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# **Lecture 3:**

# **Discrete Fourier Transform**

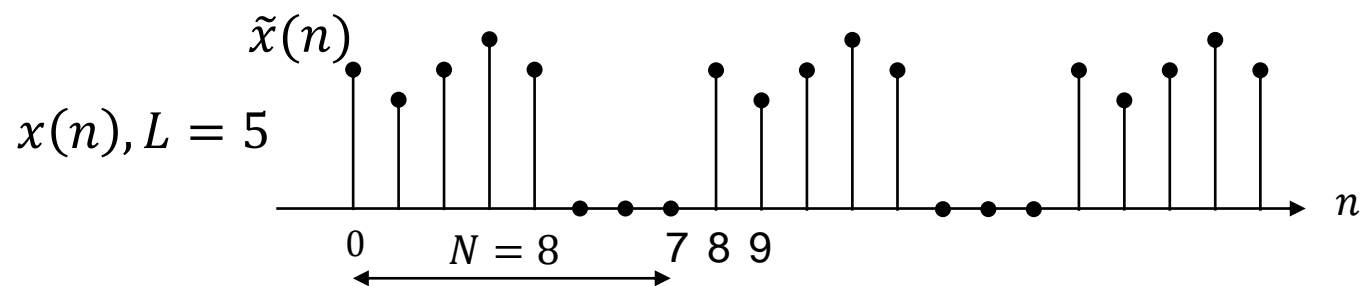
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# DFT

## Background

- ❖ Consider  $x[n] \longleftrightarrow X(\theta)$  defined for all  $\theta \in [-\pi, \pi]$
- ❖ Do we have to know  $X(\theta)$  for all  $\theta \in [-\pi, \pi]$  to recover  $x[n]$  ?
- ❖ For a finite-duration sequence  $x(n)$ ,  $0 \leq n \leq L-1$

$$\tilde{x}(n) = \sum_{k=-\infty}^{\infty} x(n-kN) \text{ where } N \geq L \quad \boxed{\tilde{x}(n) = x(n \bmod N)}$$



$$\Rightarrow \tilde{X}(\theta) = cX(\theta_k) = X(k), \quad \theta_k = \frac{2\pi}{N}k, \text{ where } 0 \leq k \leq N-1$$

**Note)** Both  $x(n)$  and  $X(k)$  are discrete

# DFT

## Background

For  $x(n)$  with a finite duration  $L$ , i.e.,  $x(n)$ ,  $0 \leq n \leq L-1$

$$X(e^{j\theta}) = \sum_{n=-\infty}^{\infty} x(n)e^{-j\theta n}$$

If we evaluate it at  $\theta = \frac{2\pi k}{N}$ ,  $N \geq L$ , to obtain

$$X(e^{j\theta})|_{\theta = k \cdot (2\pi/N)} = X(e^{j2\pi k/N}) = X\left(\frac{2\pi}{N}k\right)$$

$$\Rightarrow X(e^{j2\pi k/N}) = \sum_{n=-\infty}^{\infty} x(n)e^{-j2\pi kn/N}, \quad k = 0, 1, \dots, N-1$$

$$X(e^{j2\pi k/N}) = \dots + \sum_{n=-N}^{-1} x(n)e^{-j2\pi kn/N} + \sum_{n=0}^{N-1} x(n)e^{-j2\pi kn/N} + \sum_{n=N}^{2N-1} x(n)e^{-j2\pi kn/N} + \dots$$

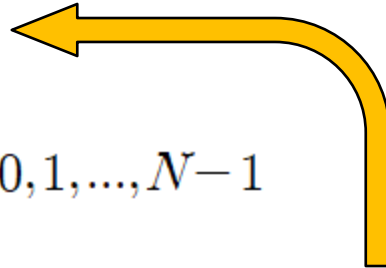
$(l = -1) \qquad (l = 0) \qquad (l = 1)$

$$= \sum_{l=-\infty}^{\infty} \sum_{n=lN}^{lN+N-1} x(n)e^{-j2\pi kn/N}$$

# DFT

By the change of variable,  $n = n - lN$

$$X(e^{j2\pi k/N}) = \sum_{n=0}^{N-1} \left[ \sum_{l=-\infty}^{\infty} x(n - lN) e^{-j2\pi k(n - lN)/N} \right]$$

$$= \sum_{n=0}^{N-1} \left[ \sum_{l=-\infty}^{\infty} x(n - lN) e^{-j2\pi kn/N} \right] \text{ for } k = 0, 1, \dots, N-1$$


$$\tilde{x}(n) = \sum_{l=-\infty}^{\infty} x(n - lN)$$

Again, repeated sequence of  $x(n)$  yield sampled spectrum  $X(k)$

When  $N \geq L$



$$X(k) = X(e^{j(2\pi/N)k}) = \sum_{n=0}^{N-1} x(n) e^{j(2\pi/N)kn}, \text{ for } 0 \leq k \leq N-1$$

$$x(n) = \frac{1}{N} \sum_{k=0}^{N-1} X(k) e^{j(2\pi/N)kn}, \text{ for } 0 \leq n \leq N-1 \quad \text{where } X(k) = X(e^{j2\pi/N}k)$$

**Remember:  $X(k)$  is the DTFT of  $\tilde{x}(n)$**

## Definition

### ❖ N-point DFT

$$X(k) = \sum_{n=0}^{N-1} x(n) e^{-j\frac{2\pi}{N}kn} = \sum_{n=0}^{N-1} x(n) W_N^{-kn}$$
$$0 \leq k, n \leq N-1$$

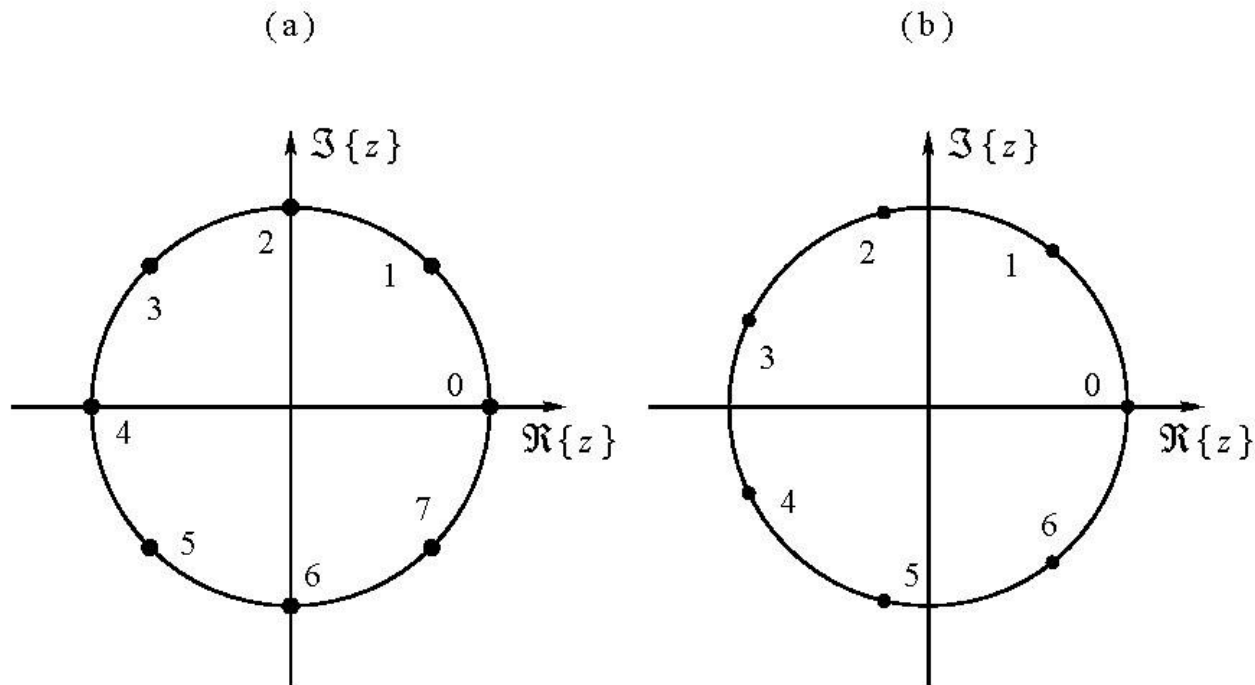
Where

$$W_N = e^{j2\pi/N} \Rightarrow W_N^0 = W_N^N = 1, \quad W_N^n = W_N^{n \pm mN}$$
$$X(k) = DTFT \left[ \tilde{x}(n) = \sum_{m=-\infty}^{\infty} x(n + mN) \right]$$
$$X(k) = X(k + N)$$

