11 (wifi), PPP 1. PHYSICAL

-> transmission over communication link/medium

### INTERNET

## "nuts and bolts" view

- millions of interconnected end-systems, each running network applications
- extremely redundant/robust
- interconnected ISPs, network of networks
- loosely hierarchical

#### service view

- communication infrastructure that enables distributed applications
- reliable data delivery, or "best effort"
- provides application programming interface

Ly ex. an API specifies how to ask for specific information

is multiple and systems in different places

network edge - hosts (clients and severs)

access betwerks - communication links, what physically connects an end-system to the first router

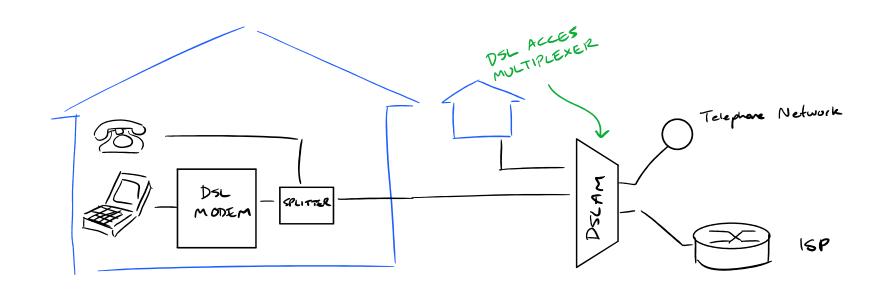
network core - interconnected routers, retworks of retworks

edge router - connects residential access nets/school or company access retworks / mobile access retworks to internet

## DSL ACCESS NETS (digital subscriber line)

- -uses existing telephone lines infrastructure to provide internet
- proves only need \$-4kHz to provide good signal
- remaining frequency range is used for internet data

La separated into uplanding/downloading ranges N2.5 Mbps ~ 25Mbps



BOTH DSL & CABLE NEED MODEMS (DATA MODULATION) TO SEPARATE CHANNELS/FREQUENCIES.

## CABLE ACCESS NETWORK

- uses frequency division multiplexing

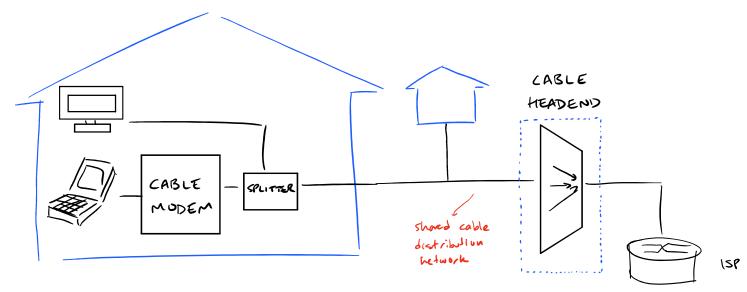
  (> different channels are transmitted
  through different bands
- lowe frequency channels 1-6 transmit video
- higher frequency channels 7-8 transmit data
- highest frequency channel a is control channel
- cable share access network to headend

  (5) better because non-used bandwidth can be
  used by other people on your network

  (5) worse because less-secure

~30Mbps 3/1, 2Mbps 0/1

access to certain



### WIRELESS ACCESS NETWORK

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- connects and system to route via access point

- LANS

-> with building ( ~looft)

-> 802.11 b/g/n (wifi)

-> 11,54,72 Mbps

- wide-area wireless access

- provided by cellular operator

- 1-10 Mbps

- 3G, 4G LTE, Wiman
```

### HOSTS

Sending packets of data

1. takes message from application

2. breaks it into packets of length L bits

3. transmits packets through access betweek at fransmission rate, R

transmission delay = time needed to send bits into the link = \frac{L(\text{bits})}{R(\text{bits}\text{fee})}

not the time to physically transmit the signal through

# LINK - what lies between transmitter and recieves

Guided media -> signals propagate in solid media (copper, fiber, coaxial) Twisted Pair (TP) - two insulated copper wires - twisted to minimize woise - much cheaper than optical - entegory 5: 100 Mbps, IGbps Ethernet - category 6: 106 bps Coaxial cable - two concentric copper conductors - bidirectional - broadband Ly multiple channels on cable Fibe optic cable -glass fiber carrying light pulses, each pulse a bit -very high speed, 10-100 Gbps - low error rate immune to electromagnetic noise/interference Unquited media -> signals propagate freely (radio) Padio - signal carried in electromagnetic spectrum - no physical wire - bidirectional