# DFT

## Definition

$$\diamond W_N = e^{j2\pi/N}$$

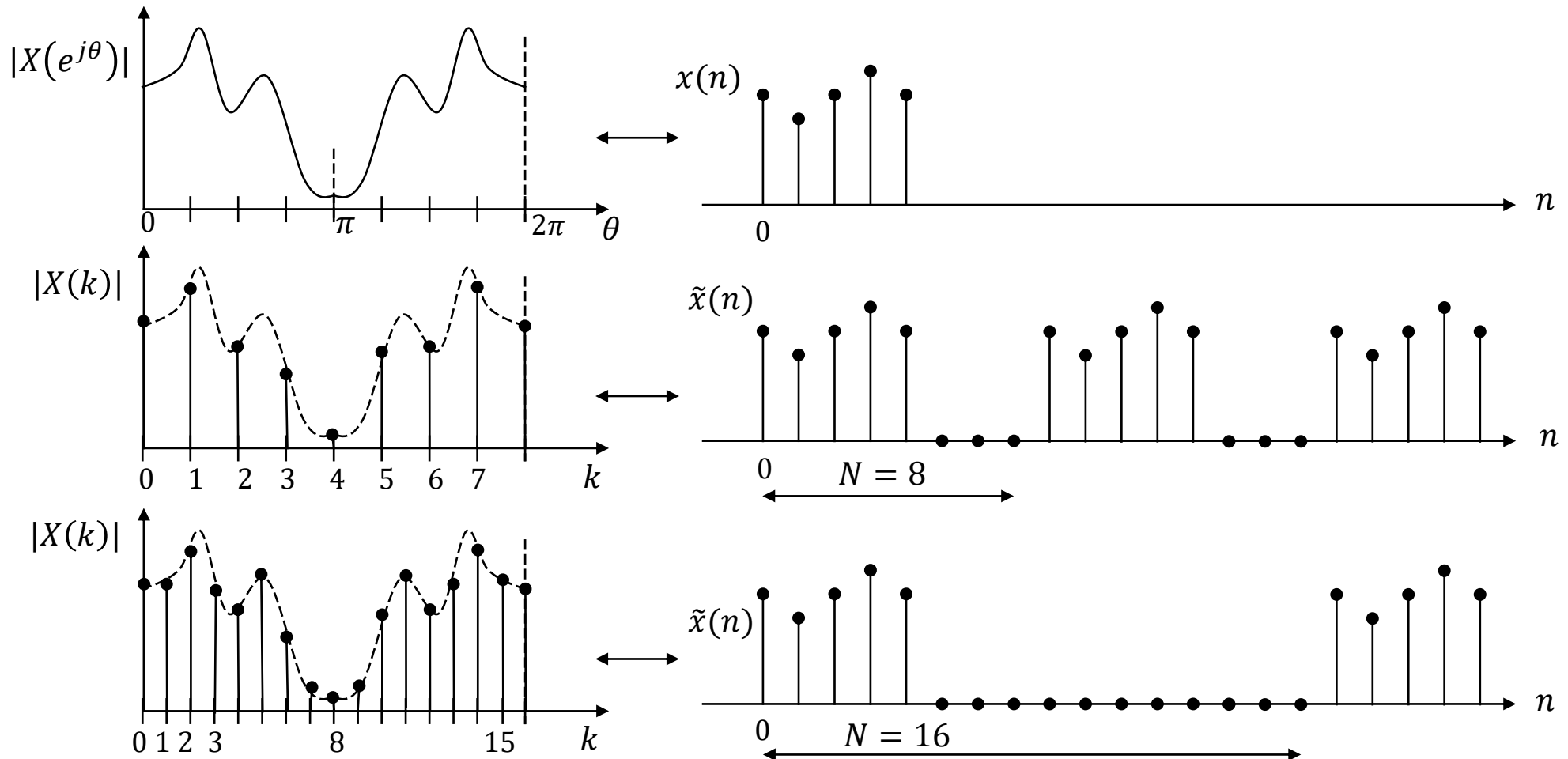


**Figure 4.2** The sequence  $W_N^n$  in the complex plane: (a) even  $N$ ; (b) odd  $N$ . The numbers indicate the values of  $n$ .

# DFT

## Definition

### ❖ N-point DFT



## Some properties of DFT and its kernel function

$$\sum_{n=0}^{N-1} W_N^{kn} = N\delta[k \bmod N] = N\delta(k - mN) = \begin{cases} N, & (k \bmod N) = 0 \\ 0, & (k \bmod N) \neq 0 \end{cases}$$

*Proof)*

$$k = mN \quad \sum_{n=0}^{N-1} W_N^{kn} = \sum_{n=0}^{N-1} W_N^{nmN} = \sum_{n=0}^{N-1} 1 = N$$

$$k \neq mN \quad \sum_{n=0}^{N-1} W_N^{kn} = \frac{W_N^{Nk} - 1}{W_N^k - 1} = \frac{1 - 1}{W_N^k - 1} = 0.$$

*Examples)*

$$1) \quad x[n] = \begin{cases} 1, & n=0 \\ 0, & 1 \leq n \leq N-1 \end{cases} \longleftrightarrow X[k] = 1, \quad 0 \leq k \leq N-1$$

$$2) \quad x[n] = 1, \quad 0 \leq n \leq N-1$$

$$\longleftrightarrow X[k] = \sum_{n=0}^{N-1} W_N^{-kn} = N \cdot \delta(k \bmod N) = \begin{cases} N, & k=0 \\ 0, & 1 \leq k \leq N-1 \end{cases}$$



# DFT

## Inverse DFT

$$x[n] = \frac{1}{N} \sum_{k=0}^{N-1} X[k] W_N^{kn} \quad ; \quad n, k \leq N-1$$

$$\begin{aligned} \frac{1}{N} \sum_{k=0}^{N-1} X[k] W_N^{kn} &= \frac{1}{N} \sum_{k=0}^{N-1} \left[ \sum_{m=0}^{N-1} x[m] W_N^{-km} \right] W_N^{kn} = \frac{1}{N} \sum_{m=0}^{N-1} x[m] \left[ \sum_{k=0}^{N-1} W_N^{(n-m)k} \right] \\ &= \frac{1}{N} \sum_{m=0}^{N-1} x[m] N \delta[(n-m) \bmod N] = x[n] \end{aligned}$$

## Things to know about DFT

- ❖ We use DFT to examine frequency spectrum of a finite length, **L**, discrete sequence (digital sequence)
- ❖ When taking its **N-DFT**, we assume that the sequence is periodic
- ❖ Hence, DFT length **N** must be chosen larger than **L** to avoid aliasing in time domain
- ❖ Frequency aliasing occurs if  $x(n)$ , or  $x(t)$ , is not band-limited

# DFT

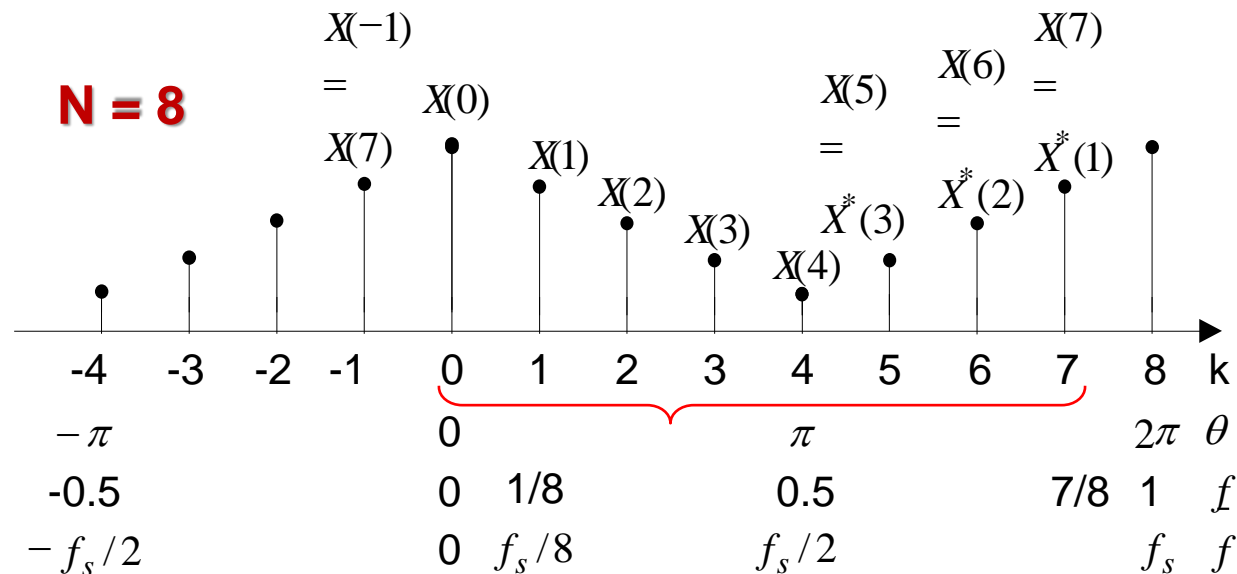
## Things to know about DFT

❖ **Analog frequency:**  $k \rightarrow f_k = k \cdot \delta_f = k \cdot \frac{f_s}{N}$  ,  $\delta_f = f_s/N$

❖ **Digital frequency:**

$$k \rightarrow \underline{f_k} = f_k' = \frac{k}{N} \quad , \quad \Theta_k = k \cdot \delta_\theta = 2\pi \cdot k/N \quad , \quad \delta_\theta = 2\pi/N$$

$$X(k) = X(k+N) \Rightarrow X(-k) = X(N-k) = X^*(k)$$



# Matrix Representation of DFT

## DFT Matrix of dimension N

$$F_N = \begin{bmatrix} W_N^0 & W_N^0 & W_N^0 & \dots & W_N^0 \\ W_N^0 & W_N^{-1} & W_N^{-2} & \dots & W_N^{-(N-1)} \\ W_N^0 & W_N^{-2} & W_N^{-4} & \dots & W_N^{-2(N-1)} \\ \vdots & \vdots & \vdots & \dots & \vdots \\ W_N^0 & W_N^{-(N-1)} & W_N^{-2(N-1)} & \dots & W_N^{-(N-1)^2} \end{bmatrix} \quad X[k] = \sum_{n=0}^{N-1} x[n] W_N^{kn}$$
$$x_N = \begin{bmatrix} x[0] \\ x[1] \\ \vdots \\ x[N-1] \end{bmatrix} \quad X_N = \begin{bmatrix} X[0] \\ X[1] \\ \vdots \\ X[N-1] \end{bmatrix}$$

## Matrix-form representation of DFT

$$X_N = F_N \cdot x_N$$

$$x_N = N^{-1} F_N^* X_N$$

# Matrix Representation of DFT

$$X[k] = \sum_{n=0}^{N-1} W_N^{-nk} x[n]$$

*Ex) case for  $N = 4$*

$$X[k] = W_4^{-0k} x[0] + W_4^{-1k} x[1] + W_4^{-2k} x[2] + W_4^{-3k} x[3]$$

*For  $X[k = 1]$*

$$\begin{bmatrix} X[0] \\ \textcolor{red}{X[1]} \\ X[2] \\ X[3] \end{bmatrix} = \begin{bmatrix} W_4^{-0 \cdot 0} & W_4^{-1 \cdot 0} & W_4^{-2 \cdot 0} & W_4^{-3 \cdot 0} \\ \textcolor{red}{W_4^{-0 \cdot 1}} & \textcolor{orange}{W_4^{-1 \cdot 1}} & \textcolor{green}{W_4^{-2 \cdot 1}} & \textcolor{blue}{W_4^{-3 \cdot 1}} \\ W_4^{-0 \cdot 2} & W_4^{-1 \cdot 2} & W_4^{-2 \cdot 2} & W_4^{-3 \cdot 2} \\ W_4^{-0 \cdot 3} & W_4^{-1 \cdot 3} & W_4^{-2 \cdot 3} & W_4^{-3 \cdot 3} \end{bmatrix} \begin{bmatrix} \textcolor{red}{x[0]} \\ \textcolor{orange}{x[1]} \\ \textcolor{green}{x[2]} \\ \textcolor{blue}{x[3]} \end{bmatrix}$$

# Matrix Representation of DFT

## Properties of DFT matrix

$$\diamond W_N^{-nk} = 1 \quad \text{for } n=0 \text{ and/or } k=0$$

$$\diamond W_N^{-kn} = W_N^{-nk} \quad [W_N^{nm}]^* = W_N^{nm}$$

$$\diamond F_N \cdot F_N^{T*} = F_N \cdot F_N^+ = NI_N$$

$$[F_N \cdot F_N^+]_{kl} = \sum_{n=0}^{N-1} W_N^{kn} \{W_N^{nl}\}^{*T} = \sum_{n=0}^{N-1} W_N^{kn} W_N^{ln} = \sum_{n=0}^{N-1} W_N^{(l-k)n} = N\delta[(l-k) \bmod N]$$

$$\diamond N^{-1/2}F_N \text{ is unitary: A symmetric unitary matrix: Normalized DFT matrix}$$

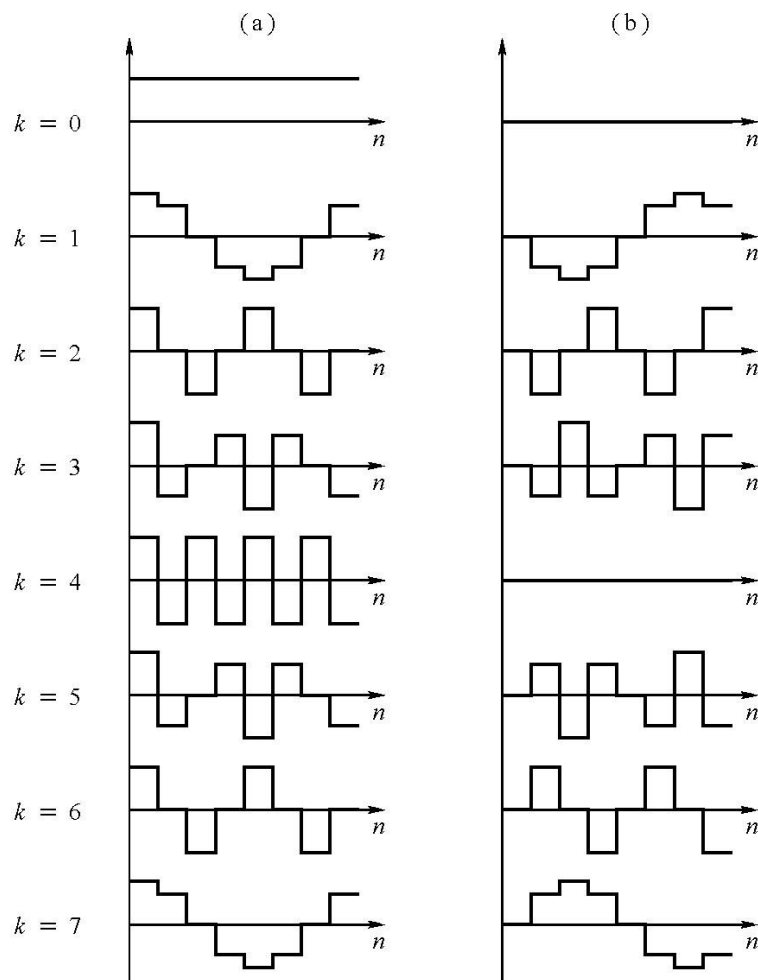
**Note:** A complex  $N \times N$  matrix is called a unitary matrix if  $Q_N \cdot Q_N^{T*} = I_N$

**If it is real, it is called an orthonormal matrix.**

# Matrix Representation of DFT

## Properties of DFT matrix

- ❖ DFT basis vectors (orthogonal)



**Figure 4.5** The DFT basis vectors for  $N = 8$ : (a) real part; (b) imaginary part.  
EEE4175 Introduction to Digital Signal Processing

# Properties of the DFT

## Linearity and Periodicity

$$z(n) = ax(n) + by(n) \Leftrightarrow Z(k) = aX(k) + bY(k)$$

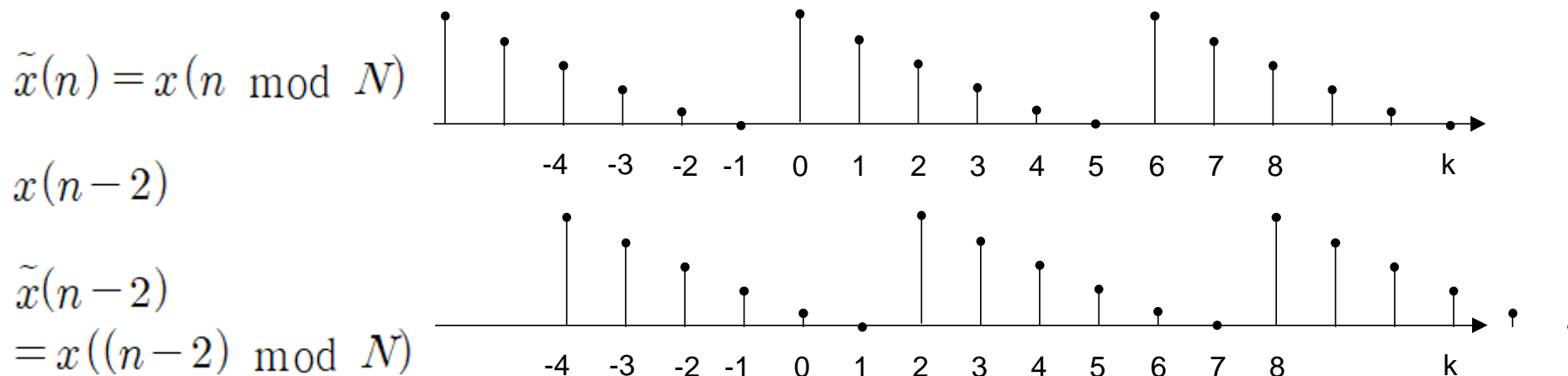
$$X(k) = X(k + N)$$

## (Circular) Time Shift

$$\tilde{y}(n) = y(n) = \tilde{x}(n - m) = x[(n - m) \bmod N] \Leftrightarrow Y(k) = W_N^{-km} X(k)$$

(pf) See pp. 102 or can be proven graphically.

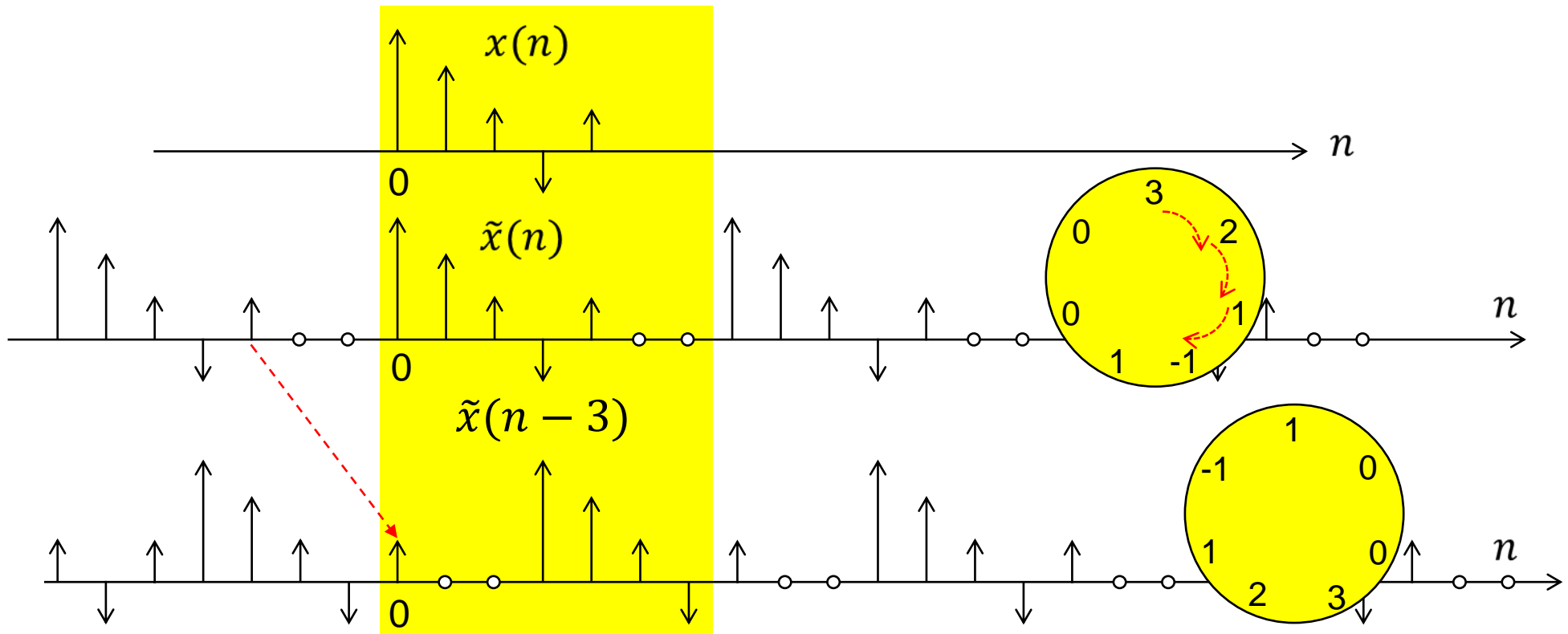
**Note: As far as DFT is concerned, any shift is considered to be circular !!**





# Properties of the DFT

## 7-Point Circular Shift



# Properties of the DFT

## (Circular) frequency shift

$$W_N^{mn} x(n) \Leftrightarrow \tilde{X}(k-m) = X[(k-m) \bmod N]$$

## Complex conjugate

$$x^*(n) \Leftrightarrow X^*[(N-k) \bmod N] = X^*(N-k)$$

$$DFT[x^*(n)] = \sum_{n=0}^{N-1} x^*(n) W_N^{-nk} = \left\{ \sum_{n=0}^{N-1} x(n) W_N^{nk} \right\}^* = X^*(-k) = X^*(N-k)$$

## Symmetry properties

❖ If  $x(n)$  is real-valued, then  $X(k) = X^*(-k) = X^*(N-k)$  or  $X^*(k) = X(-k) = X(N-k)$

- $\text{Re}[X(k)], |X(k)|$  : Even symmetric
- $\text{Im}[X(k)], \angle X(k)$  : Odd symmetric