```
- propagation environment effected by:
     obstruction by objects
     interference
 - terrestrial microwave, N 45 Mbps
 - LAN (wifi), 11 Mbps, 54 Mbps
 - wide-area (cellular), 3G ~3 Mbps
 - satellite, lebps - 45mbps
      1> 270 mac end-to-end delay
 - bluetooth
```

from source to destination

forwarding - moving an incoming packet from a router's input to the appropriate router output (link)

## NETWORK COPE

\_\_\_\_\_ mush of interconnected packet switches/routers

PACKET SUITCHING (on-demand resource allocation)

- how most of the interest transmits data - hosts break messages into packets

La packets are forwarded from one router to the

rext across links from source to destination weach packet transmitted at Isl link capacity

- store and forward

it is forwarded to the next router (: delay)

-simple

- very good for resource shaving, but no guarantee of bandwidth

## Lo can become congested, but can support more users

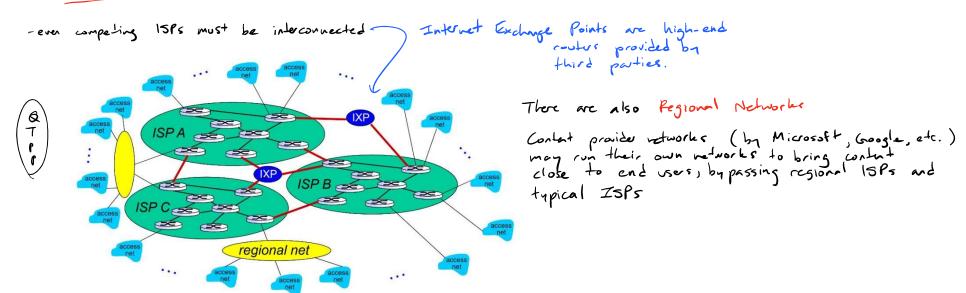
## CIRCUIT SWITCHING (reserved resource allocation)

- you said signals to set up an actual direct circuit from the saider to the recieve
- end-to-end resources allocated to and reserved for communication b/w source & destination
- has dedicated resources and guaranteed performance for each user
- however, no sharing of resources of vaused circuits are wasted while inactive

FDM (frequency division multiplexing)
frequency spectrum of a link is divided among the connections (users)

TDM (time division multiplexing)
time is divided into slots (grouped in frames) and the network
dedicates one slot per frame to each connection

# ISPS



## LOSS

if router queue becomes full due to congestion, any newly arriving packets will be lost life so, the router may send a signed back to the soder notifying them

# DELAY (packet-switching)

Total model delay (end-to-end) consists of 4 delays:

- 1. QUEUING DELAY
  - packet arrival rate exceeds output rate
  - Inly one packet may be processed at a time

  - -depends on route congestion Lap, where R is link transmission rate, L is padet length, a is average partet arival rate bits/second bits

2 must be <1 or you are dropping packets

- 2 PROCESSING DELAY
  - roster has to read data to know where to forward it (determine output link)
  - does data integrity check for bit errors
- 3. TRANSMISSION DELAY
  - -time it takes for router to send each packet
  - 4 , where L is the packet length & R is the transmission rate
  - if not a congested router, this takes the longest
- 4 PROPAGATION DELAT
  - -time to actually transmit the electrical signal through the medium
  - -d/s, where d'is the length of the link and s is the propagation speed in medium

throughput - rate at which bits can be transferred b/w sender to recieve, bits/second
- bottlenecks usually happen close to edge, like wifi
- the smallest transmission rate in a path of links a packet takes