# Properties of the DFT

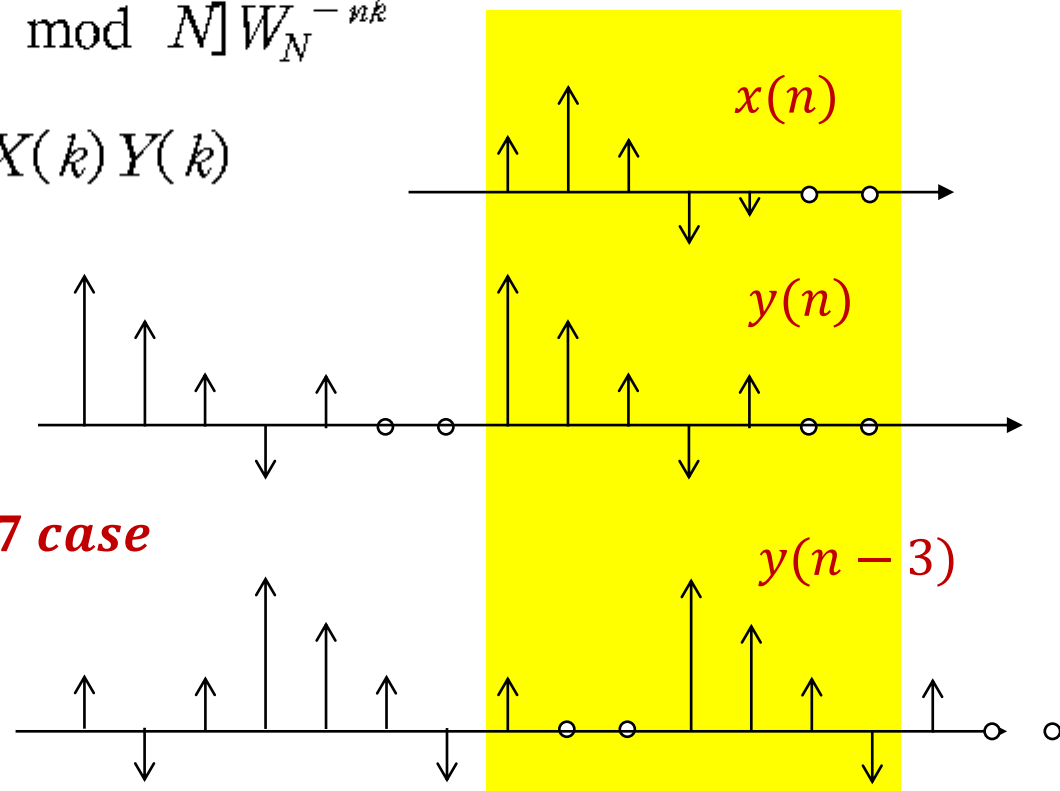
## Circular convolution in time domain

$$z(n) = x(n) * y(n) \Leftrightarrow Z(k) = X(k) Y(k) \quad \boxed{z(n) = x(n) \circledast y(n)}$$

$$\begin{aligned} Z(k) &= \sum_{n=0}^{N-1} z(n) W_N^{-nk} = \sum_{n=0}^{N-1} \left\{ \sum_{m=0}^{N-1} x(m) y[(n-m) \bmod N] \right\} \cdot W_N^{-nk} \\ &= \sum_{m=0}^{N-1} x(m) \sum_{n=0}^{N-1} y[(n-m) \bmod N] W_N^{-nk} \\ &= \sum_{m=0}^{N-1} x(m) W_N^{-mk} Y(k) = X(k) Y(k) \end{aligned}$$

Refer to 16 page.

***N = 7 case***



# Properties of the DFT

## Multiplication in time domain

$$z(n) = x(n)y(n) \Leftrightarrow Z(k) = \frac{1}{N} X(k) \circledast Y(k)$$

## Circular correlation

$$x(n) \circledast y^*(-n) \Leftrightarrow X(k) Y^*(k)$$

## Parseval's theorem

$$\sum_{n=0}^{N-1} x(n) y^*(n) \Leftrightarrow \frac{1}{N} \sum_{k=0}^{N-1} X(k) Y^*(k)$$

$$\sum_{n=0}^{N-1} |x(n)|^2 \Leftrightarrow \frac{1}{N} \sum_{k=0}^{N-1} |X(k)|^2$$

# Properties of the DFT

## Zero padding

$$x(n) \text{ vs. } x_a(n) \quad x_a(n) = \begin{cases} x(n), & 0 \leq n \leq N-1 \\ 0, & N \leq n \leq M-1 \end{cases}$$

❖  $X(k)$  vs.  $X_a(k)$

$$X(k): k=0,1,2,\dots,N-1 \quad k \rightarrow f_k = k \cdot \delta_f = k \cdot \frac{f_s}{N}, \quad \delta_f = f_s/N$$
$$\Theta_k = k \cdot \delta_\Theta = 2\pi \cdot k/N, \quad \delta_\Theta = 2\pi/N$$

$$X_a(k): k=0,1,2,\dots,M-1 \quad k \rightarrow f_k = k \cdot \underline{\delta}_f = k \cdot \frac{f_s}{M}, \quad \underline{\delta}_f = f_s/M$$
$$\Theta'_k = k \cdot \underline{\delta}_\Theta = 2\pi \cdot k/M, \quad \underline{\delta}_\Theta = 2\pi/M$$

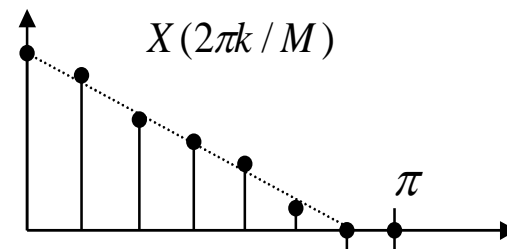
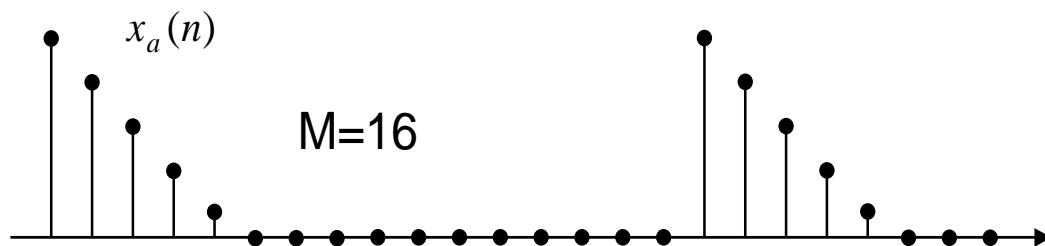
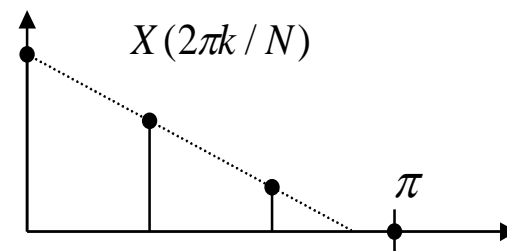
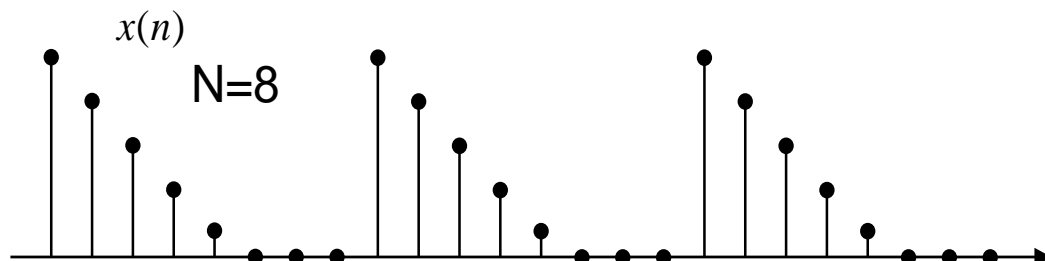
$$X_a(k) = \sum_{n=0}^{M-1} x_a(n) \exp(-j2\pi kn/M) = \sum_{n=0}^{N-1} x(n) \exp(-j2\pi kn/M) = X(\Theta'_k)$$

where  $\Theta'_k = 2\pi k/M, \quad k=0,1,2,\dots,M-1$

$$\underline{\delta}_f = \frac{N}{M} \delta_f \quad \Theta'_k = \frac{N}{M} \Theta_k$$

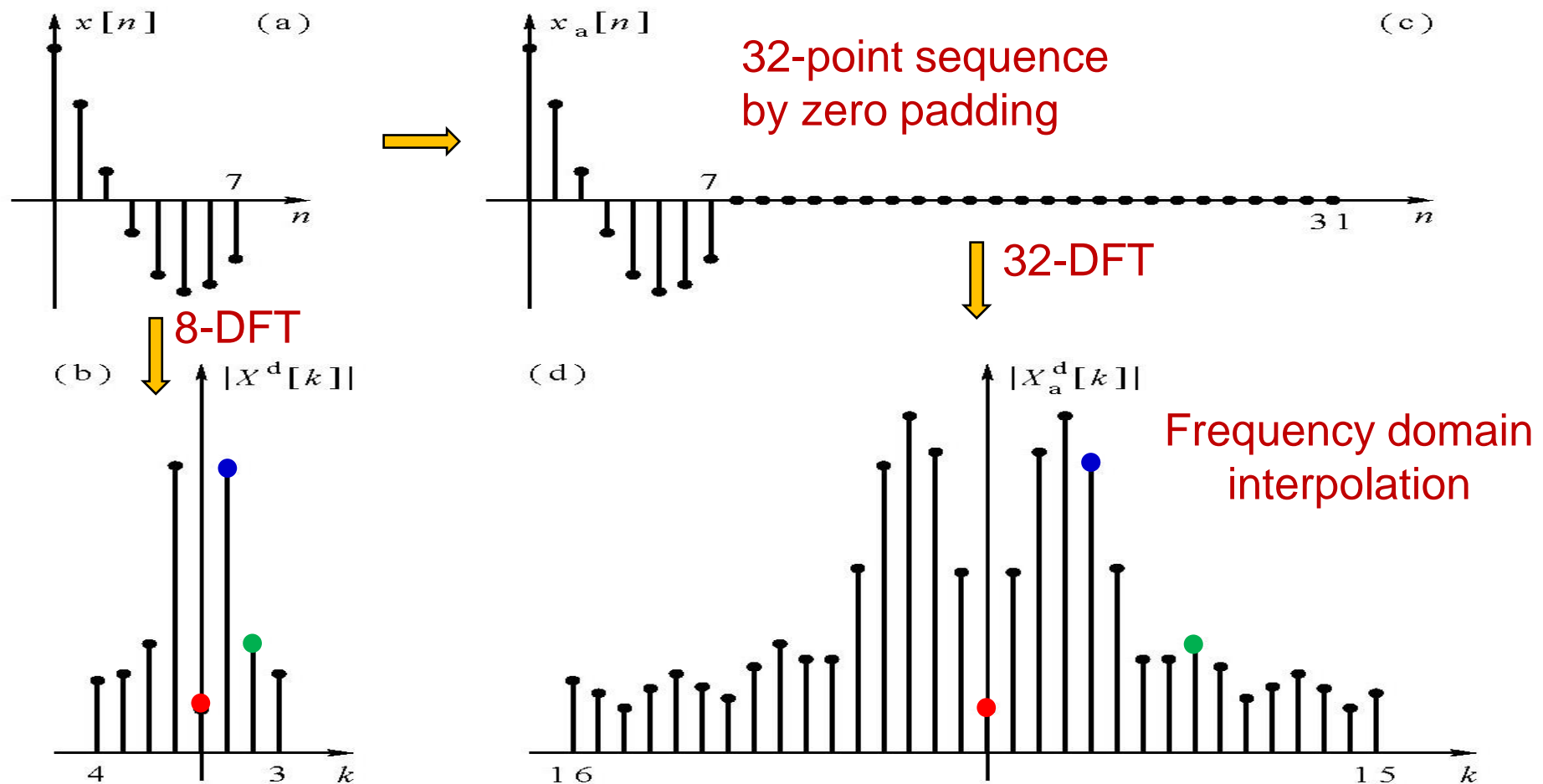
# Properties of the DFT

## Zero padding in time domain improves frequency resolution



# Properties of the DFT

Zero padding in time domain improves frequency resolution



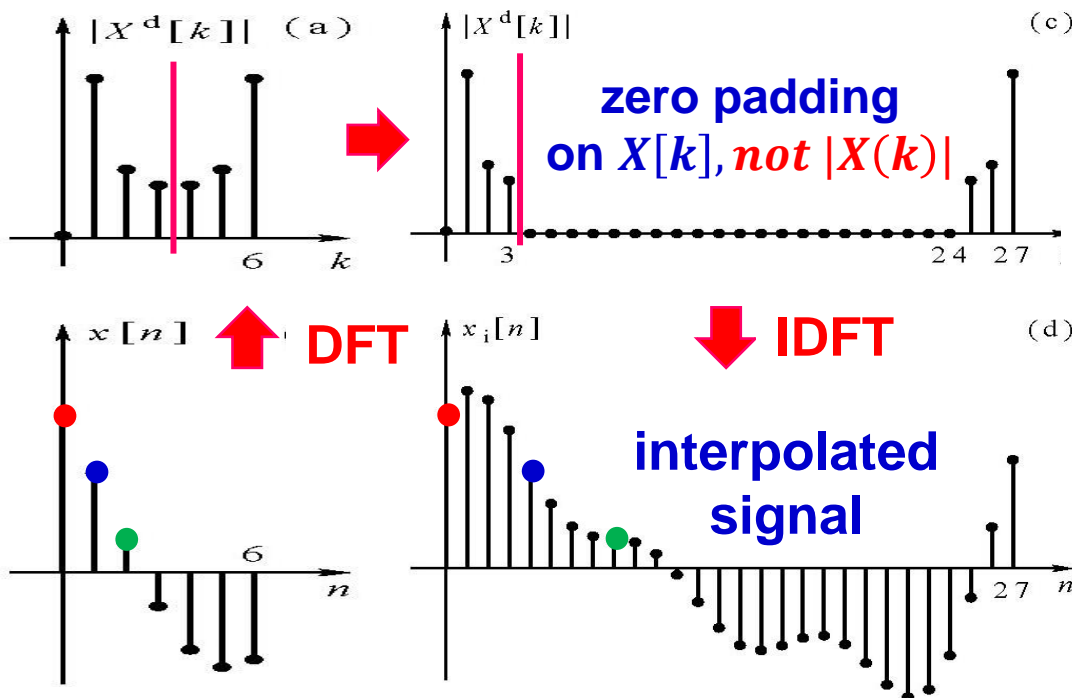
# Properties of the DFT

## Zero padding in frequency domain (time domain interpolation)

$$X_i(k) = \begin{cases} LX(k), & 0 \leq k \leq (N-1)/2 \\ LX(k - M + N), & M - (N-1)/2 \leq k \leq M-1 \\ 0 & \text{otherwise.} \end{cases} \quad \text{for odd } N$$

$L$  : Interpolation factor,  $M = LN$

Zero padding in freq. domain  $\rightarrow$  Time domain interpolation



$$x_i(n) = \frac{1}{M} \sum_{k=0}^{M-1} X_i(k) W_N^{nk}, \quad 0 \leq n \leq M-1$$

$$= \frac{1}{N} \sum_{m=0}^{N-1} x(m) \frac{\sin[\pi(n - mL)/L]}{\sin[\pi(n - mL)/M]}$$

$$x_i(kL) = x(k)$$

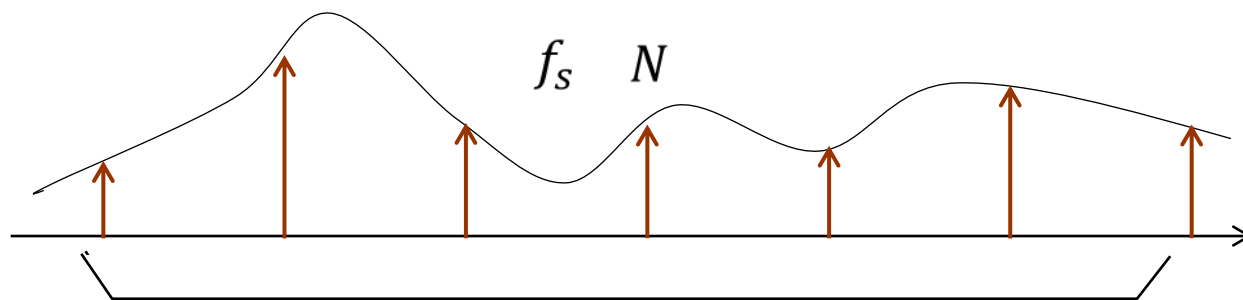
:  $L$ -fold interpolation.



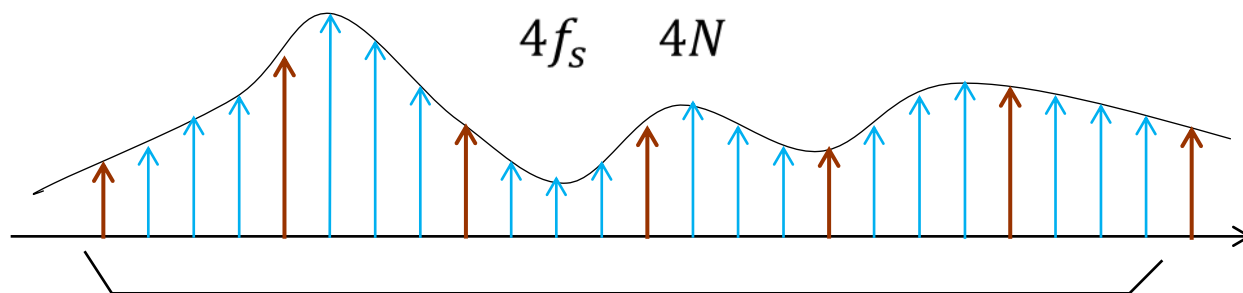
# Properties of the DFT

## Zero padding in frequency domain (time domain interpolation)

*How are the DFT spectra of  $x(t)$  different when sampled at  $f_s$  and  $4f_s$*

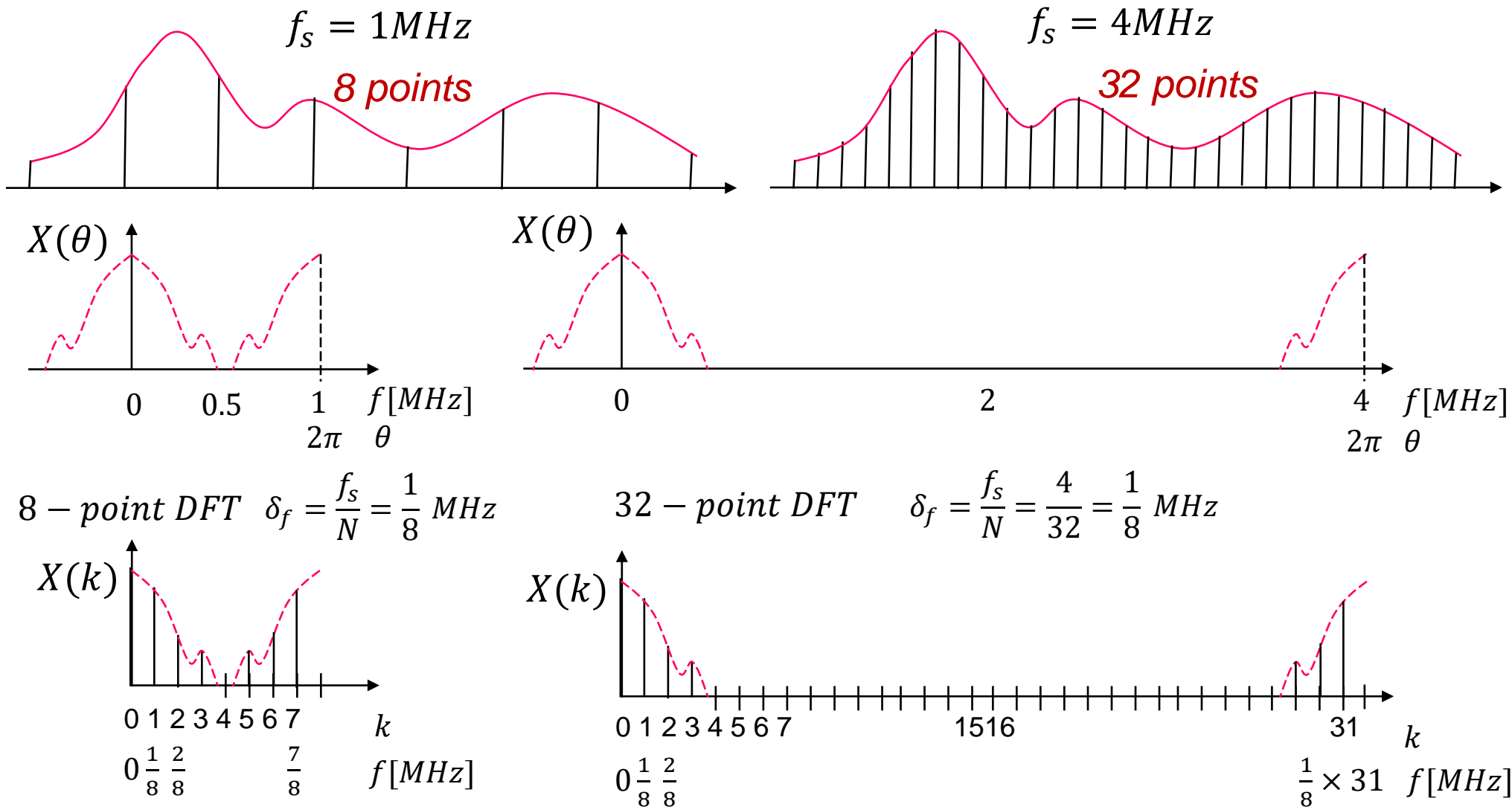


$$\delta_f = f_s / N$$



$$\delta'_f = \frac{4f_s}{4N} = \frac{f_s}{N} = \delta_f$$

# Properties of the DFT



**This is exactly what is obtained by zero padding**

# Circular Convolution

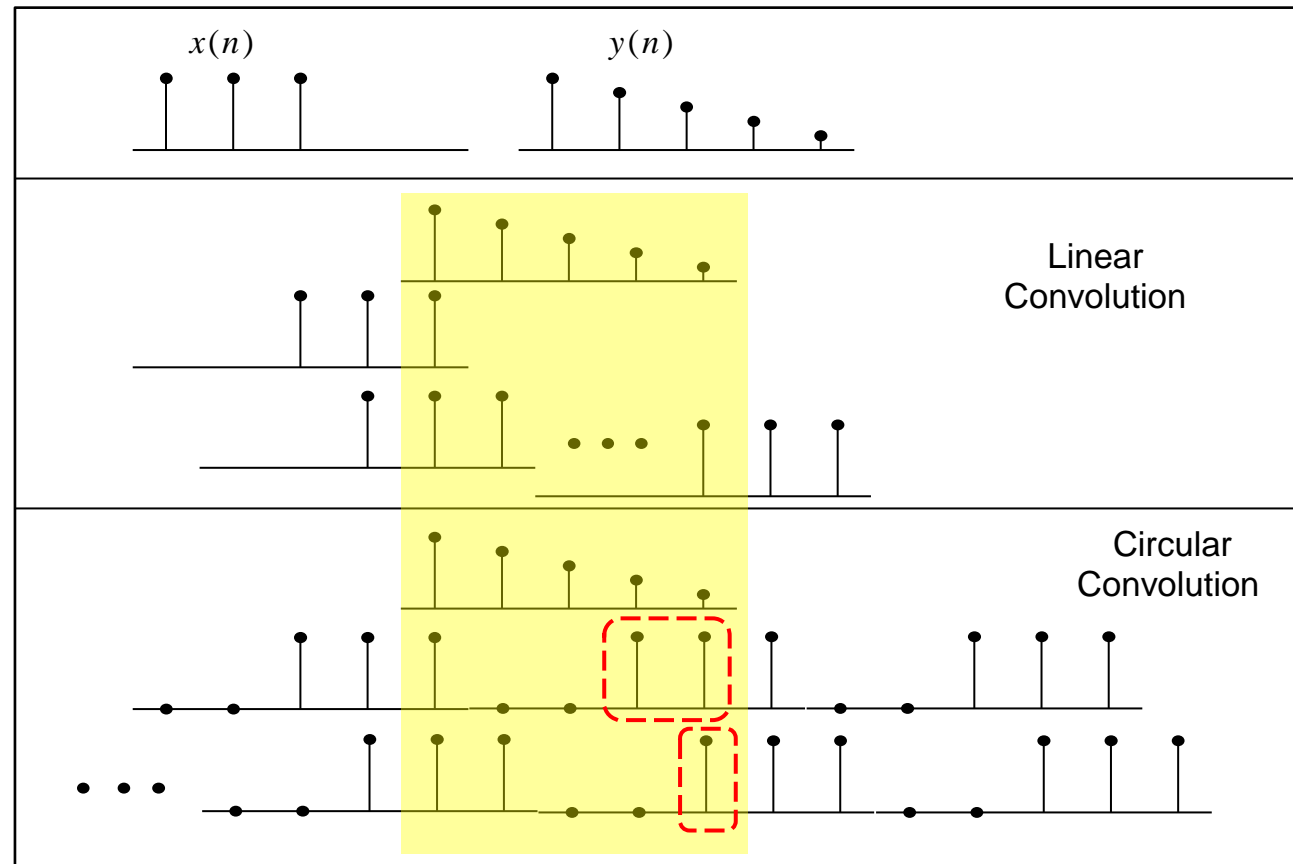
## Definition

- ❖  $x(n)$  and  $y(n)$ : Length  $N$  ( $n=0, 1, 2, \dots, N-1$ )
- ❖ Circular (or periodic) convolution

$$z(n) = x(n) \circledast y(n)$$

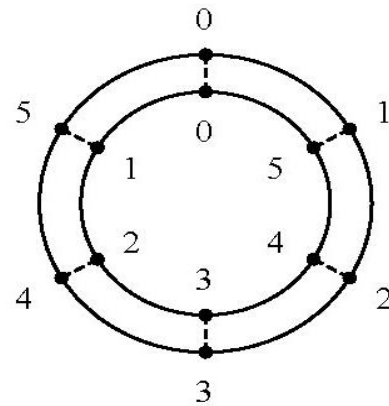
$$= \sum_{m=0}^{N-1} x[m]y[(n-m) \bmod N]$$

Generally,  
linear and circular  
correlations give  
different results

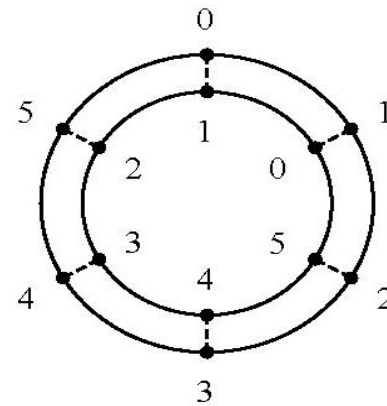


# Circular Convolution

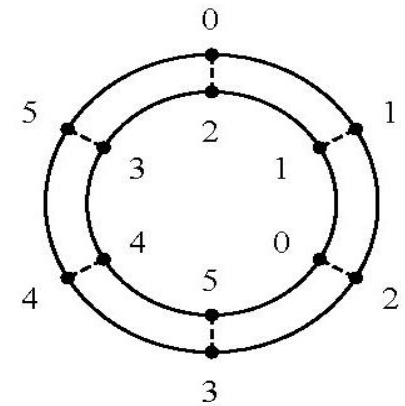
## Graphical representation



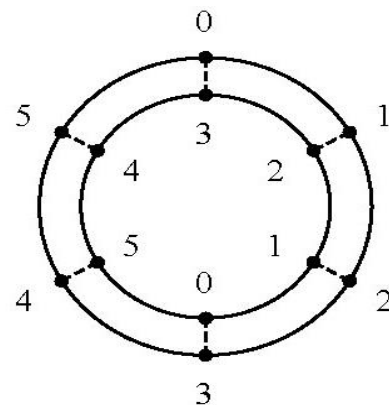
$$n = 0, \quad z[n] = 3 \ 5$$



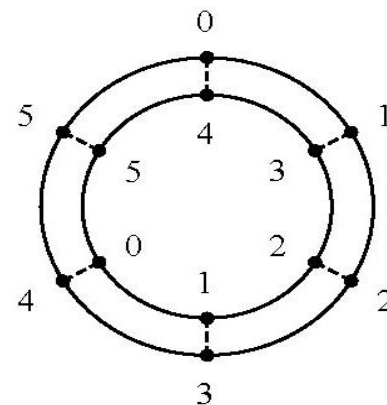
$$n = 1, \quad z[n] = 4 \ 4$$



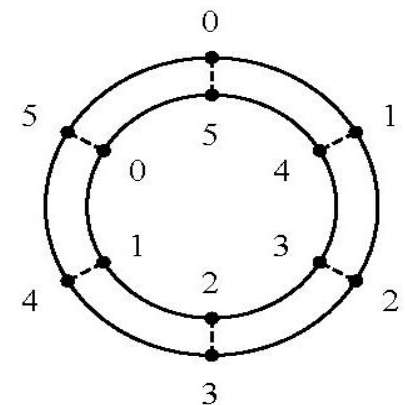
$$n = 2, \quad z[n] = 4 \ 7$$



$$n = 3, \quad z[n] = 4 \ 4$$



$$n = 4, \quad z[n] = 3 \ 5$$



$$n = 5, \quad z[n] = 2 \ 0$$

$$x(n) \circledast y(n) = y(n) \circledast x(n), \quad (x(n) \circledast y(n)) \circledast z(n) = x(n) \circledast (y(n) \circledast z(n)).$$

# Linear Convolution via Circular Convolution

## Case 1: Two finite sequences

❖ **Linear convolution**  $x_3(n) = x_1(n) * x_2(n)$  **where**

$$\begin{cases} x_1(n): \text{Length } N_1 \ (n=0, 1, \dots, N_1-1) \\ x_2(n): \text{Length } N_2 \ (n=0, 1, \dots, N_2-1) \end{cases}$$

$$\Rightarrow \text{Length of } x_3(n) = N_1 + N_2 - 1$$

❖ **Circular convolution**  $\tilde{x}_3(n) = \tilde{x}_1(n) * \tilde{x}_2(n)$  **where**

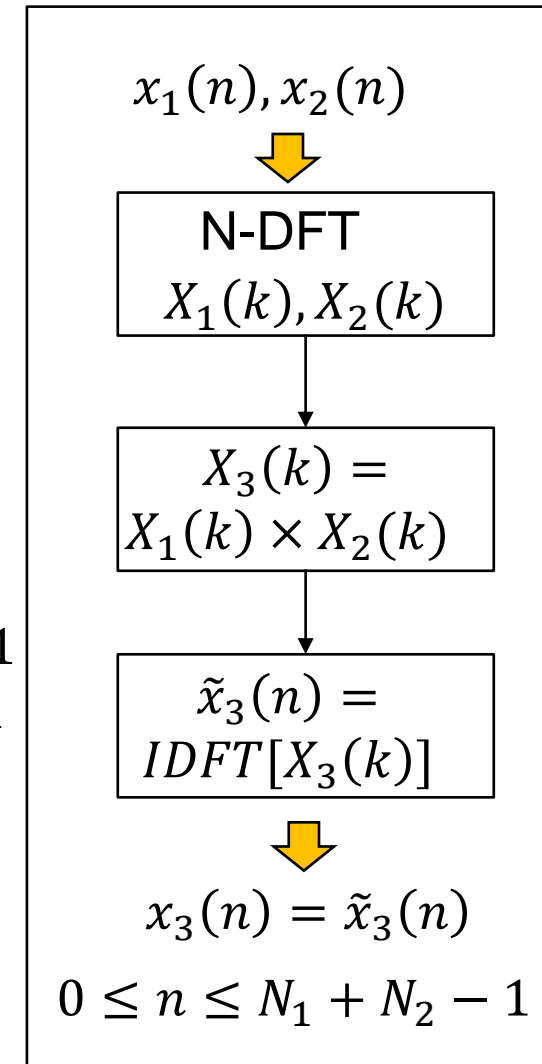
$$\tilde{x}_3(n) = x_3(n) \text{ for } 0 \leq n \leq N-1 \ (N \geq N_1 + N_2 - 1)$$

$$\tilde{x}_1(n) = \begin{cases} x_1(n), & 0 \leq n \leq N_1 - 1 \\ 0, & N_1 \leq n \leq N - 1 \end{cases} \quad \tilde{x}_2(n) = \begin{cases} x_2(n), & 0 \leq n \leq N_2 - 1 \\ 0, & N_2 \leq n \leq N - 1 \end{cases}$$

$$x_3(n) = \tilde{x}_3(n) \quad (0 \leq n \leq N-1)$$

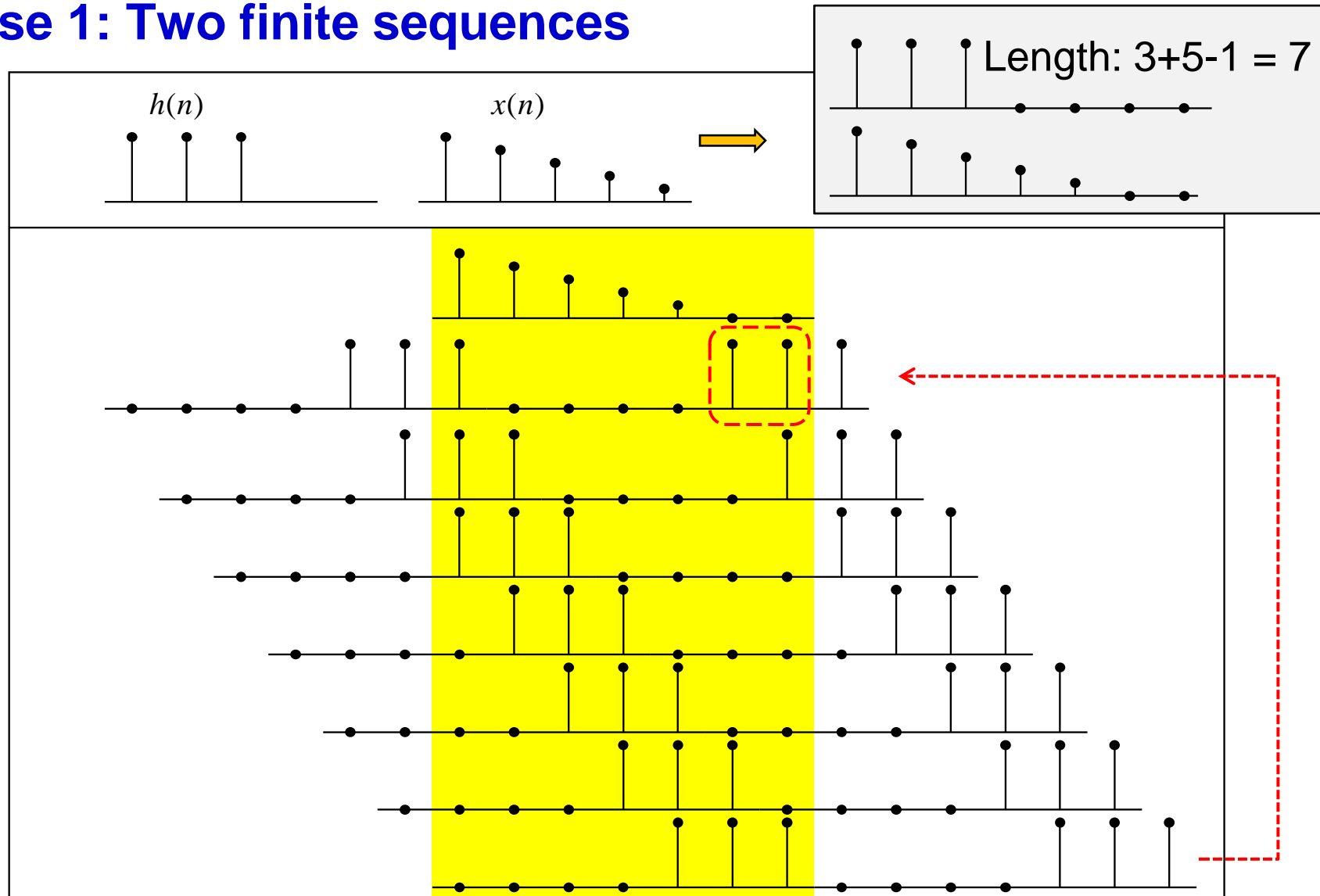
$$= \text{IDFT}\{X_1(k) \cdot X_2(k)\}$$

L. conv  
=  
C. conv



# Linear Convolution via Circular Convolution

## Case 1: Two finite sequences



# Linear Convolution via Circular Convolution

## Example

$$a = [1 \ 1 \ 1 \ 1]; \ a(n) = \{1, 1, 1, 1\}$$

$$b = [-0.5 \ 1 \ -0.5]; \ b(n) = \{-0.5, 1, -0.5\}$$

$$c_{ab} = \text{conv}(a, b) = a * b = \{-0.5 \ 0.5 \ 0 \ 0 \ 0.5 \ -0.5\}$$

$$cc_{ab4} = \text{cconv}(a, b, 4) = a \circledast b = \{0 \ 0 \ 0 \ 0\}$$

$$cc_{ab5} = \text{cconv}(a, b, 5) = a \circledast b = \{-1 \ 0.5 \ 0 \ 0 \ 0.5\}$$

$$cc_{ab6} = \text{cconv}(a, b, 6) = a \circledast b = \{-0.5 \ 0.5 \ 0 \ 0 \ 0.5 \ -0.5\}$$

$$cc_{ab7} = \text{cconv}(a, b, 7) = a \circledast b = \{-0.5 \ 0.5 \ 0 \ 0 \ 0.5 \ -0.5 \ 0\}$$

$$a_{dft4} = \text{fft}(a, 4) = \{4 \ 0 \ 0 \ 0\}$$

$$b_{dft4} = \text{fft}(b, 4) = \{0 \ -j \ -2 \ +j\}$$

$$c_{ab4} = \text{ifft}(a_{dft4} .* b_{dft4}, 4) = \{0, \quad 0, \quad 0, \quad 0\}$$

$$c_{ab6} = \text{ifft}(\text{fft}(a, 6) .* \text{fft}(b, 6), 6) = \{-0.5 \ 0.5 \ 0 \ 0 \ 0.5 \ -0.5\}$$

# Linear Convolution via Circular Convolution

## Case 2: One finite sequence and one infinite sequence

### ❖ Overlap-add method

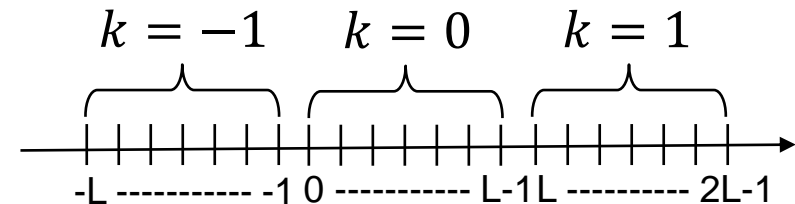
$$y(n) = h(n) * x(n) \quad \text{Len}[h(n)] = M, \quad \text{Len}[x(n)] = \infty \text{ or very long}$$

$\Rightarrow$  Block filtering technique

- Decompose  $x(n)$  into a sum of sections

$$x_k(n) = \begin{cases} x(n), & kL \leq n \leq (k+1)L - 1 \\ 0, & \text{otherwise} \end{cases}$$

$$\Rightarrow x(n) = \sum_{k=-\infty}^{\infty} x_k(n)$$

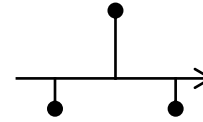
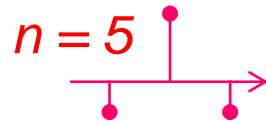
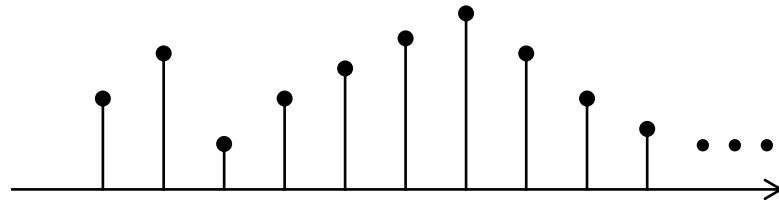


$$y(n) = \sum_{k=-\infty}^{\infty} y_k(n) = \sum_{k=-\infty}^{\infty} x_k(n) * h(n)$$

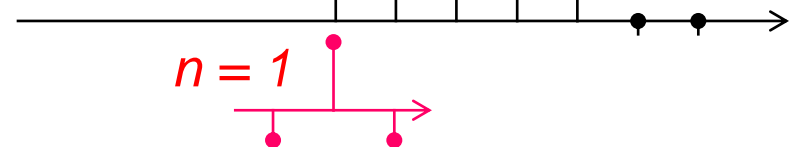
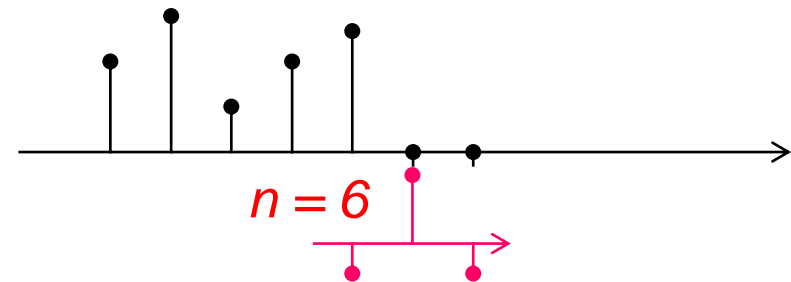
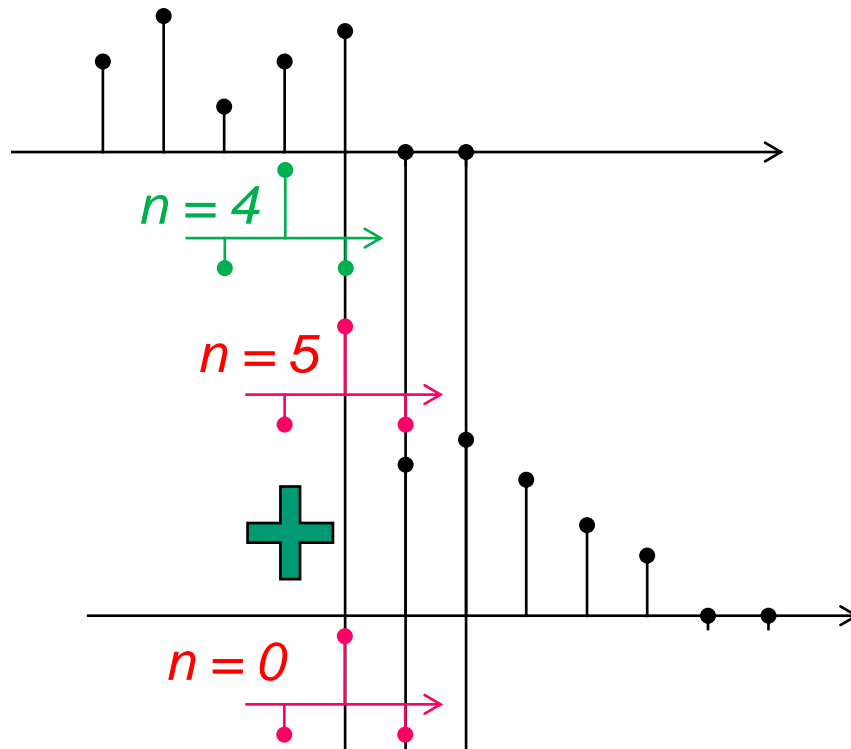
$y_k(n)$  can be obtained using DFT ( $N = L + M - 1$ )



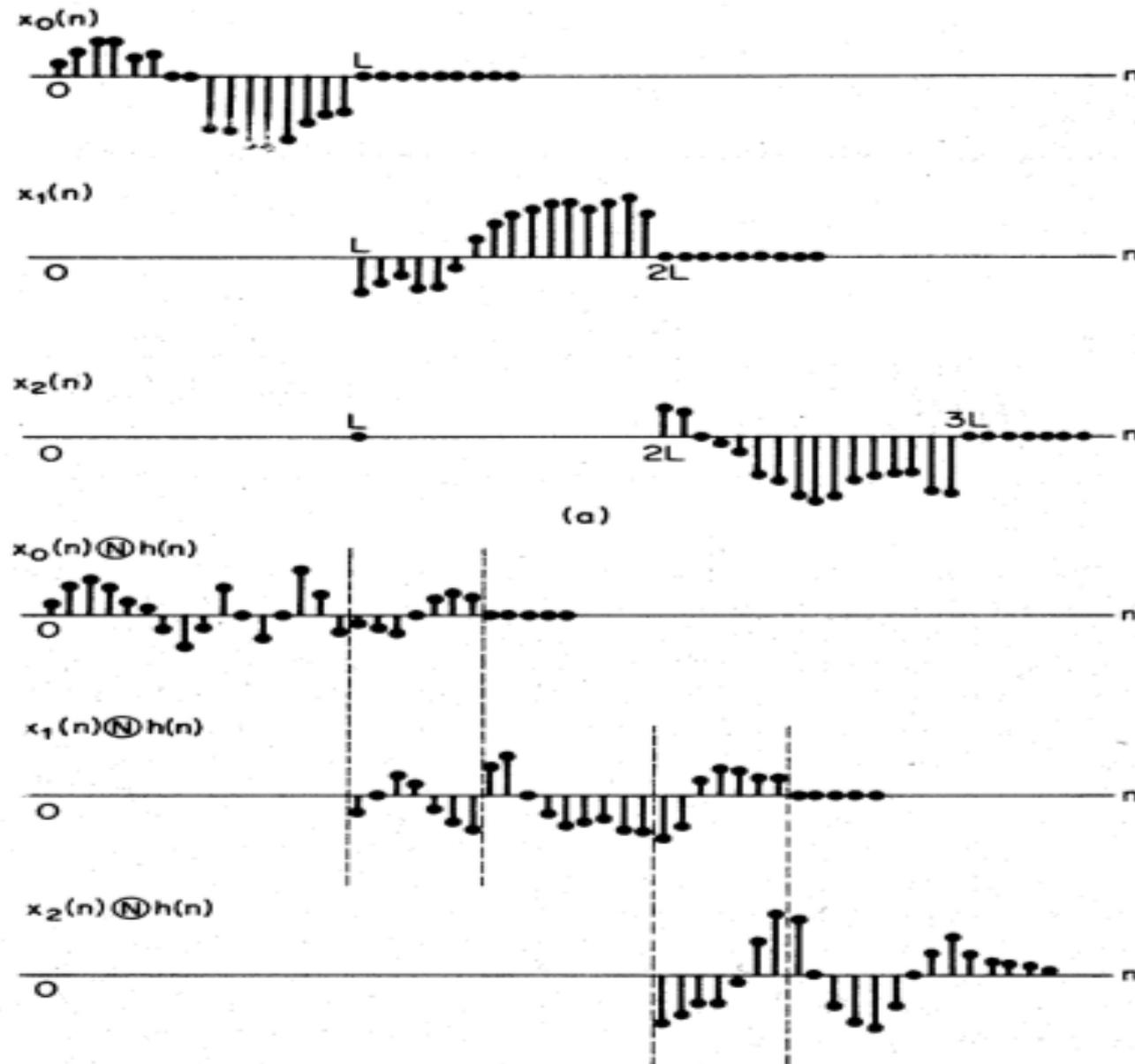
# Linear Convolution via Circular Convolution



$$\begin{aligned} L &= 5 \\ M &= 3 \\ N &= L + M - 1 = 7 \end{aligned}$$



# Linear Convolution via Circular Convolution



# Linear Convolution via Circular Convolution

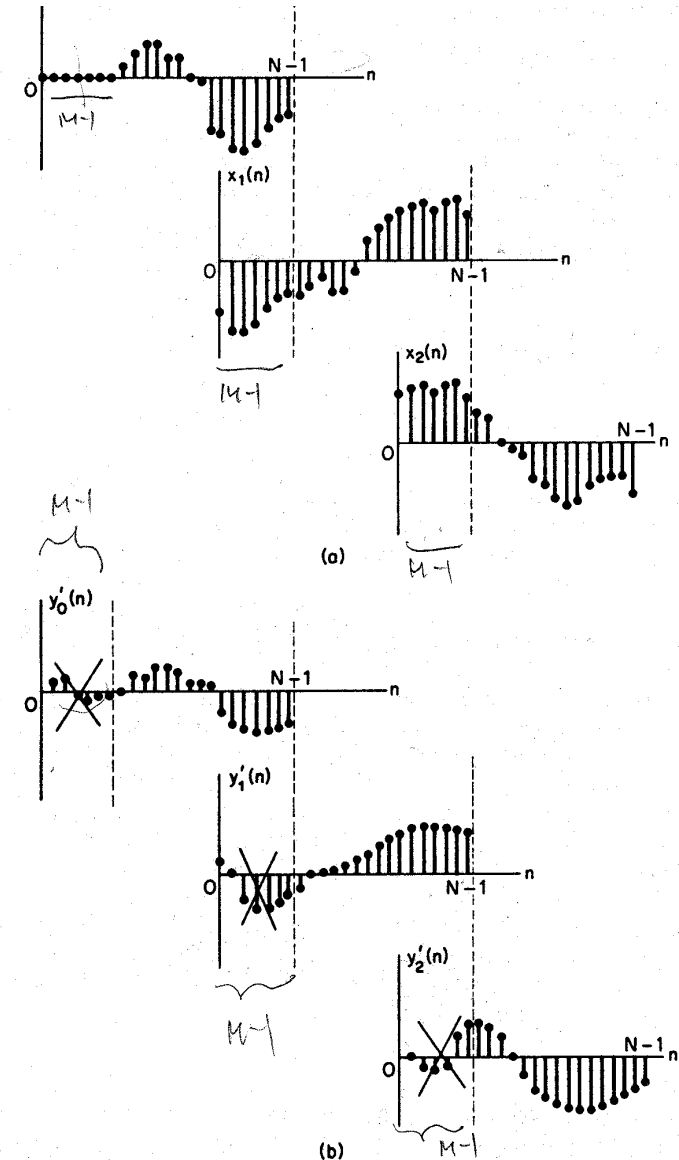
## Case 2: One finite sequence and one infinite sequence

### ❖ Overlap-save method

$$x_k(n) = x(n + k(N - M + 1)) \quad 0 \leq n \leq N - 1$$

$$y(n) = \sum_{k=0}^{\infty} y_k(n - k(N - M + 1))$$

$$\text{where } y_k(n) = \begin{cases} y'_k(n), & M-1 \leq n \leq N-1 \\ 0, & \text{otherwise} \end{cases}$$



# Discrete Cosine Transform (DCT)

## Backgrounds

### ❖ Discrete-time version of the Fourier Cosine Series

- Real valued transform for real signals
- Linear transform:  $\underline{X}_N = \underline{C}_N \cdot \underline{x}_N$
- Extend the signal symmetrically around the origin in **4 different ways**.

## Type-I DCT

❖ **N samples**  $\Rightarrow$  **2N-2 samples**

$$x_1(n) = \begin{cases} \sqrt{2}x(n), & n=0, N-1 \\ x(n) & 1 \leq n \leq N-2 \\ x(2N-n-2) & N \leq n \leq 2N-3 \end{cases}$$

